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bo cuadrado $25 \times 25 \mathrm{~mm}$, espesor 1.5 mm
ametros no especificados: 10 mm


bo cuadrado $25 \times 25 \mathrm{~mm}$, espesor 1.5 mm
ametros no especificados: 10 mm







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| 13 | 1 | Union Asiento-Apoyapie-Actuador2 |  | Acero A36 |  |  |  |  |  |
| :---: | :---: | :--- | :--- | :--- | :--- | :---: | :---: | :---: | :---: |
| 12 | 1 | Actuador LA23-150 |  |  |  |  |  |  |  |
| 11 | 1 | Actuador LA23-100 |  |  |  |  |  |  |  |
| 10 | 1 | Union Asiento-Espaldar- |  | Acero A36 |  |  |  |  |  |
| 9 | 1 | Union Asiento-Espaldar-Actuador1 |  | Acero A36 |  |  |  |  |  |
| 8 | 3 | Soporte Firgelli MB3 |  |  |  |  |  |  |  |
| 7 | 1 | Union Asiento-Apoyapie-Actuador2 |  | Acero A36 |  |  |  |  |  |
| 6 | 1 | Tabla para Glúteos |  | Madera |  |  |  |  |  |
| 5 | 1 | Almohadilla para glúteos |  |  |  |  |  |  |  |
| 4 | 1 | Soporte Asiento-Actuador4 |  | Acero A36 |  |  |  |  |  |
| 3 | 2 | Soporte de eje |  | Acero A36 |  |  |  |  |  |
| 2 | 2 | Union Asiento-Apoyapie |  | Acero A36 |  |  |  |  |  |
| 1 | 1 | Asiento |  | Acero A36 |  |  |  |  |  |
| POS. | CANT. | DESCRIPCIÓN |  |  |  |  | NORMA | MATERIAL | OBSERVACIONES |
| PONTIFICIA UNIVERSIDAD CATÓLICA DEL PERU |  |  |  |  |  |  |  |  |  |




(6) $๑ \infty$










## Features

- High Performance, Low Power Atme ${ }^{\circledR}$ AVR ${ }^{\circledR}$ 8-Bit Microcontroller
- Advanced RISC Architecture
- 135 Powerful Instructions - Most Single Clock Cycle Execution
- $32 \times 8$ General Purpose Working Registers
- Fully Static Operation
- Up to 16 MIPS Throughput at 16 MHz
- On-Chip 2-cycle Multiplier
- High Endurance Non-volatile Memory Segments
- 64K/128K/256KBytes of In-System Self-Programmable Flash
- 4Kbytes EEPROM
- 8Kbytes Internal SRAM
- Write/Erase Cycles:10,000 Flash/100,000 EEPROM
- Data retention: 20 years at $85^{\circ} \mathrm{C} / 100$ years at $25^{\circ} \mathrm{C}$
- Optional Boot Code Section with Independent Lock Bits
- In-System Programming by On-chip Boot Program
- True Read-While-Write Operation
- Programming Lock for Software Security
- Endurance: Up to 64 Kbytes Optional External Memory Space
- Atmel ${ }^{\circledR}$ QTouch ${ }^{\circledR}$ library support
- Capacitive touch buttons, sliders and wheels
- QTouch and QMatrix acquisition
- Up to 64 sense channels
- JTAG (IEEE ${ }^{\circledR}$ std. 1149.1 compliant) Interface
- Boundary-scan Capabilities According to the JTAG Standard
- Extensive On-chip Debug Support
- Programming of Flash, EEPROM, Fuses, and Lock Bits through the JTAG Interface
- Peripheral Features
- Two 8-bit Timer/Counters with Separate Prescaler and Compare Mode
- Four 16-bit Timer/Counter with Separate Prescaler, Compare- and Capture Mode
- Real Time Counter with Separate Oscillator
- Four 8-bit PWM Channels
- Six/Twelve PWM Channels with Programmable Resolution from 2 to 16 Bits (ATmega1281/2561, ATmega640/1280/2560)
- Output Compare Modulator
- 8/16-channel, 10-bit ADC (ATmega1281/2561, ATmega640/1280/2560)
- Two/Four Programmable Serial USART (ATmega1281/2561, ATmega640/1280/2560)
- Master/Slave SPI Serial Interface
- Byte Oriented 2-wire Serial Interface
- Programmable Watchdog Timer with Separate On-chip Oscillator
- On-chip Analog Comparator
- Interrupt and Wake-up on Pin Change
- Special Microcontroller Features
- Power-on Reset and Programmable Brown-out Detection
- Internal Calibrated Oscillator
- External and Internal Interrupt Sources
- Six Sleep Modes: Idle, ADC Noise Reduction, Power-save, Power-down, Standby, and Extended Standby
- I/O and Packages
- 54/86 Programmable I/O Lines (ATmega1281/2561, ATmega640/1280/2560)
- 64-pad QFN/MLF, 64-lead TQFP (ATmega1281/2561)
- 100-lead TQFP, 100-ball CBGA (ATmega640/1280/2560)
- RoHS/Fully Green
- Temperature Range:
- $-40^{\circ} \mathrm{C}$ to $85^{\circ} \mathrm{C}$ Industrial
- Ultra-Low Power Consumption
- Active Mode: 1MHz, 1.8V: $500 \mu \mathrm{~A}$
- Power-down Mode: $0.1 \mu \mathrm{~A}$ at 1.8 V
- Speed Grade:
- ATmega640V/ATmega1280V/ATmega1281V:
$\cdot 0-4 \mathrm{MHz}$ @ $1.8 \mathrm{~V}-5.5 \mathrm{~V}, 0-8 \mathrm{MHz} @ 2.7 \mathrm{~V}-5.5 \mathrm{~V}$
- ATmega2560V/ATmega2561V:
$\cdot 0-2 \mathrm{MHz} @ 1.8 \mathrm{~V}-5.5 \mathrm{~V}, 0-8 \mathrm{MHz} @ 2.7 \mathrm{~V}-5.5 \mathrm{~V}$
- ATmega640/ATmega1280/ATmega1281:
$\bullet 0-8 \mathrm{MHz} @ 2.7 \mathrm{~V}-5.5 \mathrm{~V}, 0-16 \mathrm{MHz} @ 4.5 \mathrm{~V}-5.5 \mathrm{~V}$
ATmega2560/ATmega2561:
$\cdot 0-16 \mathrm{MHz}$ @ $4.5 \mathrm{~V}-5.5 \mathrm{~V}$


## 1. Pin Configurations

Figure 1-1. TQFP-pinout ATmega640/1280/2560


Figure 1-2. CBGA-pinout ATmega640/1280/2560


Table 1-1. CBGA-pinout ATmega640/1280/2560

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ | $\mathbf{8}$ | $\mathbf{9}$ | $\mathbf{1 0}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{A}$ | GND | AREF | PF0 | PF2 | PF5 | PK0 | PK3 | PK6 | GND | VCC |
| $\mathbf{B}$ | AVCC | PG5 | PF1 | PF3 | PF6 | PK1 | PK4 | PK7 | PA0 | PA2 |
| $\mathbf{C}$ | PE2 | PE0 | PE1 | PF4 | PF7 | PK2 | PK5 | PJ7 | PA1 | PA3 |
| $\mathbf{D ~}$ | PE3 | PE4 | PE5 | PE6 | PH2 | PA4 | PA5 | PA6 | PA7 | PG2 |
| $\mathbf{E}$ | PE7 | PH0 | PH1 | PH3 | PH5 | PJ6 | PJ5 | PJ4 | PJ3 | PJ2 |
| $\mathbf{F}$ | VCC | PH4 | PH6 | PB0 | PL4 | PD1 | PJ1 | PJ0 | PC7 | GND |
| $\mathbf{G}$ | GND | PB1 | PB2 | PB5 | PL2 | PD0 | PD5 | PC5 | PC6 | VCC |
| $\mathbf{H}$ | PB3 | PB4 | RESET | PL1 | PL3 | PL7 | PD4 | PC4 | PC3 | PC2 |
| $\mathbf{J ~}$ | PH7 | PG3 | PB6 | PL0 | XTAL2 | PL6 | PD3 | PC1 | PC0 | PG1 |
| $\mathbf{K}$ | PB7 | PG4 | VCC | GND | XTAL1 | PL5 | PD2 | PD6 | PD7 | PG0 |

Note: $\quad$ The functions for each pin is the same as for the 100 pin packages shown in Figure 1-1 on page 2.

Figure 1-3. Pinout ATmega1281/2561


Note: The large center pad underneath the QFN/MLF package is made of metal and internally connected to GND. It should be soldered or glued to the board to ensure good mechanical stability. If the center pad is left unconnected, the package might loosen from the board.

## 2. Overview

The ATmega640/1280/1281/2560/2561 is a low-power CMOS 8-bit microcontroller based on the AVR enhanced RISC architecture. By executing powerful instructions in a single clock cycle, the ATmega640/1280/1281/2560/2561 achieves throughputs approaching 1 MIPS per MHz allowing the system designer to optimize power consumption versus processing speed.

### 2.1 Block Diagram

Figure 2-1. Block Diagram


The Atmel ${ }^{\circledR}$ AVR $^{\circledR}$ core combines a rich instruction set with 32 general purpose working registers. All the 32 registers are directly connected to the Arithmetic Logic Unit (ALU), allowing two independent registers to be accessed in one single instruction executed in one clock cycle. The resulting architecture is more code efficient while achieving throughputs up to ten times faster than conventional CISC microcontrollers.

The ATmega640/1280/1281/2560/2561 provides the following features: $64 \mathrm{~K} / 128 \mathrm{~K} / 256 \mathrm{~K}$ bytes of In-System Programmable Flash with Read-While-Write capabilities, 4Kbytes EEPROM, 8Kbytes SRAM, 54/86 general purpose I/O lines, 32 general purpose working registers, Real Time Counter (RTC), six flexible Timer/Counters with compare modes and PWM, four USARTs, a byte oriented 2-wire Serial Interface, a 16-channel, 10-bit ADC with optional differential input stage with programmable gain, programmable Watchdog Timer with Internal Oscillator, an SPI serial port, IEEE ${ }^{\circledR}$ std. 1149.1 compliant JTAG test interface, also used for accessing the On-chip Debug system and programming and six software selectable power saving modes. The Idle mode stops the CPU while allowing the SRAM, Timer/Counters, SPI port, and interrupt system to continue functioning. The Power-down mode saves the register contents but freezes the Oscillator, disabling all other chip functions until the next interrupt or Hardware Reset. In Power-save mode, the asynchronous timer continues to run, allowing the user to maintain a timer base while the rest of the device is sleeping. The ADC Noise Reduction mode stops the CPU and all I/O modules except Asynchronous Timer and ADC, to minimize switching noise during ADC conversions. In Standby mode, the Crystal/Resonator Oscillator is running while the rest of the device is sleeping. This allows very fast start-up combined with low power consumption. In Extended Standby mode, both the main Oscillator and the Asynchronous Timer continue to run.

Atmel offers the QTouch ${ }^{\circledR}$ library for embedding capacitive touch buttons, sliders and wheels functionality into AVR microcontrollers. The patented charge-transfer signal acquisition offersrobust sensing and includes fully debounced reporting of touch keys and includes Adjacent Key Suppression ${ }^{\circledR}$ (AKS ${ }^{\circledR}$ ) technology for unambiguous detection of key events. The easy-to-use QTouch Suite toolchain allows you to explore, develop and debug your own touch applications.

The device is manufactured using the Atmel high-density nonvolatile memory technology. The On-chip ISP Flash allows the program memory to be reprogrammed in-system through an SPI serial interface, by a conventional nonvolatile memory programmer, or by an On-chip Boot program running on the AVR core. The boot program can use any interface to download the application program in the application Flash memory. Software in the Boot Flash section will continue to run while the Application Flash section is updated, providing true Read-While-Write operation. By combining an 8 -bit RISC CPU with In-System Self-Programmable Flash on a monolithic chip, the Atmel ATmega640/1280/1281/2560/2561 is a powerful microcontroller that provides a highly flexible and cost effective solution to many embedded control applications.

The ATmega640/1280/1281/2560/2561 AVR is supported with a full suite of program and system development tools including: C compilers, macro assemblers, program debugger/simulators, in-circuit emulators, and evaluation kits.

### 2.2 Comparison Between ATmega1281/2561 and ATmega640/1280/2560

Each device in the ATmega640/1280/1281/2560/2561 family differs only in memory size and number of pins. Table 2-1 summarizes the different configurations for the six devices.

Table 2-1. Configuration Summary

| Device | Flash | EEPROM | RAM | General <br> Purpose I/O pins | 16 bits resolution <br> PWM channels | Serial <br> USARTs | ADC <br> Channels |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ATmega640 | 64 KB | 4 KB | 8 KB | 86 | 12 | 4 | 16 |
| ATmega1280 | 128 KB | 4 KB | 8 KB | 86 | 12 | 4 | 16 |
| ATmega1281 | 128 KB | 4 KB | 8 KB | 54 | 6 | 2 | 8 |
| ATmega2560 | 256 KB | 4 KB | 8 KB | 86 | 12 | 4 | 16 |
| ATmega2561 | 256 KB | 4 KB | 8 KB | 54 | 6 | 2 | 8 |

### 2.3 Pin Descriptions

### 2.3.1 VCC

Digital supply voltage.

### 2.3.2 GND

Ground.

### 2.3.3 Port A (PA7..PA0)

Port A is an 8-bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port A output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port A pins that are externally pulled low will source current if the pull-up resistors are activated. The Port A pins are tri-stated when a reset condition becomes active, even if the clock is not running.

Port A also serves the functions of various special features of the ATmega640/1280/1281/2560/2561 as listed on page 75.

### 2.3.4 Port B (PB7..PB0)

Port B is an 8-bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port B output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port B pins that are externally pulled low will source current if the pull-up resistors are activated. The Port B pins are tri-stated when a reset condition becomes active, even if the clock is not running.

Port B has better driving capabilities than the other ports.
Port B also serves the functions of various special features of the ATmega640/1280/1281/2560/2561 as listed on page 76.

### 2.3.5 Port C (PC7..PC0)

Port C is an 8 -bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port C output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port C pins that are externally pulled low will source current if the pull-up resistors are activated. The Port $C$ pins are tri-stated when a reset condition becomes active, even if the clock is not running.

Port C also serves the functions of special features of the ATmega640/1280/1281/2560/2561 as listed on page 79.

### 2.3.6 Port D (PD7..PDO)

Port D is an 8-bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port D output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port D pins that are externally pulled low will source current if the pull-up resistors are activated. The Port D pins are tri-stated when a reset condition becomes active, even if the clock is not running.
Port D also serves the functions of various special features of the ATmega640/1280/1281/2560/2561 as listed on page 80.

### 2.3.7 Port E (PE7..PEO)

Port E is an 8-bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port E output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port E pins that are externally pulled low will source current if the pull-up resistors are activated. The Port E pins are tri-stated when a reset condition becomes active, even if the clock is not running.
Port E also serves the functions of various special features of the ATmega640/1280/1281/2560/2561 as listed on page 82.

### 2.3.8 Port F (PF7..PF0)

Port $F$ serves as analog inputs to the $A / D$ Converter.
Port F also serves as an 8-bit bi-directional I/O port, if the A/D Converter is not used. Port pins can provide internal pull-up resistors (selected for each bit). The Port F output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port F pins that are externally pulled low will source current if the pull-up resistors are activated. The Port F pins are tri-stated when a reset condition becomes active, even if the clock is not running. If the JTAG interface is enabled, the pull-up resistors on pins PF7(TDI), PF5(TMS), and PF4(TCK) will be activated even if a reset occurs.

Port F also serves the functions of the JTAG interface.

### 2.3.9 Port G (PG5..PG0)

Port G is a 6-bit I/O port with internal pull-up resistors (selected for each bit). The Port G output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port $G$ pins that are externally pulled low will source current if the pull-up resistors are activated. The Port $G$ pins are tri-stated when a reset condition becomes active, even if the clock is not running.

Port $G$ also serves the functions of various special features of the ATmega640/1280/1281/2560/2561 as listed on page 86.

### 2.3.10 Port H (PH7..PH0)

Port H is a 8-bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port H output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port H pins that are externally pulled low will source current if the pull-up resistors are activated. The Port H pins are tri-stated when a reset condition becomes active, even if the clock is not running.
Port H also serves the functions of various special features of the ATmega640/1280/2560 as listed on page 88.

### 2.3.11 Port J (PJ7..PJO)

Port J is a 8 -bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port J output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port J pins that are externally pulled low will source current if the pull-up resistors are activated. The Port J pins are tri-stated when a reset condition becomes active, even if the clock is not running. Port J also serves the functions of various special features of the ATmega640/1280/2560 as listed on page 90.

### 2.3.12 Port K (PK7..PK0)

Port $K$ serves as analog inputs to the A/D Converter.
Port K is a 8-bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port K output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port K pins that are externally pulled low will source current if the pull-up resistors are activated. The Port K pins are tri-stated when a reset condition becomes active, even if the clock is not running.

Port K also serves the functions of various special features of the ATmega640/1280/2560 as listed on page 92.

### 2.3.13 Port L (PL7..PLO)

Port L is a 8-bit bi-directional I/O port with internal pull-up resistors (selected for each bit). The Port L output buffers have symmetrical drive characteristics with both high sink and source capability. As inputs, Port L pins that are externally pulled low will source current if the pull-up resistors are activated. The Port L pins are tri-stated when a reset condition becomes active, even if the clock is not running.
Port L also serves the functions of various special features of the ATmega640/1280/2560 as listed on page 94.

### 2.3.14 $\overline{\text { RESET }}$

Reset input. A low level on this pin for longer than the minimum pulse length will generate a reset, even if the clock is not running. The minimum pulse length is given in "System and Reset Characteristics" on page 360. Shorter pulses are not guaranteed to generate a reset.

### 2.3.15 XTAL1

Input to the inverting Oscillator amplifier and input to the internal clock operating circuit.

### 2.3.16 XTAL2

Output from the inverting Oscillator amplifier.

### 2.3.17 AVCC

AVCC is the supply voltage pin for Port $F$ and the $A / D$ Converter. It should be externally connected to $\mathrm{V}_{\mathrm{Cc}}$, even if the ADC is not used. If the ADC is used, it should be connected to $\mathrm{V}_{\mathrm{CC}}$ through a low-pass filter.

### 2.3.18 AREF

This is the analog reference pin for the A/D Converter.

## 3. Resources

A comprehensive set of development tools and application notes, and datasheets are available for download on http://www.atmel.com/avr.

## 4. About Code Examples

This documentation contains simple code examples that briefly show how to use various parts of the device. Be aware that not all C compiler vendors include bit definitions in the header files and interrupt handling in C is compiler dependent. Confirm with the C compiler documentation for more details.

These code examples assume that the part specific header file is included before compilation. For I/O registers located in extended I/O map, "IN", "OUT", "SBIS", "SBIC", "CBI", and "SBI" instructions must be replaced with instructions that allow access to extended I/O. Typically "LDS" and "STS" combined with "SBRS", "SBRC", "SBR", and "CBR".

## 5. Data Retention

Reliability Qualification results show that the projected data retention failure rate is much less than 1 ppm over 20 years at $85^{\circ} \mathrm{C}$ or 100 years at $25^{\circ} \mathrm{C}$.

## 6. Capacitive touch sensing

The Atmel ${ }^{\circledR}$ QTouch ${ }^{\circledR}$ Library provides a simple to use solution to realize touch sensitive interfaces on most Atmel AVR ${ }^{\circledR}$ microcontrollers. The QTouch Library includes support for the QTouch and QMatrix acquisition methods.
Touch sensing can be added to any application by linking the appropriate Atmel QTouch Library for the AVR Microcontroller. This is done by using a simple set of APIs to define the touch channels and sensors, and then calling the touch sensing APl's to retrieve the channel information and determine the touch sensor states.

The QTouch Library is FREE and downloadable from the Atmel website at the following location: www.atmel.com/qtouchlibrary. For implementation details and other information, refer to the Atmel QTouch Library User Guide - also available for download from the Atmel website.

## 7. AVR CPU Core

### 7.1 Introduction

This section discusses the AVR core architecture in general. The main function of the CPU core is to ensure correct program execution. The CPU must therefore be able to access memories, perform calculations, control peripherals, and handle interrupts.

### 7.2 Architectural Overview

Figure 7-1. Block Diagram of the AVR Architecture


In order to maximize performance and parallelism, the AVR uses a Harvard architecture - with separate memories and buses for program and data. Instructions in the program memory are executed with a single level pipelining. While one instruction is being executed, the next instruction is pre-fetched from the program memory. This concept enables instructions to be executed in every clock cycle. The program memory is In-System Reprogrammable Flash memory.

The fast-access Register File contains $32 \times 8$-bit general purpose working registers with a single clock cycle access time. This allows single-cycle Arithmetic Logic Unit (ALU) operation. In a typical ALU operation, two oper-
ands are output from the Register File, the operation is executed, and the result is stored back in the Register File - in one clock cycle.

Six of the 32 registers can be used as three 16-bit indirect address register pointers for Data Space addressing enabling efficient address calculations. One of the these address pointers can also be used as an address pointer for look up tables in Flash program memory. These added function registers are the 16 -bit X -, Y -, and Z -register, described later in this section.

The ALU supports arithmetic and logic operations between registers or between a constant and a register. Single register operations can also be executed in the ALU. After an arithmetic operation, the Status Register is updated to reflect information about the result of the operation.

Program flow is provided by conditional and unconditional jump and call instructions, able to directly address the whole address space. Most AVR instructions have a single 16 -bit word format. Every program memory address contains a 16-bit or 32-bit instruction.

Program Flash memory space is divided in two sections, the Boot Program section and the Application Program section. Both sections have dedicated Lock bits for write and read/write protection. The SPM instruction that writes into the Application Flash memory section must reside in the Boot Program section.

During interrupts and subroutine calls, the return address Program Counter (PC) is stored on the Stack. The Stack is effectively allocated in the general data SRAM, and consequently the Stack size is only limited by the total SRAM size and the usage of the SRAM. All user programs must initialize the SP in the Reset routine (before subroutines or interrupts are executed). The Stack Pointer (SP) is read/write accessible in the I/O space. The data SRAM can easily be accessed through the five different addressing modes supported in the AVR architecture.

The memory spaces in the AVR architecture are all linear and regular memory maps.
A flexible interrupt module has its control registers in the I/O space with an additional Global Interrupt Enable bit in the Status Register. All interrupts have a separate Interrupt Vector in the Interrupt Vector table. The interrupts have priority in accordance with their Interrupt Vector position. The lower the Interrupt Vector address, the higher the priority.
The I/O memory space contains 64 addresses for CPU peripheral functions as Control Registers, SPI, and other I/O functions. The I/O Memory can be accessed directly, or as the Data Space locations following those of the Register File, 0x20-0x5F. In addition, the ATmega640/1280/1281/2560/2561 has Extended I/O space from 0x60$0 x 1 F F$ in SRAM where only the ST/STS/STD and LD/LDS/LDD instructions can be used.

### 7.3 ALU - Arithmetic Logic Unit

The high-performance AVR ALU operates in direct connection with all the 32 general purpose working registers. Within a single clock cycle, arithmetic operations between general purpose registers or between a register and an immediate are executed. The ALU operations are divided into three main categories - arithmetic, logical, and bitfunctions. Some implementations of the architecture also provide a powerful multiplier supporting both signed/unsigned multiplication and fractional format. See the "Instruction Set Summary" on page 404 for a detailed description.

### 7.4 Status Register

The Status Register contains information about the result of the most recently executed arithmetic instruction. This information can be used for altering program flow in order to perform conditional operations. Note that the Status Register is updated after all ALU operations, as specified in the "Instruction Set Summary" on page 404. This will in many cases remove the need for using the dedicated compare instructions, resulting in faster and more compact code.

The Status Register is not automatically stored when entering an interrupt routine and restored when returning from an interrupt. This must be handled by software.

### 7.4.1 SREG - AVR Status Register

The AVR Status Register - SREG - is defined as:

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x3F (0x5F) | I | T | H | S | V | N | Z | C |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## - Bit 7 - I: Global Interrupt Enable

The Global Interrupt Enable bit must be set for the interrupts to be enabled. The individual interrupt enable control is then performed in separate control registers. If the Global Interrupt Enable Register is cleared, none of the interrupts are enabled independent of the individual interrupt enable settings. The I-bit is cleared by hardware after an interrupt has occurred, and is set by the RETI instruction to enable subsequent interrupts. The I-bit can also be set and cleared by the application with the SEI and CLI instructions, as described in the "Instruction Set Summary" on page 404.

## - Bit 6 - T: Bit Copy Storage

The Bit Copy instructions BLD (Bit LoaD) and BST (Bit STore) use the T-bit as source or destination for the operated bit. A bit from a register in the Register File can be copied into T by the BST instruction, and a bit in T can be copied into a bit in a register in the Register File by the BLD instruction.

## - Bit 5 - H: Half Carry Flag

The Half Carry Flag H indicates a Half Carry in some arithmetic operations. Half Carry Is useful in BCD arithmetic. See the "Instruction Set Summary" on page 404 for detailed information.

- Bit 4 - S: Sign Bit, $\mathrm{S}=\mathrm{N} \oplus \mathrm{V}$

The S-bit is always an exclusive or between the Negative Flag N and the Two's Complement Overflow Flag V. See the "Instruction Set Summary" on page 404 for detailed information.

## - Bit 3 - V: Two's Complement Overflow Flag

The Two's Complement Overflow Flag V supports two's complement arithmetics. See the "Instruction Set Summary" on page 404 for detailed information.

## - Bit 2 - N: Negative Flag

The Negative Flag N indicates a negative result in an arithmetic or logic operation. See the "Instruction Set Summary" on page 404 for detailed information.

## - Bit 1 - Z: Zero Flag

The Zero Flag Z indicates a zero result in an arithmetic or logic operation. See the "Instruction Set Summary" on page 404 for detailed information.

## - Bit 0 - C: Carry Flag

The Carry Flag C indicates a carry in an arithmetic or logic operation. See the "Instruction Set Summary" on page 404 for detailed information.

### 7.5 General Purpose Register File

The Register File is optimized for the AVR Enhanced RISC instruction set. In order to achieve the required performance and flexibility, the following input/output schemes are supported by the Register File:

- One 8 -bit output operand and one 8 -bit result input
- Two 8-bit output operands and one 8-bit result input
- Two 8-bit output operands and one 16-bit result input
- One 16 -bit output operand and one 16 -bit result input

Figure 7-2 shows the structure of the 32 general purpose working registers in the CPU.

Figure 7-2. AVR CPU General Purpose Working Registers

|  | 7 | Addr. |  |
| :---: | :---: | :---: | :---: |
|  | R0 | $0 \times 00$ |  |
|  | R1 | $0 \times 01$ |  |
|  | R2 | 0x02 |  |
|  | ... |  |  |
|  | R13 | 0x0D |  |
| General | R14 | 0x0E |  |
| Purpose | R15 | 0x0F |  |
| Working | R16 | 0x10 |  |
| Registers | R17 | $0 \times 11$ |  |
|  | $\ldots$ |  |  |
|  | R26 | $0 \times 1 \mathrm{~A}$ | X-register Low Byte |
|  | R27 | 0x1B | X-register High Byte |
|  | R28 | 0x1C | Y-register Low Byte |
|  | R29 | 0x1D | Y-register High Byte |
|  | R30 | 0x1E | Z-register Low Byte |
|  | R31 | 0x1F | Z-register High Byte |

Most of the instructions operating on the Register File have direct access to all registers, and most of them are single cycle instructions.

As shown in Figure 7-2, each register is also assigned a data memory address, mapping them directly into the first 32 locations of the user Data Space. Although not being physically implemented as SRAM locations, this memory organization provides great flexibility in access of the registers, as the X -, Y - and Z -pointer registers can be set to index any register in the file.

### 7.5.1 The X-register, Y-register, and Z-register

The registers R26..R31 have some added functions to their general purpose usage. These registers are 16-bit address pointers for indirect addressing of the data space. The three indirect address registers $\mathrm{X}, \mathrm{Y}$, and Z are defined as described in Figure 7-3.

Figure 7-3. $\quad$ The $\mathrm{X}-$, Y -, and Z-registers


In the different addressing modes these address registers have functions as fixed displacement, automatic increment, and automatic decrement (see the "Instruction Set Summary" on page 404 for details).

### 7.6 Stack Pointer

The Stack is mainly used for storing temporary data, for storing local variables and for storing return addresses after interrupts and subroutine calls. The Stack Pointer Register always points to the top of the Stack. Note that the Stack is implemented as growing from higher memory locations to lower memory locations. This implies that a Stack PUSH command decreases the Stack Pointer.

The Stack Pointer points to the data SRAM Stack area where the Subroutine and Interrupt Stacks are located. This Stack space in the data SRAM must be defined by the program before any subroutine calls are executed or interrupts are enabled. The Stack Pointer must be set to point above 0x0200. The initial value of the stack pointer is the last address of the internal SRAM. The Stack Pointer is decremented by one when data is pushed onto the Stack with the PUSH instruction, and it is decremented by two for ATmega640/1280/1281 and three for ATmega2560/2561 when the return address is pushed onto the Stack with subroutine call or interrupt. The Stack Pointer is incremented by one when data is popped from the Stack with the POP instruction, and it is incremented by two for ATmega640/1280/1281 and three for ATmega2560/2561 when data is popped from the Stack with return from subroutine RET or return from interrupt RETI.
The AVR Stack Pointer is implemented as two 8-bit registers in the I/O space. The number of bits actually used is implementation dependent. Note that the data space in some implementations of the AVR architecture is so small that only SPL is needed. In this case, the SPH Register will not be present.

| Bit | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \times 3 \mathrm{E}$ (0x5E) | SP15 | SP14 | SP13 | SP12 | SP11 | SP10 | SP9 | SP8 |
| 0x3D (0x5D) | SP7 | SP6 | SP5 | SP4 | SP3 | SP2 | SP1 | SP0 |
|  | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
|  | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 |
|  | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |

### 7.6.1 RAMPZ - Extended Z-pointer Register for ELPM/SPM

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | RAMPZ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x3B (0x5B) | RAMPZ7 | RAMPZ6 | RAMPZ5 | RAMPZ4 | RAMPZ3 | RAMPZ2 | RAMPZ1 | RAMPZ0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

For ELPM/SPM instructions, the Z-pointer is a concatenation of RAMPZ, ZH, and ZL, as shown in Figure 7-4. Note that LPM is not affected by the RAMPZ setting.

Figure 7-4. The Z-pointer used by ELPM and SPM


The actual number of bits is implementation dependent. Unused bits in an implementation will always read as zero. For compatibility with future devices, be sure to write these bits to zero.

### 7.6.2 EIND - Extended Indirect Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | EIND |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x3C (0x5C) | EIND7 | EIND6 | EIND5 | EIND4 | EIND3 | EIND2 | EIND1 | EIND0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

For EICALL/EIJMP instructions, the Indirect-pointer to the subroutine/routine is a concatenation of EIND, ZH, and ZL, as shown in Figure 7-5. Note that ICALL and IJMP are not affected by the EIND setting.

Figure 7-5. The Indirect-pointer used by EICALL and EIJMP


The actual number of bits is implementation dependent. Unused bits in an implementation will always read as zero. For compatibility with future devices, be sure to write these bits to zero.

### 7.7 Instruction Execution Timing

This section describes the general access timing concepts for instruction execution. The AVR CPU is driven by the CPU clock clk $_{\text {CPU }}$, directly generated from the selected clock source for the chip. No internal clock division is used.

Figure 7-6 on page 17 shows the parallel instruction fetches and instruction executions enabled by the Harvard architecture and the fast-access Register File concept. This is the basic pipelining concept to obtain up to 1 MIPS per MHz with the corresponding unique results for functions per cost, functions per clocks, and functions per power-unit.

Figure 7-6. The Parallel Instruction Fetches and Instruction Executions


Figure 7-7 shows the internal timing concept for the Register File. In a single clock cycle an ALU operation using two register operands is executed, and the result is stored back to the destination register.

Figure 7-7. Single Cycle ALU Operation


### 7.8 Reset and Interrupt Handling

The AVR provides several different interrupt sources. These interrupts and the separate Reset Vector each have a separate program vector in the program memory space. All interrupts are assigned individual enable bits which must be written logic one together with the Global Interrupt Enable bit in the Status Register in order to enable the interrupt. Depending on the Program Counter value, interrupts may be automatically disabled when Boot Lock bits BLB02 or BLB12 are programmed. This feature improves software security. See the section "Memory Programming" on page 325 for details.

The lowest addresses in the program memory space are by default defined as the Reset and Interrupt Vectors. The complete list of vectors is shown in "Interrupts" on page 101. The list also determines the priority levels of the different interrupts. The lower the address the higher is the priority level. RESET has the highest priority, and next is INTO - the External Interrupt Request 0 . The Interrupt Vectors can be moved to the start of the Boot Flash section by setting the IVSEL bit in the MCU Control Register (MCUCR). Refer to "Interrupts" on page 101 for more information. The Reset Vector can also be moved to the start of the Boot Flash section by programming the BOOTRST Fuse, see "Memory Programming" on page 325.
When an interrupt occurs, the Global Interrupt Enable I-bit is cleared and all interrupts are disabled. The user software can write logic one to the l-bit to enable nested interrupts. All enabled interrupts can then interrupt the current interrupt routine. The I-bit is automatically set when a Return from Interrupt instruction - RETI - is executed.

There are basically two types of interrupts. The first type is triggered by an event that sets the Interrupt Flag. For these interrupts, the Program Counter is vectored to the actual Interrupt Vector in order to execute the interrupt handling routine, and hardware clears the corresponding Interrupt Flag. Interrupt Flags can also be cleared by writing a logic one to the flag bit position(s) to be cleared. If an interrupt condition occurs while the corresponding interrupt enable bit is cleared, the Interrupt Flag will be set and remembered until the interrupt is enabled, or the
flag is cleared by software. Similarly, if one or more interrupt conditions occur while the Global Interrupt Enable bit is cleared, the corresponding Interrupt Flag(s) will be set and remembered until the Global Interrupt Enable bit is set, and will then be executed by order of priority.

The second type of interrupts will trigger as long as the interrupt condition is present. These interrupts do not necessarily have Interrupt Flags. If the interrupt condition disappears before the interrupt is enabled, the interrupt will not be triggered.

When the AVR exits from an interrupt, it will always return to the main program and execute one more instruction before any pending interrupt is served.

Note that the Status Register is not automatically stored when entering an interrupt routine, nor restored when returning from an interrupt routine. This must be handled by software.

When using the CLI instruction to disable interrupts, the interrupts will be immediately disabled. No interrupt will be executed after the CLI instruction, even if it occurs simultaneously with the CLI instruction. The following example shows how this can be used to avoid interrupts during the timed EEPROM write sequence.


When using the SEI instruction to enable interrupts, the instruction following SEI will be executed before any pending interrupts, as shown in this example.

```
Assembly Code Example
    sei ; set Global Interrupt Enable
    sleep; enter sleep, waiting for interrupt
    ; note: will enter sleep before any pending
    ; interrupt(s)
C Code Example
    __enable_interrupt(); /* set Global Interrupt Enable */
    __sleep(); /* enter sleep, waiting for interrupt */
    /* note: will enter sleep before any pending interrupt(s) */
```


### 7.8.1 Interrupt Response Time

The interrupt execution response for all the enabled AVR interrupts is five clock cycles minimum. After five clock cycles the program vector address for the actual interrupt handling routine is executed. During these five clock cycle period, the Program Counter is pushed onto the Stack. The vector is normally a jump to the interrupt routine, and this jump takes three clock cycles. If an interrupt occurs during execution of a multi-cycle instruction, this instruction is completed before the interrupt is served. If an interrupt occurs when the MCU is in sleep mode, the interrupt execution response time is increased by five clock cycles. This increase comes in addition to the start-up time from the selected sleep mode.
A return from an interrupt handling routine takes five clock cycles. During these five clock cycles, the Program Counter (three bytes) is popped back from the Stack, the Stack Pointer is incremented by three, and the I-bit in SREG is set.

## 8. AVR Memories

This section describes the different memories in the ATmega640/1280/1281/2560/2561. The AVR architecture has two main memory spaces, the Data Memory and the Program Memory space. In addition, the ATmega640/1280/1281/2560/2561 features an EEPROM Memory for data storage. All three memory spaces are linear and regular.

### 8.1 In-System Reprogrammable Flash Program Memory

The ATmega640/1280/1281/2560/2561 contains 64K/128K/256K bytes On-chip In-System Reprogrammable Flash memory for program storage, see Figure 8-1. Since all AVR instructions are 16 bit or 32 bit wide, the Flash is organized as $32 \mathrm{~K} / 64 \mathrm{~K} / 128 \mathrm{~K} \times 16$. For software security, the Flash Program memory space is divided into two sections, Boot Program section and Application Program section.

The Flash memory has an endurance of at least 10,000 write/erase cycles. The ATmega640/1280/1281/2560/2561 Program Counter (PC) is $15 / 16 / 17$ bits wide, thus addressing the $32 \mathrm{~K} / 64 \mathrm{~K} / 128 \mathrm{~K}$ program memory locations. The operation of Boot Program section and associated Boot Lock bits for software protection are described in detail in "Boot Loader Support - Read-While-Write Self-Programming" on page 310. "Memory Programming" on page 325 contains a detailed description on Flash data serial downloading using the SPI pins or the JTAG interface.
Constant tables can be allocated within the entire program memory address space (see the LPM - Load Program Memory instruction description and ELPM - Extended Load Program Memory instruction description).

Timing diagrams for instruction fetch and execution are presented in "Instruction Execution Timing" on page 16.
Figure 8-1. Program Flash Memory Map


### 8.2 SRAM Data Memory

Figure $8-2$ on page 22 shows how the ATmega640/1280/1281/2560/2561 SRAM Memory is organized.
The ATmega640/1280/1281/2560/2561 is a complex microcontroller with more peripheral units than can be supported within the 64 location reserved in the Opcode for the IN and OUT instructions. For the Extended I/O space from \$060-\$1FF in SRAM, only the ST/STS/STD and LD/LDS/LDD instructions can be used.
The first 4,608/8,704 Data Memory locations address both the Register File, the I/O Memory, Extended I/O Memory, and the internal data SRAM. The first 32 locations address the Register file, the next 64 location the standard I/O Memory, then 416 locations of Extended I/O memory and the next 8,192 locations address the internal data SRAM.

An optional external data SRAM can be used with the ATmega640/1280/1281/2560/2561. This SRAM will occupy an area in the remaining address locations in the 64 K address space. This area starts at the address following the internal SRAM. The Register file, I/O, Extended I/O and Internal SRAM occupies the lowest 4,608/8,704 bytes, so when using 64 Kbytes ( 65,536 bytes) of External Memory, $60,478 / 56,832$ Bytes of External Memory are available. See "External Memory Interface" on page 27 for details on how to take advantage of the external memory map.

When the addresses accessing the SRAM memory space exceeds the internal data memory locations, the external data SRAM is accessed using the same instructions as for the internal data memory access. When the internal data memories are accessed, the read and write strobe pins ( $\overline{\mathrm{PGO}}$ and $\overline{\mathrm{PG} 1}$ ) are inactive during the whole access cycle. External SRAM operation is enabled by setting the SRE bit in the XMCRA Register.

Accessing external SRAM takes one additional clock cycle per byte compared to access of the internal SRAM. This means that the commands LD, ST, LDS, STS, LDD, STD, PUSH, and POP take one additional clock cycle. If the Stack is placed in external SRAM, interrupts, subroutine calls and returns take three clock cycles extra because the three-byte program counter is pushed and popped, and external memory access does not take advantage of the internal pipe-line memory access. When external SRAM interface is used with wait-state, onebyte external access takes two, three, or four additional clock cycles for one, two, and three wait-states respectively. Interrupts, subroutine calls and returns will need five, seven, or nine clock cycles more than specified in the instruction set manual for one, two, and three wait-states.

The five different addressing modes for the data memory cover: Direct, Indirect with Displacement, Indirect, Indirect with Pre-decrement, and Indirect with Post-increment. In the Register file, registers R26 to R31 feature the indirect addressing pointer registers.

The direct addressing reaches the entire data space.
The Indirect with Displacement mode reaches 63 address locations from the base address given by the Y -register or Z-register.

When using register indirect addressing modes with automatic pre-decrement and post-increment, the address registers $\mathrm{X}, \mathrm{Y}$, and Z are decremented or incremented.

The 32 general purpose working registers, 64 I/O registers, and the 4,196/8,192 bytes of internal data SRAM in the ATmega640/1280/1281/2560/2561 are all accessible through all these addressing modes. The Register File is described in "General Purpose Register File" on page 13.

Figure 8-2. Data Memory Map
Address (HEX)
$0-1 F$
$20-5 F$
$60-1 \mathrm{FF}$

200
21FF
2200

|  |
| :---: |
| 32 Registers |
| 64 I/O Registers |
| 416 External I/O Registers |
| Internal SRAM |
| $(8192 \times 8)$ |

FFFF

### 8.2.1 Data Memory Access Times

This section describes the general access timing concepts for internal memory access. The internal data SRAM access is performed in two $\mathrm{clk}_{\mathrm{CPU}}$ cycles as described in Figure 8-3.

Figure 8-3. On-chip Data SRAM Access Cycles


### 8.3 EEPROM Data Memory

The ATmega640/1280/1281/2560/2561 contains 4Kbytes of data EEPROM memory. It is organized as a separate data space, in which single bytes can be read and written. The EEPROM has an endurance of at least 100,000 write/erase cycles. The access between the EEPROM and the CPU is described in the following, specifying the EEPROM Address Registers, the EEPROM Data Register, and the EEPROM Control Register.

For a detailed description of SPI, JTAG and Parallel data downloading to the EEPROM, see "Serial Downloading" on page 338, "Programming via the JTAG Interface" on page 342, and "Programming the EEPROM" on page 333 respectively.

### 8.3.1 EEPROM Read/Write Access

The EEPROM Access Registers are accessible in the I/O space, see "Register Description" on page 34.
The write access time for the EEPROM is given in Table 8-1. A self-timing function, however, lets the user software detect when the next byte can be written. If the user code contains instructions that write the EEPROM, some precautions must be taken. In heavily filtered power supplies, $\mathrm{V}_{\mathrm{CC}}$ is likely to rise or fall slowly on power-up/down. This causes the device for some period of time to run at a voltage lower than specified as minimum for the clock frequency used. See "Preventing EEPROM Corruption" on page 25. for details on how to avoid problems in these situations.

In order to prevent unintentional EEPROM writes, a specific write procedure must be followed. See the description of the EEPROM Control Register for details on this; "Register Description" on page 34.

When the EEPROM is read, the CPU is halted for four clock cycles before the next instruction is executed. When the EEPROM is written, the CPU is halted for two clock cycles before the next instruction is executed.

The calibrated Oscillator is used to time the EEPROM accesses. Table 8-1 lists the typical programming time for EEPROM access from the CPU.

Table 8-1. EEPROM Programming Time

| Symbol | Number of Calibrated RC Oscillator Cycles | Typ Programming Time |
| :--- | :---: | :---: |
| EEPROM write (from CPU) | 26,368 | 3.3 ms |

The following code examples show one assembly and one C function for writing to the EEPROM. The examples assume that interrupts are controlled (for example by disabling interrupts globally) so that no interrupts will occur during execution of these functions. The examples also assume that no Flash Boot Loader is present in the software. If such code is present, the EEPROM write function must also wait for any ongoing SPM command to finish.

```
Assembly Code Example \({ }^{(1)}\)
    EEPROM_write:
    ; Wait for completion of previous write
    sbic EECR, EEPE
    rjmp EEPROM_write
    ; Set up address (r18:r17) in address register
    out EEARH, r18
    out EEARL, r17
    ; Write data (r16) to Data Register
    out EEDR,r16
    ; Write logical one to EEMPE
    sbi EECR,EEMPE
    ; Start eeprom write by setting EEPE
    sbi EECR,EEPE
    ret
C Code Example \({ }^{(1)}\)
void EEPROM_write(unsigned int uiAddress, unsigned char ucData)
\{
    /* Wait for completion of previous write */
    while(EECR \& ( \(1 \ll E E P E)\) )
        ;
    /* Set up address and Data Registers */
    EEAR = uiAddress;
    EEDR = ucData;
    /* Write logical one to EEMPE */
    EECR |= ( \(1 \ll\) EEMPE);
    /* Start eeprom write by setting EEPE */
    EECR |= ( \(1 \ll \mathrm{EEPE})\);
    \}
```

Note: 1. See "About Code Examples" on page 10.

The next code examples show assembly and C functions for reading the EEPROM. The examples assume that interrupts are controlled so that no interrupts will occur during execution of these functions.

```
Assembly Code Example }\mp@subsup{}{}{(1)
    EEPROM_read:
    ; Wait for completion of previous write
    sbic EECR,EEPE
    rjcmp EEPROM_read
    ; Set up address (r18:r17) in address register
    out EEARH, r18
    out EEARL, r17
    ; Start eeprom read by writing EERE
    sbi EECR,EERE
    ; Read data from Data Register
    in r16,EEDR
    ret
C Code Example (1)
    unsigned char EEPROM_read(unsigned int uiAddress)
    {
    /* Wait for completion of previous write */
    while(EECR & (1<<EEPE))
        ;
    /* Set up address register */
    EEAR = uiAddress;
    /* Start eeprom read by writing EERE */
    EECR |= (1<<EERE);
    /* Return data from Data Register */
    return EEDR;
    }
```

Note: 1. See "About Code Examples" on page 10.

### 8.3.2 Preventing EEPROM Corruption

During periods of low $\mathrm{V}_{\mathrm{CC}}$, the EEPROM data can be corrupted because the supply voltage is too low for the CPU and the EEPROM to operate properly. These issues are the same as for board level systems using EEPROM, and the same design solutions should be applied.
An EEPROM data corruption can be caused by two situations when the voltage is too low. First, a regular write sequence to the EEPROM requires a minimum voltage to operate correctly. Secondly, the CPU itself can execute instructions incorrectly, if the supply voltage is too low.

EEPROM data corruption can easily be avoided by following this design recommendation:
Keep the AVR RESET active (low) during periods of insufficient power supply voltage. This can be done by enabling the internal Brown-out Detector (BOD). If the detection level of the internal BOD does not match the needed detection level, an external low $\mathrm{V}_{\mathrm{CC}}$ reset Protection circuit can be used. If a reset occurs while a write operation is in progress, the write operation will be completed provided that the power supply voltage is sufficient.

### 8.4 I/O Memory

The I/O space definition of the ATmega640/1280/1281/2560/2561 is shown in "Register Summary" on page 399.
All ATmega640/1280/1281/2560/2561 I/Os and peripherals are placed in the I/O space. All I/O locations may be accessed by the LD/LDS/LDD and ST/STS/STD instructions, transferring data between the 32 general purpose working registers and the I/O space. I/O Registers within the address range $0 \times 00-0 \times 1 \mathrm{~F}$ are directly bit-accessible using the SBI and CBI instructions. In these registers, the value of single bits can be checked by using the SBIS and SBIC instructions. Refer to the "Instruction Set Summary" on page 404 for more details. When using the I/O specific commands IN and OUT, the I/O addresses $0 \times 00-0 \times 3 \mathrm{~F}$ must be used. When addressing I/O Registers as data space using LD and ST instructions, $0 \times 20$ must be added to these addresses. The ATmega640/1280/1281/2560/2561 is a complex microcontroller with more peripheral units than can be supported within the 64 location reserved in Opcode for the IN and OUT instructions. For the Extended I/O space from 0x600x1FF in SRAM, only the ST/STS/STD and LD/LDS/LDD instructions can be used.

For compatibility with future devices, reserved bits should be written to zero if accessed. Reserved I/O memory addresses should never be written.

Some of the Status Flags are cleared by writing a logical one to them. Note that, unlike most other AVRs, the CBI and SBI instructions will only operate on the specified bit, and can therefore be used on registers containing such Status Flags. The CBI and SBI instructions work with registers $0 \times 00$ to $0 \times 1 \mathrm{~F}$ only.
The I/O and peripherals control registers are explained in later sections.

### 8.4.1 General Purpose I/O Registers

The ATmega640/1280/1281/2560/2561 contains three General Purpose I/O Registers. These registers can be used for storing any information, and they are particularly useful for storing global variables and Status Flags. General Purpose I/O Registers within the address range $0 \times 00-0 \times 1 \mathrm{~F}$ are directly bit-accessible using the SBI, CBI, SBIS, and SBIC instructions. See "Register Description" on page 34.

## 9. External Memory Interface

With all the features the External Memory Interface provides, it is well suited to operate as an interface to memory devices such as External SRAM and Flash, and peripherals such as LCD-display, A/D, and D/A. The main features are:

- Four different wait-state settings (including no wait-state)
- Independent wait-state setting for different External Memory sectors (configurable sector size)
- The number of bits dedicated to address high byte is selectable
- Bus keepers on data lines to minimize current consumption (optional)


### 9.1 Overview

When the eXternal MEMory (XMEM) is enabled, address space outside the internal SRAM becomes available using the dedicated External Memory pins (see Figure 1-3 on page 4, Table 13-3 on page 75, Table 13-9 on page 79, and Table 13-21 on page 86). The memory configuration is shown in Figure 9-1.

Figure 9-1. External Memory with Sector Select
Memory Configuration A


### 9.1.1 Using the External Memory Interface

The interface consists of:

- AD7:0: Multiplexed low-order address bus and data bus

A15:8: High-order address bus (configurable number of bits)

- ALE: Address latch enable
- $\overline{\mathrm{RD}}$ : Read strobe
- $\overline{\mathrm{WR}}:$ Write strobe

The control bits for the External Memory Interface are located in two registers, the External Memory Control Register A - XMCRA, and the External Memory Control Register B - XMCRB.
When the XMEM interface is enabled, the XMEM interface will override the setting in the data direction registers that corresponds to the ports dedicated to the XMEM interface. For details about the port override, see the alternate functions in section "//O-Ports" on page 67. The XMEM interface will auto-detect whether an access is internal or external. If the access is external, the XMEM interface will output address, data, and the control signals on the ports according to Figure 9-3 on page 29 (this figure shows the wave forms without wait-states). When ALE goes from high-to-low, there is a valid address on AD7:0. ALE is low during a data transfer. When the XMEM interface is enabled, also an internal access will cause activity on address, data and ALE ports, but the $\overline{\mathrm{RD}}$ and $\overline{\mathrm{WR}}$ strobes will not toggle during internal access. When the External Memory Interface is disabled, the normal pin and data direction settings are used. Note that when the XMEM interface is disabled, the address space above the internal SRAM boundary is not mapped into the internal SRAM. Figure 9-2 illustrates how to connect an external SRAM to the AVR using an octal latch (typically " $74 \times 573$ " or equivalent) which is transparent when $G$ is high.

### 9.1.2 Address Latch Requirements

Due to the high-speed operation of the XRAM interface, the address latch must be selected with care for system frequencies above $8 \mathrm{MHz} @ 4 \mathrm{~V}$ and $4 \mathrm{MHz} @ 2.7 \mathrm{~V}$. When operating at conditions above these frequencies, the typical old style 74 HC series latch becomes inadequate. The External Memory Interface is designed in compliance to the 74AHC series latch. However, most latches can be used as long they comply with the main timing parameters. The main parameters for the address latch are:

- D to Q propagation delay ( $t_{\text {PD }}$ )
- Data setup time before G low ( $\mathrm{t}_{\mathrm{su}}$ )
- Data (address) hold time after G low ( TH )

The External Memory Interface is designed to guaranty minimum address hold time after $G$ is asserted low of $t_{h}=$ 5ns. Refer to thaxx_ld/tLAXX_st in "External Data Memory Timing" Tables 31-11 through Tables 31-18 on pages 367 - 370. The D-to-Q propagation delay ( $\mathrm{t}_{\text {PD }}$ ) must be taken into consideration when calculating the access time requirement of the external component. The data setup time before $G$ low ( $\mathrm{t}_{\mathrm{su}}$ ) must not exceed address valid to ALE low ( $\mathrm{t}_{\text {AVLLC }}$ ) minus PCB wiring delay (dependent on the capacitive load).

Figure 9-2. External SRAM Connected to the AVR


### 9.1.3 Pull-up and Bus-keeper

The pull-ups on the AD7:0 ports may be activated if the corresponding Port register is written to one. To reduce power consumption in sleep mode, it is recommended to disable the pull-ups by writing the Port register to zero before entering sleep.
The XMEM interface also provides a bus-keeper on the AD7:0 lines. The bus-keeper can be disabled and enabled in software as described in "XMCRB - External Memory Control Register B" on page 38. When enabled, the buskeeper will keep the previous value on the AD7:0 bus while these lines are tri-stated by the XMEM interface.

### 9.1.4 Timing

External Memory devices have different timing requirements. To meet these requirements, the XMEM interface provides four different wait-states as shown in Table 9-3 on page 37. It is important to consider the timing specification of the External Memory device before selecting the wait-state. The most important parameters are the access time for the external memory compared to the set-up requirement. The access time for the External Memory is defined to be the time from receiving the chip select/address until the data of this address actually is driven on the bus. The access time cannot exceed the time from the ALE pulse must be asserted low until data is stable during a read sequence (see $\mathrm{t}_{\text {LLRL }}+\mathrm{t}_{\text {RLRH }}-\mathrm{t}_{\text {DVRH }}$ in Tables 31-11 through Tables 31-18 on pages 367-370). The different wait-states are set up in software. As an additional feature, it is possible to divide the external memory space in two sectors with individual wait-state settings. This makes it possible to connect two different memory devices with different timing requirements to the same XMEM interface. For XMEM interface timing details, refer to Table 31-11 on page 367 to Table 31-18 on page 370 and Figure 31-9 on page 370 to Figure 31-12 on page 372 in the "External Data Memory Timing" on page 367.

Note that the XMEM interface is asynchronous and that the waveforms in the following figures are related to the internal system clock. The skew between the internal and external clock (XTAL1) is not guarantied (varies between devices temperature, and supply voltage). Consequently, the XMEM interface is not suited for synchronous operation.

Figure 9-3. External Data Memory Cycles without Wait-state $(\mathrm{SRWn} 1=0 \text { and } \mathrm{SRWnO}=0)^{(1)}$


Note: 1. SRWn1 = SRW11 (upper sector) or SRW01 (lower sector), SRWn0 = SRW10 (upper sector) or SRW00 (lower sector). The ALE pulse in period T4 is only present if the next instruction accesses the RAM (internal or external).

Figure 9-4. External Data Memory Cycles with SRWn1 $=0$ and SRWn0 $=1^{(1)}$


Note: 1. SRWn1 = SRW11 (upper sector) or SRW01 (lower sector), SRWn0 = SRW10 (upper sector) or SRW00 (lower sector).
The ALE pulse in period T5 is only present if the next instruction accesses the RAM (internal or external).
Figure 9-5. External Data Memory Cycles with SRWn1 $=1$ and SRWn0 $=0^{(1)}$


Note: 1. SRWn1 = SRW11 (upper sector) or SRW01 (lower sector), SRWn0 $=$ SRW10 (upper sector) or SRW00 (lower sector).
The ALE pulse in period T6 is only present if the next instruction accesses the RAM (internal or external).

Figure 9-6. External Data Memory Cycles with SRWn1 $=1$ and SRWn0 $=1^{(1)}$


Note: 1. SRWn1 = SRW11 (upper sector) or SRW01 (lower sector), SRWn0 = SRW10 (upper sector) or SRW00 (lower sector).
The ALE pulse in period T7 is only present if the next instruction accesses the RAM (internal or external).

### 9.1.5 Using all Locations of External Memory Smaller than 64Kbytes

Since the external memory is mapped after the internal memory as shown in Figure 9-1 on page 27, the external memory is not addressed when addressing the first 8,704 bytes of data space. It may appear that the first 8,704 bytes of the external memory are inaccessible (external memory addresses $0 \times 0000$ to $0 \times 21 F F$ ). However, when connecting an external memory smaller than 64Kbytes, for example 32Kbytes, these locations are easily accessed simply by addressing from address $0 \times 8000$ to 0xA1FF. Since the External Memory Address bit A15 is not connected to the external memory, addresses $0 \times 8000$ to 0xA1FF will appear as addresses $0 \times 0000$ to $0 \times 21 \mathrm{FF}$ for the external memory. Addressing above address 0xA1FF is not recommended, since this will address an external memory location that is already accessed by another (lower) address. To the Application software, the external 32 Kbytes memory will appear as one linear 32Kbytes address space from 0x2200 to 0xA1FF. This is illustrated in Figure 9-7.

Figure 9-7. Address Map with 32Kbytes External Memory


### 9.1.6 Using all 64Kbytes Locations of External Memory

Since the External Memory is mapped after the Internal Memory as shown in Figure 9-1 on page 27, only 56Kbytes of External Memory is available by default (address space $0 \times 0000$ to $0 \times 21 \mathrm{FF}$ is reserved for internal memory). However, it is possible to take advantage of the entire External Memory by masking the higher address bits to zero. This can be done by using the XMMn bits and control by software the most significant bits of the address. By setting Port C to output $0 \times 00$, and releasing the most significant bits for normal Port Pin operation, the Memory Interface will address 0x0000-0x2FFF. See the following code examples.

Care must be exercised using this option as most of the memory is masked away.

```
Assembly Code Example }\mp@subsup{}{}{(1)
    OFFSET is defined to 0x4000 to ensure
    external memory access
    Configure Port C (address high byte) to
    output 0x00 when the pins are released
    for normal Port Pin operation
ldi r16, 0xFF
out DDRC, r16
ldi r16, 0x00
out PORTC, r16
; release PC7:6
ldi r16, (1<<XMM1)
sts XMCRB, r16
; write 0xAA to address 0x0001 of external
; memory
ldi r16, 0xaa
sts 0x0001+OFFSET, r16
; re-enable PC7:6 for external memory
ldi r16, (0<<XMM1)
sts XMCRB, r16
; store 0x55 to address (OFFSET + 1) of
; external memory
ldi r16, 0x55
sts 0x0001+OFFSET, r16
C Code Example \({ }^{(1)}\)
#define OFFSET 0x4000
void XRAM_example(void)
{
    unsigned char *p = (unsigned char *) (OFFSET + 1);
    DDRC = 0xFF;
    PORTC = 0x00;
    XMCRB = (1<<XMM1);
    *p = Oxaa;
    XMCRB = 0x00;
    *p = 0x55;
}
```

Note: 1. See "About Code Examples" on page 10.

### 9.2 Register Description

### 9.2.1 EEPROM registers

### 9.2.1.1 EEARH and EEARL - The EEPROM Address Register

| Bit | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | EEARH <br> EEARL |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x22 (0x42) | - | - | - | - | EEAR11 | EEAR10 | EEAR9 | EEAR8 |  |
| 0x21 (0x41) | EEAR7 | EEAR6 | EEAR5 | EEAR4 | EEAR3 | EEAR2 | EEAR1 | EEAR0 |  |
|  | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| Read/Write | R | R | R | R | R/W | R/W | R/W | R/W |  |
|  | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | X | X | X | X |  |
|  | X | X | X | X | X | X | X | X |  |

## - Bits 15:12 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

## - Bits 11:0 - EEAR8:0: EEPROM Address

The EEPROM Address Registers - EEARH and EEARL specify the EEPROM address in the 4Kbytes EEPROM space. The EEPROM data bytes are addressed linearly between 0 and 4096. The initial value of EEAR is undefined. A proper value must be written before the EEPROM may be accessed.

### 9.2.1.2 EEDR - The EEPROM Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | EEDR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x20 (0x40) | MSB |  |  |  |  |  |  | LSB |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7:0 - EEDR7:0: EEPROM Data

For the EEPROM write operation, the EEDR Register contains the data to be written to the EEPROM in the address given by the EEAR Register. For the EEPROM read operation, the EEDR contains the data read out from the EEPROM at the address given by EEAR.

### 9.2.1.3

EECR - The EEPROM Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | EECR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x1F (0x3F) | - | - | EEPM1 | EEPMO | EERIE | EEMPE | EEPE | EERE |  |
| Read/Write | R | R | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | X | X | 0 | 0 | X | 0 |  |

## - Bits 7:6 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

## - Bits 5, 4 - EEPM1 and EEPMO: EEPROM Programming Mode Bits

The EEPROM Programming mode bit setting defines which programming action that will be triggered when writing EEPE. It is possible to program data in one atomic operation (erase the old value and program the new value) or to split the Erase and Write operations in two different operations. The Programming times for the different modes are shown in Table 9-1 on page 35. While EEPE is set, any write to EEPMn will be ignored. During reset, the EEPMn bits will be reset to ObOO unless the EEPROM is busy programming.

Table 9-1. EEPROM Mode Bits

| EEPM1 | EEPM0 | Programming Time | Operation |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 3.4 ms | 1.8 ms |
| 0 | 1 | 1.8 ms | Erase and Write in one operation (Atomic Operation) |
| 1 | 0 | - | Wrase only |
| 1 | 1 | Reserved for future use |  |

## - Bit 3 - EERIE: EEPROM Ready Interrupt Enable

Writing EERIE to one enables the EEPROM Ready Interrupt if the I bit in SREG is set. Writing EERIE to zero disables the interrupt. The EEPROM Ready interrupt generates a constant interrupt when EEPE is cleared.

## - Bit 2 - EEMPE: EEPROM Master Programming Enable

The EEMPE bit determines whether setting EEPE to one causes the EEPROM to be written. When EEMPE is set, setting EEPE within four clock cycles will write data to the EEPROM at the selected address If EEMPE is zero, setting EEPE will have no effect. When EEMPE has been written to one by software, hardware clears the bit to zero after four clock cycles. See the description of the EEPE bit for an EEPROM write procedure.

## - Bit 1 - EEPE: EEPROM Programming Enable

The EEPROM Write Enable Signal EEPE is the write strobe to the EEPROM. When address and data are correctly set up, the EEPE bit must be written to one to write the value into the EEPROM. The EEMPE bit must be written to one before a logical one is written to EEPE, otherwise no EEPROM write takes place. The following procedure should be followed when writing the EEPROM (the order of steps 3 and 4 is not essential):

1. Wait until EEPE becomes zero.
2. Wait until SPMEN in SPMCSR becomes zero.
3. Write new EEPROM address to EEAR (optional).
4. Write new EEPROM data to EEDR (optional).
5. Write a logical one to the EEMPE bit while writing a zero to EEPE in EECR.
6. Within four clock cycles after setting EEMPE, write a logical one to EEPE.

The EEPROM can not be programmed during a CPU write to the Flash memory. The software must check that the Flash programming is completed before initiating a new EEPROM write. Step 2 is only relevant if the software contains a Boot Loader allowing the CPU to program the Flash. If the Flash is never being updated by the CPU, step 2 can be omitted. See "Memory Programming" on page 325 for details about Boot programming.

Caution: An interrupt between step 5 and step 6 will make the write cycle fail, since the EEPROM Master Write Enable will time-out. If an interrupt routine accessing the EEPROM is interrupting another EEPROM access, the EEAR or EEDR Register will be modified, causing the interrupted EEPROM access to fail. It is recommended to have the Global Interrupt Flag cleared during all the steps to avoid these problems.

When the write access time has elapsed, the EEPE bit is cleared by hardware. The user software can poll this bit and wait for a zero before writing the next byte. When EEPE has been set, the CPU is halted for two cycles before the next instruction is executed.

## - Bit 0 - EERE: EEPROM Read Enable

The EEPROM Read Enable Signal EERE is the read strobe to the EEPROM. When the correct address is set up in the EEAR Register, the EERE bit must be written to a logic one to trigger the EEPROM read. The EEPROM read access takes one instruction, and the requested data is available immediately. When the EEPROM is read, the CPU is halted for four cycles before the next instruction is executed.

The user should poll the EEPE bit before starting the read operation. If a write operation is in progress, it is neither possible to read the EEPROM, nor to change the EEAR Register.

### 9.3 General Purpose registers

### 9.3.1 GPIOR2 - General Purpose I/O Register 2

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | GPIOR2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x2B (0x4B) | MSB |  |  |  |  |  |  | LSB |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 9.3.2 GPIOR1 - General Purpose I/O Register 1

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | GPIOR1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x2A (0x4A) | MSB |  |  |  |  |  |  | LSB |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 9.3.3 GPIORO - General Purpose I/O Register 0

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | GPIOR0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x1E (0x3E) | MSB |  |  |  |  |  |  | LSB |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 9.4 External Memory registers

### 9.4.1 XMCRA - External Memory Control Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | XMCRA |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| "(0x74)" | SRE | SRL2 | SRL1 | SRLO | SRW11 | SRW10 | SRW01 | SRW00 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - SRE: External SRAM/XMEM Enable

Writing SRE to one enables the External Memory Interface.The pin functions AD7:0, A15:8, ALE, $\overline{\mathrm{WR}}$, and $\overline{\mathrm{RD}}$ are activated as the alternate pin functions. The SRE bit overrides any pin direction settings in the respective data direction registers. Writing SRE to zero, disables the External Memory Interface and the normal pin and data direction settings are used.

## - Bit 6:4 - SRL2:0: Wait-state Sector Limit

It is possible to configure different wait-states for different External Memory addresses. The external memory address space can be divided in two sectors that have separate wait-state bits. The SRL2, SRL1, and SRLO bits select the split of the sectors, see Table 9-2 on page 37 and Figure 9-1 on page 27. By default, the SRL2, SRL1, and SRLO bits are set to zero and the entire external memory address space is treated as one sector. When the entire SRAM address space is configured as one sector, the wait-states are configured by the SRW11 and SRW10 bits.

Table 9-2. Sector limits with different settings of SRL2:0

| SRL2 | SRL1 | SRLO | Sector Limits |
| :---: | :---: | :---: | :---: |
| 0 | 0 | x | Lower sector $=\mathrm{N} / \mathrm{A}$ <br> Upper sector $=0 \times 2200-0 \times F F F F$ |
| 0 | 1 | 0 | Lower sector $=0 \times 2200-0 \times 3 F F F$ <br> Upper sector $=0 \times 4000-0 \times F F F F$ |
| 0 | 1 | 1 | Lower sector $=0 \times 2200-0 \times 5$ FFF <br> Upper sector $=0 \times 6000-0 x F F F F$ |
| 1 | 0 | 0 | Lower sector $=0 \times 2200-0 \times 7 F F F$ <br> Upper sector $=0 \times 8000-0 x F F F F$ |
| 1 | 0 | 1 | $\begin{aligned} & \text { Lower sector }=0 \times 2200-0 \times 9 F F F \\ & \text { Upper sector }=0 \times A 000-0 x F F F F \end{aligned}$ |
| 1 | 1 | 0 | $\begin{aligned} & \text { Lower sector }=0 \times 2200-0 \times B F F F \\ & \text { Upper sector }=0 \times C 000-0 x F F F F \end{aligned}$ |
| 1 | 1 | 1 | $\begin{aligned} & \text { Lower sector }=0 \times 2200-0 \times D F F F \\ & \text { Upper sector }=0 \times E 000-0 x F F F F \end{aligned}$ |

## - Bit 3:2 - SRW11, SRW10: Wait-state Select Bits for Upper Sector

The SRW11 and SRW10 bits control the number of wait-states for the upper sector of the external memory address space, see Table 9-3.

- Bit 1:0 - SRW01, SRW00: Wait-state Select Bits for Lower Sector

The SRW01 and SRW00 bits control the number of wait-states for the lower sector of the external memory address space, see Table 9-3.

Table 9-3. Wait States ${ }^{(1)}$

| SRWn1 | SRWn0 | Wait States |
| :---: | :---: | :---: |
| 0 | 0 | No wait-states |
| 0 | 1 | Wait one cycle during read/write strobe |
| 1 | 0 | Wait two cycles during read/write strobe |
| 1 | 1 | Wait two cycles during read/write and wait one cycle before driving out new address |

Note: 1. $\mathrm{n}=0$ or 1 (lower/upper sector).
For further details of the timing and wait-states of the External Memory Interface, see Figure 9-3 on page 29 through Figure 9-6 on page 31 for how the setting of the SRW bits affects the timing.

### 9.4.2 XMCRB - External Memory Control Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | XMCRB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x75) | XMBK | - | - | - | - | XMM2 | XMM1 | XMM0 |  |
| Read/Write | R/W | R | R | R | R | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7- XMBK: External Memory Bus-keeper Enable

Writing XMBK to one enables the bus keeper on the AD7:0 lines. When the bus keeper is enabled, AD7:0 will keep the last driven value on the lines even if the XMEM interface has tri-stated the lines. Writing XMBK to zero disables the bus keeper. XMBK is not qualified with SRE, so even if the XMEM interface is disabled, the bus keepers are still activated as long as XMBK is one.

- Bit 6:3-Res: Reserved Bits

These bits are reserved and will always read as zero. When writing to this address location, write these bits to zero for compatibility with future devices.

## - Bit 2:0 - XMM2, XMM1, XMM0: External Memory High Mask

When the External Memory is enabled, all Port C pins are default used for the high address byte. If the full 60 Kbytes address space is not required to access the External Memory, some, or all, Port C pins can be released for normal Port Pin function as described in Table 9-4. As described in "Using all 64Kbytes Locations of External Memory" on page 32, it is possible to use the XMMn bits to access all 64Kbytes locations of the External Memory.

Table 9-4. Port C Pins Released as Normal Port Pins when the External Memory is Enabled

| XMM2 | XMM1 | XMM0 | \# Bits for External Memory Address | Released Port Pins |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 8 (Full 56Kbytes space) | None |
| 0 | 0 | 1 | 7 | PC7 |
| 0 | 1 | 0 | 6 | PC7 - PC6 |
| 0 | 1 | 1 | 5 | PC7 - PC5 |
| 1 | 0 | 0 | 4 | PC7 - PC4 |
| 1 | 0 | 1 | 3 | PC7 - PC3 |
| 1 | 1 | 0 | 2 | PC7 - PC2 |
| 1 | 1 | 1 | No Address high bits | Full Port C |

## 10. System Clock and Clock Options

This section describes the clock options for the AVR microcontroller.

### 10.1 Overview

Figure 10-1 presents the principal clock systems in the AVR and their distribution. All of the clocks need not be active at a given time. In order to reduce power consumption, the clocks to modules not being used can be halted by using different sleep modes, as described in "Power Management and Sleep Modes" on page 50. The clock systems are detailed below.

Figure 10-1. Clock Distribution.


### 10.2 Clock Systems and their Distribution

### 10.2.1 CPU Clock - clk $_{\text {cPU }}$

The CPU clock is routed to parts of the system concerned with operation of the AVR core. Examples of such modules are the General Purpose Register File, the Status Register and the data memory holding the Stack Pointer. Halting the CPU clock inhibits the core from performing general operations and calculations.

### 10.2.2 I/O Clock - clk $_{\text {/O }}$

The I/O clock is used by the majority of the I/O modules, like Timer/Counters, SPI, and USART. The I/O clock is also used by the External Interrupt module, but note that some external interrupts are detected by asynchronous logic, allowing such interrupts to be detected even if the I/O clock is halted. Also note that start condition detection in the USI module is carried out asynchronously when $\mathrm{clk}_{1 / \mathrm{O}}$ is halted, TWI address recognition in all sleep modes.

### 10.2.3 Flash Clock - clk $_{\text {FLASH }}$

The Flash clock controls operation of the Flash interface. The Flash clock is usually active simultaneously with the CPU clock.

### 10.2.4 Asynchronous Timer Clock - clk $_{\text {ASY }}$

The Asynchronous Timer clock allows the Asynchronous Timer/Counter to be clocked directly from an external clock or an external 32 kHz clock crystal. The dedicated clock domain allows using this Timer/Counter as a realtime counter even when the device is in sleep mode.

### 10.2.5 ADC Clock - clk ${ }_{\text {ADC }}$

The ADC is provided with a dedicated clock domain. This allows halting the CPU and I/O clocks in order to reduce noise generated by digital circuitry. This gives more accurate ADC conversion results.

### 10.3 Clock Sources

The device has the following clock source options, selectable by Flash Fuse bits as shown below. The clock from the selected source is input to the AVR clock generator, and routed to the appropriate modules.

Table 10-1. Device Clocking Options Select ${ }^{(1)}$

| Device Clocking Option | CKSEL3:0 |
| :---: | :---: |
| Low Power Crystal Oscillator | $1111-1000$ |
| Full Swing Crystal Oscillator | $0111-0110$ |
| Low Frequency Crystal Oscillator | $0101-0100$ |
| Internal 128kHz RC Oscillator | 0011 |
| Calibrated Internal RC Oscillator | 0010 |
| External Clock | 0000 |
| Reserved | 0001 |

Note: 1. For all fuses " 1 " means unprogrammed while " 0 " means programmed.

### 10.3.1 Default Clock Source

The device is shipped with internal RC oscillator at 8.0 MHz and with the fuse CKDIV8 programmed, resulting in 1.0 MHz system clock. The startup time is set to maximum and time-out period enabled. (CKSEL = "0010", SUT = " 10 ", CKDIV8 = "0"). The default setting ensures that all users can make their desired clock source setting using any available programming interface.

### 10.3.2 Clock Start-up Sequence

Any clock source needs a sufficient $\mathrm{V}_{\mathrm{CC}}$ to start oscillating and a minimum number of oscillating cycles before it can be considered stable.

To ensure sufficient $\mathrm{V}_{\mathrm{CC}}$, the device issues an internal reset with a time-out delay ( $\mathrm{t}_{\text {TOut }}$ ) after the device reset is released by all other reset sources. "On-chip Debug System" on page 53 describes the start conditions for the internal reset. The delay ( $\mathrm{t}_{\text {TOUT }}$ ) is timed from the Watchdog Oscillator and the number of cycles in the delay is set by the SUTx and CKSELx fuse bits. The selectable delays are shown in Table 10-2 on page 41. The frequency of the Watchdog Oscillator is voltage dependent as shown in "Typical Characteristics" on page 373.

Table 10-2. Number of Watchdog Oscillator Cycles

| Typical Time-out $\left(\mathrm{V}_{\mathrm{cc}}=\mathbf{5 . 0 V}\right)$ | Typical Time-out $\left(\mathrm{V}_{\mathrm{cc}}=\mathbf{3 . 0 V}\right)$ | Number of Cycles |
| :---: | :---: | :---: |
| 0 ms | 0 ms | 0 |
| 4.1 ms | 4.3 ms | 512 |
| 65 ms | 69 ms | $8 \mathrm{~K}(8,192)$ |

Main purpose of the delay is to keep the AVR in reset until it is supplied with minimum $\mathrm{V}_{\mathrm{CC}}$. The delay will not monitor the actual voltage and it will be required to select a delay longer than the $V_{C C}$ rise time. If this is not possible, an internal or external Brown-Out Detection circuit should be used. A BOD circuit will ensure sufficient $V_{C C}$ before it releases the reset, and the time-out delay can be disabled. Disabling the time-out delay without utilizing a BrownOut Detection circuit is not recommended.

The oscillator is required to oscillate for a minimum number of cycles before the clock is considered stable. An internal ripple counter monitors the oscillator output clock, and keeps the internal reset active for a given number of clock cycles. The reset is then released and the device will start to execute. The recommended oscillator start-up time is dependent on the clock type, and varies from 6 cycles for an externally applied clock to 32K cycles for a low frequency crystal.

The start-up sequence for the clock includes both the time-out delay and the start-up time when the device starts up from reset. When starting up from Power-save or Power-down mode, $\mathrm{V}_{\mathrm{CC}}$ is assumed to be at a sufficient level and only the start-up time is included.

### 10.4 Low Power Crystal Oscillator

Pins XTAL1 and XTAL2 are input and output, respectively, of an inverting amplifier which can be configured for use as an On-chip Oscillator, as shown in Figure 10-2. Either a quartz crystal or a ceramic resonator may be used.

This Crystal Oscillator is a low power oscillator, with reduced voltage swing on the XTAL2 output. It gives the lowest power consumption, but is not capable of driving other clock inputs, and may be more susceptible to noise in noisy environments. In these cases, refer to the "Full Swing Crystal Oscillator" on page 42.

C1 and C2 should always be equal for both crystals and resonators. The optimal value of the capacitors depends on the crystal or resonator in use, the amount of stray capacitance, and the electromagnetic noise of the environment. Some initial guidelines for choosing capacitors for use with crystals are given in Table 10-3 on page 42 . For ceramic resonators, the capacitor values given by the manufacturer should be used.

Figure 10-2. Crystal Oscillator Connections


XTAL2

XTAL1

GND

The Low Power Oscillator can operate in three different modes, each optimized for a specific frequency range. The operating mode is selected by the fuses CKSEL3:1 as shown in Table 10-3 on page 42.

Table 10-3. Low Power Crystal Oscillator Operating Modes ${ }^{(3)}$

| Frequency Range [MHz] | CKSEL3:1 $^{(1)}$ | Recommended Range for Capacitors C1 and C2 [pF] |
| :---: | :---: | :---: |
| $0.4-0.9$ | $100^{(2)}$ | - |
| $0.9-3.0$ | 101 | $12-22$ |
| $3.0-8.0$ | 110 | $12-22$ |
| $8.0-16.0^{(4)}$ | 111 | $12-22$ |

Notes: 1. This is the recommended CKSEL settings for the different frequency ranges.
2. This option should not be used with crystals, only with ceramic resonators.
3. If 8 MHz frequency exceeds the specification of the device (depends on $\mathrm{V}_{\mathrm{cc}}$ ), the CKDIV8 Fuse can be programmed in order to divide the internal frequency by 8 . It must be ensured that the resulting divided clock meets the frequency specification of the device.
4. Maximum frequency when using ceramic oscillator is 10 MHz .

The CKSEL0 Fuse together with the SUT1:0 Fuses select the start-up times as shown in Table 10-4.
Table 10-4. Start-up Times for the Low Power Crystal Oscillator Clock Selection

| Oscillator Source / Power Conditions | Start-up Time from Power- <br> down and Power-save | Additional Delay <br> from Reset <br> $\left.\mathbf{( V}_{\mathbf{C C}}=5.0 \mathrm{~V}\right)$ | CKSEL0 | SUT1:0 |
| :--- | :---: | :---: | :---: | :---: |
| Ceramic resonator, fast rising power | 258 CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}^{(1)}$ | 0 | 00 |
| Ceramic resonator, slowly rising power | 258 CK | $14 \mathrm{CK}+65 \mathrm{~ms}^{(1)}$ | 0 | 01 |
| Ceramic resonator, BOD enabled | 1 KCK | $14 \mathrm{CK}^{(2)}$ | 0 | 10 |
| Ceramic resonator, fast rising power | 1 KCK | $14 \mathrm{CK}+4.1 \mathrm{~ms}^{(2)}$ | 0 | 11 |
| Ceramic resonator, slowly rising power | 1 KCK | $14 \mathrm{CK}+65 \mathrm{~ms}^{(2)}$ | 1 | 00 |
| Crystal Oscillator, BOD enabled | 16 KCK | 14 CK | 1 | 01 |
| Crystal Oscillator, fast rising power | 16 KCK | $14 \mathrm{CK}+4.1 \mathrm{~ms}$ | 1 | 10 |
| Crystal Oscillator, slowly rising power | 16 KCK | $14 \mathrm{CK}+65 \mathrm{~ms}$ | 1 | 11 |

Notes: 1. These options should only be used when not operating close to the maximum frequency of the device, and only if frequency stability at start-up is not important for the application. These options are not suitable for crystals.
2. These options are intended for use with ceramic resonators and will ensure frequency stability at start-up. They can also be used with crystals when not operating close to the maximum frequency of the device, and if frequency stability at start-up is not important for the application.

### 10.5 Full Swing Crystal Oscillator

Pins XTAL1 and XTAL2 are input and output, respectively, of an inverting amplifier which can be configured for use as an On-chip Oscillator, as shown in Figure 10-2 on page 41. Either a quartz crystal or a ceramic resonator may be used.

This Crystal Oscillator is a full swing oscillator, with rail-to-rail swing on the XTAL2 output. This is useful for driving other clock inputs and in noisy environments. The current consumption is higher than the "Low Power Crystal Oscillator" on page 41. Note that the Full Swing Crystal Oscillator will only operate for $\mathrm{V}_{\mathrm{CC}}=2.7-5.5$ volts.
C1 and C2 should always be equal for both crystals and resonators. The optimal value of the capacitors depends on the crystal or resonator in use, the amount of stray capacitance, and the electromagnetic noise of the environment. Some initial guidelines for choosing capacitors for use with crystals are given in Table 10-6 on page 43. For ceramic resonators, the capacitor values given by the manufacturer should be used.

The operating mode is selected by the fuse CKSEL3:1 as shown in Table 10-5 on page 43.

Table 10-5. Full Swing Crystal Oscillator operating modes ${ }^{(1)}$

| Frequency Range [MHz] | CKSEL3:1 | Recommended Range for Capacitors C1 and C2 [pF] |
| :---: | :---: | :---: |
| $0.4-16$ | 011 | $12-22$ |

Note: 1. If 8 MHz frequency exceeds the specification of the device (depends on $\mathrm{V}_{\mathrm{CC}}$ ), the CKDIV8 Fuse can be programmed in order to divide the internal frequency by 8 . It must be ensured that the resulting divided clock meets the frequency specification of the device.

Table 10-6. Start-up Times for the Full Swing Crystal Oscillator Clock Selection

| Oscillator Source / Power Conditions | Start-up Time from Powerdown and Power-save | Additional Delay from Reset $\left(\mathrm{V}_{\mathrm{cc}}=5.0 \mathrm{~V}\right)$ | CKSELO | SUT1:0 |
| :---: | :---: | :---: | :---: | :---: |
| Ceramic resonator, fast rising power | 258 CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}^{(1)}$ | 0 | 00 |
| Ceramic resonator, slowly rising power | 258 CK | $14 \mathrm{CK}+65 \mathrm{~ms}^{(1)}$ | 0 | 01 |
| Ceramic resonator, BOD enabled | 1K CK | $14 \mathrm{CK}^{(2)}$ | 0 | 10 |
| Ceramic resonator, fast rising power | 1K CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}^{(2)}$ | 0 | 11 |
| Ceramic resonator, slowly rising power | 1K CK | $14 \mathrm{CK}+65 \mathrm{~ms}^{(2)}$ | 1 | 00 |
| Crystal Oscillator, BOD enabled | 16K CK | 14CK | 1 | 01 |
| Crystal Oscillator, fast rising power | 16K CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}$ | 1 | 10 |
| Crystal Oscillator, slowly rising power | 16K CK | 14CK + 65ms | 1 | 11 |

Notes: 1. These options should only be used when not operating close to the maximum frequency of the device, and only if frequency stability at start-up is not important for the application. These options are not suitable for crystals.
2. These options are intended for use with ceramic resonators and will ensure frequency stability at start-up. They can also be used with crystals when not operating close to the maximum frequency of the device, and if frequency stability at start-up is not important for the application.

### 10.6 Low Frequency Crystal Oscillator

The device can utilize a 32.768 kHz watch crystal as clock source by a dedicated Low Frequency Crystal Oscillator. The crystal should be connected as shown in Figure 10-3 on page 44. When this Oscillator is selected, start-up times are determined by the SUT Fuses and CKSELO as shown in Table 10-8 on page 44.
The Low-Frequency Crystal Oscillator provides an internal load capacitance, see Table 10-7 at each XTAL/TOSC pin.

Table 10-7. Capacitance for Low frequency oscillator

| Device | 32kHz oscillator | Cap (Xtal1/Tosc1) | Cap (Xtal2/Tosc2) |
| :--- | :--- | :--- | :--- |
| ATmega640/1280/1281/2560/2561 | System Osc. | 18 pF | $8 p F$ |
|  | Timer Osc. | $6 p F$ | $6 p F$ |

The capacitance $(\mathrm{Ce}+\mathrm{Ci})$ needed at each XTAL/TOSC pin can be calculated by using:

$$
C e+C i=2{ }^{*} C L-C_{s}
$$

where:
Ce - is optional external capacitors as described in Figure 10-3.
Ci - is the pin capacitance in Table 10-7 on page 43.
CL - is the load capacitance for a 32.768 kHz crystal specified by the crystal vendor.
$\mathrm{C}_{\mathrm{S}}$ - is the total stray capacitance for one XTAL/TOSC pin.
Crystals specifying load capacitance (CL) higher than the ones given in the Table 10-7 on page 43, require external capacitors applied as described in Figure 10-3.

Figure 10-3. Crystal Oscillator Connections


To find suitable load capacitance for a 32.768 kHz crystal, consult the crystal datasheet.
When this oscillator is selected, start-up times are determined by the SUT Fuses and CKSELO as shown in Table 10-8.

Table 10-8. Start-up times for the low frequency crystal oscillator clock selection

| Power Conditions | Start-up Time from Power-down and Power-save | Additional Delay from Reset $\left(\mathrm{V}_{\mathrm{cc}}=5.0 \mathrm{~V}\right)$ | CKSELO | SUT1:0 |
| :---: | :---: | :---: | :---: | :---: |
| BOD enabled | 1K CK | $14 \mathrm{CK}^{(1)}$ | 0 | 00 |
| Fast rising power | 1K CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}^{(1)}$ | 0 | 01 |
| Slowly rising power | 1K CK | $14 \mathrm{CK}+65 \mathrm{~ms}^{(1)}$ | 0 | 10 |
| Reserved |  |  | 0 | 11 |
| BOD enabled | 32 K CK | 14CK | 1 | 00 |
| Fast rising power | 32K CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}$ | 1 | 01 |
| Slowly rising power | 32K CK | 14CK + 65ms | 1 | 10 |
| Reserved |  |  | 1 | 11 |

Note: 1. These options should only be used if frequency stability at start-up is not important for the application.

### 10.7 Calibrated Internal RC Oscillator

By default, the Internal RC Oscillator provides an approximate 8 MHz clock. Though voltage and temperature dependent, this clock can be very accurately calibrated by the user. See Table 31-1 on page 359 and "Internal Oscillator Speed" on page 392 for more details. The device is shipped with the CKDIV8 Fuse programmed. See "System Clock Prescaler" on page 47 for more details.

This clock may be selected as the system clock by programming the CKSEL Fuses as shown in Table 10-9. If selected, it will operate with no external components. During reset, hardware loads the pre-programmed calibration value into the OSCCAL Register and thereby automatically calibrates the RC Oscillator. The accuracy of this calibration is shown as Factory calibration in Table 31-1 on page 359.

By changing the OSCCAL register from SW, see "OSCCAL - Oscillator Calibration Register" on page 48, it is possible to get a higher calibration accuracy than by using the factory calibration. The accuracy of this calibration is shown as User calibration in Table 31-1 on page 359.

When this Oscillator is used as the chip clock, the Watchdog Oscillator will still be used for the Watchdog Timer and for the Reset Time-out. For more information on the pre-programmed calibration value, see the section "Calibration Byte" on page 328.

Table 10-9. Internal Calibrated RC Oscillator Operating Modes ${ }^{(1)(2)}$

| Frequency Range [MHz] | CKSEL3:0 |
| :---: | :---: |
| $7.3-8.1$ | 0010 |

Notes: 1. The device is shipped with this option selected.
2. If 8 MHz frequency exceeds the specification of the device (depends on $\mathrm{V}_{\mathrm{CC}}$ ), the CKDIV8 Fuse can be programmed in order to divide the internal frequency by 8.
When this Oscillator is selected, start-up times are determined by the SUT Fuses as shown in Table 10-10.
Table 10-10. Start-up times for the internal calibrated RC Oscillator clock selection

| Power Conditions | Start-up Time from Power-down and <br> Power-save | Additional Delay from Reset <br> $\left(\mathbf{V}_{\mathbf{c c}}=\mathbf{5 . 0 V}\right)$ | SUT1:0 |
| :---: | :---: | :---: | :---: |
| BOD enabled | 6 CK | 14 CK | 00 |
| Fast rising power | 6 CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}$ | 01 |
| Slowly rising power | 6 CK | $14 \mathrm{CK}+65 \mathrm{~ms}^{(1)}$ | 10 |
| Reserved |  |  |  |

Note: 1. The device is shipped with this option selected.

### 10.8 128kHz Internal Oscillator

The 128 kHz internal Oscillator is a low power Oscillator providing a clock of 128 kHz . The frequency is nominal at 3 V and $25^{\circ} \mathrm{C}$. This clock may be select as the system clock by programming the CKSEL Fuses to " 11 " as shown in Table 10-11.

Table 10-11. 128 kHz Internal Oscillator Operating Modes ${ }^{(1)}$

| Nominal Frequency | CKSEL3:0 |
| :---: | :---: |
| 128 kHz | 0011 |

Note: 1. Note that the 128 kHz oscillator is a very low power clock source, and is not designed for high accuracy.
When this clock source is selected, start-up times are determined by the SUT Fuses as shown in Table 10-12 on page 46.

Table 10-12. Start-up Times for the 128 kHz Internal Oscillator

| Power Conditions | Start-up Time from Power-down and <br> Power-save | Additional Delay from Reset | SUT1:0 |
| :---: | :---: | :---: | :---: |
| BOD enabled | 6 CK | 14 CK | 00 |
| Fast rising power | 6 CK | $14 \mathrm{CK}+4 \mathrm{~ms}$ | 01 |
| Slowly rising power | 6 CK | $14 \mathrm{CK}+64 \mathrm{~ms}$ | 10 |
| Reserved |  |  |  |

### 10.9 External Clock

To drive the device from an external clock source, XTAL1 should be driven as shown in Figure 10-4. To run the device on an external clock, the CKSEL Fuses must be programmed to "0000".

Figure 10-4. External Clock Drive Configuration


When this clock source is selected, start-up times are determined by the SUT Fuses as shown in Table 10-15 on page 49.

Table 10-13. Crystal Oscillator Clock Frequency

| Nominal Frequency | CKSEL3:0 |
| :---: | :---: |
| $0-16 \mathrm{MHz}$ | 0000 |

Table 10-14. Start-up Times for the External Clock Selection

| Power Conditions | Start-up Time from Power-down and <br> Power-save | Additional Delay from Reset <br> $\left(\mathbf{V}_{\mathbf{c c}}=5.0 \mathrm{~V}\right)$ | SUT1:0 |
| :---: | :---: | :---: | :---: |
| BOD enabled | 6 CK | 14 CK | 00 |
| Fast rising power | 6 CK | $14 \mathrm{CK}+4.1 \mathrm{~ms}$ | 01 |
| Slowly rising power | 6 CK | $14 \mathrm{CK}+65 \mathrm{~ms}$ | 10 |
| Reserved |  |  |  |

When applying an external clock, it is required to avoid sudden changes in the applied clock frequency to ensure stable operation of the MCU. A variation in frequency of more than $2 \%$ from one clock cycle to the next can lead to unpredictable behavior. If changes of more than $2 \%$ is required, ensure that the MCU is kept in Reset during the changes.

Note that the System Clock Prescaler can be used to implement run-time changes of the internal clock frequency while still ensuring stable operation. Refer to "System Clock Prescaler" on page 47 for details.

### 10.10 Clock Output Buffer

The device can output the system clock on the CLKO pin. To enable the output, the CKOUT Fuse has to be programmed. This mode is suitable when the chip clock is used to drive other circuits on the system. The clock also will be output during reset, and the normal operation of I/O pin will be overridden when the fuse is programmed. Any clock source, including the internal RC Oscillator, can be selected when the clock is output on CLKO. If the System Clock Prescaler is used, it is the divided system clock that is output.

### 10.11 Timer/Counter Oscillator

The device can operate its Timer/Counter2 from an external 32.768 kHz watch crystal or a external clock source. See Figure 10-2 on page 41 for crystal connection.
Applying an external clock source to TOSC1 requires EXCLK in the ASSR Register written to logic one. See "Asynchronous Operation of Timer/Counter2" on page 179 for further description on selecting external clock as input instead of a 32 kHz crystal.

### 10.12 System Clock Prescaler

The ATmega640/1280/1281/2560/2561 has a system clock prescaler, and the system clock can be divided by setting the "CLKPR - Clock Prescale Register" on page 48. This feature can be used to decrease the system clock frequency and the power consumption when the requirement for processing power is low. This can be used with all clock source options, and it will affect the clock frequency of the CPU and all synchronous peripherals. $\mathrm{clk}_{1 / 0}$, $\mathrm{Clk}_{\text {ADC }}, \mathrm{clk}_{\text {CPU }}$, and $\mathrm{clk}_{\text {FLASH }}$ are divided by a factor as shown in Table 10-15 on page 49.

When switching between prescaler settings, the System Clock Prescaler ensures that no glitches occurs in the clock system. It also ensures that no intermediate frequency is higher than neither the clock frequency corresponding to the previous setting, nor the clock frequency corresponding to the new setting.
The ripple counter that implements the prescaler runs at the frequency of the undivided clock, which may be faster than the CPU's clock frequency. Hence, it is not possible to determine the state of the prescaler - even if it were readable, and the exact time it takes to switch from one clock division to the other cannot be exactly predicted. From the time the CLKPS values are written, it takes between T1 + T2 and T1 $+2 \times$ T2 before the new clock frequency is active. In this interval, 2 active clock edges are produced. Here, T1 is the previous clock period, and T2 is the period corresponding to the new prescaler setting.

To avoid unintentional changes of clock frequency, a special write procedure must be followed to change the CLKPS bits:

Write the Clock Prescaler Change Enable (CLKPCE) bit to one and all other bits in CLKPR to zero.
Within four cycles, write the desired value to CLKPS while writing a zero to CLKPCE.
Interrupts must be disabled when changing prescaler setting to make sure the write procedure is not interrupted.

### 10.13 Register Description

### 10.13.1 OSCCAL - Oscillator Calibration Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | OSCCAL |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x66) | CAL7 | CAL6 | CAL5 | CAL4 | CAL3 | CAL2 | CAL1 | CALO |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value |  |  |  | Spec | ration V |  |  |  |  |

## - Bits 7:0 - CAL7:0: Oscillator Calibration Value

The Oscillator Calibration Register is used to trim the Calibrated Internal RC Oscillator to remove process variations from the oscillator frequency. A pre-programmed calibration value is automatically written to this register during chip reset, giving the Factory calibrated frequency as specified in Table 31-1 on page 359. The application software can write this register to change the oscillator frequency. The oscillator can be calibrated to frequencies as specified in Table 31-1 on page 359. Calibration outside that range is not guaranteed.

Note that this oscillator is used to time EEPROM and Flash write accesses, and these write times will be affected accordingly. If the EEPROM or Flash are written, do not calibrate to more than 8.8 MHz . Otherwise, the EEPROM or Flash write may fail.

The CAL7 bit determines the range of operation for the oscillator. Setting this bit to 0 gives the lowest frequency range, setting this bit to 1 gives the highest frequency range. The two frequency ranges are overlapping, in other words a setting of OSCCAL $=0 \times 7 F$ gives a higher frequency than OSCCAL $=0 \times 80$.

The CAL6.. 0 bits are used to tune the frequency within the selected range. A setting of $0 \times 00$ gives the lowest frequency in that range, and a setting of 0x7F gives the highest frequency in the range.

### 10.13.2 CLKPR - Clock Prescale Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | CLKPR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x61) | CLKPCE | - | - | - | CLKPS3 | CLKPS2 | CLKPS1 | CLKPSO |  |
| Read/Write | R/W | R | R | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 |  | See Bit | cription |  |  |

## - Bit 7 - CLKPCE: Clock Prescaler Change Enable

The CLKPCE bit must be written to logic one to enable change of the CLKPS bits. The CLKPCE bit is only updated when the other bits in CLKPR are simultaneously written to zero. CLKPCE is cleared by hardware four cycles after it is written or when CLKPS bits are written. Rewriting the CLKPCE bit within this time-out period does neither extend the time-out period, nor clear the CLKPCE bit.

- Bits 3:0 - CLKPS3:0: Clock Prescaler Select Bits 3-0

These bits define the division factor between the selected clock source and the internal system clock. These bits can be written run-time to vary the clock frequency to suit the application requirements. As the divider divides the master clock input to the MCU, the speed of all synchronous peripherals is reduced when a division factor is used. The division factors are given in Table 10-15 on page 49.

The CKDIV8 Fuse determines the initial value of the CLKPS bits. If CKDIV8 is unprogrammed, the CLKPS bits will be reset to " 0000 ". If CKDIV8 is programmed, CLKPS bits are reset to "0011", giving a division factor of 8 at start up. This feature should be used if the selected clock source has a higher frequency than the maximum frequency of the device at the present operating conditions. Note that any value can be written to the CLKPS bits regardless of the CKDIV8 Fuse setting. The Application software must ensure that a sufficient division factor is chosen if the selected clock source has a higher frequency than the maximum frequency of the device at the present operating conditions. The device is shipped with the CKDIV8 Fuse programmed.

Table 10-15. Clock Prescaler Select

| CLKPS3 | CLKPS2 | CLKPS1 | CLKPS0 | Clock Division Factor |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 1 | 2 |
| 0 | 0 | 1 | 0 | 4 |
| 0 | 0 | 1 | 1 | 8 |
| 0 | 1 | 0 | 0 | 16 |
| 0 | 1 | 0 | 1 | 32 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 0 | 0 | 0 | 128 |
| 1 | 0 | 1 | 0 | Reserved |
| 1 | 1 | 0 | 1 | Reserved |
| 1 | 1 | 0 | 0 | Reserved |
| 1 | 1 | 1 | 1 | Reserved |
| 1 | 1 | 1 | 0 | Reserved |
| 1 | 0 | 1 | Reserved |  |
| 1 | 0 | 1 | Reserved |  |

## 11. Power Management and Sleep Modes

Sleep modes enable the application to shut down unused modules in the MCU, thereby saving power. The AVR provides various sleep modes allowing the user to tailor the power consumption to the application's requirements.

### 11.1 Sleep Modes

Figure 10-1 on page 39 presents the different clock systems in the ATmega640/1280/1281/2560/2561, and their distribution. The figure is helpful in selecting an appropriate sleep mode. Table 11-1 shows the different sleep modes and their wake-up sources.

Table 11-1. Active Clock Domains and Wake-up Sources in the Different Sleep Modes.

|  | Active Clock Domains |  |  |  |  | Oscillators |  | Wake-up Sources |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Sleep Mode | $\begin{aligned} & \overrightarrow{0} 0 \\ & \underline{0} \\ & \hline 0 \end{aligned}$ |  | $\frac{0}{\grave{0}}$ | $\begin{gathered} 0 \\ \frac{0}{4} \\ \frac{x^{2}}{} \end{gathered}$ | $\begin{gathered} x_{n}^{n} \\ \frac{x^{2}}{0} \end{gathered}$ |  |  |  |  | $\begin{aligned} & \text { N } \\ & \text { í } \\ & \text { in } \end{aligned}$ |  | U |  | O <br> ¢ <br> ¢ <br> 0 |
| Idle |  |  | X | X | X | X | $\mathrm{X}^{(2)}$ | X | X | X | X | X | X | X |
| ADCNRM |  |  |  | X | X | X | $\mathrm{X}^{(2)}$ | $\mathrm{X}^{(3)}$ | X | $\mathrm{X}^{(2)}$ | X | X | X |  |
| Power-down |  |  |  |  |  |  |  | $X^{(3)}$ | X |  |  |  | X |  |
| Power-save |  |  |  |  | X |  | $\mathrm{X}^{(2)}$ | $\mathrm{X}^{(3)}$ | X | X |  |  | X |  |
| Standby ${ }^{(1)}$ |  |  |  |  |  | X |  | $X^{(3)}$ | X |  |  |  | X |  |
| Extended Standby |  |  |  |  | $X^{(2)}$ | X | $\mathrm{X}^{(2)}$ | $\mathrm{X}^{(3)}$ | X | X |  |  | X |  |

Note: 1. Only recommended with external crystal or resonator selected as clock source.
2. If Timer/Counter2 is running in asynchronous mode.
3. For INT7:4, only level interrupt.

To enter any of the sleep modes, the SE bit in "SMCR - Sleep Mode Control Register" on page 54 must be written to logic one and a SLEEP instruction must be executed. The SM2, SM1, and SM0 bits in the SMCR Register select which sleep mode will be activated by the SLEEP instruction. See Table 11-2 on page 54 for a summary.

If an enabled interrupt occurs while the MCU is in a sleep mode, the MCU wakes up. The MCU is then halted for four cycles in addition to the start-up time, executes the interrupt routine, and resumes execution from the instruction following SLEEP. The contents of the Register File and SRAM are unaltered when the device wakes up from sleep. If a reset occurs during sleep mode, the MCU wakes up and executes from the Reset Vector.

### 11.2 Idle Mode

When the SM2:0 bits are written to 000, the SLEEP instruction makes the MCU enter Idle mode, stopping the CPU but allowing the SPI, USART, Analog Comparator, ADC, 2-wire Serial Interface, Timer/Counters, Watchdog, and the interrupt system to continue operating. This sleep mode basically halts $\mathrm{clk}_{\mathrm{CPU}}$ and $\mathrm{clk}_{\text {FLASH }}$, while allowing the other clocks to run.

Idle mode enables the MCU to wake up from external triggered interrupts as well as internal ones like the Timer Overflow and USART Transmit Complete interrupts. If wake-up from the Analog Comparator interrupt is not required, the Analog Comparator can be powered down by setting the ACD bit in the Analog Comparator Control and Status Register - ACSR. This will reduce power consumption in Idle mode. If the ADC is enabled, a conversion starts automatically when this mode is entered.

### 11.3 ADC Noise Reduction Mode

When the SM2:0 bits are written to 001, the SLEEP instruction makes the MCU enter ADC Noise Reduction mode, stopping the CPU but allowing the ADC, the external interrupts, 2-wire Serial Interface address match, Timer/Counter2 and the Watchdog to continue operating (if enabled). This sleep mode basically halts clkl/O, clkCPU, and clkFLASH, while allowing the other clocks to run.
This improves the noise environment for the ADC, enabling higher resolution measurements. If the ADC is enabled, a conversion starts automatically when this mode is entered. Apart form the ADC Conversion Complete interrupt, only an External Reset, a Watchdog System Reset, a Watchdog interrupt, a Brown-out Reset, a 2-wire serial interface interrupt, a Timer/Counter2 interrupt, an SPM/EEPROM ready interrupt, an external level interrupt on INT7:4 or a pin change interrupt can wakeup the MCU from ADC Noise Reduction mode.

### 11.4 Power-down Mode

When the SM2:0 bits are written to 010, the SLEEP instruction makes the MCU enter Power-down mode. In this mode, the external Oscillator is stopped, while the external interrupts, the 2-wire Serial Interface, and the Watchdog continue operating (if enabled). Only an External Reset, a Watchdog Reset, a Brown-out Reset, 2-wire Serial Interface address match, an external level interrupt on INT7:4, an external interrupt on INT3:0, or a pin change interrupt can wake up the MCU. This sleep mode basically halts all generated clocks, allowing operation of asynchronous modules only.
Note that if a level triggered interrupt is used for wake-up from Power-down mode, the changed level must be held for some time to wake up the MCU. Refer to "External Interrupts" on page 109 for details.
When waking up from Power-down mode, there is a delay from the wake-up condition occurs until the wake-up becomes effective. This allows the clock to restart and become stable after having been stopped. The wake-up period is defined by the same CKSEL Fuses that define the Reset Time-out period, as described in "Clock Sources" on page 40.

### 11.5 Power-save Mode

When the SM2:0 bits are written to 011, the SLEEP instruction makes the MCU enter Power-save mode. This mode is identical to Power-down, with one exception:
If Timer/Counter2 is enabled, it will keep running during sleep. The device can wake up from either Timer Overflow or Output Compare event from Timer/Counter2 if the corresponding Timer/Counter2 interrupt enable bits are set in TIMSK2, and the Global Interrupt Enable bit in SREG is set. If Timer/Counter2 is not running, Power-down mode is recommended instead of Power-save mode.

The Timer/Counter2 can be clocked both synchronously and asynchronously in Power-save mode. If the Timer/Counter2 is not using the asynchronous clock, the Timer/Counter Oscillator is stopped during sleep. If the Timer/Counter2 is not using the synchronous clock, the clock source is stopped during sleep. Note that even if the synchronous clock is running in Power-save, this clock is only available for the Timer/Counter2.

### 11.6 Standby Mode

When the SM2:0 bits are 110 and an external crystal/resonator clock option is selected, the SLEEP instruction makes the MCU enter Standby mode. This mode is identical to Power-down with the exception that the Oscillator is kept running. From Standby mode, the device wakes up in six clock cycles.

### 11.7 Extended Standby Mode

When the SM2:0 bits are 111 and an external crystal/resonator clock option is selected, the SLEEP instruction makes the MCU enter Extended Standby mode. This mode is identical to Power-save mode with the exception that the Oscillator is kept running. From Extended Standby mode, the device wakes up in six clock cycles.

### 11.8 Power Reduction Register

The Power Reduction Register (PRR), see "PRR0 - Power Reduction Register 0" on page 55 and "PRR1 - Power Reduction Register 1" on page 56, provides a method for stopping the clock to individual peripherals to reduce power consumption.
Note that when the clock for a peripheral is stopped, then:

- The current state of the peripheral is frozen
- The associated registers can not be read or written
- Resources used by the peripherals (for example I/O pin, etc.) will remain occupied

The peripheral should in most cases be disabled before stopping the clock. Waking up a module, which is done by cleaning the bit in PRR, puts the module in the same state as before shutdown. Module shutdown can be used in Idle mode or Active mode to significantly reduce the overall power consumption. See "Power-down Supply Current" on page 380 for examples. In all other sleep modes, the clock is already stopped.

### 11.9 Minimizing Power Consumption

There are several issues to consider when trying to minimize the power consumption in an AVR controlled system. In general, sleep modes should be used as much as possible, and the sleep mode should be selected so that as few as possible of the device's functions are operating. All functions not needed should be disabled. In particular, the following modules may need special consideration when trying to achieve the lowest possible power consumption.

### 11.9.1 Analog to Digital Converter

If enabled, the ADC will be enabled in all sleep modes. To save power, the ADC should be disabled before entering any sleep mode. When the ADC is turned off and on again, the next conversion will be an extended conversion. Refer to "ADC - Analog to Digital Converter" on page 268 for details on ADC operation.

### 11.9.2 Analog Comparator

When entering Idle mode, the Analog Comparator should be disabled if not used. When entering ADC Noise Reduction mode, the Analog Comparator should be disabled. In other sleep modes, the Analog Comparator is automatically disabled. However, if the Analog Comparator is set up to use the Internal Voltage Reference as input, the Analog Comparator should be disabled in all sleep modes. Otherwise, the Internal Voltage Reference will be enabled, independent of sleep mode. Refer to "AC - Analog Comparator" on page 265 for details on how to configure the Analog Comparator.

### 11.9.3 Brown-out Detector

If the Brown-out Detector is not needed by the application, this module should be turned off. If the Brown-out Detector is enabled by the BODLEVEL Fuses, it will be enabled in all sleep modes, and hence, always consume power. In the deeper sleep modes, this will contribute significantly to the total current consumption. Refer to "Brown-out Detection" on page 59 for details on how to configure the Brown-out Detector.

### 11.9.4 Internal Voltage Reference

The Internal Voltage Reference will be enabled when needed by the Brown-out Detection, the Analog Comparator or the ADC. If these modules are disabled as described in the sections above, the internal voltage reference will be disabled and it will not be consuming power. When turned on again, the user must allow the reference to start up before the output is used. If the reference is kept on in sleep mode, the output can be used immediately. Refer to "Internal Voltage Reference" on page 60 for details on the start-up time.

### 11.9.5 Watchdog Timer

If the Watchdog Timer is not needed in the application, the module should be turned off. If the Watchdog Timer is enabled, it will be enabled in all sleep modes, and hence, always consume power. In the deeper sleep modes, this will contribute significantly to the total current consumption. Refer to "Interrupts" on page 101 for details on how to configure the Watchdog Timer.

### 11.9.6 Port Pins

When entering a sleep mode, all port pins should be configured to use minimum power. The most important is then to ensure that no pins drive resistive loads. In sleep modes where both the I/O clock ( $\mathrm{Clk}_{1 / 0}$ ) and the ADC clock ( $\mathrm{Cl}_{\mathrm{ADC}}$ ) are stopped, the input buffers of the device will be disabled. This ensures that no power is consumed by the input logic when not needed. In some cases, the input logic is needed for detecting wake-up conditions, and it will then be enabled. Refer to the section "Digital Input Enable and Sleep Modes" on page 71 for details on which pins are enabled. If the input buffer is enabled and the input signal is left floating or have an analog signal level close to $\mathrm{V}_{\mathrm{CC}} / 2$, the input buffer will use excessive power.

For analog input pins, the digital input buffer should be disabled at all times. An analog signal level close to $\mathrm{V}_{\mathrm{CC}} / 2$ on an input pin can cause significant current even in active mode. Digital input buffers can be disabled by writing to the Digital Input Disable Registers (DIDR2, DIDR1 and DIDR0). Refer to "DIDR2 - Digital Input Disable Register 2" on page 288, "DIDR1 - Digital Input Disable Register 1" on page 267, and "DIDR0 - Digital Input Disable Register 0 " on page 287 for details.

### 11.9.7 On-chip Debug System

If the On-chip debug system is enabled by the OCDEN Fuse and the chip enters sleep mode, the main clock source is enabled, and hence, always consumes power. In the deeper sleep modes, this will contribute significantly to the total current consumption.

There are three alternative ways to disable the OCD system:

- Disable the OCDEN Fuse
- Disable the JTAGEN Fuse
- Write one to the JTD bit in MCUCR


### 11.10 Register Description

### 11.10.1 SMCR - Sleep Mode Control Register

The Sleep Mode Control Register contains control bits for power management.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | SMCR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x33 (0x53) | - | - | - | - | SM2 | SM1 | SM0 | SE |  |
| Read/Write | R | R | R | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

- Bits 3, 2, 1 - SM2:0: Sleep Mode Select Bits 2, 1, and 0

These bits select between the five available sleep modes as shown in Table 11-2.
Table 11-2. Sleep Mode Select

| SM2 | SM1 | SM0 | Sleep Mode |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | Idle |
| 0 | 0 | 1 | ADC Noise Reduction |
| 0 | 1 | 0 | Power-down |
| 0 | 1 | 1 | Power-save |
| 1 | 0 | 0 | Reserved |
| 1 | 0 | 1 | Reserved |
| 1 | 1 | 0 | Standby ${ }^{(1)}$ |
| 1 | 1 | Extended Standby ${ }^{(1)}$ |  |

Note: 1. Standby modes are only recommended for use with external crystals or resonators.

## - Bit 1 - SE: Sleep Enable

The SE bit must be written to logic one to make the MCU enter the sleep mode when the SLEEP instruction is executed. To avoid the MCU entering the sleep mode unless it is the programmer's purpose, it is recommended to write the Sleep Enable (SE) bit to one just before the execution of the SLEEP instruction and to clear it immediately after waking up.

### 11.10.2 PRRO - Power Reduction Register 0

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PRR0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x64) | PRTWI | PRTIM2 | PRTIMO | - | PRTIM1 | PRSPI | PRUSART0 | PRADC |  |
| Read/Write | R/W | R/W | R/W | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

- Bit 7 - PRTWI: Power Reduction TWI

Writing a logic one to this bit shuts down the TWI by stopping the clock to the module. When waking up the TWI again, the TWI should be re initialized to ensure proper operation.

## - Bit 6 - PRTIM2: Power Reduction Timer/Counter2

Writing a logic one to this bit shuts down the Timer/Counter2 module in synchronous mode (AS2 is 0). When the Timer/Counter2 is enabled, operation will continue like before the shutdown.

## - Bit 5 - PRTIMO: Power Reduction Timer/Counter0

Writing a logic one to this bit shuts down the Timer/Counter0 module. When the Timer/Counter0 is enabled, operation will continue like before the shutdown.

- Bit 4-Res: Reserved bit

This bit is reserved bit and will always read as zero.

- Bit 3 - PRTIM1: Power Reduction Timer/Counter1

Writing a logic one to this bit shuts down the Timer/Counter1 module. When the Timer/Counter1 is enabled, operation will continue like before the shutdown.

## - Bit 2 - PRSPI: Power Reduction Serial Peripheral Interface

Writing a logic one to this bit shuts down the Serial Peripheral Interface by stopping the clock to the module. When waking up the SPI again, the SPI should be re initialized to ensure proper operation.

## - Bit 1 - PRUSARTO: Power Reduction USARTO

Writing a logic one to this bit shuts down the USARTO by stopping the clock to the module. When waking up the USARTO again, the USARTO should be re initialized to ensure proper operation.

- Bit 0 - PRADC: Power Reduction ADC

Writing a logic one to this bit shuts down the ADC. The ADC must be disabled before shut down. The analog comparator cannot use the ADC input MUX when the ADC is shut down.

### 11.10.3 PRR1 - Power Reduction Register 1

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PRR1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x65) | - | - | PRTIM5 | PRTIM4 | PRTIM3 | PRUSART3 | PRUSART2 | PRUSART1 |  |
| Read/Write | R | R | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

- Bit 7:6-Res: Reserved bits

These bits are reserved and will always read as zero.

- Bit 5 - PRTIM5: Power Reduction Timer/Counter5

Writing a logic one to this bit shuts down the Timer/Counter5 module. When the Timer/Counter5 is enabled, operation will continue like before the shutdown.

## - Bit 4 - PRTIM4: Power Reduction Timer/Counter4

Writing a logic one to this bit shuts down the Timer/Counter4 module. When the Timer/Counter4 is enabled, operation will continue like before the shutdown.

- Bit 3 - PRTIM3: Power Reduction Timer/Counter3

Writing a logic one to this bit shuts down the Timer/Counter3 module. When the Timer/Counter3 is enabled, operation will continue like before the shutdown.

- Bit 2 - PRUSART3: Power Reduction USART3

Writing a logic one to this bit shuts down the USART3 by stopping the clock to the module. When waking up the USART3 again, the USART3 should be re initialized to ensure proper operation.

## - Bit 1 - PRUSART2: Power Reduction USART2

Writing a logic one to this bit shuts down the USART2 by stopping the clock to the module. When waking up the USART2 again, the USART2 should be re initialized to ensure proper operation.

- Bit 0 - PRUSART1: Power Reduction USART1

Writing a logic one to this bit shuts down the USART1 by stopping the clock to the module. When waking up the USART1 again, the USART1 should be re initialized to ensure proper operation.

## 12. System Control and Reset

### 12.1 Resetting the AVR

During reset, all I/O Registers are set to their initial values, and the program starts execution from the Reset Vector. The instruction placed at the Reset Vector must be a JMP - Absolute Jump - instruction to the reset handling routine. If the program never enables an interrupt source, the Interrupt Vectors are not used, and regular program code can be placed at these locations. This is also the case if the Reset Vector is in the Application section while the Interrupt Vectors are in the Boot section or vice versa. The circuit diagram in Figure 12-1 on page 58 shows the reset logic. "System and Reset Characteristics" on page 360 defines the electrical parameters of the reset circuitry.

The I/O ports of the AVR are immediately reset to their initial state when a reset source goes active. This does not require any clock source to be running.
After all reset sources have gone inactive, a delay counter is invoked, stretching the internal reset. This allows the power to reach a stable level before normal operation starts. The time-out period of the delay counter is defined by the user through the SUT and CKSEL Fuses. The different selections for the delay period are presented in "Clock Sources" on page 40.

### 12.2 Reset Sources

The ATmega640/1280/1281/2560/2561 has five sources of reset:

- Power-on Reset. The MCU is reset when the supply voltage is below the Power-on Reset threshold ( $\mathrm{V}_{\text {POT }}$ )
- External Reset. The MCU is reset when a low level is present on the RESET pin for longer than the minimum pulse length
- Watchdog Reset. The MCU is reset when the Watchdog Timer period expires and the Watchdog is enabled
- Brown-out Reset. The MCU is reset when the supply voltage $\mathrm{AV}_{\mathrm{CC}}$ is below the Brown-out Reset threshold $\left(\mathrm{V}_{\mathrm{BOT}}\right)$ and the Brown-out Detector is enabled
- JTAG AVR Reset. The MCU is reset as long as there is a logic one in the Reset Register, one of the scan chains of the JTAG system. Refer to the section "IEEE 1149.1 (JTAG) Boundary-scan" on page 295 for details

Figure 12-1. Reset Logic


### 12.2.1 Power-on Reset

A Power-on Reset (POR) pulse is generated by an On-chip detection circuit. The detection level is defined in "System and Reset Characteristics" on page 360 . The POR is activated whenever $\mathrm{V}_{\mathrm{CC}}$ is below the detection level. The POR circuit can be used to trigger the start-up Reset, as well as to detect a failure in supply voltage.
A Power-on Reset (POR) circuit ensures that the device is reset from Power-on. Reaching the Power-on Reset threshold voltage invokes the delay counter, which determines how long the device is kept in RESET after $\mathrm{V}_{\mathrm{CC}}$ rise. The RESET signal is activated again, without any delay, when $\mathrm{V}_{\mathrm{CC}}$ decreases below the detection level.

Figure 12-2. MCU Start-up, $\overline{\operatorname{RESET}}$ Tied to $\mathrm{V}_{\mathrm{CC}}$


Figure 12-3. MCU Start-up, RESET Extended Externally


### 12.2.2 External Reset

An External Reset is generated by a low level on the $\overline{\text { RESET }}$ pin. Reset pulses longer than the minimum pulse width (see "System and Reset Characteristics" on page 360) will generate a reset, even if the clock is not running. Shorter pulses are not guaranteed to generate a reset. When the applied signal reaches the Reset Threshold Voltage $-\mathrm{V}_{\text {RST }}$ - on its positive edge, the delay counter starts the MCU after the Time-out period $-\mathrm{t}_{\text {TOUT }}$ - has expired.

Figure 12-4. External Reset During Operation


### 12.2.3 Brown-out Detection

ATmega640/1280/1281/2560/2561 has an On-chip Brown-out Detection (BOD) circuit for monitoring the $\mathrm{AV}_{\text {CC }}$ level during operation by comparing it to a fixed trigger level. The trigger level for the BOD can be selected by the BODLEVEL Fuses. The trigger level has a hysteresis to ensure spike free Brown-out Detection. The hysteresis on the detection level should be interpreted as $\mathrm{V}_{\mathrm{BOT}+}=\mathrm{V}_{\mathrm{BOT}}+\mathrm{V}_{\mathrm{HYST}} / 2$ and $\mathrm{V}_{\mathrm{BOT}-}=\mathrm{V}_{\mathrm{BOT}}-\mathrm{V}_{\mathrm{HYST}} / 2$.

When the BOD is enabled, and $A V_{C C}$ decreases to a value below the trigger level ( $\mathrm{V}_{\text {BOT- }}$ in Figure 12-5 on page 60 ), the Brown-out Reset is immediately activated. When $\mathrm{AV}_{\mathrm{CC}}$ increases above the trigger level ( $\mathrm{V}_{\mathrm{BOT}}$ in Figure 12-5 on page 60), the delay counter starts the MCU after the Time-out period $\mathrm{t}_{\text {Tout }}$ has expired.

The BOD circuit will only detect a drop in $\mathrm{AV}_{\mathrm{CC}}$ if the voltage stays below the trigger level for longer than $\mathrm{t}_{\mathrm{BOD}}$ given in "System and Reset Characteristics" on page 360.

Figure 12-5. Brown-out Reset During Operation


### 12.2.4 Watchdog Reset

When the Watchdog times out, it will generate a short reset pulse of one CK cycle duration. On the falling edge of this pulse, the delay timer starts counting the Time-out period $\mathrm{t}_{\text {TOUT }}$. See "Watchdog Timer" on page 53. for details on operation of the Watchdog Timer.

Figure 12-6. Watchdog Reset During Operation


### 12.3 Internal Voltage Reference

ATmega640/1280/1281/2560/2561 features an internal bandgap reference. This reference is used for Brown-out Detection, and it can be used as an input to the Analog Comparator or the ADC.

### 12.3.1 Voltage Reference Enable Signals and Start-up Time

The voltage reference has a start-up time that may influence the way it should be used. The start-up time is given in "System and Reset Characteristics" on page 360. To save power, the reference is not always turned on. The reference is on during the following situations:

1. When the BOD is enabled (by programming the BODLEVEL [2:0] Fuse).
2. When the bandgap reference is connected to the Analog Comparator (by setting the ACBG bit in ACSR).
3. When the ADC is enabled.

Thus, when the BOD is not enabled, after setting the ACBG bit or enabling the ADC, the user must always allow the reference to start up before the output from the Analog Comparator or ADC is used. To reduce power consumption in Power-down mode, the user can avoid the three conditions above to ensure that the reference is turned off before entering Power-down mode.

### 12.4 Watchdog Timer

### 12.4.1 Features

- Clocked from separate On-chip Oscillator
- Three Operating modes
- Interrupt
- System Reset
- Interrupt and System Reset
- Selectable Time-out period from 16 ms to 8 s
- Possible Hardware fuse Watchdog always on (WDTON) for fail-safe mode

Figure 12-7. Watchdog Timer

### 12.4.2 Overview

ATmega640/1280/1281/2560/2561 has an Enhanced Watchdog Timer (WDT). The WDT is a timer counting cycles of a separate on-chip 128 kHz oscillator. The WDT gives an interrupt or a system reset when the counter reaches a given time-out value. In normal operation mode, it is required that the system uses the WDR - Watchdog Timer Reset - instruction to restart the counter before the time-out value is reached. If the system doesn't restart the counter, an interrupt or system reset will be issued.
In Interrupt mode, the WDT gives an interrupt when the timer expires. This interrupt can be used to wake the device from sleep-modes, and also as a general system timer. One example is to limit the maximum time allowed for certain operations, giving an interrupt when the operation has run longer than expected. In System Reset mode, the WDT gives a reset when the timer expires. This is typically used to prevent system hang-up in case of runaway code. The third mode, Interrupt and System Reset mode, combines the other two modes by first giving an interrupt and then switch to System Reset mode. This mode will for instance allow a safe shutdown by saving critical parameters before a system reset.

The Watchdog always on (WDTON) fuse, if programmed, will force the Watchdog Timer to System Reset mode. With the fuse programmed the System Reset mode bit (WDE) and Interrupt mode bit (WDIE) are locked to 1 and 0 respectively. To further ensure program security, alterations to the Watchdog set-up must follow timed sequences. The sequence for clearing WDE and changing time-out configuration is as follows:

1. In the same operation, write a logic one to the Watchdog change enable bit (WDCE) and WDE. A logic one must be written to WDE regardless of the previous value of the WDE bit.
2. Within the next four clock cycles, write the WDE and Watchdog prescaler bits (WDP) as desired, but with the WDCE bit cleared. This must be done in one operation.
The following code example shows one assembly and one C function for turning off the Watchdog Timer. The example assumes that interrupts are controlled (for example by disabling interrupts globally) so that no interrupts will occur during the execution of these functions.
```
Assembly Code Example \({ }^{(1)}\)
    WDT_off:
        ; Turn off global interrupt
    cli
    ; Reset Watchdog Timer
    wdr
    ; Clear WDRF in MCUSR
    in r16, MCUSR
    andi r16, (0xff \& ( \(0 \ll\) WDRF) )
    out MCUSR, r16
    ; Write logical one to WDCE and WDE
    ; Keep old prescaler setting to prevent unintentional time-out
    1di r16, WDTCSR
    ori r16, ( \(1 \ll\) WDCE) | ( \(1 \ll\) WDE)
    sts WDTCSR, r16
    ; Turn off WDT
    ldi r16, ( \(0 \ll\) WDE)
    sts WDTCSR, r16
    ; Turn on global interrupt
    sei
    ret
C Code Example \({ }^{(1)}\)
    void WDT_off(void)
    \{
    __disable_interrupt();
    __watchdog_reset();
    /* Clear WDRF in MCUSR */
    MCUSR \(\&=\sim(1 \ll W D R F) ;\)
    /* Write logical one to WDCE and WDE */
    /* Keep old prescaler setting to prevent unintentional time-out
    */
    WDTCSR | \(=(1 \ll\) WDCE \() \mid \quad(1 \ll\) WDE \() ;\)
    /* Turn off WDT */
    WDTCSR \(=0 x 00\);
    __enable_interrupt();
    \}
```

Note: 1. The example code assumes that the part specific header file is included.
2. If the Watchdog is accidentally enabled, for example by a runaway pointer or brown-out condition, the device will be reset and the Watchdog Timer will stay enabled. If the code is not set up to handle the Watchdog, this might lead to an eternal loop of time-out resets. To avoid this situation, the application software should always clear the Watchdog System Reset Flag (WDRF) and the WDE control bit in the initialisation routine, even if the Watchdog is not in use.
The following code example shows one assembly and one C function for changing the time-out value of the Watchdog Timer.

```
Assembly Code Example }\mp@subsup{}{}{(1)
    WDT_Prescaler_Change:
        ; Turn off global interrupt
        cli
        ; Reset Watchdog Timer
        wdr
        ; Start timed sequence
        in r16, WDTCSR
        ori r16, (1<<WDCE) | (1<<WDE)
    out WDTCSR, r16
    ; -- Got four cycles to set the new values from here -
    ; Set new prescaler(time-out) value = 64K cycles (~0.5 s)
    ldi r16, (1<<WDE) | (1<<WDP2) | (1<<WDP0)
    out WDTCSR, r16
    ; -- Finished setting new values, used 2 cycles
    ; Turn on global interrupt
    sei
    ret
C Code Example (2)
    void WDT_Prescaler_Change(void)
    {
        __disable_interrupt();
        __watchdog_reset();
    /* Start timed equence */
    WDTCSR | = (1<<WDCE) | ( }1<<\mathrm{ WDE);
    /* Set new prescaler(time-out) value = 64K cycles (~0.5 s) */
    WDTCSR = (1<<WDE) | ( 1<<WDP2) | ( 1<<WDPO);
    __enable_interrupt();
    }
```

Notes: 1. The example code assumes that the part specific header file is included.
2. The Watchdog Timer should be reset before any change of the WDP bits, since a change in the WDP bits can result in a time-out when switching to a shorter time-out period.

### 12.5 Register Description

### 12.5.1 MCUSR - MCU Status Register

The MCU Status Register provides information on which reset source caused an MCU reset.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | MCUSR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x35 (0x55) | - | - | - | JTRF | WDRF | BORF | EXTRF | PORF |  |
| Read/Write | R | R | R | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 |  |  | it Desc |  |  |  |

## - Bit 4 - JTRF: JTAG Reset Flag

This bit is set if a reset is being caused by a logic one in the JTAG Reset Register selected by the JTAG instruction AVR_RESET. This bit is reset by a Power-on Reset, or by writing a logic zero to the flag.

## - Bit 3 - WDRF: Watchdog Reset Flag

This bit is set if a Watchdog Reset occurs. The bit is reset by a Power-on Reset, or by writing a logic zero to the flag.

## - Bit 2 - BORF: Brown-out Reset Flag

This bit is set if a Brown-out Reset occurs. The bit is reset by a Power-on Reset, or by writing a logic zero to the flag.

- Bit 1 - EXTRF: External Reset Flag

This bit is set if an External Reset occurs. The bit is reset by a Power-on Reset, or by writing a logic zero to the flag.

## - Bit 0 - PORF: Power-on Reset Flag

This bit is set if a Power-on Reset occurs. The bit is reset only by writing a logic zero to the flag.
To make use of the Reset Flags to identify a reset condition, the user should read and then Reset the MCUSR as early as possible in the program. If the register is cleared before another reset occurs, the source of the reset can be found by examining the Reset Flags.

### 12.5.2 WDTCSR - Watchdog Timer Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | WDTCSR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x60) | WDIF | WDIE | WDP3 | WDCE | WDE | WDP2 | WDP1 | WDP0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | X | 0 | 0 | 0 |  |

- Bit 7 - WDIF: Watchdog Interrupt Flag

This bit is set when a time-out occurs in the Watchdog Timer and the Watchdog Timer is configured for interrupt. WDIF is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, WDIF is cleared by writing a logic one to the flag. When the I-bit in SREG and WDIE are set, the Watchdog Time-out Interrupt is executed.

## - Bit 6 - WDIE: Watchdog Interrupt Enable

When this bit is written to one and the I-bit in the Status Register is set, the Watchdog Interrupt is enabled. If WDE is cleared in combination with this setting, the Watchdog Timer is in Interrupt Mode, and the corresponding interrupt is executed if time-out in the Watchdog Timer occurs.

If WDE is set, the Watchdog Timer is in Interrupt and System Reset Mode. The first time-out in the Watchdog Timer will set WDIF. Executing the corresponding interrupt vector will clear WDIE and WDIF automatically by hardware (the Watchdog goes to System Reset Mode). This is useful for keeping the Watchdog Timer security while using the interrupt. To stay in Interrupt and System Reset Mode, WDIE must be set after each interrupt. This should however not be done within the interrupt service routine itself, as this might compromise the safety-function of the Watchdog System Reset mode. If the interrupt is not executed before the next time-out, a System Reset will be applied.

Table 12-1. Watchdog Timer Configuration

| WDTON $^{(1)}$ | WDE | WDIE | Mode | Action on Time-out |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | 0 | Stopped | None |
| 1 | 0 | 1 | Interrupt Mode | Interrupt |
| 1 | 1 | 0 | System Reset Mode | Reset |
| 1 | 1 | 1 | Interrupt and System Reset Mode | Interrupt, then go to System Reset Mode |
| 0 | x | x | System Reset Mode | Reset |

Note: 1. WDTON Fuse set to "0" means programmed and "1" means unprogrammed.

## - Bit 4 - WDCE: Watchdog Change Enable

This bit is used in timed sequences for changing WDE and prescaler bits. To clear the WDE bit, and/or change the prescaler bits, WDCE must be set.

Once written to one, hardware will clear WDCE after four clock cycles.

## - Bit 3 - WDE: Watchdog System Reset Enable

WDE is overridden by WDRF in MCUSR. This means that WDE is always set when WDRF is set. To clear WDE, WDRF must be cleared first. This feature ensures multiple resets during conditions causing failure, and a safe start-up after the failure.

## - Bit 5, 2:0 - WDP3:0: Watchdog Timer Prescaler 3, 2, 1 and 0

The WDP3:0 bits determine the Watchdog Timer prescaling when the Watchdog Timer is running. The different prescaling values and their corresponding time-out periods are shown in Table 12-2 on page 66.

Table 12-2. Watchdog Timer Prescale Select

| WDP3 | WDP2 | WDP1 | WDP0 | Number of WDT Oscillator Cycles | Typical Time-out at V $_{\text {cc }}=5.0 \mathrm{~V}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | $2 \mathrm{~K}(2048)$ cycles | 16 ms |
| 0 | 0 | 0 | 1 | $4 \mathrm{~K}(4096)$ cycles | 32 ms |
| 0 | 0 | 1 | 0 | $8 \mathrm{~K}(8192)$ cycles | 64 ms |
| 0 | 0 | 1 | 1 | $16 \mathrm{~K}(16384)$ cycles | 0.125 s |
| 0 | 1 | 0 | 0 | $32 \mathrm{~K}(32768)$ cycles | 0.25 s |
| 0 | 1 | 0 | 1 | $64 \mathrm{~K}(65536)$ cycles | 0.5 s |
| 0 | 1 | 1 | 0 | $128 \mathrm{~K}(131072)$ cycles | 1.0 s |
| 0 | 1 | 1 | 1 | $256 \mathrm{~K}(262144)$ cycles | 2.0 s |
| 1 | 0 | 0 | 0 | $512 \mathrm{~K}(524288)$ cycles | 4.0 s |
| 1 | 0 | 0 | 1 | $1024 \mathrm{~K}(1048576)$ cycles | 8.0 s |
| 1 | 0 | 1 | 0 |  | Reserved |
| 1 | 0 | 1 | 1 |  |  |
| 1 | 1 | 0 | 0 |  |  |
| 1 | 1 | 0 | 1 |  |  |
| 1 | 1 | 1 | 0 |  |  |
| 1 | 1 |  |  |  |  |

## 13. I/O-Ports

### 13.1 Introduction

All AVR ports have true Read-Modify-Write functionality when used as general digital I/O ports. This means that the direction of one port pin can be changed without unintentionally changing the direction of any other pin with the SBI and CBI instructions. The same applies when changing drive value (if configured as output) or enabling/disabling of pull-up resistors (if configured as input). Each output buffer has symmetrical drive characteristics with both high sink and source capability. The pin driver is strong enough to drive LED displays directly. All port pins have individually selectable pull-up resistors with a supply-voltage invariant resistance. All I/O pins have protection diodes to both $\mathrm{V}_{\mathrm{CC}}$ and Ground as indicated in Figure 13-1. Refer to "Electrical Characteristics" on page 355 for a complete list of parameters.

Figure 13-1. I/O Pin Equivalent Schematic


All registers and bit references in this section are written in general form. A lower case " $x$ " represents the numbering letter for the port, and a lower case " $n$ " represents the bit number. However, when using the register or bit defines in a program, the precise form must be used. For example, PORTB3 for bit no. 3 in Port B, here documented generally as PORTxn. The physical I/O Registers and bit locations are listed in "Table 13-34 and Table 1335 relates the alternate functions of Port $L$ to the overriding signals shown in Figure 13-5 on page 73." on page 95.
Three I/O memory address locations are allocated for each port, one each for the Data Register - PORTx, Data Direction Register - DDRx, and the Port Input Pins - PINx. The Port Input Pins I/O location is read only, while the Data Register and the Data Direction Register are read/write. However, writing a logic one to a bit in the PINx Register, will result in a toggle in the corresponding bit in the Data Register. In addition, the Pull-up Disable - PUD bit in MCUCR disables the pull-up function for all pins in all ports when set.

Using the I/O port as General Digital I/O is described in "Ports as General Digital I/O" on page 68. Most port pins are multiplexed with alternate functions for the peripheral features on the device. How each alternate function interferes with the port pin is described in "Alternate Port Functions" on page 72. Refer to the individual module sections for a full description of the alternate functions.

Note that enabling the alternate function of some of the port pins does not affect the use of the other pins in the port as general digital I/O.

### 13.2 Ports as General Digital I/O

The ports are bi-directional I/O ports with optional internal pull-ups. Figure 13-2 shows a functional description of one I/O-port pin, here generically called Pxn.

Figure 13-2. General Digital I/O ${ }^{(1)}$


Note: 1. WRx, WPx, WDx, RRx, RPx, and RDx are common to all pins within the same port. $\mathrm{clk}_{1 / 0}$, SLEEP, and PUD are common to all ports.

### 13.2.1 Configuring the Pin

Each port pin consists of three register bits: DDxn, PORTxn, and PINxn. As shown in "Table 13-34 and Table 1335 relates the alternate functions of Port L to the overriding signals shown in Figure 13-5 on page 73." on page 95, the DDxn bits are accessed at the DDRx I/O address, the PORTxn bits at the PORTx I/O address, and the PINxn bits at the PINx I/O address.

The DDxn bit in the DDRx Register selects the direction of this pin. If DDxn is written logic one, Pxn is configured as an output pin. If DDxn is written logic zero, Pxn is configured as an input pin.
If PORTxn is written logic one when the pin is configured as an input pin, the pull-up resistor is activated. To switch the pull-up resistor off, PORTxn has to be written logic zero or the pin has to be configured as an output pin. The port pins are tri-stated when reset condition becomes active, even if no clocks are running.
If PORTxn is written logic one when the pin is configured as an output pin, the port pin is driven high (one). If PORTxn is written logic zero when the pin is configured as an output pin, the port pin is driven low (zero).

### 13.2.2 Toggling the Pin

Writing a logic one to PINxn toggles the value of PORTxn, independent on the value of DDRxn. Note that the SBI instruction can be used to toggle one single bit in a port.

### 13.2.3 Switching Between Input and Output

When switching between tri-state (\{DDxn, PORTxn\} = 0b00) and output high ( $\{D D x n$, PORTxn $=0 b 11$ ), an intermediate state with either pull-up enabled $\{D D x n$, PORTxn\} $=0 b 01$ ) or output low (\{DDxn, PORTxn\} $=0 b 10$ ) must occur. Normally, the pull-up enabled state is fully acceptable, as a high-impedant environment will not notice the difference between a strong high driver and a pull-up. If this is not the case, the PUD bit in the MCUCR Register can be set to disable all pull-ups in all ports.

Switching between input with pull-up and output low generates the same problem. The user must use either the tristate (\{DDxn, PORTxn\} = 0b00) or the output high state (\{DDxn, PORTxn\} = 0b11) as an intermediate step.
Table 13-1 summarizes the control signals for the pin value.
Table 13-1. Port Pin Configurations

| $\begin{aligned} & \text { ᄃ } \\ & \text { O} \end{aligned}$ |  |  | I/O | Pull-up | Comment |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | X | Input | No | Tri-state (Hi-Z) |
| 0 | 1 | 0 | Input | Yes | Pxn will source current if ext. pulled low |
| 0 | 1 | 1 | Input | No | Tri-state (Hi-Z) |
| 1 | 0 | X | Output | No | Output Low (Sink) |
| 1 | 1 | X | Output | No | Output High (Source) |

### 13.2.4 Reading the Pin Value

Independent of the setting of Data Direction bit DDxn, the port pin can be read through the PINxn Register bit. As shown in Figure 13-2 on page 68, the PINxn Register bit and the preceding latch constitute a synchronizer. This is needed to avoid metastability if the physical pin changes value near the edge of the internal clock, but it also introduces a delay. Figure 13-3 on page 69 shows a timing diagram of the synchronization when reading an externally applied pin value. The maximum and minimum propagation delays are denoted $t_{p d, \max }$ and $t_{p d, \min }$ respectively.

Figure 13-3. Synchronization when Reading an Externally Applied Pin value


Consider the clock period starting shortly after the first falling edge of the system clock. The latch is closed when the clock is low, and goes transparent when the clock is high, as indicated by the shaded region of the "SYNC

LATCH" signal. The signal value is latched when the system clock goes low. It is clocked into the PINxn Register at the succeeding positive clock edge. As indicated by the two arrows tpd,max and tpd,min, a single signal transition on the pin will be delayed between $1 / 2$ and $11 / 2$ system clock period depending upon the time of assertion.

When reading back a software assigned pin value, a nop instruction must be inserted as indicated in Figure 13-4. The out instruction sets the "SYNC LATCH" signal at the positive edge of the clock. In this case, the delay tpd through the synchronizer is one system clock period.

Figure 13-4. Synchronization when Reading a Software Assigned Pin Value


The following code example shows how to set port B pins 0 and 1 high, pins 2 and 3 low, and define the port pins from 4 to 7 as input with pull-ups assigned to port pins 6 and 7 . The resulting pin values are read back again, but as previously discussed, a nop instruction is included to be able to read back the value recently assigned to some of the pins.


Note: 1. For the assembly program, two temporary registers are used to minimize the time from pull-ups are set on pins 0 , 1,6 , and 7 , until the direction bits are correctly set, defining bit 2 and 3 as low and redefining bits 0 and 1 as strong high drivers.

### 13.2.5 Digital Input Enable and Sleep Modes

As shown in Figure 13-2 on page 68, the digital input signal can be clamped to ground at the input of the schmitttrigger. The signal denoted SLEEP in the figure, is set by the MCU Sleep Controller in Power-down mode, Powersave mode, and Standby mode to avoid high power consumption if some input signals are left floating, or have an analog signal level close to $\mathrm{V}_{\mathrm{CC}} / 2$.

SLEEP is overridden for port pins enabled as external interrupt pins. If the external interrupt request is not enabled, SLEEP is active also for these pins. SLEEP is also overridden by various other alternate functions as described in "Alternate Port Functions" on page 72.

If a logic high level ("one") is present on an asynchronous external interrupt pin configured as "Interrupt on Rising Edge, Falling Edge, or Any Logic Change on Pin" while the external interrupt is not enabled, the corresponding External Interrupt Flag will be set when resuming from the above mentioned Sleep mode, as the clamping in these sleep mode produces the requested logic change.

### 13.2.6 Unconnected Pins

If some pins are unused, it is recommended to ensure that these pins have a defined level. Even though most of the digital inputs are disabled in the deep sleep modes as described above, floating inputs should be avoided to reduce current consumption in all other modes where the digital inputs are enabled (Reset, Active mode and Idle mode).

The simplest method to ensure a defined level of an unused pin, is to enable the internal pull-up. In this case, the pull-up will be disabled during reset. If low power consumption during reset is important, it is recommended to use an external pull-up or pull-down. Connecting unused pins directly to $\mathrm{V}_{\mathrm{cc}}$ or GND is not recommended, since this may cause excessive currents if the pin is accidentally configured as an output.

### 13.3 Alternate Port Functions

Most port pins have alternate functions in addition to being general digital I/Os. Figure 13-5 on page 73 shows how the port pin control signals from the simplified Figure 13-2 on page 68 can be overridden by alternate functions. The overriding signals may not be present in all port pins, but the figure serves as a generic description applicable to all port pins in the AVR microcontroller family.

Figure 13-5. Alternate Port Functions ${ }^{(1)}$


Note: 1. WRx, WPx, WDx, RRx, RPx, and RDx are common to all pins within the same port. clk ${ }_{/ / 0}$, SLEEP, and PUD are common to all ports. All other signals are unique for each pin.

Table 13-2 summarizes the function of the overriding signals. The pin and port indexes from Figure 13-5 on page 73 are not shown in the succeeding tables. The overriding signals are generated internally in the modules having the alternate function.

Table 13-2. Generic Description of Overriding Signals for Alternate Functions

| Signal Name | Full Name | Description |
| :---: | :---: | :---: |
| PUOE | Pull-up Override Enable | If this signal is set, the pull-up enable is controlled by the PUOV signal. If this signal is cleared, the pull-up is enabled when \{DDxn, PORTxn, PUD\} $=$ Ob010. |
| PUOV | Pull-up Override Value | If PUOE is set, the pull-up is enabled/disabled when PUOV is set/cleared, regardless of the setting of the DDxn, PORTxn, and PUD Register bits. |
| DDOE | Data Direction Override Enable | If this signal is set, the Output Driver Enable is controlled by the DDOV signal. If this signal is cleared, the Output driver is enabled by the DDxn Register bit. |
| DDOV | Data Direction Override Value | If DDOE is set, the Output Driver is enabled/disabled when DDOV is set/cleared, regardless of the setting of the DDxn Register bit. |
| PVOE | Port Value Override Enable | If this signal is set and the Output Driver is enabled, the port value is controlled by the PVOV signal. If PVOE is cleared, and the Output Driver is enabled, the port Value is controlled by the PORTxn Register bit. |
| PVOV | Port Value Override Value | If PVOE is set, the port value is set to PVOV, regardless of the setting of the PORTxn Register bit. |
| PTOE | Port Toggle Override Enable | If PTOE is set, the PORTxn Register bit is inverted. |
| DIEOE | Digital Input Enable Override Enable | If this bit is set, the Digital Input Enable is controlled by the DIEOV signal. If this signal is cleared, the Digital Input Enable is determined by MCU state (Normal mode, sleep mode). |
| DIEOV | Digital Input Enable Override Value | If DIEOE is set, the Digital Input is enabled/disabled when DIEOV is set/cleared, regardless of the MCU state (Normal mode, sleep mode). |
| DI | Digital Input | This is the Digital Input to alternate functions. In the figure, the signal is connected to the output of the schmitt trigger but before the synchronizer. Unless the Digital Input is used as a clock source, the module with the alternate function will use its own synchronizer. |
| AIO | Analog Input/Output | This is the Analog Input/output to/from alternate functions. The signal is connected directly to the pad, and can be used bi-directionally. |

The following subsections shortly describe the alternate functions for each port, and relate the overriding signals to the alternate function. Refer to the alternate function description for further details.

### 13.3.1 Alternate Functions of Port A

The Port A has an alternate function as the address low byte and data lines for the External Memory Interface.
Table 13-3. Port A Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :--- |
| PA7 | AD7 (External memory interface address and data bit 7) |
| PA6 | AD6 (External memory interface address and data bit 6) |
| PA5 | AD5 (External memory interface address and data bit 5) |
| PA4 | AD4 (External memory interface address and data bit 4) |
| PA3 | AD3 (External memory interface address and data bit 3) |
| PA2 | AD2 (External memory interface address and data bit 2) |
| PA1 | AD1 (External memory interface address and data bit 1 ) |
| PA0 | AD0 (External memory interface address and data bit 0 ) |

Table 13-4 and Table 13-5 on page 76 relates the alternate functions of Port A to the overriding signals shown in Figure 13-5 on page 73 .

Table 13-4. Overriding Signals for Alternate Functions in PA7:PA4

| Signal Name | PA7/AD7 | PA6/AD6 | PA5/AD5 | PA4/AD4 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | SRE | SRE | SRE | SRE |
| PUOV | $\sim\left(\overline{\mathrm{WR}} \mid \mathrm{ADA}^{(1)}\right) \cdot$ PORTA7• $\overline{\text { PUD }}$ | $\begin{gathered} \sim(\overline{W R} \mid A D A) \cdot \text { PORTA6 } \\ \bullet \overline{\text { PUD }} \end{gathered}$ | $\begin{gathered} \sim(\overline{W R} \mid A D A) \cdot \text { PORTA5 } \\ \bullet \overline{P U D} \end{gathered}$ | $\begin{gathered} \sim(\overline{W R} \mid A D A) \cdot P O R T A 4 \\ \bullet \overline{P U D} \end{gathered}$ |
| DDOE | SRE | SRE | SRE | SRE |
| DDOV | $\overline{\text { WR }}$ \| ADA | $\overline{\text { WR }}$ \| ADA | $\overline{\text { WR }}$ \| ADA | $\overline{\mathrm{WR}}$ \| ADA |
| PVOE | SRE | SRE | SRE | SRE |
| PVOV | $\begin{gathered} \text { A7•ADA \| D7 OUTPUT } \\ \bullet \mathrm{WR} \end{gathered}$ | $\begin{gathered} \mathrm{A} 6 \cdot \mathrm{ADA} \mid \mathrm{D} 6 \text { OUTPUT } \\ \cdot \mathrm{WR} \end{gathered}$ | $\begin{gathered} \text { A5 • ADA \| D5 OUTPUT } \\ \bullet \mathrm{WR} \end{gathered}$ | $\begin{gathered} \mathrm{A} 4 \cdot \mathrm{ADA} \mid \mathrm{D} 4 \text { OUTPUT } \\ \bullet \mathrm{WR} \end{gathered}$ |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | 0 | 0 |
| DI | D7 INPUT | D6 INPUT | D5 INPUT | D4 INPUT |
| AIO | - | - | - | - |

Note: 1. ADA is short for ADdress Active and represents the time when address is output. See "External Memory Interface" on page 27 for details.

Table 13-5. Overriding Signals for Alternate Functions in PA3:PAO

| Signal Name | PA3/AD3 | PA2/AD2 | PA1/AD1 | PAO/ADO |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | SRE | SRE | SRE | SRE |
| PUOV | $\begin{gathered} \sim(\overline{\mathrm{WR}} \mid \mathrm{ADA}) \bullet \text { PORTA3 } \\ \bullet \cdot \overline{\mathrm{PUD}} \end{gathered}$ | $\begin{gathered} \sim(\overline{W R} \mid A D A) \cdot \text { PORTA2 } \\ \bullet \cdot \overline{\text { PUD }} \end{gathered}$ | $\begin{gathered} \sim(\overline{\mathrm{WR}} \mid \mathrm{ADA}) \cdot \text { PORTA1 } \\ \bullet \overline{\mathrm{PUD}} \end{gathered}$ | $\begin{gathered} \sim(\overline{\mathrm{WR}} \mid \mathrm{ADA}) \cdot \mathrm{PORTAO} \\ \bullet \overline{\mathrm{PUD}} \end{gathered}$ |
| DDOE | SRE | SRE | SRE | SRE |
| DDOV | $\overline{\text { WR }}$ \| ADA | $\overline{\text { WR }}$ A ADA | $\overline{\text { WR }}$ \| ADA | $\overline{\text { WR }}$ \| ADA |
| PVOE | SRE | SRE | SRE | SRE |
| PVOV | $\begin{gathered} \mathrm{A} 3 \cdot \mathrm{ADA} \mid \mathrm{D} 3 \text { OUTPUT } \\ \cdot \mathrm{WR} \end{gathered}$ | $\begin{gathered} \mathrm{A} 2 \cdot \mathrm{ADA} \mid \mathrm{D} 2 \text { OUTPUT } \\ \bullet \mathrm{WR} \end{gathered}$ | $\begin{gathered} \mathrm{A} 1 \cdot \mathrm{ADA} \mid \mathrm{D} 1 \text { OUTPUT } \\ \cdot \mathrm{WR} \end{gathered}$ | $\begin{gathered} \mathrm{AO} \cdot \mathrm{ADA} \mid \mathrm{DO} \text { OUTPUT } \\ \bullet \mathrm{WR} \end{gathered}$ |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | 0 | 0 |
| DI | D3 INPUT | D2 INPUT | D1 INPUT | DO INPUT |
| AIO | - | - | - | - |

### 13.3.2 Alternate Functions of Port B

The Port B pins with alternate functions are shown in Table 13-6.
Table 13-6. Port B Pins Alternate Functions

| Port Pin | Alternate Functions |
| :---: | :---: |
| PB7 | OC0A/OC1C/PCINT7 (Output Compare and PWM Output A for Timer/Counter0, Output Compare and |
| PWM Output C for Timer/Counter1 or Pin Change Interrupt 7) |  |

The alternate pin configuration is as follows:

## - OCOA/OC1C/PCINT7, Bit 7

OCOA, Output Compare Match A output: The PB7 pin can serve as an external output for the Timer/Counter0 Output Compare. The pin has to be configured as an output (DDB7 set "one") to serve this function. The OCOA pin is also the output pin for the PWM mode timer function.

OC1C, Output Compare Match C output: The PB7 pin can serve as an external output for the Timer/Counter1 Output Compare C. The pin has to be configured as an output (DDB7 set (one)) to serve this function. The OC1C pin is also the output pin for the PWM mode timer function.
PCINT7, Pin Change Interrupt source 7: The PB7 pin can serve as an external interrupt source.

## - OC1B/PCINT6, Bit 6

OC1B, Output Compare Match B output: The PB6 pin can serve as an external output for the Timer/Counter1 Output Compare B. The pin has to be configured as an output (DDB6 set (one)) to serve this function. The OC1B pin is also the output pin for the PWM mode timer function.

PCINT6, Pin Change Interrupt source 6: The PB6 pin can serve as an external interrupt source.

## - OC1A/PCINT5, Bit 5

OC1A, Output Compare Match A output: The PB5 pin can serve as an external output for the Timer/Counter1 Output Compare A. The pin has to be configured as an output (DDB5 set (one)) to serve this function. The OC1A pin is also the output pin for the PWM mode timer function.

PCINT5, Pin Change Interrupt source 5: The PB5 pin can serve as an external interrupt source.

## - OC2A/PCINT4, Bit 4

OC2A, Output Compare Match output: The PB4 pin can serve as an external output for the Timer/Counter2 Output Compare. The pin has to be configured as an output (DDB4 set (one)) to serve this function. The OC2A pin is also the output pin for the PWM mode timer function.

PCINT4, Pin Change Interrupt source 4: The PB4 pin can serve as an external interrupt source.

## - MISO/PCINT3 - Port B, Bit 3

MISO: Master Data input, Slave Data output pin for SPI channel. When the SPI is enabled as a master, this pin is configured as an input regardless of the setting of DDB3. When the SPI is enabled as a slave, the data direction of this pin is controlled by DDB3. When the pin is forced to be an input, the pull-up can still be controlled by the PORTB3 bit.

PCINT3, Pin Change Interrupt source 3: The PB3 pin can serve as an external interrupt source.

- MOSI/PCINT2 - Port B, Bit 2

MOSI: SPI Master Data output, Slave Data input for SPI channel. When the SPI is enabled as a slave, this pin is configured as an input regardless of the setting of DDB2. When the SPI is enabled as a master, the data direction of this pin is controlled by DDB2. When the pin is forced to be an input, the pull-up can still be controlled by the PORTB2 bit.

PCINT2, Pin Change Interrupt source 2: The PB2 pin can serve as an external interrupt source.

## - SCK/PCINT1 - Port B, Bit 1

SCK: Master Clock output, Slave Clock input pin for SPI channel. When the SPI is enabled as a slave, this pin is configured as an input regardless of the setting of DDB1. When the SPIO is enabled as a master, the data direction of this pin is controlled by DDB1. When the pin is forced to be an input, the pull-up can still be controlled by the PORTB1 bit.

PCINT1, Pin Change Interrupt source 1: The PB1 pin can serve as an external interrupt source.

## - SS/PCINT0 - Port B, Bit 0

$\overline{\mathrm{SS}}$ : Slave Port Select input. When the SPI is enabled as a slave, this pin is configured as an input regardless of the setting of DDBO. As a slave, the SPI is activated when this pin is driven low. When the SPI is enabled as a master, the data direction of this pin is controlled by DDBO. When the pin is forced to be an input, the pull-up can still be controlled by the PORTB0 bit.

Table 13-7 and Table 13-8 relate the alternate functions of Port $B$ to the overriding signals shown in Figure 13-5 on page 73. SPI MSTR INPUT and SPI SLAVE OUTPUT constitute the MISO signal, while MOSI is divided into SPI MSTR OUTPUT and SPI SLAVE INPUT.

PCINT0, Pin Change Interrupt source 0: The PB0 pin can serve as an external interrupt source.
Table 13-7. Overriding Signals for Alternate Functions in PB7:PB4

| Signal Name | PB7/OC0A/OC1C | PB6/OC1B | PB5/OC1A | PB4/OC2A |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | OC0/OC1C ENABLE | OC1B ENABLE | OC1A ENABLE | OC2A ENABLE |
| PVOV | OC0/OC1C | OC1B | OC1A | OC2A |
| DIEOE | PCINT7 • PCIE0 | PCINT6 • PCIE0 | PCINT5•PCIE0 | PCINT4•PCIE0 |
| DIEOV | 1 | 1 | 1 | 1 |
| DI | PCINT7 INPUT | PCINT6 INPUT | PCINT5 INPUT | PCINT4 INPUT |
| AIO | - | - | - | - |

Table 13-8. Overriding Signals for Alternate Functions in PB3:PB0

| Signal Name | PB3/MISO | PB2/MOSI | PB1/SCK | PB0/SS |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | SPE • MSTR | SPE • $\overline{\text { MSTR }}$ | SPE • $\overline{\text { SSTR }}$ | SPE• $\overline{\text { MSTR }}$ |
| PUOV | PORTB3 • $\overline{\text { PUD }}$ | PORTB2 • $\overline{\text { PUD }}$ | PORTB1 • $\overline{\text { PUD }}$ | PORTBO - $\overline{\text { PUD }}$ |
| DDOE | SPE • MSTR | SPE • MSTR | SPE • MSTR | SPE • MSTR |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | SPE • $\overline{\text { MSTR }}$ | SPE - MSTR | SPE • MSTR | 0 |
| PVOV | SPI SLAVE OUTPUT | SPI MSTR OUTPUT | SCK OUTPUT | 0 |
| DIEOE | PCINT3 • PCIE0 | PCINT2 • PCIE0 | PCINT1 • PCIE0 | PCINTO • PCIEO |
| DIEOV | 1 | 1 | 1 | 1 |
| DI | SPI MSTR INPUT PCINT3 INPUT | SPI SLAVE INPUT PCINT2 INPUT | SCK INPUT PCINT1 INPUT | SPI $\overline{\mathrm{SS}}$ PCINTO INPUT |
| AIO | - | - | - | - |

### 13.3.3 Alternate Functions of Port C

The Port C alternate function is as follows:

Table 13-9. Port C Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PC7 | A15 (External Memory interface address bit 15) |
| PC6 | A14 (External Memory interface address bit 14) |
| PC5 | A13 (External Memory interface address bit 13) |
| PC4 | A12 (External Memory interface address bit 12) |
| PC3 | A11 (External Memory interface address bit 11) |
| PC2 | A10 (External Memory interface address bit 10) |
| PC1 | A9 (External Memory interface address bit 9 ) |
| PC0 | A8 (External Memory interface address bit 8) |

Table 13-10 and Table 13-11 on page 80 relate the alternate functions of Port $C$ to the overriding signals shown in Figure 13-5 on page 73.

Table 13-10. Overriding Signals for Alternate Functions in PC7:PC4

| Signal Name | PC7/A15 | PC6/A14 | PC5/A13 | PC4/A12 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | SRE $\bullet(\mathrm{XMM}<1)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<2)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<3)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<4)$ |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | $\mathrm{SRE} \cdot(\mathrm{XMM}<1)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<2)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<3)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<4)$ |
| DDOV | 1 | 1 | 1 | 1 |
| PVOE | $\mathrm{SRE} \cdot(\mathrm{XMM}<1)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<2)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<3)$ | $\mathrm{SRE} \cdot(\mathrm{XMM}<4)$ |
| PVOV | $\mathrm{A15}$ | A 14 | A 13 | A 12 |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | 0 | 0 |
| DI | - | - | - | - |
| AIO | - | - | - | - |

Table 13-11. Overriding Signals for Alternate Functions in PC3:PC0

| Signal Name | PC3/A11 | PC2/A10 | PC1/A9 | PC0/A8 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | SRE•(XMM<5) | SRE•(XMM<6) | SRE•(XMM<7) | SRE•(XMM<7) |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | SRE - (XMM<5) | SRE - $(\mathrm{XMM}<6)$ | SRE - (XMM<7) | SRE - (XMM<7) |
| DDOV | 1 | 1 | 1 | 1 |
| PVOE | SRE - (XMM<5) | SRE - $(\mathrm{XMM}<6)$ | SRE - (XMM<7) | SRE - (XMM<7) |
| PVOV | A11 | A10 | A9 | A8 |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | 0 | 0 |
| DI | - | - | - | - |
| AIO | - | - | - | - |

### 13.3.4 Alternate Functions of Port D

The Port D pins with alternate functions are shown in Table 13-12.
Table 13-12. Port D Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PD7 | T0 (Timer/Counter0 Clock Input) |
| PD6 | T1 (Timer/Counter1 Clock Input) |
| PD5 | XCK1 (USART1 External Clock Input/Output) |
| PD4 | ICP1 (Timer/Counter1 Input Capture Trigger) |
| PD3 | $\overline{\text { INT3/TXD1 (External Interrupt3 Input or USART1 Transmit Pin) }}$ |
| PD2 | $\overline{\text { INT2/RXD1 }}$ (External Interrupt2 Input or USART1 Receive Pin) |
| PD1 | $\overline{\mathrm{INT} 1 / S D A ~(E x t e r n a l ~ I n t e r r u p t 1 ~ I n p u t ~ o r ~ T W I ~ S e r i a l ~ D A t a) ~}$ |
| PD0 | $\overline{\mathrm{INT0} / S C L ~(E x t e r n a l ~ I n t e r r u p t 0 ~ I n p u t ~ o r ~ T W I ~ S e r i a l ~ C L o c k) ~}$ |

The alternate pin configuration is as follows:

- TO - Port D, Bit 7

T0, Timer/Counter0 counter source.

## - T1 - Port D, Bit 6

T1, Timer/Counter1 counter source.

- XCK1 - Port D, Bit 5

XCK1, USART1 External clock. The Data Direction Register (DDD5) controls whether the clock is output (DDD5 set) or input (DDD5 cleared). The XCK1 pin is active only when the USART1 operates in Synchronous mode.

## - ICP1 - Port D, Bit 4

ICP1 - Input Capture Pin 1: The PD4 pin can act as an input capture pin for Timer/Counter1.

## - $\overline{\text { INT3/TXXD1 - Port D, Bit } 3}$

INT3, External Interrupt source 3: The PD3 pin can serve as an external interrupt source to the MCU.
TXD1, Transmit Data (Data output pin for the USART1). When the USART1 Transmitter is enabled, this pin is configured as an output regardless of the value of DDD3.

## - INT2/RXD1 - Port D, Bit 2

INT2, External Interrupt source 2. The PD2 pin can serve as an External Interrupt source to the MCU.
RXD1, Receive Data (Data input pin for the USART1). When the USART1 receiver is enabled this pin is configured as an input regardless of the value of DDD2. When the USART forces this pin to be an input, the pull-up can still be controlled by the PORTD2 bit.

## - $\overline{\text { INT1 }} /$ SDA - Port D, Bit 1

INT1, External Interrupt source 1. The PD1 pin can serve as an external interrupt source to the MCU.
SDA, 2-wire Serial Interface Data: When the TWEN bit in TWCR is set (one) to enable the 2-wire Serial Interface, pin PD1 is disconnected from the port and becomes the Serial Data I/O pin for the 2-wire Serial Interface. In this mode, there is a spike filter on the pin to suppress spikes shorter than 50 ns on the input signal, and the pin is driven by an open drain driver with slew-rate limitation.

## - INTO/SCL - Port D, Bit 0

INTO, External Interrupt source 0. The PDO pin can serve as an external interrupt source to the MCU.
SCL, 2-wire Serial Interface Clock: When the TWEN bit in TWCR is set (one) to enable the 2-wire Serial Interface, pin PDO is disconnected from the port and becomes the Serial Clock I/O pin for the 2-wire Serial Interface. In this mode, there is a spike filter on the pin to suppress spikes shorter than 50 ns on the input signal, and the pin is driven by an open drain driver with slew-rate limitation.

Table 13-13 on page 81 and Table 13-14 on page 82 relates the alternate functions of Port $D$ to the overriding signals shown in Figure 13-5 on page 73.

Table 13-13. Overriding Signals for Alternate Functions PD7:PD4

| Signal Name | PD7/T0 | PD6/T1 | PD5/XCK1 | PD4/ICP1 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 1 | 0 |
| DDOV | 0 | 0 | XCK1 OUTPUT ENABLE | 0 |
| PVOE | 0 | 0 | XCK1 OUTPUT ENABLE | 0 |
| PVOV | 0 | 0 | XCK1 OUTPUT | 0 |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | XCK1 INPUT | 0 |
| DI | T0 INPUT | T1 INPUT | - | ICP1 INPUT |
| AIO | - | - |  | - |

Table 13-14. Overriding Signals for Alternate Functions in PD3:PDO ${ }^{(1)}$

| Signal Name | PD3/INT3/TXD1 | PD2/INT2/RXD1 | PD1/INT1/SDA | PD0/INT0/SCL |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | TXEN1 | RXEN1 | TWEN | TWEN |
| PUOV | 0 | PORTD2• $\overline{\text { PUD }}$ | PORTD1•PUD | PORTD0• $\overline{\text { PUD }}$ |
| DDOE | TXEN1 | RXEN1 | TWEN | TWEN |
| DDOV | 1 | 0 | SDA_OUT | SCL_OUT |
| PVOE | TXEN1 | 0 | TWEN | TWEN |
| PVOV | TXD1 | 0 | 0 | 0 |
| DIEOE | INT3 ENABLE | INT2 ENABLE | INT1 ENABLE | INT0 ENABLE |
| DIEOV | 1 | 1 | 1 | 1 |
| DI | INT3 INPUT | INT2 INPUT/RXD1 | INT1 INPUT | INT0 INPUT |
| AIO | - | - | SDA INPUT | SCL INPUT |

Note: 1. When enabled, the 2-wire Serial Interface enables Slew-Rate controls on the output pins PD0 and PD1. This is not shown in this table. In addition, spike filters are connected between the AIO outputs shown in the port figure and the digital logic of the TWI module.

### 13.3.5 Alternate Functions of Port E

The Port E pins with alternate functions are shown in Table 13-15.
Table 13-15. Port E Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PE7 | (External Interrupt 7 Input, Timer/Counter3 Input Capture Trigger or Divided System Clock) |
| PE6 | INT6/ T3 |
| PE5 | (External Interrupt 6 Input or Timer/Counter3 Clock Input) |
| PE4 | INT5/OC3C |
| (External Interrupt 5 Input or Output Compare and PWM Output C for Timer/Counter3) |  |

[^0]
## - INT7/ICP3/CLKO - Port E, Bit 7

INT7, External Interrupt source 7: The PE7 pin can serve as an external interrupt source.
ICP3, Input Capture Pin 3: The PE7 pin can act as an input capture pin for Timer/Counter3.
CLKO - Divided System Clock: The divided system clock can be output on the PE7 pin. The divided system clock will be output if the CKOUT Fuse is programmed, regardless of the PORTE7 and DDE7 settings. It will also be output during reset.

- INT6/T3 - Port E, Bit 6

INT6, External Interrupt source 6: The PE6 pin can serve as an external interrupt source.
T3, Timer/Counter3 counter source.

- INT5/OC3C - Port E, Bit 5

INT5, External Interrupt source 5: The PE5 pin can serve as an External Interrupt source.
OC3C, Output Compare Match C output: The PE5 pin can serve as an External output for the Timer/Counter3 Output Compare C. The pin has to be configured as an output (DDE5 set "one") to serve this function. The OC3C pin is also the output pin for the PWM mode timer function.

## - INT4/OC3B - Port E, Bit 4

INT4, External Interrupt source 4: The PE4 pin can serve as an External Interrupt source.
OC3B, Output Compare Match B output: The PE4 pin can serve as an External output for the Timer/Counter3 Output Compare B. The pin has to be configured as an output (DDE4 set (one)) to serve this function. The OC3B pin is also the output pin for the PWM mode timer function.

- AIN1/OC3A - Port E, Bit 3

AIN1 - Analog Comparator Negative input. This pin is directly connected to the negative input of the Analog Comparator.

OC3A, Output Compare Match A output: The PE3 pin can serve as an External output for the Timer/Counter3 Output Compare A. The pin has to be configured as an output (DDE3 set "one") to serve this function. The OC3A pin is also the output pin for the PWM mode timer function.

- AINO/XCKO - Port E, Bit 2

AINO - Analog Comparator Positive input. This pin is directly connected to the positive input of the Analog Comparator.

XCKO, USARTO External clock. The Data Direction Register (DDE2) controls whether the clock is output (DDE2 set) or input (DDE2 cleared). The XCK0 pin is active only when the USART0 operates in Synchronous mode.

## - PDO/TXDO - Port E, Bit 1

PDO, SPI Serial Programming Data Output. During Serial Program Downloading, this pin is used as data output line for the ATmega1281/2561. For ATmega640/1280/2560 this function is placed on MISO.

TXDO, USARTO Transmit pin.

- PDI/RXD0/PCINT8 - Port E, Bit 0

PDI, SPI Serial Programming Data Input. During Serial Program Downloading, this pin is used as data input line for the ATmega1281/2561. For ATmega640/1280/2560 this function is placed on MOSI.

RXD0, USARTO Receive Pin. Receive Data (Data input pin for the USARTO). When the USART0 receiver is enabled this pin is configured as an input regardless of the value of DDREO. When the USARTO forces this pin to be an input, a logical one in PORTE0 will turn on the internal pull-up.

PCINT8, Pin Change Interrupt source 8: The PE0 pin can serve as an external interrupt source.
Table $13-16$ on page 84 and Table $13-17$ on page 84 relates the alternate functions of Port $E$ to the overriding signals shown in Figure 13-5 on page 73.

Table 13-16. Overriding Signals for Alternate Functions PE7:PE4

| Signal Name | PE7/INT7/ICP3 | PE6/INT6/T3 | PE5/INT5/OC3C | PE4/INT4/OC3B |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | 0 | 0 | OC3C ENABLE | OC3B ENABLE |
| PVOV | 0 | 0 | OC3C | OC3B |
| DIEOE | INT7 ENABLE | INT6 ENABLE | INT5 ENABLE | INT4 ENABLE |
| DIEOV | 1 | 1 | 1 | 1 |
| DI | INT7 INPUT/ICP3 |  |  |  |
| INPUT | INT7 INPUT/T3 INPUT | INT5 INPUT | INT4 INPUT |  |
| AIO | - | - | - | - |

Table 13-17. Overriding Signals for Alternate Functions in PE3:PE0

| Signal Name | PE3/AIN1/OC3A | PE2/AIN0/XCK0 | PE1/PDO $^{(1)}$ /TXD0 | PE0/PDI $^{(1) / R X D 0 / P C I N T 8 ~}$ |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | TXEN0 | RXEN0 |
| PUOV | 0 | 0 | 0 | PORTE0• PUD |
| DDOE | 0 | XCK0 OUTPUT <br> ENABLE | TXEN0 | RXEN0 |
| DDOV | 0 | 1 | 1 | 0 |
| PVOE | OC3B ENABLE | XCK0 OUTPUT <br> ENABLE | TXEN0 | 0 |
| PVOV | OC3B | XCK0 OUTPUT | TXD0 | 0 |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | - | PCINT8•PCIE1 |
| DI | 0 | XCK0 INPUT | 0 | RXD0 |
| PE0 | 0 | AIN0 INPUT | - | PCINT8 INPUT |
| AIO | AIN1 INPUT | 0 |  | - |

Note: 1. PDO/PDI only available at PE1/PE0 for ATmega1281/2561.

### 13.3.6 Alternate Functions of Port F

The Port F has an alternate function as analog input for the ADC as shown in Table 13-18. If some Port F pins are configured as outputs, it is essential that these do not switch when a conversion is in progress. This might corrupt the result of the conversion. If the JTAG interface is enabled, the pull-up resistors on pins PF7(TDI), PF5(TMS), and PF4(TCK) will be activated even if a Reset occurs.

Table 13-18. Port F Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PF7 | ADC7/TDI (ADC input channel 7 or JTAG Test Data Input) |
| PF6 | ADC6/TDO (ADC input channel 6 or JTAG Test Data Output) |
| PF5 | ADC5/TMS (ADC input channel 5 or JTAG Test Mode Select) |
| PF4 | ADC4/TCK (ADC input channel 4 or JTAG Test ClocK) |
| PF3 | ADC3 (ADC input channel 3) |
| PF2 | ADC2 (ADC input channel 2) |
| PF1 | ADC1 (ADC input channel 1$)$ |
| PF0 | ADC0 (ADC input channel 0$)$ |

## - TDI, ADC7 - Port F, Bit 7

ADC7, Analog to Digital Converter, Channel 7.
TDI, JTAG Test Data In: Serial input data to be shifted in to the Instruction Register or Data Register (scan chains).
When the JTAG interface is enabled, this pin can not be used as an I/O pin.

- TDO, ADC6 - Port F, Bit 6

ADC6, Analog to Digital Converter, Channel 6.
TDO, JTAG Test Data Out: Serial output data from Instruction Register or Data Register. When the JTAG interface is enabled, this pin can not be used as an I/O pin.
The TDO pin is tri-stated unless TAP states that shift out data are entered.

## - TMS, ADC5 - Port F, Bit 5

ADC5, Analog to Digital Converter, Channel 5.
TMS, JTAG Test Mode Select: This pin is used for navigating through the TAP-controller state machine. When the JTAG interface is enabled, this pin can not be used as an I/O pin.

- TCK, ADC4 - Port F, Bit 4

ADC4, Analog to Digital Converter, Channel 4.
TCK, JTAG Test Clock: JTAG operation is synchronous to TCK. When the JTAG interface is enabled, this pin can not be used as an I/O pin.

## - ADC3 - ADC0 - Port F, Bit 3:0

Analog to Digital Converter, Channel 3:0.

Table 13-19. Overriding Signals for Alternate Functions in PF7:PF4

| Signal Name | PF7/ADC7/TDI | PF6/ADC6/TDO | PF5/ADC5/TMS | PF4/ADC4/TCK |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | JTAGEN | JTAGEN | JTAGEN | JTAGEN |
| PUOV | 1 | 0 | 1 | 1 |
| DDOE | JTAGEN | JTAGEN | JTAGEN | JTAGEN |
| DDOV | 0 | SHIFT_IR + SHIFT_DR | 0 | 0 |
| PVOE | 0 | JTAGEN | 0 | 0 |
| PVOV | 0 | TDO | 0 | 0 |
| DIEOE | JTAGEN | JTAGEN | JTAGEN | JTAGEN |
| DIEOV | 0 | 0 | 0 | 0 |
| DI | - | - | - | - |
| AIO | TDI/ADC7 INPUT | ADC6 INPUT | TMS/ADC5 INPUT | TCK/ADC4 INPUT |

Table 13-20. Overriding Signals for Alternate Functions in PF3:PFO

| Signal Name | PF3/ADC3 | PF2/ADC2 | PF1/ADC1 | PF0/ADC0 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | 0 | 0 | 0 | 0 |
| PVOV | 0 | 0 | 0 | 0 |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | - | - | 0 |
| DI | - | ADC2 INPUT | ADC1 INPUT | ADC0 INPUT |
| AIO | ADC3 INPUT |  |  |  |

### 13.3.7 Alternate Functions of Port G

The Port G alternate pin configuration is as follows:
Table 13-21. Port G Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PG5 | OCOB (Output Compare and PWM Output B for Timer/Counter0) |
| PG4 | TOSC1 (RTC Oscillator Timer/Counter2) |
| PG3 | TOSC2 (RTC Oscillator Timer/Counter2) |
| PG2 | ALE (Address Latch Enable to external memory) |
| PG1 | $\overline{R D}$ (Read strobe to external memory) |
| PG0 | $\overline{W R}$ (Write strobe to external memory) |

## - OCOB - Port G, Bit 5

OCOB, Output Compare match B output: The PG5 pin can serve as an external output for the Tlmer/Counter0 Output Compare. The pin has to be configured as an output (DDG5 set) to serve this function. The OCOB pin is also the output pin for the PWM mode timer function.

- TOSC1 - Port G, Bit 4

TOSC2, Timer Oscillator pin 1: When the AS2 bit in ASSR is set (one) to enable asynchronous clocking of Timer/Counter2, pin PG4 is disconnected from the port, and becomes the input of the inverting Oscillator amplifier. In this mode, a Crystal Oscillator is connected to this pin, and the pin can not be used as an I/O pin.

## - TOSC2 - Port G, Bit 3

TOSC2, Timer Oscillator pin 2: When the AS2 bit in ASSR is set (one) to enable asynchronous clocking of Timer/Counter2, pin PG3 is disconnected from the port, and becomes the inverting output of the Oscillator amplifier. In this mode, a Crystal Oscillator is connected to this pin, and the pin can not be used as an I/O pin.

- ALE - Port G, Bit 2

ALE is the external data memory Address Latch Enable signal.

- RD - Port G, Bit 1
$\overline{\mathrm{RD}}$ is the external data memory read control strobe.
- WR - Port G, Bit 0
$\overline{W R}$ is the external data memory write control strobe.
Table 13-22 on page 87 and Table 13-23 on page 88 relates the alternate functions of Port $G$ to the overriding signals shown in Figure 13-5 on page 73.

Table 13-22. Overriding Signals for Alternate Functions in PG5:PG4

| Signal Name | - | - | PG5/OC0B | PG4/TOSC1 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | - | - | - | AS2 |
| PUOV | - | - | - | 0 |
| DDOE | - | - | - | AS2 |
| DDOV | - | - | - | 0 |
| PVOE | - | - | OC0B Enable | 0 |
| PVOV | - | - | - | 0 OBB |
| PTOE | - | - | - | - |
| DIEOE | - | - | - | AS2 |
| DIEOV | - | - | - | EXCLK |
| DI | - |  | - |  |
| AIO | - | - | T/C2 OSC INPUT |  |

Table 13-23. Overriding Signals for Alternate Functions in PG3:PG0

| Signal Name | PG3/TOSC2 | PG2/ALE/A7 | PG1/RD | PG0/WR |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | AS2• EXCLK | SRE | SRE | SRE |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | AS2• EXCLK | SRE | SRE | SRE |
| DDOV | 0 | 1 | 1 | 1 |
| PVOE | 0 | SRE | SRE | SRE |
| PVOV | 0 | ALE | - | WD |
| PTOE | - | - | 0 | - |
| DIEOE | AS2•EXCLK | 0 | 0 | 0 |
| DIEOV | 0 | - | - | 0 |
| DI | - | - | - | - |
| AIO | T/C2 OSC OUTPUT |  |  |  |

### 13.3.8 Alternate Functions of Port H

The Port H alternate pin configuration is as follows:
Table 13-24. Port H Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PH7 | T4 (Timer/Counter4 Clock Input) |
| PH6 | OC2B (Output Compare and PWM Output B for Timer/Counter2) |
| PH5 | OC4C (Output Compare and PWM Output C for Timer/Counter4) |
| PH4 | OC4B (Output Compare and PWM Output B for Timer/Counter4) |
| PH3 | OC4A (Output Compare and PWM Output A for Timer/Counter4) |
| PH2 | XCK2 (USART2 External Clock) |
| PH1 | TXD2 (USART2 Transmit Pin) |
| PH0 | RXD2 (USART2 Receive Pin) |

## - T4 - Port H, Bit 7

T4, Timer/Counter4 counter source.

- OC2B - Port H, Bit 6

OC2B, Output Compare Match B output: The PH6 pin can serve as an external output for the Timer/Counter2 Output Compare B. The pin has to be configured as an output (DDH6 set) to serve this function. The OC2B pin is also the output pin for the PWM mode timer function.

## - OC4C - Port H, Bit 5

OC4C, Output Compare Match C output: The PH5 pin can serve as an external output for the Timer/Counter4 Output Compare C. The pin has to be configured as an output (DDH5 set) to serve this function. The OC4C pin is also the output pin for the PWM mode timer function.

## - OC4B - Port H, Bit 4

OC4B, Output Compare Match B output: The PH4 pin can serve as an external output for the Timer/Counter2 Output Compare B. The pin has to be configured as an output (DDH4 set) to serve this function. The OC4B pin is also the output pin for the PWM mode timer function.

## - OC4A - Port H, Bit 3

OC4C, Output Compare Match A output: The PH3 pin can serve as an external output for the Timer/Counter4 Output Compare A. The pin has to be configured as an output (DDH3 set) to serve this function. The OC4A pin is also the output pin for the PWM mode timer function.

- XCK2 - Port H, Bit 2

XCK2, USART2 External Clock: The Data Direction Register (DDH2) controls whether the clock is output (DDH2 set) or input (DDH2 cleared). The XC2K pin is active only when the USART2 operates in synchronous mode.

- TXD2 - Port H, Bit 1

TXD2, USART2 Transmit Pin.

## - RXD2 - Port H, Bit 0

RXD2, USART2 Receive pin: Receive Data (Data input pin for the USART2). When the USART2 Receiver is enabled, this pin is configured as an input regardless of the value of DDHO. When the USART2 forces this pin to be an input, a logical on in PORTH0 will turn on the internal pull-up.

Table 13-25. Overriding Signals for Alternate Functions in PH7:PH4

| Signal Name | PH7/T4 | PH6/OC2B | PH5/OC4C | PH4/OC4B |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | 0 | OC2B ENABLE | OC4C ENABLE | OC4B ENABLE |
| PVOV | 0 | OC2B | OC4C | OC4B |
| PTOE | - | - | - | - |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | 0 | 0 |
| DI | T4 INPUT | - | 0 | 0 |
| AIO | - | - | - |  |

Table 13-26. Overriding Signals for Alternate Functions in PH3:PH0

| Signal Name | PH3/OC4A | PH2/XCK2 | PH1/TXD2 | PH0/RXD2 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | TXEN2 | RXEN2 |
| PUOV | 0 | 0 | 0 | PORTH0 • PUD |
| DDOE | 0 | XCK2 OUTPUT <br> ENABLE | TXEN2 | RXEN2 |
| DDOV | 0 | 1 | 1 | 0 |
| PVOE | OC4A ENABLE | XCK2 OUTPUT <br> ENABLE | TXEN2 | 0 |
| PVOV | OC4A | XCK2 | TXD2 | 0 |
| PTOE | - | - | - | 0 |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | - | 0 |
| DI | 0 | - | XC2K INPUT | 0 |
| AIO | - |  | 0 | 0 |

### 13.3.9 Alternate Functions of Port J

The Port J alternate pin configuration is as follows:
Table 13-27. Port J Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PJ7 | - |
| PJ6 | PCINT15 (Pin Change Interrupt 15) |
| PJ5 | PCINT14 (Pin Change Interrupt 14) |
| PJ4 | PCINT13 (Pin Change Interrupt 13) |
| PJ3 | PCINT12 (Pin Change Interrupt 12) |
| PJ2 | XCK3/PCINT11 (USART3 External Clock or Pin Change Interrupt 11) |
| PJ1 | TXD3/PCINT10 (USART3 Transmit Pin or Pin Change Interrupt 10) |
| PJ0 | RXD3/PCINT9 (USART3 Receive Pin or Pin Change Interrupt 9) |

## - PCINT15:12 - Port J, Bit 6:3

PCINT15:12, Pin Change Interrupt Source 15:12. The PJ6:3 pins can serve as External Interrupt Sources.

- XCK2/PCINT11 - Port J, Bit 2

XCK2, USART 2 External Clock. The Data Direction Register (DDJ2) controls whether the clock is output (DDJ2 set) or input (DDJ2 cleared). The XCK2 pin is active only when the USART2 operates in synchronous mode.

PCINT11, Pin Change Interrupt Source 11. The PJ2 pin can serve as External Interrupt Sources.

- TXD3/PCINT10 - Port J, Bit 1

TXD3, USART3 Transmit pin.
PCINT10, Pin Change Interrupt Source 10. The PJ1 pin can serve as External Interrupt Sources.

## - RXD3/PCINT9 - Port J, Bit 0

RXD3, USART3 Receive pin. Receive Data (Data input pin for the USART3). When the USART3 Receiver is enabled, this pin is configured as an input regardless of the value of DDJO. When the USART3 forces this pin to be an input, a logical one in PORTJO will turn on the internal pull-up.

PCINT9, Pin Change Interrupt Source 9. The PJO pin can serve as External Interrupt Sources.
Table 13-28 on page 92 and Table 13-29 on page 92 relates the alternate functions of Port J to the overriding signals shown in Figure 13-5 on page 73.

Table 13-28. Overriding Signals for Alternate Functions in PJ7:PJ4

| Signal Name | PJ7 | PJ6/ PCINT15 | PJ5/ PCINT14 | PJ4/ PCINT13 |
| :--- | :--- | :--- | :--- | :--- |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | 0 | 0 | 0 | 0 |
| PVOV | 0 | 0 | 0 | 0 |
| PTOE | - | - | - | - |
| DIEOE | 0 | PCINT15.PCIE1 | PCINT14.PCIE1 | PCINT13.PCIE1 |
| DIEOV | 0 | 1 | 1 | 1 |
| DI | 0 | PCINT15 INPUT | PCINT14 INPUT | PCINT13 INPUT |
| AIO | - | - | - | - |

Table 13-29. Overriding Signals for Alternate Functions in PJ3:PJ0

| Signal Name | PJ3/PCINT12 | PJ2/XCK3/PCINT11 | PJ1/TXD3/PCINT10 | PJ0/RXD3/PCINT9 |
| :--- | :--- | :--- | :--- | :--- |
| PUOE | 0 | 0 | TXEN3 | RXEN3 |
| PUOV | 0 | 0 | 0 | PORTJ0•PUD |
| DDOE | 0 | XCK3 OUTPUT <br> ENABLE | TXEN3 | RXEN3 |
| DDOV | 0 | 1 | 1 | 0 |
| PVOE | 0 | XCK3 OUTPUT <br> ENABLE | TXEN3 | 0 |
| PVOV | 0 | XCK3 | - | - |
| PTOE | - | PCINT11•PCIE1 | PCINT10.PCIE1 | PCINT9.PCIE1 |
| DIEOE | PCINT12•PCIE1 | 1 | 1 | 1 |
| DIEOV | 1 | PCINT11 INPUT <br> XCK3 INPUT | PCINT10 INPUT | PCINT9 INPUT RXD3 |
| DI | PCINT12 INPUT | - | - | - |
| AIO | - |  |  |  |

### 13.3.10 Alternate Functions of Port K

The Port K alternate pin configuration is as follows:
Table 13-30. Port K Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PK7 | ADC15/PCINT23 (ADC Input Channel 15 or Pin Change Interrupt 23) |
| PK6 | ADC14/PCINT22 (ADC Input Channel 14 or Pin Change Interrupt 22) |
| PK5 | ADC13/PCINT21 (ADC Input Channel 13 or Pin Change Interrupt 21) |

Table 13-30. Port K Pins Alternate Functions (Continued)

| Port Pin | Alternate Function |
| :---: | :---: |
| PK4 | ADC12/PCINT20 (ADC Input Channel 12 or Pin Change Interrupt 20) |
| PK3 | ADC11/PCINT19 (ADC Input Channel 11 or Pin Change Interrupt 19) |
| PK2 | ADC10/PCINT18 (ADC Input Channel 10 or Pin Change Interrupt 18) |
| PK1 | ADC9/PCINT17 (ADC Input Channel 9 or Pin Change Interrupt 17) |
| PK0 | ADC8 /PCINT16 (ADC Input Channel 8 or Pin Change Interrupt 16) |

- ADC15:8/PCINT23:16 - Port K, Bit 7:0

ADC15:8, Analog to Digital Converter, Channel 15-8.
PCINT23:16, Pin Change Interrupt Source 23:16. The PK7:0 pins can serve as External Interrupt Sources.
Table 13-31. Overriding Signals for Alternate Functions in PK7:PK4

| Signal Name | PK7/ADC15/PCINT23 | PK6/ADC14/PCINT22 | PK5/ADC13/PCINT21 | PK4/ADC12/PCINT20 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | 0 | 0 | 0 | 0 |
| PVOV | 0 | 0 | - | 0 |
| PTOE | - | - | 1 | - |
| DIEOE | PCINT23 • PCIE2 | PCINT22 • PCIE2 | PCINT21•PCIE2 | PCINT20 • PCIE2 |
| DIEOV | 1 | 1 | PCINT21 INPUT | PCINT20 INPUT |
| DI | PCINT23 INPUT | PCINT22 INPUT | PC13 INPUT | ADC12 INPUT |
| AIO | ADC15 INPUT | ADC14 INPUT | ADC13 |  |

Table 13-32. Overriding Signals for Alternate Functions in PK3:PK0

| Signal Name | PK3/ADC11/PCINT19 | PK2/ADC10/PCINT18 | PK1/ADC9/PCINT17 | PK0/ADC8/PCINT16 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | 0 | 0 | 0 | 0 |
| PVOV | 0 | 0 | - | 0 |
| PTOE | - | - | 1 | - |
| DIEOE | PCINT19 • PCIE2 | PCINT18 • PCIE2 | PCINT17 • PCIE2 | PCINT16•PCIE2 |
| DIEOV | 1 | 1 | PCINT17 INPUT | PCINT16 INPUT |
| DI | PCINT19 INPUT | PCINT18 INPUT | PDC9 INPUT | ADC8 INPUT |
| AIO | ADC11 INPUT | ADC10INPUT | ADP |  |

### 13.3.11 Alternate Functions of Port L

The Port L alternate pin configuration is as follows:
Table 13-33. Port L Pins Alternate Functions

| Port Pin | Alternate Function |
| :---: | :---: |
| PL7 | - |
| PL6 | - |
| PL5 | OC5C (Output Compare and PWM Output C for Timer/Counter5) |
| PL4 | OC5B (Output Compare and PWM Output B for Timer/Counter5) |
| PL3 | OC5A (Output Compare and PWM Output A for Timer/Counter5) |
| PL2 | T5 (Timer/Counter5 Clock Input) |
| PL1 | ICP5 (Timer/Counter5 Input Capture Trigger) |
| PL0 | ICP4 (Timer/Counter4 Input Capture Trigger) |

## - OC5C - Port L, Bit 5

OC5C, Output Compare Match C output: The PL5 pin can serve as an external output for the Timer/Counter5 Output Compare C. The pin has to be configured as an output (DDL5 set) to serve this function. The OC5C pin is also the output pin for the PWM mode timer function.

## - OC5B - Port L, Bit 4

OC5B, Output Compare Match B output: The PL4 pin can serve as an external output for the Timer/Counter 5 Output Compare B. The pin has to be configured as an output (DDL4 set) to serve this function. The OC5B pin is also the output pin for the PWM mode timer function.

- OC5A - Port L, Bit 3

OC5A, Output Compare Match A output: The PL3 pin can serve as an external output for the Timer/Counter 5 Output Compare A. The pin has to be configured as an output (DDL3 set) to serve this function. The OC5A pin is also the output pin for the PWM mode timer function.

## - T5 - Port L, Bit 2

T5, Timer/Counter5 counter source.

- ICP5 - Port L, Bit 1

ICP5, Input Capture Pin 5: The PL1 pin can serve as an Input Capture pin for Timer/Counter5.

- ICP4 - Port L, Bit 0

ICP4, Input Capture Pin 4: The PLO pin can serve as an Input Capture pin for Timer/Counter4.

Table 13-34 and Table 13-35 relates the alternate functions of Port $L$ to the overriding signals shown in Figure 135 on page 73 .

Table 13-34. Overriding Signals for Alternate Functions in PL7:PL4

| Signal Name | PL7 | PL6 | PL5/OC5C | PL4/OC5B |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | - | - | 0 | 0 |
| DDOV | - | - | 0 | 0 |
| PVOE | - | - | OC5C ENABLE | OC5B ENABLE |
| PVOV | - | - | 0 | OC5CB |
| PTOE | 0 | - | 0 | - |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | 0 | - | 0 |
| DI | - | - | 0 | - |
| AIO | - |  | 0 | 0 |

Table 13-35. Overriding Signals for Alternate Functions in PL3:PL0

| Signal Name | PL3/OC5A | PL2/T5 | PL1/ICP5 | PL0/ICP4 |
| :---: | :---: | :---: | :---: | :---: |
| PUOE | 0 | 0 | 0 | 0 |
| PUOV | 0 | 0 | 0 | 0 |
| DDOE | 0 | 0 | 0 | 0 |
| DDOV | 0 | 0 | 0 | 0 |
| PVOE | OC5A ENABLE | 0 | 0 | 0 |
| PVOV | OC5A | 0 | 0 | 0 |
| PTOE | - | - | 0 | - |
| DIEOE | 0 | 0 | 0 | 0 |
| DIEOV | 0 | T5 INPUT | - | ICP5 INPUT |

### 13.4 Register Description for I/O-Ports

### 13.4.1 MCUCR - MCU Control Register

Bit
$0 \times 35$ ( $0 \times 55$ )
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| JTD | - | - | PUD | - | - | IVSEL | IVCE |
| R/W | $R$ | $R$ | $R / W$ | $R$ | $R$ | $R / W$ | $R / W$ |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

- Bit 4 - PUD: Pull-up Disable

When this bit is written to one, the I/O ports pull-up resistors are disabled even if the DDxn and PORTxn Registers are configured to enable the pull-up resistor (\{DDxn, PORTxn\} $=0 b 01$ ). See "Configuring the Pin" on page 68 for more details about this feature.

### 13.4.2 PORTA - Port A Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PORTA |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x02 (0x22) | PORTA7 | PORTA6 | PORTA5 | PORTA4 | PORTA3 | PORTA2 | PORTA1 | PORTA0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 13.4.3 DDRA - Port A Data Direction Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | DDRA |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x01 (0x21) | DDA7 | DDA6 | DDA5 | DDA4 | DDA3 | DDA2 | DDA1 | DDA0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

13.4.4 PINA - Port A Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0x00 (0x20) | PINA7 | PINA6 | PINA5 | PINA4 | PINA3 | PINA2 | PINA1 | PINA0 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |  |

### 13.4.5 PORTB - Port B Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PORTB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x05 (0x25) | PORTB7 | PORTB6 | PORTB5 | PORTB4 | PORTB3 | PORTB2 | PORTB1 | PORTB0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 13.4.6 DDRB - Port B Data Direction Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x04 (0x24) | DDB7 | DDB6 | DDB5 | DDB4 | DDB3 | DDB2 | DDB1 | DDB0 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.7 PINB - Port B Input Pins Address

Bit
$0 \times 03$ (0x23)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PINB7 | PINB6 | PINB5 | PINB4 | PINB3 | PINB2 | PINB1 | PINB0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |

### 13.4.8 PORTC - Port C Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PORTC |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x08 (0x28) | PORTC7 | PORTC6 | PORTC5 | PORTC4 | PORTC3 | PORTC2 | PORTC1 | PORTC0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 13.4.9 DDRC - Port C Data Direction Register

Bit
0x07 (0x27)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DDC7 | DDC6 | DDC5 | DDC4 | DDC3 | DDC2 | DDC1 | DDC0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.10 PINC- Port C Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \times 06(0 \times 26)$ | PINC7 | PINC6 | PINC5 | PINC4 | PINC3 | PINC2 | PINC1 | PINC0 |  |
|  | Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |  |

### 13.4.11 PORTD - Port D Data Register

Bit
0x0B (0x2B)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PORTD7 | PORTD6 | PORTD5 | PORTD4 | PORTD3 | PORTD2 | PORTD1 | PORTD0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.12 DDRD - Port D Data Direction Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0x0A (0x2A) | DDD7 | DDD6 | DDD5 | DDD4 | DDD3 | DDD2 | DDD1 | DDD0 |
|  | Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 13.4.13 PIND - Port D Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PIND |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x09 (0x29) | PIND7 | PIND6 | PIND5 | PIND4 | PIND3 | PIND2 | PIND1 | PIND0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |  |

### 13.4.14 PORTE - Port E Data Register

Bit
0x0E (0x2E)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PORTE7 | PORTE6 | PORTE5 | PORTE4 | PORTE3 | PORTE2 | PORTE1 | PORTE0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.15 DDRE - Port E Data Direction Register

Bit
0x0D (0x2D)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DDE7 | DDE6 | DDE5 | DDE4 | DDE3 | DDE2 | DDE1 | DDE0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.16 PINE - Port E Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x0C (0x2C) | PINE7 | PINE6 | PINE5 | PINE4 | PINE3 | PINE2 | PINE1 | PINE0 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |

PINE

### 13.4.17 PORTF - Port F Data Register

Bit
0x11 (0x31)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PORTF7 | PORTF6 | PORTF5 | PORTF4 | PORTF3 | PORTF2 | PORTF1 | PORTF0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.18 DDRF - Port F Data Direction Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \times 10(0 \times 30)$ | DDF7 | DDF6 | DDF5 | DDF4 | DDF3 | DDF2 | DDF1 | DDF0 |  |
|  | Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

DDRF
13.4.19 PINF - Port F Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x0F (0x2F) | PINF7 | PINF6 | PINF5 | PINF4 | PINF3 | PINF2 | PINF1 | PINFO |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |

PINF

### 13.4.20 PORTG - Port G Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x14 (0x34) | - | - | PORTG5 | PORTG4 | PORTG3 | PORTG2 | PORTG1 | PORTG0 |
| Read/Write | R | R | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

PORTG

### 13.4.21 DDRG - Port G Data Direction Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x13 (0x33) | - | - | DDG5 | DDG4 | DDG3 | DDG2 | DDG1 | DDG0 |
| Read/Write | R | R | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.22 PING - Port G Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | PING |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \times 12(0 \times 32)$ | - | - | PING5 | PING4 | PING3 | PING2 | PING1 | PING0 |
| Read/Write | R | R | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | N/A | N/A | N/A | N/A | N/A | N/A |

### 13.4.23 PORTH - Port H Data Register

Bit
(0x102)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PORTH7 | PORTH6 | PORTH5 | PORTH4 | PORTH3 | PORTH2 | PORTH1 | PORTH0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## PORTH

### 13.4.24 DDRH - Port H Data Direction Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | DDRH |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x101) | DDH7 | DDH6 | DDH5 | DDH4 | DDH3 | DDH2 | DDH1 | DDH0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 13.4.25 PINH - Port H Input Pins Address

Bit
(0x100)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PINH5 | PINH5 | PINH5 | PINH4 | PINH3 | PINGH | PINH1 | PINH0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| N/A | N/A | N/A | N/A | N/A | N/A | R/W |  |
| N/A | N/A |  |  |  |  |  |  |

### 13.4.26 PORTJ - Port J Data Register

Bit
(0x105)
Read/Write Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PORTJ7 | PORTJ6 | PORTJ5 | PORTJ4 | PORTJ3 | PORTJ2 | PORTJ1 | PORTJ0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

13.4.27 DDRJ - Port J Data Direction Register

Bit
(0x104)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DDJ7 | DDJ6 | DDJ5 | DDJ4 | DDJ3 | DDJ2 | DDJ1 | DDJ0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.28 PINJ - Port J Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x103) | PINJ5 | PINJ5 | PINJ5 | PINJ4 | PINJ3 | PINGJ | PINJ1 | PINJO |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |

PINJ

### 13.4.29 PORTK - Port K Data Register

Bit
(0x108)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PORTK7 | PORTK6 | PORTK5 | PORTK4 | PORTK3 | PORTK2 | PORTK1 | PORTK0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.30 DDRK - Port K Data Direction Register

Bit
(0x107)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| DDK7 | DDK6 | DDK5 | DDK4 | DDK3 | DDK2 | DDK1 | DDK0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

DDRK

### 13.4.31 PINK - Port K Input Pins Address

Bit
(0x106)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PINK5 | PINK5 | PINK5 | PINK4 | PINK3 | PINGK | PINK1 | PINK0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |

PINK

### 13.4.32 PORTL - Port L Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x10B) | PORTL7 | PORTL6 | PORTL5 | PORTL4 | PORTL3 | PORTL2 | PORTL1 | PORTLO |
|  | Pead/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W

PORTL
13.4.33 DDRL - Port L Data Direction Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x10A) | DDL7 | DDL6 | DDL5 | DDL4 | DDL3 | DDL2 | DDL1 | DDL0 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 13.4.34 PINL - Port L Input Pins Address

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x109) | PINL5 | PINL5 | PINL5 | PINL4 | PINL3 | PINGL | PINL1 | PINLO |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | N/A | N/A | N/A | N/A | N/A | N/A | N/A | N/A |

PINL

## 14. Interrupts

This section describes the specifics of the interrupt handling as performed in ATmega640/1280/1281/2560/2561. For a general explanation of the AVR interrupt handling, refer to "Reset and Interrupt Handling" on page 17.

### 14.1 Interrupt Vectors in ATmega640/1280/1281/2560/2561

Table 14-1. Reset and Interrupt Vectors

| Vector No. | Program Address ${ }^{(2)}$ | Source | Interrupt Definition |
| :---: | :---: | :---: | :---: |
| 1 | \$0000 ${ }^{(1)}$ | RESET | External Pin, Power-on Reset, Brown-out Reset, Watchdog Reset, and JTAG AVR Reset |
| 2 | \$0002 | INTO | External Interrupt Request 0 |
| 3 | \$0004 | INT1 | External Interrupt Request 1 |
| 4 | \$0006 | INT2 | External Interrupt Request 2 |
| 5 | \$0008 | INT3 | External Interrupt Request 3 |
| 6 | \$000A | INT4 | External Interrupt Request 4 |
| 7 | \$000C | INT5 | External Interrupt Request 5 |
| 8 | \$000E | INT6 | External Interrupt Request 6 |
| 9 | \$0010 | INT7 | External Interrupt Request 7 |
| 10 | \$0012 | PCINTO | Pin Change Interrupt Request 0 |
| 11 | \$0014 | PCINT1 | Pin Change Interrupt Request 1 |
| 12 | \$0016 ${ }^{(3)}$ | PCINT2 | Pin Change Interrupt Request 2 |
| 13 | \$0018 | WDT | Watchdog Time-out Interrupt |
| 14 | \$001A | TIMER2 COMPA | Timer/Counter2 Compare Match A |
| 15 | \$001C | TIMER2 COMPB | Timer/Counter2 Compare Match B |
| 16 | \$001E | TIMER2 OVF | Timer/Counter2 Overflow |
| 17 | \$0020 | TIMER1 CAPT | Timer/Counter1 Capture Event |
| 18 | \$0022 | TIMER1 COMPA | Timer/Counter1 Compare Match A |
| 19 | \$0024 | TIMER1 COMPB | Timer/Counter1 Compare Match B |
| 20 | \$0026 | TIMER1 COMPC | Timer/Counter1 Compare Match C |
| 21 | \$0028 | TIMER1 OVF | Timer/Counter1 Overflow |
| 22 | \$002A | TIMER0 COMPA | Timer/Counter0 Compare Match A |
| 23 | \$002C | TIMER0 COMPB | Timer/Counter0 Compare match B |
| 24 | \$002E | TIMERO OVF | Timer/Counter0 Overflow |
| 25 | \$0030 | SPI, STC | SPI Serial Transfer Complete |
| 26 | \$0032 | USARTO RX | USARTO Rx Complete |
| 27 | \$0034 | USARTO UDRE | USARTO Data Register Empty |
| 28 | \$0036 | USARTO TX | USART0 Tx Complete |
| 29 | \$0038 | ANALOG COMP | Analog Comparator |
| 30 | \$003A | ADC | ADC Conversion Complete |

Table 14-1. Reset and Interrupt Vectors (Continued)

| Vector No. | Program Address ${ }^{(2)}$ | Source | Interrupt Definition |
| :---: | :---: | :---: | :---: |
| 31 | \$003C | EE READY | EEPROM Ready |
| 32 | \$003E | TIMER3 CAPT | Timer/Counter3 Capture Event |
| 33 | \$0040 | TIMER3 COMPA | Timer/Counter3 Compare Match A |
| 34 | \$0042 | TIMER3 COMPB | Timer/Counter3 Compare Match B |
| 35 | \$0044 | TIMER3 COMPC | Timer/Counter3 Compare Match C |
| 36 | \$0046 | TIMER3 OVF | Timer/Counter3 Overflow |
| 37 | \$0048 | USART1 RX | USART1 Rx Complete |
| 38 | \$004A | USART1 UDRE | USART1 Data Register Empty |
| 39 | \$004C | USART1 TX | USART1 Tx Complete |
| 40 | \$004E | TWI | 2-wire Serial Interface |
| 41 | \$0050 | SPM READY | Store Program Memory Ready |
| 42 | \$0052 ${ }^{(3)}$ | TIMER4 CAPT | Timer/Counter4 Capture Event |
| 43 | \$0054 | TIMER4 COMPA | Timer/Counter4 Compare Match A |
| 44 | \$0056 | TIMER4 COMPB | Timer/Counter4 Compare Match B |
| 45 | \$0058 | TIMER4 COMPC | Timer/Counter4 Compare Match C |
| 46 | \$005A | TIMER4 OVF | Timer/Counter4 Overflow |
| 47 | \$005C ${ }^{(3)}$ | TIMER5 CAPT | Timer/Counter5 Capture Event |
| 48 | \$005E | TIMER5 COMPA | Timer/Counter5 Compare Match A |
| 49 | \$0060 | TIMER5 COMPB | Timer/Counter5 Compare Match B |
| 50 | \$0062 | TIMER5 COMPC | Timer/Counter5 Compare Match C |
| 51 | \$0064 | TIMER5 OVF | Timer/Counter5 Overflow |
| 52 | \$0066 ${ }^{(3)}$ | USART2 RX | USART2 Rx Complete |
| 53 | \$0068 ${ }^{(3)}$ | USART2 UDRE | USART2 Data Register Empty |
| 54 | \$006A ${ }^{(3)}$ | USART2 TX | USART2 Tx Complete |
| 55 | \$006C ${ }^{(3)}$ | USART3 RX | USART3 Rx Complete |
| 56 | \$006E ${ }^{(3)}$ | USART3 UDRE | USART3 Data Register Empty |
| 57 | \$0070 ${ }^{(3)}$ | USART3 TX | USART3 Tx Complete |

Notes: 1. When the BOOTRST Fuse is programmed, the device will jump to the Boot Loader address at reset, see "Memory Programming" on page 325.
2. When the IVSEL bit in MCUCR is set, Interrupt Vectors will be moved to the start of the Boot Flash Section. The address of each Interrupt Vector will then be the address in this table added to the start address of the Boot Flash Section.
3. Only available in ATmega640/1280/2560.

### 14.2 Reset and Interrupt Vector placement

Table 14-2 on page 103 shows Reset and Interrupt Vectors placement for the various combinations of BOOTRST and IVSEL settings. If the program never enables an interrupt source, the Interrupt Vectors are not used, and regular program code can be placed at these locations. This is also the case if the Reset Vector is in the Application section while the Interrupt Vectors are in the Boot section or vice versa.

Table 14-2. Reset and Interrupt Vectors Placement ${ }^{(1)}$

| BOOTRST | IVSEL | Reset Address | Interrupt Vectors Start Address |
| :---: | :---: | :---: | :---: |
| 1 | 0 | $0 \times 0000$ | $0 \times 0002$ |
| 1 | 1 | $0 \times 0000$ | Boot Reset Address $+0 \times 0002$ |
| 0 | 0 | Boot Reset Address | $0 \times 0002$ |
| 0 | 1 | Boot Reset Address | Boot Reset Address $+0 \times 0002$ |

Note: 1. The Boot Reset Address is shown in Table 29-7 on page 320 through Table 29-15 on page 322. For the BOOTRST Fuse " 1 " means unprogrammed while " 0 " means programmed.
The most typical and general program setup for the Reset and Interrupt Vector Addresses in ATmega640/1280/1281/2560/2561 is:

| Addre <br> ss | Label | Code |  | Comments |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0x000 |  | jmp | RESET |  | Reset Handler |
| 0 |  |  |  |  |  |
| 0x000 |  | jmp | INTO | ; | IRQO Handler |
| 2 |  |  |  |  |  |
| 0x000 |  | jmp | INT1 | ; | IRQ1 Handler |
| 4 |  |  |  |  |  |
| 0x000 |  | jmp | INT2 | ; | IRQ2 Handler |
| 6 |  |  |  |  |  |
| 0x000 |  | jmp | INT3 | ; | IRQ3 Handler |
| 8 |  |  |  |  |  |
| 0x000 |  | jmp | INT4 | ; | IRQ4 Handler |
| A |  |  |  |  |  |
| 0x000 |  | jmp | INT5 | ; | IRQ5 Handler |
| C |  |  |  |  |  |
| 0x000 |  | jmp | INT6 | ; | IRQ6 Handler |
| E |  |  |  |  |  |
| 0x001 |  | jmp | INT7 | ; | IRQ7 Handler |
| 0 |  |  |  |  |  |
| 0x001 |  | jmp | PCINT0 | ; | PCINTO Handler |
| 2 |  |  |  |  |  |
| 0x001 |  | jmp | PCINT1 | ; | PCINT1 Handler |
| 4 |  |  |  |  |  |
| 0x001 |  | jmp | PCINT2 | ; | PCINT2 Handler |
| 6 |  |  |  |  |  |
| 0x001 |  | jmp | WDT | ; | Watchdog Timeout Handler |
| 8 |  |  |  |  |  |
| 0x001 |  | jmp | TIM2_COMPA | ; | Timer2 CompareA Handler |
| A |  |  |  |  |  |
| 0x001 |  | jmp | TIM2_COMPB | ; | Timer2 CompareB Handler |
| C |  |  |  |  |  |
| 0x001 |  | jmp | TIM2_OVF | ; | Timer2 Overflow Handler |
| E |  |  |  |  |  |
| 0x002 |  | jmp | TIM1_CAPT |  | Timer1 Capture Handler |
| 0 |  |  |  |  |  |
| 0x002 |  | jmp | TIM1_COMPA |  | Timer1 CompareA Handler |
| 2 |  |  |  |  |  |
| 0x002 |  | jmp | TIM1_COMPB |  | Timer1 CompareB Handler |
| 4 |  |  |  |  |  |
| 0x002 |  | jmp | TIM1_COMPC |  | Timer1 CompareC Handler |
| 6 |  |  |  |  |  |


| $\begin{gathered} 0 \times 002 \\ 8 \end{gathered}$ | jmp | TIM1_OVF | ; | Timer1 | Overflow Handler |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{gathered} 0 \times 002 \\ \text { A } \end{gathered}$ | jmp | TIMO_COMPA | ; | Timer0 | CompareA Handler |
| $\begin{gathered} 0 \times 002 \\ C \end{gathered}$ | jmp | TIMO_COMPB | ; | Timer0 | CompareB Handler |
| $\begin{gathered} 0 \times 002 \\ E \end{gathered}$ | jmp | TIMO_OVF | ; | Timer0 | Overflow Handler |
| $\begin{gathered} 0 \times 003 \\ 0 \end{gathered}$ | jmp | SPI_STC | ; | SPI Tran | nsfer Complete Handler |
| $\begin{gathered} 0 \times 003 \\ 2 \end{gathered}$ | jmp | USART0_RXC | ; | USART0 | RX Complete Handler |
| $\begin{gathered} 0 \times 003 \\ 4 \end{gathered}$ | jmp | USART0_UDRE | ; | USART0 | UDR Empty Handler |
| $\begin{gathered} 0 \times 003 \\ 6 \end{gathered}$ | jmp | USART0_TXC | ; | USART0 | TX Complete Handler |
| $\begin{gathered} 0 \times 003 \\ 8 \end{gathered}$ | jmp | ANA_COMP | ; | Analog | Comparator Handler |
| $\begin{gathered} 0 \times 003 \\ \mathrm{~A} \end{gathered}$ | jmp | ADC | ; | ADC Con Handler | version Complete |
| $\begin{gathered} 0 \times 003 \\ C \end{gathered}$ | jmp | EE_RDY | ; | EEPROM | Ready Handler |
| $\begin{gathered} 0 \times 003 \\ E \end{gathered}$ | jmp | TIM3_CAPT | ; | Timer3 | Capture Handler |
| $\begin{gathered} 0 \times 004 \\ 0 \end{gathered}$ | jmp | TIM3_COMPA | ; | Timer3 | CompareA Handler |
| $\begin{gathered} 0 \times 004 \\ 2 \end{gathered}$ | jmp | TIM3_COMPB | ; | Timer3 | CompareB Handler |
| $\begin{gathered} 0 \times 004 \\ 4 \end{gathered}$ | jmp | TIM3_COMPC | ; | Timer3 | CompareC Handler |
| $\begin{gathered} 0 \times 004 \\ 6 \end{gathered}$ | jmp | TIM3_OVF | ; | Timer3 | Overflow Handler |
| $\begin{gathered} 0 \times 004 \\ 8 \end{gathered}$ | jmp | USART1_RXC | ; | USART1 | RX Complete Handler |
| $\begin{gathered} 0 \times 004 \\ \mathrm{~A} \end{gathered}$ | jmp | USART1_UDRE | ; | USART1 | UDR Empty Handler |
| $\begin{gathered} 0 \times 004 \\ C \end{gathered}$ | jmp | USART1_TXC | ; | USART1 | TX Complete Handler |
| $\begin{gathered} 0 \times 004 \\ E \end{gathered}$ | jmp | TWI | ; | 2-wire | Serial Handler |
| $\begin{gathered} 0 \times 005 \\ 0 \end{gathered}$ | jmp | SPM_RDY | ; | SPM Read | dy Handler |
| $\begin{gathered} 0 \times 005 \\ 2 \end{gathered}$ | jmp | TIM4_CAPT | ; | Timer4 | Capture Handler |
| $\begin{gathered} 0 \times 005 \\ 4 \end{gathered}$ | jmp | TIM4_COMPA | ; | Timer4 | CompareA Handler |
| $\begin{gathered} 0 \times 005 \\ 6 \end{gathered}$ | jmp | TIM4_COMPB | ; | Timer4 | CompareB Handler |
| $\begin{gathered} 0 \times 005 \\ 8 \end{gathered}$ | jmp | TIM4_COMPC | ; | Timer4 | CompareC Handler |
| $\begin{gathered} 0 \times 005 \\ \text { A } \end{gathered}$ | jmp | TIM4_OVF | ; | Timer4 | Overflow Handler |
| $\begin{gathered} 0 \times 005 \\ C \end{gathered}$ | jmp | TIM5_CAPT | ; | Timer5 | Capture Handler |
| $\begin{gathered} 0 \times 005 \\ E \end{gathered}$ | jmp | TIM5_COMPA | ; | Timer5 | CompareA Handler |
| $\begin{gathered} 0 \times 006 \\ 0 \end{gathered}$ | jmp | TIM5_COMPB | ; | Timer5 | CompareB Handler |



When the BOOTRST Fuse is unprogrammed, the Boot section size set to 8Kbytes and the IVSEL bit in the MCUCR Register is set before any interrupts are enabled, the most typical and general program setup for the Reset and Interrupt Vector Addresses is:


When the BOOTRST Fuse is programmed and the Boot section size set to 8 Kbytes , the most typical and general program setup for the Reset and Interrupt Vector Addresses is:

```
Address Labels Code Comments
.org 0x0002
\begin{tabular}{llll}
\(0 x 00002\) & jmp & EXT_INT0 & ; IRQ0 Handler \\
\(0 \times 00004\) & jmp & EXT_INT1 & ; IRQ1 Handler \\
\(\ldots\). & \(\cdots\) & \(\cdots\) & \(;\) \\
\(0 x 00070\) & jmp & USART3_TXC & ; USART3 TX Complete Handler
\end{tabular}
. org 0x1F000
0x1F000 RESET: ldi r16,high(RAMEND); Main program start 
0x1F002 ldi r16,low(RAMEND)
0x1F003 out SPL,r16
0x1F004 sei ; Enable interrupts
0x1F005 <instr> xxx
```

When the BOOTRST Fuse is programmed, the Boot section size set to 8Kbytes and the IVSEL bit in the MCUCR Register is set before any interrupts are enabled, the most typical and general program setup for the Reset and Interrupt Vector Addresses is:

| Address Labels | Code |  | Comments |
| :---: | :---: | :---: | :---: |
| ; |  |  |  |
| . org 0x1F000 |  |  |  |
| $0 \times 1 F 000$ | jmp | RESET | ; Reset handler |
| 0x1F002 | jmp | EXT_INT0 | ; IRQO Handler |
| 0x1F004 | jmp | EXT_INT1 | ; IRQ1 Handler |
| $\cdots$ |  |  | ; |
| $0 \times 1 F 070$ | jmp | USART3_TXC | ; USART3 TX Complete Handler |
| ; |  |  |  |
| 0x1F072 RESET: | $1 d i$ | r16,high (RAMEND) | ; Main program start |
| 0x1F073 | out | SPH,r16 | ; Set Stack Pointer to top of RAM |
| 0x1F074 | 1di | r16, low (RAMEND) |  |
| 0x1F075 | out | SPL,r16 |  |
| 0x1F076 | sei |  | ; Enable interrupts |
| 0x1F077 | <instr> | xxx |  |

### 14.3 Moving Interrupts Between Application and Boot Section

The MCU Control Register controls the placement of the Interrupt Vector table, see Code Example below. For more details, see "Reset and Interrupt Handling" on page 17.

```
Assembly Code Example
    Move_interrupts:
    ; Get MCUCR
    in r16, MCUCR
    mov r17, r16
    ; Enable change of Interrupt Vectors
    ori r16, (1<<IVCE)
    out MCUCR, r16
    ; Move interrupts to Boot Flash section
    ori r16, (1<<IVSEL)
    out MCUCR, r17
    ret
```

C Code Example
void Move_interrupts (void)
\{
uchar temp;
/* Get MCUCR*/
temp = MCUCR;
/* Enable change of Interrupt Vectors */
MCUCR $=$ temp $\mid(1 \ll$ IVCE $) ;$
/* Move interrupts to Boot Flash section */
MCUCR $=$ temp| $(1 \ll$ IVSEL $) ;$
\}

### 14.4 Register Description

### 14.4.1 MCUCR - MCU Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | MCUCR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x35 (0x55) | JTD | - | - | PUD | - | - | IVSEL | IVCE |  |
| Read/Write | R/W | R | R | R/W | R | R | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 1 - IVSEL: Interrupt Vector Select

When the IVSEL bit is cleared (zero), the Interrupt Vectors are placed at the start of the Flash memory. When this bit is set (one), the Interrupt Vectors are moved to the beginning of the Boot Loader section of the Flash. The actual address of the start of the Boot Flash Section is determined by the BOOTSZ Fuses. Refer to the section "Memory Programming" on page 325 for details. To avoid unintentional changes of Interrupt Vector tables, a special write procedure must be followed to change the IVSEL bit (see "Moving Interrupts Between Application and Boot Section" on page 107):

1. Write the Interrupt Vector Change Enable (IVCE) bit to one.
2. Within four cycles, write the desired value to IVSEL while writing a zero to IVCE.

Interrupts will automatically be disabled while this sequence is executed. Interrupts are disabled in the cycle IVCE is set, and they remain disabled until after the instruction following the write to IVSEL. If IVSEL is not written, interrupts remain disabled for four cycles. The l-bit in the Status Register is unaffected by the automatic disabling.

Note: If Interrupt Vectors are placed in the Boot Loader section and Boot Lock bit BLB02 is programmed, interrupts are disabled while executing from the Application section. If Interrupt Vectors are placed in the Application section and Boot Lock bit BLB12 is programed, interrupts are disabled while executing from the Boot Loader section. Refer to the section "Memory Programming" on page 325 for details on Boot Lock bits.

## - Bit 0 - IVCE: Interrupt Vector Change Enable

The IVCE bit must be written to logic one to enable change of the IVSEL bit. IVCE is cleared by hardware four cycles after it is written or when IVSEL is written. Setting the IVCE bit will disable interrupts, as explained in the IVSEL description.

## 15. External Interrupts

The External Interrupts are triggered by the INT7:0 pin or any of the PCINT23:0 pins. Observe that, if enabled, the interrupts will trigger even if the INT7:0 or PCINT23:0 pins are configured as outputs. This feature provides a way of generating a software interrupt.

The Pin change interrupt PCl 2 will trigger if any enabled $\mathrm{PCINT23:16}$ pin toggles, Pin change interrupt PCl 1 if any enabled PCINT15:8 toggles and Pin change interrupts PCIO will trigger if any enabled PCINT7:0 pin toggles. PCMSK2, PCMSK1 and PCMSK0 Registers control which pins contribute to the pin change interrupts. Pin change interrupts on PCINT23:0 are detected asynchronously. This implies that these interrupts can be used for waking the part also from sleep modes other than Idle mode.

The External Interrupts can be triggered by a falling or rising edge or a low level. This is set up as indicated in the specification for the External Interrupt Control Registers - EICRA (INT3:0) and EICRB (INT7:4). When the external interrupt is enabled and is configured as level triggered, the interrupt will trigger as long as the pin is held low.
Low level interrupts and the edge interrupt on INT3:0 are detected asynchronously. This implies that these interrupts can be used for waking the part also from sleep modes other than Idle mode. The I/O clock is halted in all sleep modes except Idle mode.

Note that if a level triggered interrupt is used for wake-up from Power-down, the required level must be held long enough for the MCU to complete the wake-up to trigger the level interrupt. If the level disappears before the end of the Start-up Time, the MCU will still wake up, but no interrupt will be generated. The start-up time is defined by the SUT and CKSEL Fuses as described in "System Clock and Clock Options" on page 39.

### 15.1 Pin Change Interrupt Timing

An example of timing of a pin change interrupt is shown in Figure 15-1.
Figure 15-1. Normal pin change interrupt.


### 15.2 Register Description

### 15.2.1 EICRA - External Interrupt Control Register A

The External Interrupt Control Register A contains control bits for interrupt sense control.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | EICRA |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x69) | ISC31 | ISC30 | ISC21 | ISC20 | ISC11 | ISC10 | ISC01 | ISC00 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7:0 - ISC31, ISC30 - ISC00, ISC00: External Interrupt 3-0 Sense Control Bits

The External Interrupts 3-0 are activated by the external pins INT3:0 if the SREG I-flag and the corresponding interrupt mask in the EIMSK is set. The level and edges on the external pins that activate the interrupts are defined in Table 15-1. Edges on INT3:0 are registered asynchronously. Pulses on INT3:0 pins wider than the minimum pulse width given in Table 15-2 will generate an interrupt. Shorter pulses are not guaranteed to generate an interrupt. If low level interrupt is selected, the low level must be held until the completion of the currently executing instruction to generate an interrupt. If enabled, a level triggered interrupt will generate an interrupt request as long as the pin is held low. When changing the ISCn bit, an interrupt can occur. Therefore, it is recommended to first disable INTn by clearing its Interrupt Enable bit in the EIMSK Register. Then, the ISCn bit can be changed. Finally, the INTn interrupt flag should be cleared by writing a logical one to its Interrupt Flag bit (INTFn) in the EIFR Register before the interrupt is re-enabled.

Table 15-1. Interrupt Sense Control ${ }^{(1)}$

| ISCn1 | ISCn0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | The low level of INTn generates an interrupt request |
| 0 | 1 | Any edge of INTn generates asynchronously an interrupt request |
| 1 | 0 | The falling edge of INTn generates asynchronously an interrupt request |
| 1 | 1 | The rising edge of INTn generates asynchronously an interrupt request |

Note: 1. $\mathrm{n}=3,2$, 1 or 0 .
When changing the ISCn1/ISCn0 bits, the interrupt must be disabled by clearing its Interrupt Enable bit in the EIMSK Register. Otherwise an interrupt can occur when the bits are changed.

Table 15-2. Asynchronous External Interrupt Characteristics

| Symbol | Parameter | Condition | Min. | Typ. | Max. |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Units |  |  |  |  |  |
| $\mathrm{t}_{\mathrm{INT}}$ | Minimum pulse width for asynchronous external interrupt |  |  | 50 |  |
| n |  |  |  |  |  |

### 15.2.2 EICRB - External Interrupt Control Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x6A) | ISC71 | ISC70 | ISC61 | ISC60 | ISC51 | ISC50 | ISC41 | ISC40 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

EICRB

## - Bits 7:0 - ISC71, ISC70-ISC41, ISC40: External Interrupt 7-4 Sense Control Bits

The External Interrupts 7-4 are activated by the external pins INT7:4 if the SREG I-flag and the corresponding interrupt mask in the EIMSK is set. The level and edges on the external pins that activate the interrupts are defined in Table 15-3. The value on the INT7:4 pins are sampled before detecting edges. If edge or toggle interrupt is selected, pulses that last longer than one clock period will generate an interrupt. Shorter pulses are not guaranteed to generate an interrupt. Observe that CPU clock frequency can be lower than the XTAL frequency if the XTAL divider is enabled. If low level interrupt is selected, the low level must be held until the completion of the currently executing instruction to generate an interrupt. If enabled, a level triggered interrupt will generate an interrupt request as long as the pin is held low.

Table 15-3. Interrupt Sense Control ${ }^{(1)}$

| ISCn1 | ISCn0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | The low level of INTn generates an interrupt request |
| 0 | 1 | Any logical change on INTn generates an interrupt request |
| 1 | 0 | The falling edge between two samples of INTn generates an interrupt request |
| 1 | 1 | The rising edge between two samples of INTn generates an interrupt request |

Note: 1. $\mathrm{n}=7,6,5$ or 4 .
When changing the ISCn1/ISCn0 bits, the interrupt must be disabled by clearing its Interrupt Enable bit in the EIMSK Register. Otherwise an interrupt can occur when the bits are changed.

### 15.2.3 EIMSK - External Interrupt Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x1D (0x3D) | INT7 | INT6 | INT5 | INT4 | INT3 | INT2 | INT1 | INTO |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## - Bits 7:0-INT7:0: External Interrupt Request 7-0 Enable

When an INT7:0 bit is written to one and the I-bit in the Status Register (SREG) is set (one), the corresponding external pin interrupt is enabled. The Interrupt Sense Control bits in the External Interrupt Control Registers EICRA and EICRB - defines whether the external interrupt is activated on rising or falling edge or level sensed. Activity on any of these pins will trigger an interrupt request even if the pin is enabled as an output. This provides a way of generating a software interrupt.

### 15.2.4 EIFR - External Interrupt Flag Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | EIFR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x1C (0x3C) | INTF7 | INTF6 | INTF5 | INTF4 | INTF3 | INTF2 | INTF1 | IINTF0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7:0 - INTF7:0: External Interrupt Flags 7-0

When an edge or logic change on the INT7:0 pin triggers an interrupt request, INTF7:0 becomes set (one). If the Ibit in SREG and the corresponding interrupt enable bit, INT7:0 in EIMSK, are set (one), the MCU will jump to the interrupt vector. The flag is cleared when the interrupt routine is executed. Alternatively, the flag can be cleared by writing a logical one to it. These flags are always cleared when INT7:0 are configured as level interrupt. Note that when entering sleep mode with the INT3:0 interrupts disabled, the input buffers on these pins will be disabled. This may cause a logic change in internal signals which will set the INTF3:0 flags. See "Digital Input Enable and Sleep Modes" on page 71 for more information.

### 15.2.5 PCICR - Pin Change Interrupt Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PCICR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x68) | - | - | - | - | - | PCIE2 | PCIE1 | PCIE0 |  |
| Read/Write | R | R | R | R | R | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 2 - PCIE2: Pin Change Interrupt Enable 1

When the PCIE2 bit is set (one) and the I-bit in the Status Register (SREG) is set (one), pin change interrupt 2 is enabled. Any change on any enabled PCINT23:16 pin will cause an interrupt. The corresponding interrupt of Pin Change Interrupt Request is executed from the PCI2 Interrupt Vector. PCINT23:16 pins are enabled individually by the PCMSK2 Register.

## - Bit 1 - PCIE1: Pin Change Interrupt Enable 1

When the PCIE1 bit is set (one) and the I-bit in the Status Register (SREG) is set (one), pin change interrupt 1 is enabled. Any change on any enabled PCINT15:8 pin will cause an interrupt. The corresponding interrupt of Pin Change Interrupt Request is executed from the PCI1 Interrupt Vector. PCINT15:8 pins are enabled individually by the PCMSK1 Register.

## - Bit 0 - PCIE0: Pin Change Interrupt Enable 0

When the PCIEO bit is set (one) and the I-bit in the Status Register (SREG) is set (one), pin change interrupt 0 is enabled. Any change on any enabled PCINT7:0 pin will cause an interrupt. The corresponding interrupt of Pin Change Interrupt Request is executed from the PCIO Interrupt Vector. PCINT7:0 pins are enabled individually by the PCMSKO Register.

### 15.2.6 PCIFR - Pin Change Interrupt Flag Register



- Bit 2 - PCIF2: Pin Change Interrupt Flag 1

When a logic change on any PCINT23:16 pin triggers an interrupt request, PCIF2 becomes set (one). If the I-bit in SREG and the PCIE2 bit in PCICR are set (one), the MCU will jump to the corresponding Interrupt Vector. The flag is cleared when the interrupt routine is executed. Alternatively, the flag can be cleared by writing a logical one to it.

## - Bit 1 - PCIF1: Pin Change Interrupt Flag 1

When a logic change on any PCINT15:8 pin triggers an interrupt request, PCIF1 becomes set (one). If the l-bit in SREG and the PCIE1 bit in PCICR are set (one), the MCU will jump to the corresponding Interrupt Vector. The flag is cleared when the interrupt routine is executed. Alternatively, the flag can be cleared by writing a logical one to it.

## - Bit 0 - PCIFO: Pin Change Interrupt Flag 0

When a logic change on any PCINT7:0 pin triggers an interrupt request, PCIFO becomes set (one). If the l-bit in SREG and the PCIEO bit in PCICR are set (one), the MCU will jump to the corresponding Interrupt Vector. The flag is cleared when the interrupt routine is executed. Alternatively, the flag can be cleared by writing a logical one to it.

### 15.2.7 PCMSK2 - Pin Change Mask Register 2

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PCMSK2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x6D) | PCINT23 | PCINT22 | PCINT21 | PCINT20 | PCINT19 | PCINT18 | PCINT17 | PCINT16 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

- Bit 7:0 - PCINT23:16: Pin Change Enable Mask 23:16

Each PCINT23:16-bit selects whether pin change interrupt is enabled on the corresponding I/O pin. If PCINT23:16 is set and the PCIE2 bit in PCICR is set, pin change interrupt is enabled on the corresponding I/O pin. If PCINT23:16 is cleared, pin change interrupt on the corresponding I/O pin is disabled.

### 15.2.8 PCMSK1 - Pin Change Mask Register 1

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PCMSK1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x6C) | PCINT15 | PCINT14 | PCINT13 | PCINT12 | PCINT11 | PCINT10 | PCINT9 | PCINT8 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7:0 - PCINT15:8: Pin Change Enable Mask 15:8

Each PCINT15:8-bit selects whether pin change interrupt is enabled on the corresponding I/O pin. If PCINT15:8 is set and the PCIE1 bit in PCICR is set, pin change interrupt is enabled on the corresponding I/O pin. If PCINT15:8 is cleared, pin change interrupt on the corresponding I/O pin is disabled.

### 15.2.9 PCMSKO - Pin Change Mask Register 0

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | PCMSK0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x6B) | PCINT7 | PCINT6 | PCINT5 | PCINT4 | PCINT3 | PCINT2 | PCINT1 | PCINT0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

- Bit 7:0 - PCINT7:0: Pin Change Enable Mask 7:0

Each PCINT7:0 bit selects whether pin change interrupt is enabled on the corresponding I/O pin. If PCINT7:0 is set and the PCIEO bit in PCICR is set, pin change interrupt is enabled on the corresponding I/O pin. If PCINT7:0 is cleared, pin change interrupt on the corresponding I/O pin is disabled.

## 16. 8-bit Timer/Counter0 with PWM

### 16.1 Features

- Two Independent Output Compare Units
- Double Buffered Output Compare Registers
- Clear Timer on Compare Match (Auto Reload)
- Glitch Free, Phase Correct Pulse Width Modulator (PWM)
- Variable PWM Period
- Frequency Generator
- Three Independent Interrupt Sources (TOVO, OCFOA, and OCFOB)


### 16.2 Overview

Timer/Counter0 is a general purpose 8-bit Timer/Counter module, with two independent Output Compare Units, and with PWM support. It allows accurate program execution timing (event management) and wave generation.

A simplified block diagram of the 8 -bit Timer/Counter is shown in Figure 16-1. For the actual placement of I/O pins, refer to "TQFP-pinout ATmega640/1280/2560" on page 2. CPU accessible I/O Registers, including I/O bits and I/O pins, are shown in bold. The device-specific I/O Register and bit locations are listed in the "Register Description" on page 126.

Figure 16-1. 8-bit Timer/Counter Block Diagram


### 16.2.1 Registers

The Timer/Counter (TCNTO) and Output Compare Registers (OCROA and OCROB) are 8-bit registers. Interrupt request (abbreviated to Int.Req. in the figure) signals are all visible in the Timer Interrupt Flag Register (TIFRO). All interrupts are individually masked with the Timer Interrupt Mask Register (TIMSK0). TIFRO and TIMSK0 are not shown in the figure.
The Timer/Counter can be clocked internally, via the prescaler, or by an external clock source on the T0 pin. The Clock Select logic block controls which clock source and edge the Timer/Counter uses to increment (or decrement) its value. The Timer/Counter is inactive when no clock source is selected. The output from the Clock Select logic is referred to as the timer clock ( $\mathrm{Clk}_{\mathrm{T}}$ ).

The double buffered Output Compare Registers (OCROA and OCROB) are compared with the Timer/Counter value at all times. The result of the compare can be used by the Waveform Generator to generate a PWM or variable frequency output on the Output Compare pins (OCOA and OCOB). See "Output Compare Unit" on page 117. for details. The Compare Match event will also set the Compare Flag (OCFOA or OCFOB) which can be used to generate an Output Compare interrupt request.

### 16.2.2 Definitions

Many register and bit references in this section are written in general form. A lower case " n " replaces the Timer/Counter number, in this case 0. A lower case " $x$ " replaces the Output Compare Unit, in this case Compare Unit A or Compare Unit B. However, when using the register or bit defines in a program, the precise form must be used, that is, TCNTO for accessing Timer/Counter0 counter value and so on.
The definitions in Table 16-1 are also used extensively throughout the document.
Table 16-1. Definitions

| BOTTOM | The counter reaches the BOTTOM when it becomes 0x00. |
| :--- | :--- |
| MAX | The counter reaches its MAXimum when it becomes 0xFF (decimal 255). <br> TOP |
| The counter reaches the TOP when it becomes equal to the highest value in the count <br> sequence. The TOP value can be assigned to be the fixed value 0xFF (MAX) or the value stored <br> in the OCROA Register. The assignment is dependent on the mode of operation. |  |

### 16.3 Timer/Counter Clock Sources

The Timer/Counter can be clocked by an internal or an external clock source. The clock source is selected by the Clock Select logic which is controlled by the Clock Select (CS02:0) bits located in the Timer/Counter Control Register (TCCROB). For details on clock sources and prescaler, see "Timer/Counter 0, 1, 3, 4, and 5 Prescaler" on page 164.

### 16.4 Counter Unit

The main part of the 8 -bit Timer/Counter is the programmable bi-directional counter unit. Figure 16-2 on page 117 shows a block diagram of the counter and its surroundings.

Figure 16-2. Counter Unit Block Diagram


Signal description (internal signals):

| count | Increment or decrement TCNTO by 1. |
| :--- | :--- |
| direction | Select between increment and decrement. |
| clear | Clear TCNTO (set all bits to zero). |
| $\mathbf{c l k}_{\mathbf{T n}_{\mathrm{n}}}$ | Timer/Counter clock, referred to as clk ${ }_{\text {TO }}$ in the following. |
| top | Signalize that TCNT0 has reached maximum value. |
| bottom | Signalize that TCNTO has reached minimum value (zero). |

Depending of the mode of operation used, the counter is cleared, incremented, or decremented at each timer clock $\left(\mathrm{clk}_{\mathrm{T} 0}\right) . \mathrm{clk}_{\mathrm{T} 0}$ can be generated from an external or internal clock source, selected by the Clock Select bits (CS02:0). When no clock source is selected $(\operatorname{CSO2:0}=0)$ the timer is stopped. However, the TCNT0 value can be accessed by the CPU, regardless of whether $\mathrm{Cl}_{\mathrm{T} 0}$ is present or not. A CPU write overrides (has priority over) all counter clear or count operations.

The counting sequence is determined by the setting of the WGM01 and WGM00 bits located in the Timer/Counter Control Register (TCCROA) and the WGM02 bit located in the Timer/Counter Control Register B (TCCROB). There are close connections between how the counter behaves (counts) and how waveforms are generated on the Output Compare outputs OCOA and OCOB. For more details about advanced counting sequences and waveform generation, see "Modes of Operation" on page 120.

The Timer/Counter Overflow Flag (TOVO) is set according to the mode of operation selected by the WGM02:0 bits. TOVO can be used for generating a CPU interrupt.

### 16.5 Output Compare Unit

The 8-bit comparator continuously compares TCNT0 with the Output Compare Registers (OCROA and OCROB). Whenever TCNTO equals OCROA or OCROB, the comparator signals a match. A match will set the Output Compare Flag (OCFOA or OCFOB) at the next timer clock cycle. If the corresponding interrupt is enabled, the Output Compare Flag generates an Output Compare interrupt. The Output Compare Flag is automatically cleared when the interrupt is executed. Alternatively, the flag can be cleared by software by writing a logical one to its I/O bit location. The Waveform Generator uses the match signal to generate an output according to operating mode set by the WGM02:0 bits and Compare Output mode (COM0x1:0) bits. The maximum and bottom signals are used by the Waveform Generator for handling the special cases of the extreme values in some modes of operation ("Modes of Operation" on page 120).
Figure $16-3$ on page 118 shows a block diagram of the Output Compare unit.

Figure 16-3. Output Compare Unit, Block Diagram


The OCROx Registers are double buffered when using any of the Pulse Width Modulation (PWM) modes. For the normal and Clear Timer on Compare (CTC) modes of operation, the double buffering is disabled. The double buffering synchronizes the update of the OCROx Compare Registers to either top or bottom of the counting sequence. The synchronization prevents the occurrence of odd-length, non-symmetrical PWM pulses, thereby making the output glitch-free.

The OCROx Register access may seem complex, but this is not case. When the double buffering is enabled, the CPU has access to the OCROx Buffer Register, and if double buffering is disabled the CPU will access the OCR0x directly.

### 16.5.1 Force Output Compare

In non-PWM waveform generation modes, the match output of the comparator can be forced by writing a one to the Force Output Compare (FOCOx) bit. Forcing Compare Match will not set the OCFOx Flag or reload/clear the timer, but the OC0x pin will be updated as if a real Compare Match had occurred (the COM0x1:0 bits settings define whether the OCOx pin is set, cleared or toggled).

### 16.5.2 Compare Match Blocking by TCNTO Write

All CPU write operations to the TCNTO Register will block any Compare Match that occur in the next timer clock cycle, even when the timer is stopped. This feature allows OCROx to be initialized to the same value as TCNTO without triggering an interrupt when the Timer/Counter clock is enabled.

### 16.5.3 Using the Output Compare Unit

Since writing TCNTO in any mode of operation will block all Compare Matches for one timer clock cycle, there are risks involved when changing TCNTO when using the Output Compare Unit, independently of whether the Timer/Counter is running or not. If the value written to TCNT0 equals the OCROx value, the Compare Match will be missed, resulting in incorrect waveform generation. Similarly, do not write the TCNTO value equal to BOTTOM when the counter is down-counting.

The setup of the OCOx should be performed before setting the Data Direction Register for the port pin to output. The easiest way of setting the OC0x value is to use the Force Output Compare (FOCOx) strobe bits in Normal mode. The OC0x Registers keep their values even when changing between Waveform Generation modes.

Be aware that the COM0x1:0 bits are not double buffered together with the compare value. Changing the COM0x1:0 bits will take effect immediately.

### 16.6 Compare Match Output Unit

The Compare Output mode (COM0x1:0) bits have two functions. The Waveform Generator uses the COM0x1:0 bits for defining the Output Compare (OC0x) state at the next Compare Match. Also, the COM0x1:0 bits control the OC0x pin output source. Figure 16-4 shows a simplified schematic of the logic affected by the COM0x1:0 bit setting. The I/O Registers, I/O bits, and I/O pins in the figure are shown in bold. Only the parts of the general I/O Port Control Registers (DDR and PORT) that are affected by the COM0x1:0 bits are shown. When referring to the OC0x state, the reference is for the internal OC0x Register, not the OC0x pin. If a system reset occur, the OC0x Register is reset to " 0 ".

Figure 16-4. Compare Match Output Unit, Schematic


The general I/O port function is overridden by the Output Compare (OCOx) from the Waveform Generator if either of the COM0x1:0 bits are set. However, the OC0x pin direction (input or output) is still controlled by the Data Direction Register (DDR) for the port pin. The Data Direction Register bit for the OC0x pin (DDR_OC0x) must be set as output before the OC0x value is visible on the pin. The port override function is independent of the Waveform Generation mode.

The design of the Output Compare pin logic allows initialization of the OC0x state before the output is enabled. Note that some COM0x1:0 bit settings are reserved for certain modes of operation. See "Register Description" on page 126.

### 16.6.1 Compare Output Mode and Waveform Generation

The Waveform Generator uses the COM0x1:0 bits differently in Normal, CTC, and PWM modes. For all modes, setting the COM0x1:0 = 0 tells the Waveform Generator that no action on the OC0x Register is to be performed on the next Compare Match. For compare output actions in the non-PWM modes refer to Table 16-2 on page 126. For fast PWM mode, refer to Table 16-3 on page 126, and for phase correct PWM refer to Table 16-4 on page 127.

A change of the COM0x1:0 bits state will have effect at the first Compare Match after the bits are written. For nonPWM modes, the action can be forced to have immediate effect by using the FOCOx strobe bits.

### 16.7 Modes of Operation

The mode of operation, that is, the behavior of the Timer/Counter and the Output Compare pins, is defined by the combination of the Waveform Generation mode (WGM02:0) and Compare Output mode (COM0x1:0) bits. The Compare Output mode bits do not affect the counting sequence, while the Waveform Generation mode bits do. The COM0x1:0 bits control whether the PWM output generated should be inverted or not (inverted or non-inverted PWM). For non-PWM modes the COM0x1:0 bits control whether the output should be set, cleared, or toggled at a Compare Match. See "Compare Match Output Unit" on page 143.

For detailed timing information see "Timer/Counter Timing Diagrams" on page 124.

### 16.7.1 Normal Mode

The simplest mode of operation is the Normal mode ( $\mathrm{WGMO2:0}=0$ ). In this mode the counting direction is always up (incrementing), and no counter clear is performed. The counter simply overruns when it passes its maximum 8bit value (TOP = 0xFF) and then restarts from the bottom ( $0 \times 00$ ). In normal operation the Timer/Counter Overflow Flag (TOVO) will be set in the same timer clock cycle as the TCNT0 becomes zero. The TOV0 Flag in this case behaves like a ninth bit, except that it is only set, not cleared. However, combined with the timer overflow interrupt that automatically clears the TOVO Flag, the timer resolution can be increased by software. There are no special cases to consider in the Normal mode, a new counter value can be written anytime.

The Output Compare Unit can be used to generate interrupts at some given time. Using the Output Compare to generate waveforms in Normal mode is not recommended, since this will occupy too much of the CPU time.

### 16.7.2 Clear Timer on Compare Match (CTC) Mode

In Clear Timer on Compare or CTC mode (WGM02:0 = 2), the OCROA Register is used to manipulate the counter resolution. In CTC mode the counter is cleared to zero when the counter value (TCNTO) matches the OCROA. The OCROA defines the top value for the counter, hence also its resolution. This mode allows greater control of the Compare Match output frequency. It also simplifies the operation of counting external events.

The timing diagram for the CTC mode is shown in Figure 16-5. The counter value (TCNT0) increases until a Compare Match occurs between TCNTO and OCROA, and then counter (TCNTO) is cleared.

Figure 16-5. CTC Mode, Timing Diagram


An interrupt can be generated each time the counter value reaches the TOP value by using the OCFOA Flag. If the interrupt is enabled, the interrupt handler routine can be used for updating the TOP value. However, changing TOP to a value close to BOTTOM when the counter is running with none or a low prescaler value must be done with care since the CTC mode does not have the double buffering feature. If the new value written to OCROA is lower than the current value of TCNTO, the counter will miss the Compare Match. The counter will then have to count to its maximum value ( $0 x F F$ ) and wrap around starting at $0 \times 00$ before the Compare Match can occur.

For generating a waveform output in CTC mode, the OCOA output can be set to toggle its logical level on each Compare Match by setting the Compare Output mode bits to toggle mode (COMOA1:0 = 1). The OCOA value will not be visible on the port pin unless the data direction for the pin is set to output. The waveform generated will have a maximum frequency of $f_{\mathrm{OCO}}=\mathrm{f}_{\mathrm{clk} \_/ 0} / 2$ when OCROA is set to zero ( $0 \times 00$ ). The waveform frequency is defined by the following equation:

$$
f_{O C n x}=\frac{f_{\mathrm{clk}_{\_} \_/ O}}{2 \cdot N \cdot(1+O C R n x)}
$$

The $N$ variable represents the prescale factor ( $1,8,64,256$, or 1024).
As for the Normal mode of operation, the TOVO Flag is set in the same timer clock cycle that the counter counts from MAX to 0x00.

### 16.7.3 Fast PWM Mode

The fast Pulse Width Modulation or fast PWM mode (WGM02:0 $=3$ or 7 ) provides a high frequency PWM waveform generation option. The fast PWM differs from the other PWM option by its single-slope operation. The counter counts from BOTTOM to TOP then restarts from BOTTOM. TOP is defined as $0 x F F$ when WGM2:0 $=3$, and OCROA when WGM2:0 $=7$. In non-inverting Compare Output mode, the Output Compare (OCOX) is cleared on the Compare Match between TCNTO and OCR0x, and set at BOTTOM. In inverting Compare Output mode, the output is set on Compare Match and cleared at BOTTOM. Due to the single-slope operation, the operating frequency of the fast PWM mode can be twice as high as the phase correct PWM mode that use dual-slope operation. This high frequency makes the fast PWM mode well suited for power regulation, rectification, and DAC applications. High frequency allows physically small sized external components (coils, capacitors), and therefore reduces total system cost.

In fast PWM mode, the counter is incremented until the counter value matches the TOP value. The counter is then cleared at the following timer clock cycle. The timing diagram for the fast PWM mode is shown in Figure 16-6. The TCNTO value is in the timing diagram shown as a histogram for illustrating the single-slope operation. The diagram includes non-inverted and inverted PWM outputs. The small horizontal line marks on the TCNTO slopes represent Compare Matches between OCROx and TCNTO.

Figure 16-6. Fast PWM Mode, Timing Diagram


The Timer/Counter Overflow Flag (TOVO) is set each time the counter reaches TOP. If the interrupt is enabled, the interrupt handler routine can be used for updating the compare value.
In fast PWM mode, the compare unit allows generation of PWM waveforms on the OC0x pins. Setting the COM0x1:0 bits to two will produce a non-inverted PWM and an inverted PWM output can be generated by setting the COM0x1:0 to three: Setting the COM0A1:0 bits to one allows the OC0A pin to toggle on Compare Matches if the WGM02 bit is set. This option is not available for the OCOB pin (see Table 16-3 on page 126). The actual OC0x value will only be visible on the port pin if the data direction for the port pin is set as output. The PWM waveform is generated by setting (or clearing) the OCOx Register at the Compare Match between OCROx and TCNTO, and clearing (or setting) the OC0x Register at the timer clock cycle the counter is cleared (changes from TOP to BOTTOM).

The PWM frequency for the output can be calculated by the following equation:

$$
f_{O C n \times P W M}=\frac{f_{\text {clk }} 1 / \mathrm{O}}{N \cdot 256}
$$

The $N$ variable represents the prescale factor ( $1,8,64,256$, or 1024).
The extreme values for the OCROA Register represents special cases when generating a PWM waveform output in the fast PWM mode. If the OCROA is set equal to BOTTOM, the output will be a narrow spike for each MAX+1 timer clock cycle. Setting the OCROA equal to MAX will result in a constantly high or low output (depending on the polarity of the output set by the COM0A1:0 bits).

A frequency (with $50 \%$ duty cycle) waveform output in fast PWM mode can be achieved by setting OCOx to toggle its logical level on each Compare Match (COM0x1:0 $=1$ ). The waveform generated will have a maximum frequency of $f_{\mathrm{OCO}}=\mathrm{f}_{\text {clk } \_/ 1 / 2} / 2$ when OCROA is set to zero. This feature is similar to the OCOA toggle in CTC mode, except the double buffer feature of the Output Compare unit is enabled in the fast PWM mode.

### 16.7.4 Phase Correct PWM Mode

The phase correct PWM mode (WGM02:0 $=1$ or 5 ) provides a high resolution phase correct PWM waveform generation option. The phase correct PWM mode is based on a dual-slope operation. The counter counts repeatedly from BOTTOM to TOP and then from TOP to BOTTOM. TOP is defined as $0 x F F$ when WGM2:0 $=1$, and OCROA when WGM2:0 $=5$. In non-inverting Compare Output mode, the Output Compare ( $O C O x$ ) is cleared on the Compare Match between TCNTO and OCROx while upcounting, and set on the Compare Match while down-counting. In inverting Output Compare mode, the operation is inverted. The dual-slope operation has lower maximum operation frequency than single slope operation. However, due to the symmetric feature of the dual-slope PWM modes, these modes are preferred for motor control applications.
In phase correct PWM mode the counter is incremented until the counter value matches TOP. When the counter reaches TOP, it changes the count direction. The TCNT0 value will be equal to TOP for one timer clock cycle. The timing diagram for the phase correct PWM mode is shown on Figure 16-7 on page 123. The TCNT0 value is in the timing diagram shown as a histogram for illustrating the dual-slope operation. The diagram includes non-inverted and inverted PWM outputs. The small horizontal line marks on the TCNT0 slopes represent Compare Matches between OCR0x and TCNTO.

Figure 16-7. Phase Correct PWM Mode, Timing Diagram


The Timer/Counter Overflow Flag (TOVO) is set each time the counter reaches BOTTOM. The Interrupt Flag can be used to generate an interrupt each time the counter reaches the BOTTOM value.

In phase correct PWM mode, the compare unit allows generation of PWM waveforms on the OC0x pins. Setting the COM0x1:0 bits to two will produce a non-inverted PWM. An inverted PWM output can be generated by setting the COM0x1:0 to three: Setting the COMOAO bits to one allows the OCOA pin to toggle on Compare Matches if the WGM02 bit is set. This option is not available for the OCOB pin (see Table 16-4 on page 127). The actual OC0x value will only be visible on the port pin if the data direction for the port pin is set as output. The PWM waveform is generated by clearing (or setting) the OC0x Register at the Compare Match between OCROx and TCNTO when the counter increments, and setting (or clearing) the OCOx Register at Compare Match between OCROx and TCNTO when the counter decrements. The PWM frequency for the output when using phase correct PWM can be calculated by the following equation:

$$
f_{O C n x P C P W M}=\frac{f_{\text {clk } 1 / O}}{N \cdot 510}
$$

The $N$ variable represents the prescale factor ( $1,8,64,256$, or 1024).
The extreme values for the OCROA Register represent special cases when generating a PWM waveform output in the phase correct PWM mode. If the OCROA is set equal to BOTTOM, the output will be continuously low and if set equal to MAX the output will be continuously high for non-inverted PWM mode. For inverted PWM the output will have the opposite logic values.

At the very start of period 2 in Figure 16-7 OCnx has a transition from high to low even though there is no Compare Match. The point of this transition is to guarantee symmetry around BOTTOM. There are two cases that give a transition without Compare Match.

- OCROA changes its value from MAX, like in Figure 16-7. When the OCROA value is MAX the OCn pin value is the same as the result of a down-counting Compare Match. To ensure symmetry around BOTTOM the OCn value at MAX must correspond to the result of an up-counting Compare Match.
- The timer starts counting from a value higher than the one in OCROA, and for that reason misses the Compare Match and hence the OCn change that would have happened on the way up.


### 16.8 Timer/Counter Timing Diagrams

The Timer/Counter is a synchronous design and the timer clock ( $\mathrm{clk}_{\mathrm{T}}$ ) is therefore shown as a clock enable signal in the following figures. The figures include information on when Interrupt Flags are set. Figure 16-8 contains timing data for basic Timer/Counter operation. The figure shows the count sequence close to the MAX value in all modes other than phase correct PWM mode.

Figure 16-8. Timer/Counter Timing Diagram, no Prescaling


Figure 16-9 shows the same timing data, but with the prescaler enabled.
Figure 16-9. Timer/Counter Timing Diagram, with Prescaler ( $\mathrm{f}_{\text {clk_// }} / 8$ )


Figure 16-10 shows the setting of OCFOB in all modes and OCFOA in all modes except CTC mode and PWM mode, where OCROA is TOP.

Figure 16-10. Timer/Counter Timing Diagram, Setting of OCFOx, with Prescaler ( $f_{\text {clk_/// }} / 8$ )


Figure $16-11$ on page 125 shows the setting of OCFOA and the clearing of TCNT0 in CTC mode and fast PWM mode where OCROA is TOP.

Figure 16-11. Timer/Counter Timing Diagram, Clear Timer on Compare Match mode, with Prescaler ( $\mathrm{f}_{\mathrm{clk} \_/ / 0} / 8$ )


### 16.9 Register Description

### 16.9.1 TCCROA - Timer/Counter Control Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCR0A |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x24 (0x44) | COMOA1 | COMOAO | COM0B1 | COMOBO | - | - | WGM01 | WGM00 |  |
| Read/Write | R/W | R/W | R/W | R/W | R | R | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7:6 - COMOA1:0: Compare Match Output A Mode

These bits control the Output Compare pin (OCOA) behavior. If one or both of the COMOA1:0 bits are set, the OCOA output overrides the normal port functionality of the I/O pin it is connected to. However, note that the Data Direction Register (DDR) bit corresponding to the OCOA pin must be set in order to enable the output driver.

When OCOA is connected to the pin, the function of the COMOA1:0 bits depends on the WGM02:0 bit setting. Table 16-2 shows the COM0A1:0 bit functionality when the WGM02:0 bits are set to a normal or CTC mode (nonPWM).

Table 16-2. Compare Output Mode, non-PWM Mode

| COMOA1 | COMOAO | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OCOA disconnected |
| 0 | 1 | Toggle OCOA on Compare Match |
| 1 | 0 | Clear OCOA on Compare Match |
| 1 | 1 | Set OCOA on Compare Match |

Table 16-3 shows the COM0A1:0 bit functionality when the WGM01:0 bits are set to fast PWM mode.
Table 16-3. Compare Output Mode, Fast PWM Mode ${ }^{(1)}$

| COMOA1 | COMOAO | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC0A disconnected |
| 0 | 1 | WGMO2 $=0:$ Normal Port Operation, OCOA Disconnected <br> WGM02 = 1: Toggle OCOA on Compare Match |
| 1 | 0 | Clear OCOA on Compare Match, set OC0A at BOTTOM |
| (non-inverting mode) |  |  |

Note: 1. A special case occurs when OCROA equals TOP and COMOA1 is set. In this case, the Compare Match is ignored, but the set or clear is done at BOTTOM. See "Fast PWM Mode" on page 121 for more details.
Table 16-4 on page 127 shows the COMOA1:0 bit functionality when the WGM02:0 bits are set to phase correct PWM mode.

Table 16-4. Compare Output Mode, Phase Correct PWM Mode ${ }^{(1)}$

| COMOA1 | COMOAO | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC0A disconnected |
| 0 | 1 | WGMO2 = 0: Normal Port Operation, OCOA Disconnected <br> WGM02 = 1: Toggle OCOA on Compare Match |
| 1 | 0 | Clear OCOA on Compare Match when up-counting. Set OCOA on Compare Match when |
| down-counting |  |  |

Note: 1. A special case occurs when OCROA equals TOP and COMOA1 is set. In this case, the Compare Match is ignored, but the set or clear is done at TOP. See "Phase Correct PWM Mode" on page 122 for more details.

## - Bits 5:4 - COM0B1:0: Compare Match Output B Mode

These bits control the Output Compare pin (OCOB) behavior. If one or both of the COM0B1:0 bits are set, the OCOB output overrides the normal port functionality of the I/O pin it is connected to. However, note that the Data Direction Register (DDR) bit corresponding to the OCOB pin must be set in order to enable the output driver.

When OCOB is connected to the pin, the function of the COM0B1:0 bits depends on the WGM02:0 bit setting. Table 16-5 shows the COM0B1:0 bit functionality when the WGM02:0 bits are set to a normal or CTC mode (nonPWM).

Table 16-5. Compare Output Mode, non-PWM Mode

| COMOB1 | COMOB0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC0B disconnected |
| 0 | 1 | Toggle OCOB on Compare Match |
| 1 | 0 | Clear OCOB on Compare Match |
| 1 | 1 | Set OCOB on Compare Match |

Table 16-6 shows the COM0B1:0 bit functionality when the WGM02:0 bits are set to fast PWM mode.
Table 16-6. Compare Output Mode, Fast PWM Mode ${ }^{(1)}$

| COMOB1 | COMOB0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OCOB disconnected |
| 0 | 1 | Reserved |
| 1 | 0 | Clear OCOB on Compare Match, set OCOB at BOTTOM |
| (non-inverting mode) |  |  |

Note: 1. A special case occurs when OCROB equals TOP and COMOB1 is set. In this case, the Compare Match is ignored, but the set or clear is done at BOTTOM. See "Fast PWM Mode" on page 121 for more details.

Table 16-7 shows the COMOB1:0 bit functionality when the WGM02:0 bits are set to phase correct PWM mode.
Table 16-7. $\quad$ Compare Output Mode, Phase Correct PWM Mode ${ }^{(1)}$

| COMOB1 | COMOB0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OCOB disconnected |
| 0 | 1 | Reserved |
| 1 | 0 | Clear OCOB on Compare Match when up-counting. Set OCOB on Compare Match when |
| down-counting |  |  |

Note: 1. A special case occurs when OCROB equals TOP and COMOB1 is set. In this case, the Compare Match is ignored, but the set or clear is done at TOP. See "Phase Correct PWM Mode" on page 122 for more details.

## - Bits 3, 2 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

- Bits 1:0 - WGM01:0: Waveform Generation Mode

Combined with the WGM02 bit found in the TCCR0B Register, these bits control the counting sequence of the counter, the source for maximum (TOP) counter value, and what type of waveform generation to be used, see Table 16-8. Modes of operation supported by the Timer/Counter unit are: Normal mode (counter), Clear Timer on Compare Match (CTC) mode, and two types of Pulse Width Modulation (PWM) modes (see "Modes of Operation" on page 144).

Table 16-8. Waveform Generation Mode Bit Description

| Mode | WGM2 | WGM1 | WGM0 | Timer/Counter Mode of Operation | TOP | Update of OCRx at | TOV Flag Set on ${ }^{(1)(2)}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | Normal | 0xFF | Immediate | MAX |
| 1 | 0 | 0 | 1 | PWM, Phase Correct | 0xFF | TOP | BOTTOM |
| 2 | 0 | 1 | 0 | CTC | OCRA | Immediate | MAX |
| 3 | 0 | 1 | 1 | Fast PWM | 0xFF | TOP | MAX |
| 4 | 1 | 0 | 0 | Reserved | - | - | - |
| 5 | 1 | 0 | 1 | PWM, Phase Correct | OCRA | TOP | BOTTOM |
| 6 | 1 | 1 | 0 | Reserved | - | - | - |
| 7 | 1 | 1 | 1 | Fast PWM | OCRA | BOTTOM | TOP |

Note: 1. MAX $=0 x F F$
2. $\mathrm{BOTTOM}=0 \times 00$

### 16.9.2 TCCROB - Timer/Counter Control Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCROB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x25 (0x45) | FOC0A | FOC0B | - | - | WGM02 | CS02 | CS01 | CSOO |  |
| Read/Write | W | W | R | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - FOCOA: Force Output Compare A

The FOCOA bit is only active when the WGM bits specify a non-PWM mode.
However, for ensuring compatibility with future devices, this bit must be set to zero when TCCROB is written when operating in PWM mode. When writing a logical one to the FOCOA bit, an immediate Compare Match is forced on the Waveform Generation unit. The OCOA output is changed according to its COM0A1:0 bits setting. Note that the FOCOA bit is implemented as a strobe. Therefore it is the value present in the COM0A1:0 bits that determines the effect of the forced compare.

A FOCOA strobe will not generate any interrupt, nor will it clear the timer in CTC mode using OCROA as TOP.
The FOCOA bit is always read as zero.

## - Bit 6 - FOCOB: Force Output Compare B

The FOCOB bit is only active when the WGM bits specify a non-PWM mode.
However, for ensuring compatibility with future devices, this bit must be set to zero when TCCROB is written when operating in PWM mode. When writing a logical one to the FOCOB bit, an immediate Compare Match is forced on the Waveform Generation unit. The OCOB output is changed according to its COMOB1:0 bits setting. Note that the FOCOB bit is implemented as a strobe. Therefore it is the value present in the COMOB1:0 bits that determines the effect of the forced compare.

A FOCOB strobe will not generate any interrupt, nor will it clear the timer in CTC mode using OCROB as TOP.
The FOCOB bit is always read as zero.

- Bits 5:4 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

## - Bit 3 - WGMO2: Waveform Generation Mode

See the description in the "TCCROA - Timer/Counter Control Register A" on page 126.

- Bits 2:0 - CS02:0: Clock Select

The three Clock Select bits select the clock source to be used by the Timer/Counter, see Table 16-9 on page 130.

Table 16-9. Clock Select Bit Description

| CS02 | CS01 | CSO0 | Description |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | No clock source (Timer/Counter stopped) |
| 0 | 0 | 1 | $\mathrm{ckk}_{1 / 0} /$ (No prescaling) |
| 0 | 1 | 0 | $\mathrm{clk}_{1 / 0} / 8$ (From prescaler) |
| 0 | 1 | 1 | $\mathrm{clk}_{/ / \mathrm{O}} / 64$ (From prescaler) |
| 1 | 0 | 0 | $\mathrm{clk}_{1 / 0} / 256$ (From prescaler) |
| 1 | 0 | 1 | $\mathrm{clk}_{1 / 0} / 1024$ (From prescaler) |
| 1 | 1 | 0 | External clock source on T0 pin. Clock on falling edge |
| 1 | 1 | 1 | External clock source on T0 pin. Clock on rising edge |

If external pin modes are used for the Timer/CounterO, transitions on the T0 pin will clock the counter even if the pin is configured as an output. This feature allows software control of the counting.

### 16.9.3 TCNTO - Timer/Counter Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCNTO |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x26 (0x46) | TCNT0[7:0] |  |  |  |  |  |  |  |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

The Timer/Counter Register gives direct access, both for read and write operations, to the Timer/Counter unit 8-bit counter. Writing to the TCNT0 Register blocks (removes) the Compare Match on the following timer clock. Modifying the counter (TCNTO) while the counter is running, introduces a risk of missing a Compare Match between TCNTO and the OCROx Registers.

### 16.9.4 OCROA - Output Compare Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | OCROA |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x27 (0x47) | OCROA[7:0] |  |  |  |  |  |  |  |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

The Output Compare Register A contains an 8-bit value that is continuously compared with the counter value (TCNTO). A match can be used to generate an Output Compare interrupt, or to generate a waveform output on the OCOA pin.

### 16.9.5 OCROB - Output Compare Register B



The Output Compare Register B contains an 8-bit value that is continuously compared with the counter value (TCNTO). A match can be used to generate an Output Compare interrupt, or to generate a waveform output on the OCOB pin.

### 16.9.6 TIMSK0 - Timer/Counter Interrupt Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TIMSK0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x6E) | - | - | - | - | - | OCIE0B | OCIEOA | TOIE0 |  |
| Read/Write | R | R | R | R | R | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7:3, 0 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

## - Bit 2 - OCIEOB: Timer/Counter Output Compare Match B Interrupt Enable

When the OCIEOB bit is written to one, and the I-bit in the Status Register is set, the Timer/Counter Compare Match B interrupt is enabled. The corresponding interrupt is executed if a Compare Match in Timer/Counter occurs, that is, when the OCFOB bit is set in the Timer/Counter Interrupt Flag Register - TIFRO.

- Bit 1 - OCIEOA: Timer/Counter0 Output Compare Match A Interrupt Enable

When the OCIEOA bit is written to one, and the I-bit in the Status Register is set, the Timer/Counter0 Compare Match A interrupt is enabled. The corresponding interrupt is executed if a Compare Match in Timer/Counter0 occurs, that is, when the OCFOA bit is set in the Timer/Counter 0 Interrupt Flag Register - TIFRO.

## - Bit 0 - TOIEO: Timer/Counter0 Overflow Interrupt Enable

When the TOIEO bit is written to one, and the I-bit in the Status Register is set, the Timer/Counter0 Overflow interrupt is enabled. The corresponding interrupt is executed if an overflow in Timer/Counter0 occurs, that is, when the TOVO bit is set in the Timer/Counter 0 Interrupt Flag Register - TIFRO.

### 16.9.7 TIFRO - Timer/Counter 0 Interrupt Flag Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x15 (0x35) | - | - | - | - | - | OCFOB | OCFOA | TOVO |
| Read/Write | R | R | R | R | R | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## - Bits 7:3, 0 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

- Bit 2 - OCFOB: Timer/Counter 0 Output Compare B Match Flag

The OCFOB bit is set when a Compare Match occurs between the Timer/Counter and the data in OCROB - Output Compare RegisterO B. OCFOB is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, OCFOB is cleared by writing a logic one to the flag. When the I-bit in SREG, OCIEOB (Timer/Counter Compare B Match Interrupt Enable), and OCFOB are set, the Timer/Counter Compare Match Interrupt is executed.

## - Bit 1 - OCFOA: Timer/Counter 0 Output Compare A Match Flag

The OCFOA bit is set when a Compare Match occurs between the Timer/Counter0 and the data in OCROA - Output Compare Register0. OCFOA is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, OCFOA is cleared by writing a logic one to the flag. When the I-bit in SREG, OCIEOA (Timer/Counter0 Compare Match Interrupt Enable), and OCFOA are set, the Timer/Counter0 Compare Match Interrupt is executed.

## - Bit 0 - TOVO: Timer/CounterO Overflow Flag

The bit TOVO is set when an overflow occurs in Timer/Counter0. TOVO is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, TOVO is cleared by writing a logic one to the flag. When the

SREG I-bit, TOIE0 (Timer/Counter0 Overflow Interrupt Enable), and TOVO are set, the Timer/CounterO Overflow interrupt is executed.

The setting of this flag is dependent of the WGM02:0 bit setting. Refer to Table 16-8, "Waveform Generation Mode Bit Description" on page 128.

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## 17. 16-bit Timer/Counter (Timer/Counter 1, 3, 4, and 5)

### 17.1 Features

- True 16-bit Design (that is, allows 16-bit PWM)
- Three independent Output Compare Units
- Double Buffered Output Compare Registers
- One Input Capture Unit
- Input Capture Noise Canceler
- Clear Timer on Compare Match (Auto Reload)
- Glitch-free, Phase Correct Pulse Width Modulator (PWM)
- Variable PWM Period
- Frequency Generator
- External Event Counter
- Twenty independent interrupt sources (TOV1, OCF1A, OCF1B, OCF1C, ICF1, TOV3, OCF3A, OCF3B, OCF3C, ICF3, TOV4, OCF4A, OCF4B, OCF4C, ICF4, TOV5, OCF5A, OCF5B, OCF5C, and ICF5)


### 17.2 Overview

The 16-bit Timer/Counter unit allows accurate program execution timing (event management), wave generation, and signal timing measurement.

Most register and bit references in this section are written in general form. A lower case "n" replaces the Timer/Counter number, and a lower case " $x$ " replaces the Output Compare unit channel. However, when using the register or bit defines in a program, the precise form must be used, that is, TCNT1 for accessing Timer/Counter1 counter value and so on.

A simplified block diagram of the 16-bit Timer/Counter is shown in Figure 17-1 on page 134. For the actual placement of I/O pins, see "TQFP-pinout ATmega640/1280/2560" on page 2 and "Pinout ATmega1281/2561" on page 4. CPU accessible I/O Registers, including I/O bits and I/O pins, are shown in bold. The device-specific I/O Register and bit locations are listed in the "Register Description" on page 154.

The Power Reduction Timer/Counter1 bit, PRTIM1, in "PRR0 - Power Reduction Register 0" on page 55 must be written to zero to enable Timer/Counter1 module.

The Power Reduction Timer/Counter3 bit, PRTIM3, in "PRR1 - Power Reduction Register 1" on page 56 must be written to zero to enable Timer/Counter3 module.

The Power Reduction Timer/Counter4 bit, PRTIM4, in "PRR1 - Power Reduction Register 1" on page 56 must be written to zero to enable Timer/Counter4 module.

The Power Reduction Timer/Counter5 bit, PRTIM5, in "PRR1 - Power Reduction Register 1" on page 56 must be written to zero to enable Timer/Counter5 module.

Timer/Counter4 and Timer/Counter5 only have full functionality in the ATmega640/1280/2560. Input capture and output compare are not available in the ATmega1281/2561.

Figure 17-1. 16 -bit Timer/Counter Block Diagram ${ }^{(1)}$


Note: 1. Refer to Figure 1-1 on page 2, Table 13-5 on page 76, and Table 13-11 on page 80 for Timer/Counter1 and 3 and 3 pin placement and description.

### 17.2.1 Registers

The Timer/Counter (TCNTn), Output Compare Registers (OCRnA/B/C), and Input Capture Register (ICRn) are all 16 -bit registers. Special procedures must be followed when accessing the 16-bit registers. These procedures are described in the section "Accessing 16-bit Registers" on page 135. The Timer/Counter Control Registers (TCCRnA/B/C) are 8-bit registers and have no CPU access restrictions. Interrupt requests (shorten as Int.Req.) signals are all visible in the Timer Interrupt Flag Register (TIFRn). All interrupts are individually masked with the Timer Interrupt Mask Register (TIMSKn). TIFRn and TIMSKn are not shown in the figure since these registers are shared by other timer units.
The Timer/Counter can be clocked internally, via the prescaler, or by an external clock source on the Tn pin. The Clock Select logic block controls which clock source and edge the Timer/Counter uses to increment (or decrement)
its value. The Timer/Counter is inactive when no clock source is selected. The output from the clock select logic is referred to as the timer clock $\left(\mathrm{clk}_{\mathrm{Tn}}\right)$.

The double buffered Output Compare Registers (OCRnA/B/C) are compared with the Timer/Counter value at all time. The result of the compare can be used by the Waveform Generator to generate a PWM or variable frequency output on the Output Compare pin (OCnA/B/C). See "Output Compare Units" on page 141. The compare match event will also set the Compare Match Flag (OCFnA/B/C) which can be used to generate an Output Compare interrupt request.

The Input Capture Register can capture the Timer/Counter value at a given external (edge triggered) event on either the Input Capture pin (ICPn) or on the Analog Comparator pins (see "AC - Analog Comparator" on page 265). The Input Capture unit includes a digital filtering unit (Noise Canceler) for reducing the chance of capturing noise spikes.

The TOP value, or maximum Timer/Counter value, can in some modes of operation be defined by either the OCRnA Register, the ICRn Register, or by a set of fixed values. When using OCRnA as TOP value in a PWM mode, the OCRnA Register can not be used for generating a PWM output. However, the TOP value will in this case be double buffered allowing the TOP value to be changed in run time. If a fixed TOP value is required, the ICRn Register can be used as an alternative, freeing the OCRnA to be used as PWM output.

### 17.2.2 Definitions

The following definitions are used extensively throughout the document:
Table 17-1. Definitions

| BOTTOM | The counter reaches the BOTTOM when it becomes 0x0000. <br> MAX |
| :--- | :--- |
| The counter reaches its MAXimum when it becomes 0xFFFF (decimal 65535). |  |
| The counter reaches the TOP when it becomes equal to the highest value in the count |  |
| sequence. The TOP value can be assigned to be one of the fixed values: 0x00FF, Ox01FF, or |  |
| 0x03FF, or to the value stored in the OCRnA or ICRn Register. The assignment is dependent of |  |
| the mode of operation. |  |

### 17.3 Accessing 16-bit Registers

The TCNTn, OCRnA/B/C, and ICRn are 16 -bit registers that can be accessed by the AVR CPU via the 8 -bit data bus. The 16 -bit register must be byte accessed using two read or write operations. Each 16 -bit timer has a single 8 bit register for temporary storing of the high byte of the 16 -bit access. The same Temporary Register is shared between all 16-bit registers within each 16-bit timer. Accessing the low byte triggers the 16-bit read or write operation. When the low byte of a 16 -bit register is written by the CPU, the high byte stored in the Temporary Register, and the low byte written are both copied into the 16 -bit register in the same clock cycle. When the low byte of a 16bit register is read by the CPU, the high byte of the 16-bit register is copied into the Temporary Register in the same clock cycle as the low byte is read.

Not all 16-bit accesses uses the Temporary Register for the high byte. Reading the OCRnA/B/C 16-bit registers does not involve using the Temporary Register.
To do a 16-bit write, the high byte must be written before the low byte. For a 16 -bit read, the low byte must be read before the high byte.

The following code examples show how to access the 16-bit timer registers assuming that no interrupts updates the temporary register. The same principle can be used directly for accessing the OCRnA/B/C and ICRn Registers. Note that when using " $C$ ", the compiler handles the 16 -bit access.


Note: 1. See "About Code Examples" on page 10.
The assembly code example returns the TCNTn value in the r17:r16 register pair.
It is important to notice that accessing 16-bit registers are atomic operations. If an interrupt occurs between the two instructions accessing the 16 -bit register, and the interrupt code updates the temporary register by accessing the same or any other of the 16 -bit Timer Registers, then the result of the access outside the interrupt will be corrupted. Therefore, when both the main code and the interrupt code update the temporary register, the main code must disable the interrupts during the 16 -bit access.

The following code examples show how to do an atomic read of the TCNTn Register contents. Reading any of the OCRnA/B/C or ICRn Registers can be done by using the same principle.

```
Assembly Code Example \({ }^{(1)}\)
    TIM16_ReadTCNTn:
    ; Save global interrupt flag
    in r18,SREG
    ; Disable interrupts
    cli
    ; Read TCNTn into r17:r16
    in r16, TCNTnL
    in \(r 17\), TCNTnH
    ; Restore global interrupt flag
    out SREG, r18
    ret
C Code Example \({ }^{(1)}\)
unsigned int TIM16_ReadTCNTn( void )
\{
    unsigned char sreg;
    unsigned int i;
    /* Save global interrupt flag */
    sreg = SREG;
    /* Disable interrupts */
    __disable_interrupt();
    /* Read TCNTn into i */
    i \(=\) TCNTn;
    /* Restore global interrupt flag */
    SREG = sreg;
    return i;
\}
```

Note: 1. See "About Code Examples" on page 10.
The assembly code example returns the TCNTn value in the r17:r16 register pair.

The following code examples show how to do an atomic write of the TCNTn Register contents. Writing any of the OCRnA/B/C or ICRn Registers can be done by using the same principle.

```
Assembly Code Example }\mp@subsup{}{}{(1)
    TIM16_WriteTCNTn:
    ; Save global interrupt flag
    in r18,SREG
    ; Disable interrupts
    cli
    ; Set TCNTn to r17:r16
    out TCNTnH,r17
    out TCNTnL,r16
    ; Restore global interrupt flag
    out SREG,r18
    ret
C Code Example }\mp@subsup{}{}{(1)
    void TIM16_WriteTCNTn( unsigned int i )
    {
    unsigned char sreg;
    unsigned int i;
    /* Save global interrupt flag */
    sreg = SREG;
    /* Disable interrupts */
    __disable_interrupt();
    /* Set TCNTn to i */
    TCNTn = i;
    /* Restore global interrupt flag */
    SREG = sreg;
    }
```

Note: 1. See "About Code Examples" on page 10.
The assembly code example requires that the r17:r16 register pair contains the value to be written to TCNTn.

### 17.3.1 Reusing the Temporary High Byte Register

If writing to more than one 16-bit register where the high byte is the same for all registers written, then the high byte only needs to be written once. However, note that the same rule of atomic operation described previously also applies in this case.

### 17.4 Timer/Counter Clock Sources

The Timer/Counter can be clocked by an internal or an external clock source. The clock source is selected by the Clock Select logic which is controlled by the Clock Select (CSn2:0) bits located in the Timer/Counter control Register $B$ (TCCRnB). For details on clock sources and prescaler, see "Timer/Counter 0, 1, 3, 4, and 5 Prescaler" on page 164.

### 17.5 Counter Unit

The main part of the 16-bit Timer/Counter is the programmable 16-bit bi-directional counter unit. Figure 17-2 shows a block diagram of the counter and its surroundings.

Figure 17-2. Counter Unit Block Diagram


Signal description (internal signals):

| Count | Increment or decrement TCNTn by 1. |
| :--- | :--- |
| Direction | Select between increment and decrement. |
| Clear | Clear TCNTn (set all bits to zero). |
| clk $_{\text {Tn }}$ | Timer/Counter clock. |
| TOP | Signalize that TCNTn has reached maximum value. |
| BOTTOM | Signalize that TCNTn has reached minimum value (zero). |

The 16-bit counter is mapped into two 8-bit I/O memory locations: Counter High (TCNTnH) containing the upper eight bits of the counter, and Counter Low (TCNTnL) containing the lower eight bits. The TCNTnH Register can only be indirectly accessed by the CPU. When the CPU does an access to the TCNTnH I/O location, the CPU accesses the high byte temporary register (TEMP). The temporary register is updated with the TCNTnH value when the TCNTnL is read, and TCNTnH is updated with the temporary register value when TCNTnL is written. This allows the CPU to read or write the entire 16 -bit counter value within one clock cycle via the 8 -bit data bus. It is important to notice that there are special cases of writing to the TCNTn Register when the counter is counting that will give unpredictable results. The special cases are described in the sections where they are of importance.

Depending on the mode of operation used, the counter is cleared, incremented, or decremented at each timer clock $\left(\mathrm{clk}_{\mathrm{T}_{\mathrm{n}}}\right)$. The $\mathrm{ck}_{\mathrm{T}_{\mathrm{n}}}$ can be generated from an external or internal clock source, selected by the Clock Select bits (CSn2:0). When no clock source is selected (CSn2:0 $=0$ ) the timer is stopped. However, the TCNTn value can be accessed by the CPU, independent of whether $\mathrm{clk}_{\mathrm{T}_{\mathrm{n}}}$ is present or not. A CPU write overrides (has priority over) all counter clear or count operations.
The counting sequence is determined by the setting of the Waveform Generation mode bits (WGMn3:0) located in the Timer/Counter Control Registers A and B (TCCRnA and TCCRnB). There are close connections between how the counter behaves (counts) and how waveforms are generated on the Output Compare outputs OCnx. For more details about advanced counting sequences and waveform generation, see "Modes of Operation" on page 144.

The Timer/Counter Overflow Flag (TOVn) is set according to the mode of operation selected by the WGMn3:0 bits. TOVn can be used for generating a CPU interrupt.

### 17.6 Input Capture Unit

The Timer/Counter incorporates an input capture unit that can capture external events and give them a time-stamp indicating time of occurrence. The external signal indicating an event, or multiple events, can be applied via the ICPn pin or alternatively, for the Timer/Counter1 only, via the Analog Comparator unit. The time-stamps can then be used to calculate frequency, duty-cycle, and other features of the signal applied. Alternatively the time-stamps can be used for creating a log of the events.

The Input Capture unit is illustrated by the block diagram shown in Figure 17-3. The elements of the block diagram that are not directly a part of the input capture unit are gray shaded. The small " n " in register and bit names indicates the Timer/Counter number.

Figure 17-3. Input Capture Unit Block Diagram


Note: $\quad$ The Analog Comparator Output (ACO) can only trigger the Timer/Counter1 ICP - not Timer/Counter3, 4 or 5.
When a change of the logic level (an event) occurs on the Input Capture Pin (ICPn), alternatively on the analog Comparator output (ACO), and this change confirms to the setting of the edge detector, a capture will be triggered. When a capture is triggered, the 16 -bit value of the counter (TCNTn) is written to the Input Capture Register (ICRn). The Input Capture Flag (ICFn) is set at the same system clock as the TCNTn value is copied into ICRn Register. If enabled (TICIEn = 1), the input capture flag generates an input capture interrupt. The ICFn flag is automatically cleared when the interrupt is executed. Alternatively the ICFn flag can be cleared by software by writing a logical one to its I/O bit location.

Reading the 16 -bit value in the Input Capture Register (ICRn) is done by first reading the low byte (ICRnL) and then the high byte (ICRnH). When the low byte is read the high byte is copied into the high byte Temporary Register (TEMP). When the CPU reads the ICRnH I/O location it will access the TEMP Register.
The ICRn Register can only be written when using a Waveform Generation mode that utilizes the ICRn Register for defining the counter's TOP value. In these cases the Waveform Generation mode (WGMn3:0) bits must be set before the TOP value can be written to the ICRn Register. When writing the ICRn Register the high byte must be written to the ICRnH I/O location before the low byte is written to ICRnL.

For more information on how to access the 16-bit registers refer to "Accessing 16-bit Registers" on page 135.

### 17.6.1 Input Capture Trigger Source

The main trigger source for the input capture unit is the Input Capture Pin (ICPn). Timer/Counter1 can alternatively use the analog comparator output as trigger source for the input capture unit. The Analog Comparator is selected as trigger source by setting the analog Comparator Input Capture (ACIC) bit in the Analog Comparator Control and Status Register (ACSR). Be aware that changing trigger source can trigger a capture. The input capture flag must therefore be cleared after the change.

Both the Input Capture Pin (ICPn) and the Analog Comparator output (ACO) inputs are sampled using the same technique as for the Tn pin (Figure 18-1 on page 164). The edge detector is also identical. However, when the noise canceler is enabled, additional logic is inserted before the edge detector, which increases the delay by four system clock cycles. Note that the input of the noise canceler and edge detector is always enabled unless the Timer/Counter is set in a Waveform Generation mode that uses ICRn to define TOP.

An input capture can be triggered by software by controlling the port of the ICPn pin.

### 17.6.2 Noise Canceler

The noise canceler improves noise immunity by using a simple digital filtering scheme. The noise canceler input is monitored over four samples, and all four must be equal for changing the output that in turn is used by the edge detector.

The noise canceler is enabled by setting the Input Capture Noise Canceler (ICNCn) bit in Timer/Counter Control Register $B$ (TCCRnB). When enabled the noise canceler introduces additional four system clock cycles of delay from a change applied to the input, to the update of the ICRn Register. The noise canceler uses the system clock and is therefore not affected by the prescaler.

### 17.6.3 Using the Input Capture Unit

The main challenge when using the Input Capture unit is to assign enough processor capacity for handling the incoming events. The time between two events is critical. If the processor has not read the captured value in the ICRn Register before the next event occurs, the ICRn will be overwritten with a new value. In this case the result of the capture will be incorrect.
When using the Input Capture interrupt, the ICRn Register should be read as early in the interrupt handler routine as possible. Even though the Input Capture interrupt has relatively high priority, the maximum interrupt response time is dependent on the maximum number of clock cycles it takes to handle any of the other interrupt requests.

Using the Input Capture unit in any mode of operation when the TOP value (resolution) is actively changed during operation, is not recommended.

Measurement of an external signal's duty cycle requires that the trigger edge is changed after each capture. Changing the edge sensing must be done as early as possible after the ICRn Register has been read. After a change of the edge, the Input Capture Flag (ICFn) must be cleared by software (writing a logical one to the I/O bit location). For measuring frequency only, the clearing of the ICFn Flag is not required (if an interrupt handler is used).

### 17.7 Output Compare Units

The 16 -bit comparator continuously compares TCNTn with the Output Compare Register (OCRnx). If TCNT equals OCRnx the comparator signals a match. A match will set the Output Compare Flag (OCFnx) at the next timer clock cycle. If enabled (OCIEnx = 1), the Output Compare Flag generates an Output Compare interrupt. The OCFnx Flag is automatically cleared when the interrupt is executed. Alternatively the OCFnx Flag can be cleared by software by writing a logical one to its I/O bit location. The Waveform Generator uses the match signal to generate an output according to operating mode set by the Waveform Generation mode (WGMn3:0) bits and Compare Output mode (COMnx1:0) bits. The TOP and BOTTOM signals are used by the Waveform Generator for handling the special cases of the extreme values in some modes of operation. See "Modes of Operation" on page 144.

A special feature of Output Compare unit A allows it to define the Timer/Counter TOP value (that is, counter resolution). In addition to the counter resolution, the TOP value defines the period time for waveforms generated by the Waveform Generator.

Figure 17-4 shows a block diagram of the Output Compare unit. The small " $n$ " in the register and bit names indicates the device number ( $n=n$ for Timer/Counter $n$ ), and the " $x$ " indicates Output Compare unit ( $A / B / C$ ). The elements of the block diagram that are not directly a part of the Output Compare unit are gray shaded.

Figure 17-4. Output Compare Unit, Block Diagram


The OCRnx Register is double buffered when using any of the twelve Pulse Width Modulation (PWM) modes. For the Normal and Clear Timer on Compare (CTC) modes of operation, the double buffering is disabled. The double buffering synchronizes the update of the OCRnx Compare Register to either TOP or BOTTOM of the counting sequence. The synchronization prevents the occurrence of odd-length, non-symmetrical PWM pulses, thereby making the output glitch-free.

The OCRnx Register access may seem complex, but this is not case. When the double buffering is enabled, the CPU has access to the OCRnx Buffer Register, and if double buffering is disabled the CPU will access the OCRnx directly. The content of the OCR1x (Buffer or Compare) Register is only changed by a write operation (the Timer/Counter does not update this register automatically as the TCNT1 and ICR1 Register). Therefore OCR1x is not read via the high byte temporary register (TEMP). However, it is a good practice to read the low byte first as when accessing other 16-bit registers. Writing the OCRnx Registers must be done via the TEMP Register since the compare of all 16 bits is done continuously. The high byte (OCRnxH) has to be written first. When the high byte I/O location is written by the CPU, the TEMP Register will be updated by the value written. Then when the low byte (OCRnxL) is written to the lower eight bits, the high byte will be copied into the upper 8 -bits of either the OCRnx buffer or OCRnx Compare Register in the same system clock cycle.

For more information of how to access the 16-bit registers refer to "Accessing 16-bit Registers" on page 135.

### 17.7.1 Force Output Compare

In non-PWM Waveform Generation modes, the match output of the comparator can be forced by writing a one to the Force Output Compare (FOCnx) bit. Forcing compare match will not set the OCFnx Flag or reload/clear the timer, but the OCnx pin will be updated as if a real compare match had occurred (the COMn1:0 bits settings define whether the OCnx pin is set, cleared or toggled).

### 17.7.2 Compare Match Blocking by TCNTn Write

All CPU writes to the TCNTn Register will block any compare match that occurs in the next timer clock cycle, even when the timer is stopped. This feature allows OCRnx to be initialized to the same value as TCNTn without triggering an interrupt when the Timer/Counter clock is enabled.

### 17.7.3 Using the Output Compare Unit

Since writing TCNTn in any mode of operation will block all compare matches for one timer clock cycle, there are risks involved when changing TCNTn when using any of the Output Compare channels, independent of whether the Timer/Counter is running or not. If the value written to TCNTn equals the OCRnx value, the compare match will be missed, resulting in incorrect waveform generation. Do not write the TCNTn equal to TOP in PWM modes with variable TOP values. The compare match for the TOP will be ignored and the counter will continue to 0xFFFF. Similarly, do not write the TCNTn value equal to BOTTOM when the counter is downcounting.
The setup of the OCnx should be performed before setting the Data Direction Register for the port pin to output. The easiest way of setting the OCnx value is to use the Force Output Compare (FOCnx) strobe bits in Normal mode. The OCnx Register keeps its value even when changing between Waveform Generation modes.
Be aware that the COMnx1:0 bits are not double buffered together with the compare value. Changing the COMnx1:0 bits will take effect immediately.

### 17.8 Compare Match Output Unit

The Compare Output mode (COMnx1:0) bits have two functions. The Waveform Generator uses the COMnx1:0 bits for defining the Output Compare (OCnx) state at the next compare match. Secondly the COMnx1:0 bits control the OCnx pin output source. Figure 17-5 on page 144 shows a simplified schematic of the logic affected by the COMnx1:0 bit setting. The I/O Registers, I/O bits, and I/O pins in the figure are shown in bold. Only the parts of the general I/O Port Control Registers (DDR and PORT) that are affected by the COMnx1:0 bits are shown. When referring to the OCnx state, the reference is for the internal OCnx Register, not the OCnx pin. If a system reset occur, the OCnx Register is reset to " 0 ".

Figure 17-5. Compare Match Output Unit, Schematic


The general I/O port function is overridden by the Output Compare (OCnx) from the Waveform Generator if either of the COMnx1:0 bits are set. However, the OCnx pin direction (input or output) is still controlled by the Data Direction Register (DDR) for the port pin. The Data Direction Register bit for the OCnx pin (DDR_OCnx) must be set as output before the OCnx value is visible on the pin. The port override function is generally independent of the Waveform Generation mode, but there are some exceptions. Refer to Table 17-3 on page 155, Table 17-4 on page 155 and Table 17-5 on page 155 for details.
The design of the Output Compare pin logic allows initialization of the OCnx state before the output is enabled. Note that some COMnx1:0 bit settings are reserved for certain modes of operation. See "Register Description" on page 154.
The COMnx1:0 bits have no effect on the Input Capture unit.

### 17.8.1 Compare Output Mode and Waveform Generation

The Waveform Generator uses the COMnx1:0 bits differently in normal, CTC, and PWM modes. For all modes, setting the COMnx1:0 $=0$ tells the Waveform Generator that no action on the OCnx Register is to be performed on the next compare match. For compare output actions in the non-PWM modes refer to Table 17-3 on page 155. For fast PWM mode refer to Table 17-4 on page 155, and for phase correct and phase and frequency correct PWM refer to Table 17-5 on page 155.

A change of the COMnx1:0 bits state will have effect at the first compare match after the bits are written. For nonPWM modes, the action can be forced to have immediate effect by using the FOCnx strobe bits.

### 17.9 Modes of Operation

The mode of operation, that is, the behavior of the Timer/Counter and the Output Compare pins, is defined by the combination of the Waveform Generation mode (WGMn3:0) and Compare Output mode (COMnx1:0) bits. The Compare Output mode bits do not affect the counting sequence, while the Waveform Generation mode bits do. The COMnx1:0 bits control whether the PWM output generated should be inverted or not (inverted or non-inverted PWM). For non-PWM modes the COMnx1:0 bits control whether the output should be set, cleared or toggle at a compare match. See "Compare Match Output Unit" on page 143.

Table 17-2. Waveform Generation Mode Bit Description ${ }^{(1)}$

| Mode | WGMn3 | WGMn2 (CTCn) | WGMn1 <br> (PWMn1) | WGMn0 (PWMn0) | Timer/Counter Mode of Operation | TOP | Update of OCRnX at | TOVn Flag Set on |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | Normal | 0xFFFF | Immediate | MAX |
| 1 | 0 | 0 | 0 | 1 | PWM, Phase Correct, 8-bit | 0x00FF | TOP | BOTTOM |
| 2 | 0 | 0 | 1 | 0 | PWM, Phase Correct, 9-bit | 0x01FF | TOP | BOTTOM |
| 3 | 0 | 0 | 1 | 1 | PWM, Phase Correct, 10-bit | 0x03FF | TOP | BOTTOM |
| 4 | 0 | 1 | 0 | 0 | CTC | OCRnA | Immediate | MAX |
| 5 | 0 | 1 | 0 | 1 | Fast PWM, 8-bit | 0x00FF | BOTTOM | TOP |
| 6 | 0 | 1 | 1 | 0 | Fast PWM, 9-bit | 0x01FF | BOTTOM | TOP |
| 7 | 0 | 1 | 1 | 1 | Fast PWM, 10-bit | 0x03FF | BOTTOM | TOP |
| 8 | 1 | 0 | 0 | 0 | PWM, Phase and Frequency Correct | ICRn | BOTTOM | BOTTOM |
| 9 | 1 | 0 | 0 | 1 | PWM,Phase and Frequency Correct | OCRnA | BOTTOM | BOTTOM |
| 10 | 1 | 0 | 1 | 0 | PWM, Phase Correct | ICRn | TOP | BOTTOM |
| 11 | 1 | 0 | 1 | 1 | PWM, Phase Correct | OCRnA | TOP | BOTTOM |
| 12 | 1 | 1 | 0 | 0 | CTC | ICRn | Immediate | MAX |
| 13 | 1 | 1 | 0 | 1 | (Reserved) | - | - | - |
| 14 | 1 | 1 | 1 | 0 | Fast PWM | ICRn | BOTTOM | TOP |
| 15 | 1 | 1 | 1 | 1 | Fast PWM | OCRnA | BOTTOM | TOP |

Note: 1. The CTCn and PWMn1:0 bit definition names are obsolete. Use the WGMn2:0 definitions. However, the functionality and location of these bits are compatible with previous versions of the timer.
For detailed timing information refer to "Timer/Counter Timing Diagrams" on page 152.

### 17.9.1 Normal Mode

The simplest mode of operation is the Normal mode (WGMn3:0 = 0). In this mode the counting direction is always up (incrementing), and no counter clear is performed. The counter simply overruns when it passes its maximum 16 -bit value (MAX $=0 \times F F F F$ ) and then restarts from the BOTTOM ( $0 \times 0000$ ). In normal operation the Timer/Counter Overflow Flag (TOVn) will be set in the same timer clock cycle as the TCNTn becomes zero. The TOVn Flag in this case behaves like a $17^{\text {th }}$ bit, except that it is only set, not cleared. However, combined with the timer overflow interrupt that automatically clears the TOVn Flag, the timer resolution can be increased by software. There are no special cases to consider in the Normal mode, a new counter value can be written anytime.

The Input Capture unit is easy to use in Normal mode. However, observe that the maximum interval between the external events must not exceed the resolution of the counter. If the interval between events are too long, the timer overflow interrupt or the prescaler must be used to extend the resolution for the capture unit.

The Output Compare units can be used to generate interrupts at some given time. Using the Output Compare to generate waveforms in Normal mode is not recommended, since this will occupy too much of the CPU time.

### 17.9.2 Clear Timer on Compare Match (CTC) Mode

In Clear Timer on Compare or CTC mode (WGMn3:0 = 4 or 12), the OCRnA or ICRn Register are used to manipulate the counter resolution. In CTC mode the counter is cleared to zero when the counter value (TCNTn) matches either the OCRnA (WGMn3:0 = 4) or the ICRn (WGMn3:0 = 12). The OCRnA or ICRn define the top value for the
counter, hence also its resolution. This mode allows greater control of the compare match output frequency. It also simplifies the operation of counting external events.

The timing diagram for the CTC mode is shown in Figure 17-6. The counter value (TCNTn) increases until a compare match occurs with either OCRnA or ICRn, and then counter (TCNTn) is cleared.

Figure 17-6. CTC Mode, Timing Diagram


An interrupt can be generated at each time the counter value reaches the TOP value by either using the OCFnA or ICFn Flag according to the register used to define the TOP value. If the interrupt is enabled, the interrupt handler routine can be used for updating the TOP value. However, changing the TOP to a value close to BOTTOM when the counter is running with none or a low prescaler value must be done with care since the CTC mode does not have the double buffering feature. If the new value written to OCRnA or ICRn is lower than the current value of TCNTn, the counter will miss the compare match. The counter will then have to count to its maximum value ( $0 x F F F F$ ) and wrap around starting at $0 \times 0000$ before the compare match can occur. In many cases this feature is not desirable. An alternative will then be to use the fast PWM mode using OCRnA for defining TOP (WGMn3:0 = 15) since the OCRnA then will be double buffered.

For generating a waveform output in CTC mode, the OCnA output can be set to toggle its logical level on each compare match by setting the Compare Output mode bits to toggle mode (COMnA1:0 = 1). The OCnA value will not be visible on the port pin unless the data direction for the pin is set to output (DDR_OCnA = 1). The waveform generated will have a maximum frequency of $\mathrm{f}_{\mathrm{OCnA}}=\mathrm{f}_{\mathrm{clk} \_/ 0} / 2$ when OCRnA is set to zero ( $0 \times 0000$ ). The waveform frequency is defined by the following equation:

$$
f_{O C n A}=\frac{f_{\text {ck }} / / / O}{2 \cdot N \cdot(1+O C R n A)}
$$

The $N$ variable represents the prescaler factor ( $1,8,64,256$, or 1024 ).
As for the Normal mode of operation, the TOVn Flag is set in the same timer clock cycle that the counter counts from MAX to $0 \times 0000$.

### 17.9.3 Fast PWM Mode

The fast Pulse Width Modulation or fast PWM mode (WGMn3:0 $=5,6,7,14$, or 15) provides a high frequency PWM waveform generation option. The fast PWM differs from the other PWM options by its single-slope operation. The counter counts from BOTTOM to TOP then restarts from BOTTOM. In non-inverting Compare Output mode, the Output Compare (OCnx) is cleared on the compare match between TCNTn and OCRnx, and set at BOTTOM. In inverting Compare Output mode output is set on compare match and cleared at BOTTOM. Due to the singleslope operation, the operating frequency of the fast PWM mode can be twice as high as the phase correct and phase and frequency correct PWM modes that use dual-slope operation. This high frequency makes the fast PWM
mode well suited for power regulation, rectification, and DAC applications. High frequency allows physically small sized external components (coils, capacitors), hence reduces total system cost.
The PWM resolution for fast PWM can be fixed to 8 -bit, 9 -bit, or 10 -bit, or defined by either ICRn or OCRnA. The minimum resolution allowed is 2 -bit (ICRn or OCRnA set to $0 \times 0003$ ), and the maximum resolution is 16 -bit (ICRn or OCRnA set to MAX). The PWM resolution in bits can be calculated by using the following equation:

$$
R_{F P W M}=\frac{\log (T O P+1)}{\log (2)}
$$

In fast PWM mode the counter is incremented until the counter value matches either one of the fixed values $0 \times 00 F F$, $0 \times 01 \mathrm{FF}$, or $0 \times 03 F F$ (WGMn3:0 $=5,6$, or 7 ), the value in ICRn (WGMn3:0 $=14$ ), or the value in OCRnA (WGMn3:0 = 15). The counter is then cleared at the following timer clock cycle. The timing diagram for the fast PWM mode is shown in Figure 17-7 on page 147. The figure shows fast PWM mode when OCRnA or ICRn is used to define TOP. The TCNTn value is in the timing diagram shown as a histogram for illustrating the single-slope operation. The diagram includes non-inverted and inverted PWM outputs. The small horizontal line marks on the TCNTn slopes represent compare matches between OCRnx and TCNTn. The OCnx Interrupt Flag will be set when a compare match occurs.

Figure 17-7. Fast PWM Mode, Timing Diagram


The Timer/Counter Overflow Flag (TOVn) is set each time the counter reaches TOP. In addition the OCnA or ICFn Flag is set at the same timer clock cycle as TOVn is set when either OCRnA or ICRn is used for defining the TOP value. If one of the interrupts are enabled, the interrupt handler routine can be used for updating the TOP and compare values.
When changing the TOP value the program must ensure that the new TOP value is higher or equal to the value of all of the Compare Registers. If the TOP value is lower than any of the Compare Registers, a compare match will never occur between the TCNTn and the OCRnx. Note that when using fixed TOP values the unused bits are masked to zero when any of the OCRnx Registers are written.
The procedure for updating ICRn differs from updating OCRnA when used for defining the TOP value. The ICRn Register is not double buffered. This means that if ICRn is changed to a low value when the counter is running with none or a low prescaler value, there is a risk that the new ICRn value written is lower than the current value of TCNTn. The result will then be that the counter will miss the compare match at the TOP value. The counter will then have to count to the MAX value ( $0 \times \mathrm{FFFFF}$ ) and wrap around starting at $0 \times 0000$ before the compare match can occur. The OCRnA Register however, is double buffered. This feature allows the OCRnA I/O location to be written anytime. When the OCRnA I/O location is written the value written will be put into the OCRnA Buffer Register. The OCRnA Compare Register will then be updated with the value in the Buffer Register at the next timer clock cycle
the TCNTn matches TOP. The update is done at the same timer clock cycle as the TCNTn is cleared and the TOVn Flag is set.

Using the ICRn Register for defining TOP works well when using fixed TOP values. By using ICRn, the OCRnA Register is free to be used for generating a PWM output on OCnA. However, if the base PWM frequency is actively changed (by changing the TOP value), using the OCRnA as TOP is clearly a better choice due to its double buffer feature.

In fast PWM mode, the compare units allow generation of PWM waveforms on the OCnx pins. Setting the COMnx1:0 bits to two will produce a non-inverted PWM and an inverted PWM output can be generated by setting the COMnx1:0 to three (see Table on page 155). The actual OCnx value will only be visible on the port pin if the data direction for the port pin is set as output (DDR_OCnx). The PWM waveform is generated by setting (or clearing) the OCnx Register at the compare match between OCRnx and TCNTn, and clearing (or setting) the OCnx Register at the timer clock cycle the counter is cleared (changes from TOP to BOTTOM).
The PWM frequency for the output can be calculated by the following equation:

$$
f_{O C n x P W M}=\frac{f_{\mathrm{ckk}} / / O}{N \cdot(1+T O P)}
$$

The N variable represents the prescaler divider (1, 8, 64, 256, or 1024).
The extreme values for the OCRnx Register represents special cases when generating a PWM waveform output in the fast PWM mode. If the OCRnx is set equal to BOTTOM ( $0 \times 0000$ ) the output will be a narrow spike for each TOP +1 timer clock cycle. Setting the OCRnx equal to TOP will result in a constant high or low output (depending on the polarity of the output set by the COMnx1:0 bits).

A frequency (with $50 \%$ duty cycle) waveform output in fast PWM mode can be achieved by setting OCnA to toggle its logical level on each compare match (COMnA1:0 $=1$ ). This applies only if OCR1A is used to define the TOP value (WGM13:0 = 15). The waveform generated will have a maximum frequency of $f_{\text {OCnA }}=f_{\text {clk_/I/ }} / 2$ when OCRnA is set to zero ( $0 \times 0000$ ). This feature is similar to the OCnA toggle in CTC mode, except the double buffer feature of the Output Compare unit is enabled in the fast PWM mode.

### 17.9.4 Phase Correct PWM Mode

The phase correct Pulse Width Modulation or phase correct PWM mode (WGMn3:0 $=1,2,3,10$, or 11) provides a high resolution phase correct PWM waveform generation option. The phase correct PWM mode is, like the phase and frequency correct PWM mode, based on a dual-slope operation. The counter counts repeatedly from BOTTOM $(0 \times 0000)$ to TOP and then from TOP to BOTTOM. In non-inverting Compare Output mode, the Output Compare (OCnx) is cleared on the compare match between TCNTn and OCRnx while upcounting, and set on the compare match while downcounting. In inverting Output Compare mode, the operation is inverted. The dual-slope operation has lower maximum operation frequency than single slope operation. However, due to the symmetric feature of the dual-slope PWM modes, these modes are preferred for motor control applications.

The PWM resolution for the phase correct PWM mode can be fixed to 8-bit, 9-bit, or 10 -bit, or defined by either ICRn or OCRnA. The minimum resolution allowed is 2-bit (ICRn or OCRnA set to 0x0003), and the maximum resolution is 16-bit (ICRn or OCRnA set to MAX). The PWM resolution in bits can be calculated by using the following equation:

$$
R_{P C P W M}=\frac{\log (T O P+1)}{\log (2)}
$$

In phase correct PWM mode the counter is incremented until the counter value matches either one of the fixed values $0 \times 00 F F$, $0 \times 01 \mathrm{FF}$, or $0 \times 03 F F$ ( $\mathrm{WGMn} 3: 0=1,2$, or 3 ), the value in ICRn (WGMn3:0 $=10$ ), or the value in OCRnA (WGMn3:0 = 11). The counter has then reached the TOP and changes the count direction. The TCNTn value will be equal to TOP for one timer clock cycle. The timing diagram for the phase correct PWM mode is shown
on Figure 17-8 on page 149. The figure shows phase correct PWM mode when OCRnA or ICRn is used to define TOP. The TCNTn value is in the timing diagram shown as a histogram for illustrating the dual-slope operation. The diagram includes non-inverted and inverted PWM outputs. The small horizontal line marks on the TCNTn slopes represent compare matches between OCRnx and TCNTn. The OCnx Interrupt Flag will be set when a compare match occurs.

Figure 17-8. Phase Correct PWM Mode, Timing Diagram


The Timer/Counter Overflow Flag (TOVn) is set each time the counter reaches BOTTOM. When either OCRnA or ICRn is used for defining the TOP value, the OCnA or ICFn Flag is set accordingly at the same timer clock cycle as the OCRnx Registers are updated with the double buffer value (at TOP). The Interrupt Flags can be used to generate an interrupt each time the counter reaches the TOP or BOTTOM value.

When changing the TOP value the program must ensure that the new TOP value is higher or equal to the value of all of the Compare Registers. If the TOP value is lower than any of the Compare Registers, a compare match will never occur between the TCNTn and the OCRnx. Note that when using fixed TOP values, the unused bits are masked to zero when any of the OCRnx Registers are written. As the third period shown in Figure 17-8 illustrates, changing the TOP actively while the Timer/Counter is running in the phase correct mode can result in an unsymmetrical output. The reason for this can be found in the time of update of the OCRnx Register. Since the OCRnx update occurs at TOP, the PWM period starts and ends at TOP. This implies that the length of the falling slope is determined by the previous TOP value, while the length of the rising slope is determined by the new TOP value. When these two values differ the two slopes of the period will differ in length. The difference in length gives the unsymmetrical result on the output.

It is recommended to use the phase and frequency correct mode instead of the phase correct mode when changing the TOP value while the Timer/Counter is running. When using a static TOP value there are practically no differences between the two modes of operation.

In phase correct PWM mode, the compare units allow generation of PWM waveforms on the OCnx pins. Setting the COMnx1:0 bits to two will produce a non-inverted PWM and an inverted PWM output can be generated by setting the COMnx1:0 to three (see Table 17-5 on page 155). The actual OCnx value will only be visible on the port pin if the data direction for the port pin is set as output (DDR_OCnx). The PWM waveform is generated by setting (or clearing) the OCnx Register at the compare match between OCRnx and TCNTn when the counter increments, and clearing (or setting) the OCnx Register at compare match between OCRnx and TCNTn when the counter dec-
rements. The PWM frequency for the output when using phase correct PWM can be calculated by the following equation:

$$
f_{O C n \times P C P W M}=\frac{f_{\text {clk }} / / O}{2 \cdot N \cdot T O P}
$$

The $N$ variable represents the prescaler divider (1, 8, 64, 256, or 1024).
The extreme values for the OCRnx Register represent special cases when generating a PWM waveform output in the phase correct PWM mode. If the OCRnx is set equal to BOTTOM the output will be continuously low and if set equal to TOP the output will be continuously high for non-inverted PWM mode. For inverted PWM the output will have the opposite logic values. If OCR1A is used to define the TOP value (WGM13:0 $=11$ ) and COM1A1:0 $=1$, the OC1A output will toggle with a $50 \%$ duty cycle.

### 17.9.5 Phase and Frequency Correct PWM Mode

The phase and frequency correct Pulse Width Modulation, or phase and frequency correct PWM mode (WGMn3:0 $=8$ or 9 ) provides a high resolution phase and frequency correct PWM waveform generation option. The phase and frequency correct PWM mode is, like the phase correct PWM mode, based on a dual-slope operation. The counter counts repeatedly from BOTTOM ( $0 \times 0000$ ) to TOP and then from TOP to BOTTOM. In non-inverting Compare Output mode, the Output Compare (OCnx) is cleared on the compare match between TCNTn and OCRnx while upcounting, and set on the compare match while downcounting. In inverting Compare Output mode, the operation is inverted. The dual-slope operation gives a lower maximum operation frequency compared to the sin-gle-slope operation. However, due to the symmetric feature of the dual-slope PWM modes, these modes are preferred for motor control applications.

The main difference between the phase correct, and the phase and frequency correct PWM mode is the time the OCRnx Register is updated by the OCRnx Buffer Register, see Figure 17-8 on page 149 and Figure 17-9 on page 151.

The PWM resolution for the phase and frequency correct PWM mode can be defined by either ICRn or OCRnA. The minimum resolution allowed is 2 -bit (ICRn or OCRnA set to $0 \times 0003$ ), and the maximum resolution is 16 -bit (ICRn or OCRnA set to MAX). The PWM resolution in bits can be calculated using the following equation:

$$
R_{P F C P W M}=\frac{\log (T O P+1)}{\log (2)}
$$

In phase and frequency correct PWM mode the counter is incremented until the counter value matches either the value in ICRn $(W G M n 3: 0=8)$, or the value in OCRnA $(W G M n 3: 0=9)$. The counter has then reached the TOP and changes the count direction. The TCNTn value will be equal to TOP for one timer clock cycle. The timing diagram for the phase correct and frequency correct PWM mode is shown on Figure 17-9 on page 151. The figure shows phase and frequency correct PWM mode when OCRnA or ICRn is used to define TOP. The TCNTn value is in the timing diagram shown as a histogram for illustrating the dual-slope operation. The diagram includes non-inverted and inverted PWM outputs. The small horizontal line marks on the TCNTn slopes represent compare matches between OCRnx and TCNTn. The OCnx Interrupt Flag will be set when a compare match occurs.

Figure 17-9. Phase and Frequency Correct PWM Mode, Timing Diagram


The Timer/Counter Overflow Flag (TOVn) is set at the same timer clock cycle as the OCRnx Registers are updated with the double buffer value (at BOTTOM). When either OCRnA or ICRn is used for defining the TOP value, the OCnA or ICFn Flag set when TCNTn has reached TOP. The Interrupt Flags can then be used to generate an interrupt each time the counter reaches the TOP or BOTTOM value.
When changing the TOP value the program must ensure that the new TOP value is higher or equal to the value of all of the Compare Registers. If the TOP value is lower than any of the Compare Registers, a compare match will never occur between the TCNTn and the OCRnx.

As Figure 17-9 shows the output generated is, in contrast to the phase correct mode, symmetrical in all periods. Since the OCRnx Registers are updated at BOTTOM, the length of the rising and the falling slopes will always be equal. This gives symmetrical output pulses and is therefore frequency correct.

Using the ICRn Register for defining TOP works well when using fixed TOP values. By using ICRn, the OCRnA Register is free to be used for generating a PWM output on OCnA. However, if the base PWM frequency is actively changed by changing the TOP value, using the OCRnA as TOP is clearly a better choice due to its double buffer feature.

In phase and frequency correct PWM mode, the compare units allow generation of PWM waveforms on the OCnx pins. Setting the COMnx1:0 bits to two will produce a non-inverted PWM and an inverted PWM output can be generated by setting the COMnx1:0 to three (see Table 17-5 on page 155). The actual OCnx value will only be visible on the port pin if the data direction for the port pin is set as output (DDR_OCnx). The PWM waveform is generated by setting (or clearing) the OCnx Register at the compare match between OCRnx and TCNTn when the counter increments, and clearing (or setting) the OCnx Register at compare match between OCRnx and TCNTn when the counter decrements. The PWM frequency for the output when using phase and frequency correct PWM can be calculated by the following equation:

$$
f_{O C n x P F C P W M}=\frac{f_{\text {clk } \_/ O}}{2 \cdot N \cdot T O P}
$$

The N variable represents the prescaler divider (1, 8, 64, 256, or 1024).
The extreme values for the OCRnx Register represents special cases when generating a PWM waveform output in the phase correct PWM mode. If the OCRnx is set equal to BOTTOM the output will be continuously low and if set equal to TOP the output will be set to high for non-inverted PWM mode. For inverted PWM the output will have the opposite logic values. If OCR1A is used to define the TOP value (WGM13:0 $=9$ ) and COM1A1:0 $=1$, the OC1A output will toggle with a $50 \%$ duty cycle.

### 17.10 Timer/Counter Timing Diagrams

The Timer/Counter is a synchronous design and the timer clock ( $\mathrm{clk}_{\mathrm{T}_{\mathrm{n}}}$ ) is therefore shown as a clock enable signal in the following figures. The figures include information on when Interrupt Flags are set, and when the OCRnx Register is updated with the OCRnx buffer value (only for modes utilizing double buffering). Figure 17-10 shows a timing diagram for the setting of OCFnx.

Figure 17-10. Timer/Counter Timing Diagram, Setting of OCFnx, no Prescaling


Figure 17-11 shows the same timing data, but with the prescaler enabled.
Figure 17-11. Timer/Counter Timing Diagram, Setting of OCFnx, with Prescaler ( $f_{\text {clk_// }} / 8$ )


Figure 17-12 shows the count sequence close to TOP in various modes. When using phase and frequency correct PWM mode the OCRnx Register is updated at BOTTOM. The timing diagrams will be the same, but TOP should be replaced by BOTTOM, TOP-1 by BOTTOM+1 and so on. The same renaming applies for modes that set the TOVn Flag at BOTTOM.

Figure 17-12. Timer/Counter Timing Diagram, no Prescaling


Figure $17-13$ shows the same timing data, but with the prescaler enabled.

Figure 17-13. Timer/Counter Timing Diagram, with Prescaler ( $\left.\mathrm{f}_{\mathrm{clk} \_} / / / \mathrm{/} / 8\right)$


### 17.11 Register Description

### 17.11.1 TCCR1A - Timer/Counter 1 Control Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x80) | COM1A1 | COM1A0 | COM1B1 | COM1B0 | COM1C1 | COM1C0 | WGM11 | WGM10 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TCCR1A

### 17.11.2 TCCR3A - Timer/Counter 3 Control Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x90) | COM3A1 | COM3A0 | COM3B1 | COM3B0 | COM3C1 | COM3C0 | WGM31 | WGM30 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## TCCR3A

### 17.11.3 TCCR4A - Timer/Counter 4 Control Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xA0) | COM4A1 | COM4A0 | COM4B1 | COM4B0 | COM4C1 | COM4C0 | WGM41 | WGM40 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TCCR4A

### 17.11.4 TCCR5A - Timer/Counter 5 Control Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x120) | COM5A1 | COM5A0 | COM5B1 | COM5B0 | COM5C1 | COM5C0 | WGM51 | WGM50 |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TCCR5A

- Bit 7:6 - COMnA1:0: Compare Output Mode for Channel A
- Bit 5:4 - COMnB1:0: Compare Output Mode for Channel B
- Bit 3:2 - COMnC1:0: Compare Output Mode for Channel C

The COMnA1:0, COMnB1:0, and COMnC1:0 control the output compare pins (OCnA, OCnB, and OCnC respectively) behavior. If one or both of the COMnA1:0 bits are written to one, the OCnA output overrides the normal port functionality of the I/O pin it is connected to. If one or both of the COMnB1:0 bits are written to one, the OCnB output overrides the normal port functionality of the I/O pin it is connected to. If one or both of the COMnC1:0 bits are written to one, the OCnC output overrides the normal port functionality of the I/O pin it is connected to. However, note that the Data Direction Register (DDR) bit corresponding to the OCnA, OCnB or OCnC pin must be set in order to enable the output driver.

When the OCnA, OCnB or OCnC is connected to the pin, the function of the COMnx1:0 bits is dependent of the WGMn3:0 bits setting. Table 17-3 on page 155 shows the COMnx1:0 bit functionality when the WGMn3:0 bits are set to a normal or a CTC mode (non-PWM).

## - Bit 1:0 - WGMn1:0: Waveform Generation Mode

Combined with the WGMn3:2 bits found in the TCCRnB Register, these bits control the counting sequence of the counter, the source for maximum (TOP) counter value, and what type of waveform generation to be used, see Table 17-2 on page 145. Modes of operation supported by the Timer/Counter unit are: Normal mode (counter), Clear Timer on Compare match (CTC) mode, and three types of Pulse Width Modulation (PWM) modes. For more information on the different modes, see "Modes of Operation" on page 144.

Table 17-3. Compare Output Mode, non-PWM

| COMnA1 <br> COMnB1 <br> COMnC1 | COMnA0 <br> COMnB0 <br> COMnC0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OCnA/OCnB/OCnC disconnected |
| 0 | 1 | Toggle OCnA/OCnB/OCnC on compare match |
| 1 | 0 | Clear OCnA/OCnB/OCnC on compare match (set output to low level) |
| 1 | 1 | Set OCnA/OCnB/OCnC on compare match (set output to high level) |

Table 17-4 shows the COMnx1:0 bit functionality when the WGMn3:0 bits are set to the fast PWM mode.
Table 17-4. Compare Output Mode, Fast PWM

| COMnA1 <br> COMnB1 <br> COMnC1 | COMnA0 <br> COMnB0 <br> COMnC0 | Description |
| :---: | :---: | :---: | (0 10 | Normal port operation, OCnA/OCnB/OCnC disconnected |
| :---: |
| 0 |

Note: A special case occurs when OCRnA/OCRnB/OCRnC equals TOP and COMnA1/COMnB1/COMnC1 is set. In this case the compare match is ignored, but the set or clear is done at BOTTOM. See "Fast PWM Mode" on page 146. for more details.
Table 17-5 shows the COMnx1:0 bit functionality when the WGMn3:0 bits are set to the phase correct and frequency correct PWM mode.

Table 17-5. Compare Output Mode, Phase Correct and Phase and Frequency Correct PWM

| COMnA1 <br> COMnB1 <br> COMnC1 | COMnA0 <br> COMnB0 <br> COMnC0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OCnA/OCnB/OCnC disconnected |
| 0 | 1 | WGM13:0 =9 or 11: Toggle OC1A on Compare Match, OC1B and OC1C disconnected (normal port <br> operation). For all other WGM1 settings, normal port operation, OC1A/OC1B/OC1C disconnected |
| 1 | 0 | Clear OCnA/OCnB/OCnC on compare match when up-counting <br> Set OCnA/OCnB/OCnC on compare match when downcounting |
| 1 | 1 | Set OCnA/OCnB/OCnC on compare match when up-counting <br> Clear OCnA/OCnB/OCnC on compare match when downcounting |

Note: A special case occurs when OCRnA/OCRnB/OCRnC equals TOP and COMnA1/COMnB1//COMnC1 is set. See "Phase Correct PWM Mode" on page 148. for more details.

### 17.11.5 TCCR1B - Timer/Counter 1 Control Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCR1B |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x81) | ICNC1 | ICES1 | - | WGM13 | WGM12 | CS12 | CS11 | CS10 |  |
| Read/Write | R/W | R/W | R | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.6 TCCR3B - Timer/Counter 3 Control Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCR3B |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x91) | ICNC3 | ICES3 | - | WGM33 | WGM32 | CS32 | CS31 | CS30 |  |
| Read/Write | R/W | R/W | R | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.7 TCCR4B - Timer/Counter 4 Control Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xA1) | ICNC4 | ICES4 | - | WGM43 | WGM42 | CS42 | CS41 | CS40 |
| Read/Write | R/W | R/W | R | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TCCR4B

### 17.11.8 TCCR5B - Timer/Counter 5 Control Register B

Bit
(0x121)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ICNC5 | ICES5 | - | WGM53 | WGM52 | CS52 | CS51 | CS50 |
| R/W | R/W | R | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TCCR5B

## - Bit 7 - ICNCn: Input Capture Noise Canceler

Setting this bit (to one) activates the Input Capture Noise Canceler. When the Noise Canceler is activated, the input from the Input Capture Pin (ICPn) is filtered. The filter function requires four successive equal valued samples of the ICPn pin for changing its output. The input capture is therefore delayed by four Oscillator cycles when the noise canceler is enabled.

## - Bit 6 - ICESn: Input Capture Edge Select

This bit selects which edge on the Input Capture Pin (ICPn) that is used to trigger a capture event. When the ICESn bit is written to zero, a falling (negative) edge is used as trigger, and when the ICESn bit is written to one, a rising (positive) edge will trigger the capture.
When a capture is triggered according to the ICESn setting, the counter value is copied into the Input Capture Register (ICRn). The event will also set the Input Capture Flag (ICFn), and this can be used to cause an Input Capture Interrupt, if this interrupt is enabled.

When the ICRn is used as TOP value (see description of the WGMn3:0 bits located in the TCCRnA and the TCCRnB Register), the ICPn is disconnected and consequently the input capture function is disabled.

- Bit 5 - Reserved Bit

This bit is reserved for future use. For ensuring compatibility with future devices, this bit must be written to zero when TCCRnB is written.

- Bit 4:3 - WGMn3:2: Waveform Generation Mode

See TCCRnA Register description.

## - Bit 2:0 - CSn2:0: Clock Select

The three clock select bits select the clock source to be used by the Timer/Counter, see Figure 17-10 and Figure 17-11 on page 152.

Table 17-6. Clock Select Bit Description

| CSn2 | CSn1 | CSn0 | Description |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | No clock source. (Timer/Counter stopped) |
| 0 | 0 | 1 | $\mathrm{clk}_{/ / /} / 1$ (No prescaling |
| 0 | 1 | 0 | $\mathrm{clk}_{1 / 0} / 8$ (From prescaler) |
| 0 | 1 | 1 | $\mathrm{clk}_{1 / 0} / 64$ (From prescaler) |
| 1 | 0 | 0 | $\mathrm{clk}_{1 / 0} / 256$ (From prescaler) |
| 1 | 0 | 1 | $\mathrm{clk}_{1 / 0} / 1024$ (From prescaler) |
| 1 | 1 | 0 | External clock source on Tn pin. Clock on falling edge |
| 1 | 1 | 1 | External clock source on Tn pin. Clock on rising edge |

If external pin modes are used for the Timer/Countern, transitions on the Tn pin will clock the counter even if the pin is configured as an output. This feature allows software control of the counting.

### 17.11.9 TCCR1C - Timer/Counter 1 Control Register C

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCR1C |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x82) | FOC1A | FOC1B | FOC1C | - | - | - | - | - |  |
| Read/Write | W | W | W | R | R | R | R | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.10 TCCR3C - Timer/Counter 3 Control Register C

Bit

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| FOC3A | FOC3B | FOC3C | - | - | - | - | - |
| W | W | W | $R$ | $R$ | $R$ | $R$ | $R$ |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TCCR3C
Read/Write
Initial Value
0
0
17.11.11 TCCR4C - Timer/Counter 4 Control Register C

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCR4C |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xA2) | FOC4A | FOC4B | FOC4C | - | - | - | - | - |  |
| Read/Write | W | W | W | R | R | R | R | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.12 TCCR5C - Timer/Counter 5 Control Register C

Bit
(0x122)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| FOC5A | FOC5B | FOC3C | - | - | - | - | - |
| W | W | W | $R$ | $R$ | $R$ | $R$ | $R$ |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## TCCR5C

- Bit 7 - FOCnA: Force Output Compare for Channel A
- Bit 6 - FOCnB: Force Output Compare for Channel B
- Bit 5 - FOCnC: Force Output Compare for Channel C

The FOCnA/FOCnB/FOCnC bits are only active when the WGMn3:0 bits specifies a non-PWM mode. When writing a logical one to the FOCnA/FOCnB/FOCnC bit, an immediate compare match is forced on the waveform generation unit. The $\mathrm{OCnA} / \mathrm{OCnB} / \mathrm{OCnC}$ output is changed according to its COMnx1:0 bits setting. Note that the

FOCnA/FOCnB/FOCnC bits are implemented as strobes. Therefore it is the value present in the COMnx1:0 bits that determine the effect of the forced compare.

A FOCnA/FOCnB/FOCnC strobe will not generate any interrupt nor will it clear the timer in Clear Timer on Compare Match (CTC) mode using OCRnA as TOP.

The FOCnA/FOCnB/FOCnB bits are always read as zero.

## - Bit 4:0 - Reserved Bits

These bits are reserved for future use. For ensuring compatibility with future devices, these bits must be written to zero when TCCRnC is written.

### 17.11.13 TCNT1H and TCNT1L - Timer/Counter 1

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x85) | TCNT1[15:8] |  |  |  |  |  |  |  | TCNT1H |
| (0x84) | TCNT1[7:0] |  |  |  |  |  |  |  | TCNT1L |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.14 TCNT3H and TCNT3L - Timer/Counter 3

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x95) | TCNT3[15:8] |  |  |  |  |  |  |  | TCNT3H |
| (0x94) | TCNT3[7:0] |  |  |  |  |  |  |  | TCNT3L |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.15 TCNT4H and TCNT4L -Timer/Counter 4



### 17.11.16 TCNT5H and TCNT5L -Timer/Counter 5

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x125) | TCNT5[15:8] |  |  |  |  |  |  |  | TCNT5H |
| (0x124) | TCNT5[7:0] |  |  |  |  |  |  |  | TCNT5L |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

The two Timer/Counter I/O locations (TCNTnH and TCNTnL, combined TCNTn) give direct access, both for read and for write operations, to the Timer/Counter unit 16 -bit counter. To ensure that both the high and low bytes are read and written simultaneously when the CPU accesses these registers, the access is performed using an 8-bit temporary High Byte Register (TEMP). This temporary register is shared by all the other 16-bit registers. See "Accessing 16-bit Registers" on page 135.

Modifying the counter (TCNTn) while the counter is running introduces a risk of missing a compare match between TCNTn and one of the OCRnx Registers.

Writing to the TCNTn Register blocks (removes) the compare match on the following timer clock for all compare units.

### 17.11.17 OCR1AH and OCR1AL - Output Compare Register 1 A

Bit
(0x89)
(0x88)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | OCR1A[15:8] |  |  |  |  |  |  |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

OCR1AH
OCR1AL

### 17.11.18 OCR1BH and OCR1BL - Output Compare Register 1 B

Bit
(0x8B)
(0x8A)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| OCR1B[15:8] |  |  |  |  |  |  |  |
| OCR1B[7:0] |  |  |  |  |  |  |  |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

OCR1BH OCR1BL
17.11.19 OCR1CH and OCR1CL - Output Compare Register 1 C

Bit
(0x8D)
(0x8C)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |

OCR1CH OCR1CL
17.11.20 OCR3AH and OCR3AL - Output Compare Register 3 A

Bit
(0x99)
$(0 \times 98)$
Read/Write

| 7 | 6 |
| :---: | :---: |
|  |  |
|  |  |
| R/W |  |
| 0 | 0 |

6

R/W
0

### 17.11.24 OCR4BH and OCR4BL - Output Compare Register 4 B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xAA) | OCR4B[15:8] |  |  |  |  |  |  |  |
| (0xAB) | OCR4B[7:0] |  |  |  |  |  |  |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

OCR4BH
OCR4BL

### 17.11.25 OCR4CH and OCR4CL -Output Compare Register 4 C

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xAD) | OCR4C[15:8] |  |  |  |  |  |  |  |
| (0xAC) | OCR4C[7:0] |  |  |  |  |  |  |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 17.11.26 OCR5AH and OCR5AL - Output Compare Register 5 A



OCR5AH
OCR5AL
17.11.27 OCR5BH and OCR5BL - Output Compare Register 5 B

17.11.28 OCR5CH and OCR5CL -Output Compare Register 5 C

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | OCR5CH |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x12D) | OCR5C[15:8] |  |  |  |  |  |  |  |  |
| (0x12C) | OCR5C[7:0] |  |  |  |  |  |  |  | OCR5CL |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

The Output Compare Registers contain a 16-bit value that is continuously compared with the counter value (TCNTn). A match can be used to generate an Output Compare interrupt, or to generate a waveform output on the OCnx pin.

The Output Compare Registers are 16-bit in size. To ensure that both the high and low bytes are written simultaneously when the CPU writes to these registers, the access is performed using an 8-bit temporary High Byte Register (TEMP). This temporary register is shared by all the other 16 -bit registers. See "Accessing 16 -bit Registers" on page 135.

### 17.11.29 ICR1H and ICR1L - Input Capture Register 1

Bit
(0x87)
(0x86)
Read/Write
Initial Value


ICR1H ICR1L

### 17.11.30 ICR3H and ICR3L - Input Capture Register 3

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x97) | ICR3[15:8] |  |  |  |  |  |  |  |
| (0x96) | ICR3[7:0] |  |  |  |  |  |  |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 17.11.31 ICR4H and ICR4L - Input Capture Register 4

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xA7) | ICR4[15:8] |  |  |  |  |  |  |  | ICR4H |
| (0xA6) | ICR4[7:0] |  |  |  |  |  |  |  | ICR4L |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.32 ICR5H and ICR5L - Input Capture Register 5

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x127) | ICR5[15:8] |  |  |  |  |  |  |  | ICR5H |
| (0x126) | ICR5[7:0] |  |  |  |  |  |  |  | ICR5L |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

The Input Capture is updated with the counter (TCNTn) value each time an event occurs on the ICPn pin (or optionally on the Analog Comparator output for Timer/Counter1). The Input Capture can be used for defining the counter TOP value.

The Input Capture Register is 16 -bit in size. To ensure that both the high and low bytes are read simultaneously when the CPU accesses these registers, the access is performed using an 8-bit temporary High Byte Register (TEMP). This temporary register is shared by all the other 16 -bit registers. See "Accessing 16 -bit Registers" on page 135.

### 17.11.33 TIMSK1 - Timer/Counter 1 Interrupt Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x6F) | - | - | ICIE1 | - | OCIE1C | OCIE1B | OCIE1A | TOIE1 |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TIMSK1

### 17.11.34 TIMSK3 - Timer/Counter 3 Interrupt Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x71) | - | - | ICIE3 | - | OCIE3C | OCIE3B | OCIE3A | TOIE3 |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 17.11.35 TIMSK4 - Timer/Counter 4 Interrupt Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TIMSK4 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x72) | - | - | ICIE4 | - | OCIE4C | OCIE4B | OCIE4A | TOIE4 |  |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 17.11.36 TIMSK5 - Timer/Counter 5 Interrupt Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TIMSK5 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x73) | - | - | ICIE5 | - | OCIE5C | OCIE5B | OCIE5A | TOIE5 |  |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 5 - ICIEn: Timer/Countern, Input Capture Interrupt Enable

When this bit is written to one, and the I-flag in the Status Register is set (interrupts globally enabled), the Timer/Countern Input Capture interrupt is enabled. The corresponding Interrupt Vector (see "Interrupts" on page 101) is executed when the ICFn Flag, located in TIFRn, is set.

## - Bit 3 - OCIEnC: Timer/Countern, Output Compare C Match Interrupt Enable

When this bit is written to one, and the I-flag in the Status Register is set (interrupts globally enabled), the Timer/Countern Output Compare C Match interrupt is enabled. The corresponding Interrupt Vector (see "Interrupts" on page 101) is executed when the OCFnC Flag, located in TIFRn, is set.

- Bit 2 - OCIEnB: Timer/Countern, Output Compare B Match Interrupt Enable

When this bit is written to one, and the I-flag in the Status Register is set (interrupts globally enabled), the Timer/Countern Output Compare B Match interrupt is enabled. The corresponding Interrupt Vector (see "Interrupts" on page 101) is executed when the OCFnB Flag, located in TIFRn, is set.

- Bit 1 - OCIEnA: Timer/Countern, Output Compare A Match Interrupt Enable

When this bit is written to one, and the I-flag in the Status Register is set (interrupts globally enabled), the Timer/Countern Output Compare A Match interrupt is enabled. The corresponding Interrupt Vector (see "Interrupts" on page 101) is executed when the OCFnA Flag, located in TIFRn, is set.

- Bit 0 - TOIEn: Timer/Countern, Overflow Interrupt Enable

When this bit is written to one, and the I-flag in the Status Register is set (interrupts globally enabled), the Timer/Countern Overflow interrupt is enabled. The corresponding Interrupt Vector (see "Interrupts" on page 101) is executed when the TOVn Flag, located in TIFRn, is set.

### 17.11.37 TIFR1 - Timer/Counter1 Interrupt Flag Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TIFR1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x16 (0x36) | - | - | ICF1 | - | OCF1C | OCF1B | OCF1A | TOV1 |  |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

17.11.38 TIFR3 - Timer/Counter3 Interrupt Flag Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x18 (0x38) | - | - | ICF3 | - | OCF3C | OCF3B | OCF3A | TOV3 |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## TIFR3

### 17.11.39 TIFR4 - Timer/Counter4 Interrupt Flag Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x19 (0x39) | - | - | ICF4 | - | OCF4C | OCF4B | OCF4A | TOV4 |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

### 17.11.40 TIFR5 - Timer/Counter5 Interrupt Flag Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \times 1 \mathrm{~A}(0 \times 3 \mathrm{~A})$ | - | - | ICF5 | - | OCF5C | OCF5B | OCF5A | TOV5 |
| Read/Write | R | R | R/W | R | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

TIFR5

- Bit 5 - ICFn: Timer/Countern, Input Capture Flag

This flag is set when a capture event occurs on the ICPn pin. When the Input Capture Register (ICRn) is set by the WGMn3:0 to be used as the TOP value, the ICFn Flag is set when the counter reaches the TOP value.
ICFn is automatically cleared when the Input Capture Interrupt Vector is executed. Alternatively, ICFn can be cleared by writing a logic one to its bit location.

- Bit 3- OCFnC: Timer/Countern, Output Compare C Match Flag

This flag is set in the timer clock cycle after the counter (TCNTn) value matches the Output Compare Register C (OCRnC).

Note that a Forced Output Compare (FOCnC) strobe will not set the OCFnC Flag.
OCFnC is automatically cleared when the Output Compare Match C Interrupt Vector is executed. Alternatively, OCFnC can be cleared by writing a logic one to its bit location.

- Bit 2 - OCFnB: Timer/Counter1, Output Compare B Match Flag

This flag is set in the timer clock cycle after the counter (TCNTn) value matches the Output Compare Register B (OCRnB).

Note that a Forced Output Compare (FOCnB) strobe will not set the OCFnB Flag.
OCFnB is automatically cleared when the Output Compare Match B Interrupt Vector is executed. Alternatively, OCFnB can be cleared by writing a logic one to its bit location.

- Bit 1 - OCF1A: Timer/Counter1, Output Compare A Match Flag

This flag is set in the timer clock cycle after the counter (TCNTn value matches the Output Compare Register A (OCRnA).

Note that a Forced Output Compare (FOCnA) strobe will not set the OCFnA Flag.
OCFnA is automatically cleared when the Output Compare Match A Interrupt Vector is executed. Alternatively, OCFnA can be cleared by writing a logic one to its bit location.

## - Bit 0 - TOVn: Timer/Countern, Overflow Flag

The setting of this flag is dependent of the WGMn3:0 bits setting. In Normal and CTC modes, the TOVn Flag is set when the timer overflows. Refer to Table 17-2 on page 145 for the TOVn Flag behavior when using another WGMn3:0 bit setting.

TOVn is automatically cleared when the Timer/Countern Overflow Interrupt Vector is executed. Alternatively, TOVn can be cleared by writing a logic one to its bit location.

## 18. Timer/Counter 0, 1, 3, 4, and 5 Prescaler

Timer/Counter 0, 1, 3, 4, and 5 share the same prescaler module, but the Timer/Counters can have different prescaler settings. The description below applies to all Timer/Counters. Tn is used as a general name, $n=0,1,3,4$, or 5.

### 18.1 Internal Clock Source

The Timer/Counter can be clocked directly by the system clock (by setting the CSn2:0 = 1). This provides the fastest operation, with a maximum Timer/Counter clock frequency equal to system clock frequency (f $\mathrm{f}_{\text {CLK_I/O }}$ ). Alternatively, one of four taps from the prescaler can be used as a clock source. The prescaled clock has a frequency of either $\mathrm{f}_{\text {CLK_//O }} / 8, \mathrm{f}_{\text {CLK_// }} / 64, \mathrm{f}_{\text {CLK_// }} / 256$, or $\mathrm{f}_{\text {CLK_// }} / 1024$.

### 18.2 Prescaler Reset

The prescaler is free running, that is, operates independently of the Clock Select logic of the Timer/Counter, and it is shared by the Timer/Counter Tn. Since the prescaler is not affected by the Timer/Counter's clock select, the state of the prescaler will have implications for situations where a prescaled clock is used. One example of prescaling artifacts occurs when the timer is enabled and clocked by the prescaler ( $6>\operatorname{CSn} 2: 0>1$ ). The number of system clock cycles from when the timer is enabled to the first count occurs can be from 1 to $\mathrm{N}+1$ system clock cycles, where $N$ equals the prescaler divisor ( $8,64,256$, or 1024 ).

It is possible to use the prescaler reset for synchronizing the Timer/Counter to program execution. However, care must be taken if the other Timer/Counter that shares the same prescaler also uses prescaling. A prescaler reset will affect the prescaler period for all Timer/Counters it is connected to.

### 18.3 External Clock Source

An external clock source applied to the Tn pin can be used as Timer/Counter clock ( $\mathrm{clk}_{\mathrm{Tn}_{\mathrm{n}}}$ ). The Tn pin is sampled once every system clock cycle by the pin synchronization logic. The synchronized (sampled) signal is then passed through the edge detector. Figure 18-1 shows a functional equivalent block diagram of the Tn synchronization and edge detector logic. The registers are clocked at the positive edge of the internal system clock (clk//0). The latch is transparent in the high period of the internal system clock.

The edge detector generates one $\mathrm{clk}_{\text {Tn }}$ pulse for each positive (CSn2:0 $=7$ ) or negative (CSn2:0 $=6$ ) edge it detects.

Figure 18-1. Tn/TO Pin Sampling


The synchronization and edge detector logic introduces a delay of 2.5 to 3.5 system clock cycles from an edge has been applied to the Tn pin to the counter is updated.

Enabling and disabling of the clock input must be done when Tn has been stable for at least one system clock cycle, otherwise it is a risk that a false Timer/Counter clock pulse is generated.

Each half period of the external clock applied must be longer than one system clock cycle to ensure correct sampling. The external clock must be guaranteed to have less than half the system clock frequency ( $\mathrm{f}_{\text {ExtIIk }}<\mathrm{f}_{\text {clk_//0 }} / 2$ ) given a $50 / 50 \%$ duty cycle. Since the edge detector uses sampling, the maximum frequency of an external clock it can detect is half the sampling frequency (Nyquist sampling theorem). However, due to variation of the system clock frequency and duty cycle caused by Oscillator source (crystal, resonator, and capacitors) tolerances, it is recommended that maximum frequency of an external clock source is less than $f_{\text {clk_//o }} / 2.5$.

An external clock source can not be prescaled.
Figure 18-2. Prescaler for synchronous Timer/Counters


### 18.4 Register Description

### 18.4.1 GTCCR - General Timer/Counter Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | GTCCR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x23 (0x43) | TSM | - | - | - | - | - | PSRASY | PSRSYNC |  |
| Read/Write | R/W | R | R | R | R | R | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - TSM: Timer/Counter Synchronization Mode

Writing the TSM bit to one activates the Timer/Counter Synchronization mode. In this mode, the value that is written to the PSRASY and PSRSYNC bits is kept, hence keeping the corresponding prescaler reset signals asserted. This ensures that the corresponding Timer/Counters are halted and can be configured to the same value without the risk of one of them advancing during configuration. When the TSM bit is written to zero, the PSRASY and PSRSYNC bits are cleared by hardware, and the Timer/Counters start counting simultaneously.

## - Bit 0 - PSRSYNC: Prescaler Reset for Synchronous Timer/Counters

When this bit is one, Timer/Counter0, Timer/Counter1, Timer/Counter3, Timer/Counter4 and Timer/Counter5 prescaler will be Reset. This bit is normally cleared immediately by hardware, except if the TSM bit is set. Note that Timer/Counter0, Timer/Counter1, Timer/Counter3, Timer/Counter4 and Timer/Counter5 share the same prescaler and a reset of this prescaler will affect all timers.

## 19. Output Compare Modulator (OCM1COA)

### 19.1 Overview

The Output Compare Modulator (OCM) allows generation of waveforms modulated with a carrier frequency. The modulator uses the outputs from the Output Compare Unit C of the 16-bit Timer/Counter1 and the Output Compare Unit of the 8-bit Timer/Counter0. For more details about these Timer/Counters see "Timer/Counter 0, 1, 3, 4, and 5 Prescaler" on page 164 and "8-bit Timer/Counter2 with PWM and Asynchronous Operation" on page 169.

Figure 19-1. Output Compare Modulator, Block Diagram


When the modulator is enabled, the two output compare channels are modulated together as shown in the block diagram (see Figure 19-1).

### 19.2 Description

The Output Compare unit 1C and Output Compare unit 2 shares the PB7 port pin for output. The outputs of the Output Compare units (OC1C and OCOA) overrides the normal PORTB7 Register when one of them is enabled (that is, when COMnx1:0 is not equal to zero). When both OC1C and OCOA are enabled at the same time, the modulator is automatically enabled.

The functional equivalent schematic of the modulator is shown on Figure 19-2. The schematic includes part of the Timer/Counter units and the port B pin 7 output driver circuit.

Figure 19-2. Output Compare Modulator, Schematic


When the modulator is enabled the type of modulation (logical AND or OR) can be selected by the PORTB7 Register. Note that the DDRB7 controls the direction of the port independent of the COMnx1:0 bit setting.

### 19.2.1 Timing example

Figure 19-3 illustrates the modulator in action. In this example the Timer/Counter1 is set to operate in fast PWM mode (non-inverted) and Timer/CounterO uses CTC waveform mode with toggle Compare Output mode (COMnx1:0 = 1).

Figure 19-3. Output Compare Modulator, Timing Diagram


In this example, Timer/Counter2 provides the carrier, while the modulating signal is generated by the Output Compare unit C of the Timer/Counter1.

The resolution of the PWM signal (OC1C) is reduced by the modulation. The reduction factor is equal to the number of system clock cycles of one period of the carrier (OCOA). In this example the resolution is reduced by a factor of two. The reason for the reduction is illustrated in Figure 19-3 at the second and third period of the PB7 output when PORTB7 equals zero. The period 2 high time is one cycle longer than the period 3 high time, but the result on the PB7 output is equal in both periods.

## 20. 8-bit Timer/Counter2 with PWM and Asynchronous Operation

Timer/Counter2 is a general purpose, single channel, 8 -bit Timer/Counter module. The main features are:

- Single Channel Counter
- Clear Timer on Compare Match (Auto Reload)
- Glitch-free, Phase Correct Pulse Width Modulator (PWM)
- Frequency Generator
- 10-bit Clock Prescaler
- Overflow and Compare Match Interrupt Sources (TOV2, OCF2A and OCF2B)
- Allows Clocking from External 32kHz Watch Crystal Independent of the I/O Clock


### 20.1 Overview

A simplified block diagram of the 8 -bit Timer/Counter is shown in Figure 17-12 on page 153 For the actual placement of I/O pins, see "Pin Configurations" on page 2. CPU accessible I/O Registers, including I/O bits and I/O pins, are shown in bold. The device-specific I/O Register and bit locations are listed in the "Register Description" on page 182.
The Power Reduction Timer/Counter2 bit, PRTIM2, in "PRR0 - Power Reduction Register 0" on page 55 must be written to zero to enable Timer/Counter2 module.

Figure 20-1. 8 -bit Timer/Counter Block Diagram


### 20.1.1 Registers

The Timer/Counter (TCNT2) and Output Compare Register (OCR2A and OCR2B) are 8-bit registers. Interrupt request (abbreviated to Int.Req.) signals are all visible in the Timer Interrupt Flag Register (TIFR2). All interrupts are individually masked with the Timer Interrupt Mask Register (TIMSK2). TIFR2 and TIMSK2 are not shown in the figure.
The Timer/Counter can be clocked internally, via the prescaler, or asynchronously clocked from the TOSC $1 / 2$ pins, as detailed later in this section. The asynchronous operation is controlled by the Asynchronous Status Register (ASSR). The Clock Select logic block controls which clock source the Timer/Counter uses to increment (or decrement) its value. The Timer/Counter is inactive when no clock source is selected. The output from the Clock Select logic is referred to as the timer clock $\left(\mathrm{Clk}_{\mathrm{T}_{2}}\right)$.

The double buffered Output Compare Register (OCR2A and OCR2B) are compared with the Timer/Counter value at all times. The result of the compare can be used by the Waveform Generator to generate a PWM or variable frequency output on the Output Compare pins (OC2A and OC2B). See "Output Compare Unit" on page 175 for details. The compare match event will also set the Compare Flag (OCF2A or OCF2B) which can be used to generate an Output Compare interrupt request.

### 20.1.2 Definitions

Many register and bit references in this document are written in general form. A lower case " n " replaces the Timer/Counter number, in this case 2. However, when using the register or bit defines in a program, the precise form must be used, that is, TCNT2 for accessing Timer/Counter2 counter value and so on.

The definitions in Table 20-1 are also used extensively throughout the section.
Table 20-1. Definitions

| BOTTOM | The counter reaches the BOTTOM when it becomes zero (0x00) |
| :--- | :--- |
| MAX | The counter reaches its MAXimum when it becomes 0xFF (decimal 255) |
| TOP | The counter reaches the TOP when it becomes equal to the highest value in the count <br> sequence. The TOP value can be assigned to be the fixed value 0xFF (MAX) or the value stored <br> in the OCR2A Register. The assignment is dependent on the mode of operation |

### 20.2 Timer/Counter Clock Sources

The Timer/Counter can be clocked by an internal synchronous or an external asynchronous clock source. The clock source $\mathrm{clk}_{\mathrm{T} 2}$ is by default equal to the MCU clock, $\mathrm{clk}_{1 / 2}$. When the AS2 bit in the ASSR Register is written to logic one, the clock source is taken from the Timer/Counter Oscillator connected to TOSC1 and TOSC2. For details on asynchronous operation, see "Asynchronous Operation of Timer/Counter2" on page 179. For details on clock sources and prescaler, see "Timer/Counter Prescaler" on page 180.

### 20.3 Counter Unit

The main part of the 8 -bit Timer/Counter is the programmable bi-directional counter unit. Figure 20-2 on page 171 shows a block diagram of the counter and its surrounding environment.

Figure 20-2. Counter Unit Block Diagram


Signal description (internal signals):
count Increment or decrement TCNT2 by 1.
direction
clear Clear TCNT2 (set all bits to zero).
$\mathbf{c l k}_{\text {Tn }} \quad$ Timer/Counter clock, referred to as $\mathrm{Clk}_{\mathrm{T} 2}$ in the following.
top Signalizes that TCNT2 has reached maximum value.
bottom
Signalizes that TCNT2 has reached minimum value (zero).
Depending on the mode of operation used, the counter is cleared, incremented, or decremented at each timer clock ( $\mathrm{clk}_{\mathrm{T} 2}$ ). $\mathrm{cl}_{\mathrm{T} 2}$ can be generated from an external or internal clock source, selected by the Clock Select bits (CS22:0). When no clock source is selected (CS22:0 = 0) the timer is stopped. However, the TCNT2 value can be accessed by the CPU, regardless of whether $\mathrm{Clk}_{\mathrm{T} 2}$ is present or not. A CPU write overrides (has priority over) all counter clear or count operations.
The counting sequence is determined by the setting of the WGM21 and WGM20 bits located in the Timer/Counter Control Register (TCCR2A) and the WGM22 located in the Timer/Counter Control Register B (TCCR2B). There are close connections between how the counter behaves (counts) and how waveforms are generated on the Output Compare outputs OC2A and OC2B. For more details about advanced counting sequences and waveform generation, see "Modes of Operation".

The Timer/Counter Overflow Flag (TOV2) is set according to the mode of operation selected by the WGM22:0 bits. TOV2 can be used for generating a CPU interrupt.

### 20.4 Modes of Operation

The mode of operation, that is, the behavior of the Timer/Counter and the Output Compare pins, is defined by the combination of the Waveform Generation mode (WGM22:0) and Compare Output mode (COM2x1:0) bits. The Compare Output mode bits do not affect the counting sequence, while the Waveform Generation mode bits do. The COM2x1:0 bits control whether the PWM output generated should be inverted or not (inverted or non-inverted PWM). For non-PWM modes the COM2x1:0 bits control whether the output should be set, cleared, or toggled at a compare match. See "Compare Match Output Unit" on page 176.

For detailed timing information refer to "Timer/Counter Timing Diagrams" on page 177.

### 20.4.1 Normal Mode

The simplest mode of operation is the Normal mode (WGM22:0 = 0). In this mode the counting direction is always up (incrementing), and no counter clear is performed. The counter simply overruns when it passes its maximum 8bit value (TOP $=0 x F F$ ) and then restarts from the bottom ( $0 \times 00$ ). In normal operation the Timer/Counter Overflow Flag (TOV2) will be set in the same timer clock cycle as the TCNT2 becomes zero. The TOV2 Flag in this case behaves like a ninth bit, except that it is only set, not cleared. However, combined with the timer overflow interrupt that automatically clears the TOV2 Flag, the timer resolution can be increased by software. There are no special cases to consider in the Normal mode, a new counter value can be written anytime.

The Output Compare unit can be used to generate interrupts at some given time. Using the Output Compare to generate waveforms in Normal mode is not recommended, since this will occupy too much of the CPU time.

### 20.4.2 Clear Timer on Compare Match (CTC) Mode

In Clear Timer on Compare or CTC mode (WGM22:0 = 2), the OCR2A Register is used to manipulate the counter resolution. In CTC mode the counter is cleared to zero when the counter value (TCNT2) matches the OCR2A. The OCR2A defines the top value for the counter, hence also its resolution. This mode allows greater control of the compare match output frequency. It also simplifies the operation of counting external events.

The timing diagram for the CTC mode is shown in Figure 20-3. The counter value (TCNT2) increases until a compare match occurs between TCNT2 and OCR2A, and then counter (TCNT2) is cleared.

Figure 20-3. CTC Mode, Timing Diagram


An interrupt can be generated each time the counter value reaches the TOP value by using the OCF2A Flag. If the interrupt is enabled, the interrupt handler routine can be used for updating the TOP value. However, changing TOP to a value close to BOTTOM when the counter is running with none or a low prescaler value must be done with care since the CTC mode does not have the double buffering feature. If the new value written to OCR2A is lower than the current value of TCNT2, the counter will miss the compare match. The counter will then have to count to its maximum value ( $0 x F F$ ) and wrap around starting at $0 x 00$ before the compare match can occur.

For generating a waveform output in CTC mode, the OC2A output can be set to toggle its logical level on each compare match by setting the Compare Output mode bits to toggle mode (COM2A1:0 = 1). The OC2A value will not be visible on the port pin unless the data direction for the pin is set to output. The waveform generated will have a maximum frequency of $\mathrm{f}_{\mathrm{OC} 2 \mathrm{~A}}=\mathrm{f}_{\mathrm{Clk} \_/ / 0} / 2$ when OCR2A is set to zero ( $0 \times 00$ ). The waveform frequency is defined by the following equation:

$$
f_{O C n x}=\frac{f_{\mathrm{clk}_{\mathrm{l}} / \mathrm{O}}}{2 \cdot N \cdot(1+O C R n x)}
$$

The $N$ variable represents the prescale factor (1, 8, 32, $64,128,256$, or 1024 ).
As for the Normal mode of operation, the TOV2 Flag is set in the same timer clock cycle that the counter counts from MAX to $0 \times 00$.

### 20.4.3 Fast PWM Mode

Figure 20-4. Fast PWM Mode, Timing Diagram


The Timer/Counter Overflow Flag (TOV2) is set each time the counter reaches TOP. If the interrupt is enabled, the interrupt handler routine can be used for updating the compare value.

In fast PWM mode, the compare unit allows generation of PWM waveforms on the OC2x pin. Setting the COM $2 \times 1: 0$ bits to two will produce a non-inverted PWM and an inverted PWM output can be generated by setting the COM2x1:0 to three. TOP is defined as $0 \times$ FF when WGM2:0 $=3$, and OCR2A when WGM2:0 $=7$ (see Table $20-3$ on page 182). The actual OC2x value will only be visible on the port pin if the data direction for the port pin is set as output. The PWM waveform is generated by setting (or clearing) the OC2x Register at the compare match between OCR2x and TCNT2, and clearing (or setting) the OC2x Register at the timer clock cycle the counter is cleared (changes from TOP to BOTTOM).

The PWM frequency for the output can be calculated by the following equation:

$$
f_{O C n x P W M}=\frac{f_{\mathrm{clk} 1 / \mathrm{O}}}{N \cdot 256}
$$

The $N$ variable represents the prescale factor ( $1,8,32,64,128,256$, or 1024 ).
The extreme values for the OCR2A Register represent special cases when generating a PWM waveform output in the fast PWM mode. If the OCR2A is set equal to BOTTOM, the output will be a narrow spike for each MAX+1 timer clock cycle. Setting the OCR2A equal to MAX will result in a constantly high or low output (depending on the polarity of the output set by the COM2A1:0 bits).

A frequency (with $50 \%$ duty cycle) waveform output in fast PWM mode can be achieved by setting OC2x to toggle its logical level on each compare match $(C O M 2 \times 1: 0=1)$. The waveform generated will have a maximum frequency of $f_{\text {oc2 } 2}=f_{\text {clk_/ } / 0} / 2$ when OCR2A is set to zero. This feature is similar to the OC2A toggle in CTC mode, except the double buffer feature of the Output Compare unit is enabled in the fast PWM mode.

### 20.4.4 Phase Correct PWM Mode

The phase correct PWM mode (WGM22:0 = 1 or 5) provides a high resolution phase correct PWM waveform generation option. The phase correct PWM mode is based on a dual-slope operation. The counter counts repeatedly from BOTTOM to TOP and then from TOP to BOTTOM. TOP is defined as $0 x F F$ when WGM22:0 $=1$, and OCR2A when MGM22:0 = 5. In non-inverting Compare Output mode, the Output Compare (OC2x) is cleared on the compare match between TCNT2 and OCR2x while upcounting, and set on the compare match while downcounting. In inverting Output Compare mode, the operation is inverted. The dual-slope operation has lower maximum operation
frequency than single slope operation. However, due to the symmetric feature of the dual-slope PWM modes, these modes are preferred for motor control applications.

In phase correct PWM mode the counter is incremented until the counter value matches TOP. When the counter reaches TOP, it changes the count direction. The TCNT2 value will be equal to TOP for one timer clock cycle. The timing diagram for the phase correct PWM mode is shown on Figure 20-5. The TCNT2 value is in the timing diagram shown as a histogram for illustrating the dual-slope operation. The diagram includes non-inverted and inverted PWM outputs. The small horizontal line marks on the TCNT2 slopes represent compare matches between OCR2x and TCNT2.

Figure 20-5. Phase Correct PWM Mode, Timing Diagram


The Timer/Counter Overflow Flag (TOV2) is set each time the counter reaches BOTTOM. The Interrupt Flag can be used to generate an interrupt each time the counter reaches the BOTTOM value.

In phase correct PWM mode, the compare unit allows generation of PWM waveforms on the OC2x pin. Setting the COM2x1:0 bits to two will produce a non-inverted PWM. An inverted PWM output can be generated by setting the COM2x1:0 to three. TOP is defined as 0xFF when WGM2:0 = 3, and OCR2A when MGM2:0 $=7$ (see Table 20-4 on page 183). The actual OC2x value will only be visible on the port pin if the data direction for the port pin is set as output. The PWM waveform is generated by clearing (or setting) the OC2x Register at the compare match between OCR2x and TCNT2 when the counter increments, and setting (or clearing) the OC2x Register at compare match between OCR2x and TCNT2 when the counter decrements. The PWM frequency for the output when using phase correct PWM can be calculated by the following equation:

$$
f_{O C n \times P C P W M}=\frac{f_{\text {clk } 1 / \mathrm{O}}}{N \cdot 510}
$$

The $N$ variable represents the prescale factor ( $1,8,32,64,128,256$, or 1024 ).
The extreme values for the OCR2A Register represent special cases when generating a PWM waveform output in the phase correct PWM mode. If the OCR2A is set equal to BOTTOM, the output will be continuously low and if set equal to MAX the output will be continuously high for non-inverted PWM mode. For inverted PWM the output will have the opposite logic values.

At the very start of period 2 in Figure 20-5 on page 174 OCnx has a transition from high to low even though there is no Compare Match. The point of this transition is to guarantee symmetry around BOTTOM. There are two cases that give a transition without Compare Match.

- OCR2A changes its value from MAX, like in Figure 20-5 on page 174. When the OCR2A value is MAX the OCn pin value is the same as the result of a down-counting compare match. To ensure symmetry around BOTTOM the OCn value at MAX must correspond to the result of an up-counting Compare Match.
- The timer starts counting from a value higher than the one in OCR2A, and for that reason misses the Compare Match and hence the OCn change that would have happened on the way up.


### 20.5 Output Compare Unit

The 8-bit comparator continuously compares TCNT2 with the Output Compare Register (OCR2A and OCR2B). Whenever TCNT2 equals OCR2A or OCR2B, the comparator signals a match. A match will set the Output Compare Flag (OCF2A or OCF2B) at the next timer clock cycle. If the corresponding interrupt is enabled, the Output Compare Flag generates an Output Compare interrupt. The Output Compare Flag is automatically cleared when the interrupt is executed. Alternatively, the Output Compare Flag can be cleared by software by writing a logical one to its I/O bit location. The Waveform Generator uses the match signal to generate an output according to operating mode set by the WGM22:0 bits and Compare Output mode (COM2x1:0) bits. The max and bottom signals are used by the Waveform Generator for handling the special cases of the extreme values in some modes of operation (see "Modes of Operation" on page 171).
Figure 20-6 shows a block diagram of the Output Compare unit.
Figure 20-6. Output Compare Unit, Block Diagram


The OCR2x Register is double buffered when using any of the Pulse Width Modulation (PWM) modes. For the Normal and Clear Timer on Compare (CTC) modes of operation, the double buffering is disabled. The double buffering synchronizes the update of the OCR2x Compare Register to either top or bottom of the counting sequence. The synchronization prevents the occurrence of odd-length, non-symmetrical PWM pulses, thereby making the output glitch-free.

The OCR2x Register access may seem complex, but this is not case. When the double buffering is enabled, the CPU has access to the OCR2x Buffer Register, and if double buffering is disabled the CPU will access the OCR2x directly.

### 20.5.1 Force Output Compare

In non-PWM waveform generation modes, the match output of the comparator can be forced by writing a one to the Force Output Compare (FOC2x) bit. Forcing compare match will not set the OCF2x Flag or reload/clear the timer, but the OC2x pin will be updated as if a real compare match had occurred (the COM $2 \times 1: 0$ bits settings define whether the OC2x pin is set, cleared or toggled).

### 20.5.2 Compare Match Blocking by TCNT2 Write

All CPU write operations to the TCNT2 Register will block any compare match that occurs in the next timer clock cycle, even when the timer is stopped. This feature allows OCR2x to be initialized to the same value as TCNT2 without triggering an interrupt when the Timer/Counter clock is enabled.

### 20.5.3 Using the Output Compare Unit

Since writing TCNT2 in any mode of operation will block all compare matches for one timer clock cycle, there are risks involved when changing TCNT2 when using the Output Compare channel, independently of whether the Timer/Counter is running or not. If the value written to TCNT2 equals the OCR2x value, the compare match will be missed, resulting in incorrect waveform generation. Similarly, do not write the TCNT2 value equal to BOTTOM when the counter is downcounting.

The setup of the OC2x should be performed before setting the Data Direction Register for the port pin to output. The easiest way of setting the OC2x value is to use the Force Output Compare (FOC2x) strobe bit in Normal mode. The OC2x Register keeps its value even when changing between Waveform Generation modes.
Be aware that the COM2×1:0 bits are not double buffered together with the compare value. Changing the COM $2 \times 1: 0$ bits will take effect immediately.

### 20.6 Compare Match Output Unit

The Compare Output mode (COM2x1:0) bits have two functions. The Waveform Generator uses the COM $2 \times 1: 0$ bits for defining the Output Compare (OC2x) state at the next compare match. Also, the COM $2 \times 1: 0$ bits control the OC2x pin output source. Figure 20-7 on page 177 shows a simplified schematic of the logic affected by the COM $2 \times 1: 0$ bit setting. The I/O Registers, I/O bits, and I/O pins in the figure are shown in bold. Only the parts of the general I/O Port Control Registers (DDR and PORT) that are affected by the COM2x1:0 bits are shown. When referring to the OC2x state, the reference is for the internal OC2x Register, not the OC2x pin.

Figure 20-7. Compare Match Output Unit, Schematic


The general I/O port function is overridden by the Output Compare (OC2x) from the Waveform Generator if either of the COM2x1:0 bits are set. However, the OC2x pin direction (input or output) is still controlled by the Data Direction Register (DDR) for the port pin. The Data Direction Register bit for the OC2x pin (DDR_OC2x) must be set as output before the OC2x value is visible on the pin. The port override function is independent of the Waveform Generation mode.

The design of the Output Compare pin logic allows initialization of the OC2x state before the output is enabled. Note that some COM2×1:0 bit settings are reserved for certain modes of operation. See "Register Description" on page 182.

### 20.6.1 Compare Output Mode and Waveform Generation

The Waveform Generator uses the COM2x1:0 bits differently in normal, CTC, and PWM modes. For all modes, setting the COM2x1:0 $=0$ tells the Waveform Generator that no action on the OC2x Register is to be performed on the next compare match. For compare output actions in the non-PWM modes refer to Table 20-5 on page 183. For fast PWM mode, refer to Table 20-6 on page 183, and for phase correct PWM refer to Table 20-7 on page 184.

A change of the COM $2 \times 1: 0$ bits state will have effect at the first compare match after the bits are written. For nonPWM modes, the action can be forced to have immediate effect by using the FOC2x strobe bits.

### 20.7 Timer/Counter Timing Diagrams

The following figures show the Timer/Counter in synchronous mode, and the timer clock $\left(\mathrm{clk}_{\mathrm{T} 2}\right)$ is therefore shown as a clock enable signal. In asynchronous mode, $\mathrm{clk}_{1 / \mathrm{O}}$ should be replaced by the Timer/Counter Oscillator clock. The figures include information on when Interrupt Flags are set. Figure 20-8 on page 178 contains timing data for basic Timer/Counter operation. The figure shows the count sequence close to the MAX value in all modes other than phase correct PWM mode.

Figure 20-8. Timer/Counter Timing Diagram, no Prescaling


Figure 20-9 shows the same timing data, but with the prescaler enabled.
Figure 20-9. Timer/Counter Timing Diagram, with Prescaler ( $\mathrm{f}_{\mathrm{clk} \_/ / 0} / 8$ )


Figure $20-10$ shows the setting of OCF2A in all modes except CTC mode.
Figure 20-10. Timer/Counter Timing Diagram, Setting of OCF2A, with Prescaler ( $\mathrm{f}_{\mathrm{clk} \_/ 1 / 0} / 8$ )


Figure 20-11 shows the setting of OCF2A and the clearing of TCNT2 in CTC mode.
Figure 20-11. Timer/Counter Timing Diagram, Clear Timer on Compare Match mode, with Prescaler ( $\mathrm{f}_{\mathrm{clk} \_/ / 0} / 8$ )


### 20.8 Asynchronous Operation of Timer/Counter2

When Timer/Counter2 operates asynchronously, some considerations must be taken.

- Warning: When switching between asynchronous and synchronous clocking of Timer/Counter2, the Timer Registers TCNT2, OCR2x, and TCCR2x might be corrupted. A safe procedure for switching clock source is:

1. Disable the Timer/Counter2 interrupts by clearing OCIE2x and TOIE2.
2. Select clock source by setting AS2 as appropriate.
3. Write new values to TCNT2, OCR2x, and TCCR2x.
4. To switch to asynchronous operation: Wait for TCN2UB, OCR2xUB, and TCR2xUB.
5. Clear the Timer/Counter2 Interrupt Flags.
6. Enable interrupts, if needed.

- The CPU main clock frequency must be more than four times the Oscillator frequency.
- When writing to one of the registers TCNT2, OCR2x, or TCCR2x, the value is transferred to a temporary register, and latched after two positive edges on TOSC1. The user should not write a new value before the contents of the temporary register have been transferred to its destination. Each of the five mentioned registers have their individual temporary register, which means that, for example, writing to TCNT2 does not disturb an OCR2x write in progress. To detect that a transfer to the destination register has taken place, the Asynchronous Status Register - ASSR has been implemented.
- When entering Power-save or ADC Noise Reduction mode after having written to TCNT2, OCR2x, or TCCR2x, the user must wait until the written register has been updated if Timer/Counter2 is used to wake up the device. Otherwise, the MCU will enter sleep mode before the changes are effective. This is particularly important if any of the Output Compare2 interrupt is used to wake up the device, since the Output Compare function is disabled during writing to OCR2x or TCNT2. If the write cycle is not finished, and the MCU enters sleep mode before the corresponding OCR2xUB bit returns to zero, the device will never receive a compare match interrupt, and the MCU will not wake up.
- If Timer/Counter2 is used to wake the device up from Power-save or ADC Noise Reduction mode, precautions must be taken if the user wants to re-enter one of these modes: The interrupt logic needs one TOSC1 cycle to be reset. If the time between wake-up and re-entering sleep mode is less than one TOSC1 cycle, the interrupt will not occur, and the device will fail to wake up. If the user is in doubt whether the time before re-entering Power-save or ADC Noise Reduction mode is sufficient, the following algorithm can be used to ensure that one TOSC1 cycle has elapsed:

1. Write a value to TCCR2x, TCNT2, or OCR2x.
2. Wait until the corresponding Update Busy Flag in ASSR returns to zero.
3. Enter Power-save or ADC Noise Reduction mode.

- When the asynchronous operation is selected, the 32.768 kHz Oscillator for Timer/Counter2 is always running, except in Power-down and Standby modes. After a Power-up Reset or wake-up from Power-down or Standby mode, the user should be aware of the fact that this Oscillator might take as long as one second to stabilize. The user is advised to wait for at least one second before using Timer/Counter2 after power-up or wake-up from Power-down or Standby mode. The contents of all Timer/Counter2 Registers must be considered lost after a wake-up from Power-down or Standby mode due to unstable clock signal upon start-up, no matter whether the Oscillator is in use or a clock signal is applied to the TOSC1 pin.
- Description of wake up from Power-save or ADC Noise Reduction mode when the timer is clocked asynchronously: When the interrupt condition is met, the wake up process is started on the following cycle of the timer clock, that is, the timer is always advanced by at least one before the processor can read the counter value. After wake-up, the MCU is halted for four cycles, it executes the interrupt routine, and resumes execution from the instruction following SLEEP.
- Reading of the TCNT2 Register shortly after wake-up from Power-save may give an incorrect result. Since TCNT2 is clocked on the asynchronous TOSC clock, reading TCNT2 must be done through a register synchronized to the internal I/O clock domain. Synchronization takes place for every rising TOSC1 edge. When waking up from Power-save mode, and the I/O clock ( $\mathrm{clk}_{/ / 0}$ ) again becomes active, TCNT2 will read as the previous value (before entering sleep) until the next rising TOSC1 edge. The phase of the TOSC clock after waking up from Power-save mode is essentially unpredictable, as it depends on the wake-up time. The recommended procedure for reading TCNT2 is thus as follows:

1. Write any value to either of the registers OCR2x or TCCR2x.
2. Wait for the corresponding Update Busy Flag to be cleared.
3. Read TCNT2.

- During asynchronous operation, the synchronization of the Interrupt Flags for the asynchronous timer takes three processor cycles plus one timer cycle. The timer is therefore advanced by at least one before the processor can read the timer value causing the setting of the Interrupt Flag. The Output Compare pin is changed on the timer clock and is not synchronized to the processor clock.


### 20.9 Timer/Counter Prescaler

Figure 20-12. Prescaler for Timer/Counter2


The clock source for Timer/Counter2 is named $\mathrm{clk}_{\mathrm{T} 2 \mathrm{~S}} . \mathrm{Clk}_{\mathrm{T} 2 \mathrm{~S}}$ is by default connected to the main system I/O clock $\mathrm{clk}_{10}$. By setting the AS2 bit in ASSR, Timer/Counter2 is asynchronously clocked from the TOSC1 pin. This enables use of Timer/Counter2 as a Real Time Counter (RTC). When AS2 is set, pins TOSC1 and TOSC2 are disconnected from Port C. A crystal can then be connected between the TOSC1 and TOSC2 pins to serve as an independent clock source for Timer/Counter2. The Oscillator is optimized for use with a 32.768 kHz crystal. By setting the EXCLK bit in the ASSR, a 32 kHz external clock can be applied. See "ASSR - Asynchronous Status Register" on page 187 for details.

For Timer/Counter2, the possible prescaled selections are: $\mathrm{clk}_{\mathrm{T} 2 \mathrm{~S}} / 8, \mathrm{clk}_{\mathrm{T} 2 S} / 32$, $\mathrm{clk}_{\mathrm{T} 2 S} / 64, \mathrm{clk}_{\mathrm{T} 2 S} / 128$, $\mathrm{clk}_{\mathrm{T} 2 S} / 256$, and $\mathrm{clk}_{\mathrm{T} 2 \mathrm{~S}} / 1024$. Additionally, $\mathrm{clk}_{\mathrm{T} 2 \mathrm{~S}}$ as well as 0 (stop) may be selected. Setting the PSRASY bit in GTCCR resets the prescaler. This allows the user to operate with a predictable prescaler.

### 20.10 Register Description

### 20.10.1 TCCR2A -Timer/Counter Control Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCR2A |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xBO) | COM2A1 | COM2A0 | COM2B1 | COM2B0 | - | - | WGM21 | WGM20 |  |
| Read/Write | R/W | R/W | R/W | R/W | R | R | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7:6 - COM2A1:0: Compare Match Output A Mode

These bits control the Output Compare pin (OC2A) behavior. If one or both of the COM2A1:0 bits are set, the OC2A output overrides the normal port functionality of the I/O pin it is connected to. However, note that the Data Direction Register (DDR) bit corresponding to the OC2A pin must be set in order to enable the output driver.

When OC2A is connected to the pin, the function of the COM2A1:0 bits depends on the WGM22:0 bit setting. Table 20-2 shows the COM2A1:0 bit functionality when the WGM22:0 bits are set to a normal or CTC mode (nonPWM).

Table 20-2. Compare Output Mode, non-PWM Mode

| COM2A1 | COM2A0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC2A disconnected |
| 0 | 1 | Toggle OC2A on Compare Match |
| 1 | 0 | Clear OC2A on Compare Match |
| 1 | 1 | Set OC2A on Compare Match |

Table 20-3 shows the COM2A1:0 bit functionality when the WGM21:0 bits are set to fast PWM mode.
Table 20-3. Compare Output Mode, Fast PWM Mode ${ }^{(1)}$

| COM2A1 | COM2A0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC2A disconnected |
| 0 | 1 | WGM22 $=0:$ Normal Port Operation, OC2A Disconnected <br> WGM22 $=1:$ Toggle OC2A on Compare Match |
| 1 | 0 | Clear OC2A on Compare Match, set OC2A at BOTTOM |
| (non-inverting mode) |  |  |

Note: 1. A special case occurs when OCR2A equals TOP and COM2A1 is set. In this case, the Compare Match is ignored, but the set or clear is done at BOTTOM. See "Fast PWM Mode" on page 173 for more details.
Table 20-4 on page 183 shows the COM2A1:0 bit functionality when the WGM22:0 bits are set to phase correct PWM mode.

Table 20-4. Compare Output Mode, Phase Correct PWM Mode ${ }^{(1)}$

| COM2A1 | COM2A0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC2A disconnected |
| 0 | 1 | WGM22 $=0:$ Normal Port Operation, OC2A Disconnected <br> WGM22 = 1: Toggle OC2A on Compare Match |
| 1 | 0 | Clear OC2A on Compare Match when up-counting <br> Set OC2A on Compare Match when down-counting |
| 1 | 1 | Set OC2A on Compare Match when up-counting <br> Clear OC2A on Compare Match when down-counting |

Note: 1. A special case occurs when OCR2A equals TOP and COM2A1 is set. In this case, the Compare Match is ignored, but the set or clear is done at TOP. See "Phase Correct PWM Mode" on page 173 for more details.

## - Bits 5:4 - COM2B1:0: Compare Match Output B Mode

These bits control the Output Compare pin (OC2B) behavior. If one or both of the COM2B1:0 bits are set, the OC2B output overrides the normal port functionality of the I/O pin it is connected to. However, note that the Data Direction Register (DDR) bit corresponding to the OC2B pin must be set in order to enable the output driver.

When OC2B is connected to the pin, the function of the COM2B1:0 bits depends on the WGM22:0 bit setting. Table 20-5 shows the COM2B1:0 bit functionality when the WGM22:0 bits are set to a normal or CTC mode (nonPWM).

Table 20-5. Compare Output Mode, non-PWM Mode

| COM2B1 | COM2B0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC2B disconnected |
| 0 | 1 | Toggle OC2B on Compare Match |
| 1 | 0 | Clear OC2B on Compare Match |
| 1 | 1 | Set OC2B on Compare Match |

Table 20-6 shows the COM2B1:0 bit functionality when the WGM22:0 bits are set to fast PWM mode.
Table 20-6. $\quad$ Compare Output Mode, Fast PWM Mode ${ }^{(1)}$

| COM2B1 | COM2B0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC2B disconnected |
| 0 | 1 | Reserved |
| 1 | 0 | Clear OC2B on Compare Match, set OC2B at BOTTOM |
| (non-inverting mode) |  |  |

Note: 1. A special case occurs when OCR2B equals TOP and COM2B1 is set. In this case, the Compare Match is ignored, but the set or clear is done at BOTTOM. See "Fast PWM Mode" on page 173 for more details.

Table 20-7 shows the COM2B1:0 bit functionality when the WGM22:0 bits are set to phase correct PWM mode.
Table 20-7. Compare Output Mode, Phase Correct PWM Mode ${ }^{(1)}$

| COM2B1 | COM2B0 | Description |
| :---: | :---: | :---: |
| 0 | 0 | Normal port operation, OC2B disconnected |
| 0 | 1 | Reserved |
| 1 | 0 | Clear OC2B on Compare Match when up-counting <br> Set OC2B on Compare Match when down-counting |
| 1 | 1 | Set OC2B on Compare Match when up-counting <br> Clear OC2B on Compare Match when down-counting |

Note: 1. A special case occurs when OCR2B equals TOP and COM2B1 is set. In this case, the Compare Match is ignored, but the set or clear is done at TOP. See "Phase Correct PWM Mode" on page 173 for more details.

## - Bits 3, 2 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

- Bits 1:0 - WGM21:0: Waveform Generation Mode

Combined with the WGM22 bit found in the TCCR2B Register, these bits control the counting sequence of the counter, the source for maximum (TOP) counter value, and what type of waveform generation to be used, see Table 20-8. Modes of operation supported by the Timer/Counter unit are: Normal mode (counter), Clear Timer on Compare Match (CTC) mode, and two types of Pulse Width Modulation (PWM) modes (see "Modes of Operation" on page 171).

Table 20-8. Waveform Generation Mode Bit Description

| Mode | WGM2 | WGM1 | WGM0 | Timer/Counter Mode of <br> Operation | TOP | Update of <br> OCRx at | TOV Flag <br> Set on |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | Normal | 0xFF | Immediate | MAX |
| 1 | 0 | 0 | 1 | PWM, Phase Correct | 0xFF | TOP | BOTTOM |
| 2 | 0 | 1 | 0 | CTC | OCRA | Immediate | MAX |
| 3 | 0 | 1 | 1 | Fast PWM | $0 x F F$ | BOTTOM | MAX |
| 4 | 1 | 0 | 0 | Reserved | - | - | - |
| 5 | 1 | 0 | 1 | PWM, Phase Correct | OCRA | TOP | BOTTOM |
| 6 | 1 | 1 | 0 | Reserved | - | - | - |
| 7 | 1 | 1 | 1 | Fast PWM | OCRA | BOTTOM | TOP |

Notes: 1. MAX $=0 x F F$.
2. BOTTOM $=0 \times 00$.

### 20.10.2 TCCR2B - Timer/Counter Control Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TCCR2B |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xB1) | FOC2A | FOC2B | - | - | WGM22 | CS22 | CS21 | CS20 |  |
| Read/Write | W | W | R | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - FOC2A: Force Output Compare A

The FOC2A bit is only active when the WGM bits specify a non-PWM mode.
However, for ensuring compatibility with future devices, this bit must be set to zero when TCCR2B is written when operating in PWM mode. When writing a logical one to the FOC2A bit, an immediate Compare Match is forced on the Waveform Generation unit. The OC2A output is changed according to its COM2A1:0 bits setting. Note that the FOC2A bit is implemented as a strobe. Therefore it is the value present in the COM2A1:0 bits that determines the effect of the forced compare.

A FOC2A strobe will not generate any interrupt, nor will it clear the timer in CTC mode using OCR2A as TOP.
The FOC2A bit is always read as zero.

- Bit 6 - FOC2B: Force Output Compare B

The FOC2B bit is only active when the WGM bits specify a non-PWM mode.
However, for ensuring compatibility with future devices, this bit must be set to zero when TCCR2B is written when operating in PWM mode. When writing a logical one to the FOC2B bit, an immediate Compare Match is forced on the Waveform Generation unit. The OC2B output is changed according to its COM2B1:0 bits setting. Note that the FOC2B bit is implemented as a strobe. Therefore it is the value present in the COM2B1:0 bits that determines the effect of the forced compare.

A FOC2B strobe will not generate any interrupt, nor will it clear the timer in CTC mode using OCR2B as TOP.
The FOC2B bit is always read as zero.

## - Bits 5:4 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

- Bit 3 - WGM22: Waveform Generation Mode

See the description in the "TCCR2A -Timer/Counter Control Register A" on page 182.

- Bit 2:0 - CS22:0: Clock Select

The three Clock Select bits select the clock source to be used by the Timer/Counter, see Table 20-9.
Table 20-9. Clock Select Bit Description

| CS22 | CS21 | CS20 | Description |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | No clock source (Timer/Counter stopped) |
| 0 | 0 | 1 | $\mathrm{clk}_{\mathrm{T} 2 \mathrm{~S}} /$ (No prescaling) |
| 0 | 1 | 0 | $\mathrm{clk}_{\mathrm{T} 2 \mathrm{~S}} / 8$ (From prescaler) |
| 0 | 1 | 1 | $\mathrm{clk}_{\mathrm{T} 2 \mathrm{~S}} / 32$ (From prescaler) |
| 1 | 0 | 0 | $\mathrm{ck}_{\mathrm{T} 2 \mathrm{~S}} / 64$ (From prescaler) |

Table 20-9. Clock Select Bit Description (Continued)

| CS22 | CS21 | CS20 | Description |
| :---: | :---: | :---: | :---: |
| 1 | 0 | 1 | $\mathrm{ck}_{\mathrm{T} 2 \mathrm{~S}} / 128$ (From prescaler) |
| 1 | 1 | 0 | $\mathrm{ck}_{\mathrm{T} 2 \mathrm{~S}} / 256$ (From prescaler) |
| 1 | 1 | 1 | $\mathrm{ck}_{\mathrm{T} 2 \mathrm{~S}} / 1024$ (From prescaler) |

If external pin modes are used for the Timer/CounterO, transitions on the TO pin will clock the counter even if the pin is configured as an output. This feature allows software control of the counting.

### 20.10.3 TCNT2 - Timer/Counter Register



The Timer/Counter Register gives direct access, both for read and write operations, to the Timer/Counter unit 8-bit counter. Writing to the TCNT2 Register blocks (removes) the Compare Match on the following timer clock. Modifying the counter (TCNT2) while the counter is running, introduces a risk of missing a Compare Match between TCNT2 and the OCR2x Registers.

### 20.10.4 OCR2A - Output Compare Register A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xB3) | OCR2A[7:0] |  |  |  |  |  |  |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

OCR2A

The Output Compare Register A contains an 8-bit value that is continuously compared with the counter value (TCNT2). A match can be used to generate an Output Compare interrupt, or to generate a waveform output on the OC2A pin.

### 20.10.5 OCR2B - Output Compare Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xB4) | OCR2B[7:0] |  |  |  |  |  |  |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

OCR2B

The Output Compare Register B contains an 8-bit value that is continuously compared with the counter value (TCNT2). A match can be used to generate an Output Compare interrupt, or to generate a waveform output on the OC2B pin.

### 20.10.6 ASSR - Asynchronous Status Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | ASSR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xB6) | - | EXCLK | AS2 | TCN2UB | OCR2AUB | OCR2BUB | TCR2AUB | TCR2BUB |  |
| Read/Write | R | R/W | R/W | R | R | R | R | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 6 - EXCLK: Enable External Clock Input

When EXCLK is written to one, and asynchronous clock is selected, the external clock input buffer is enabled and an external clock can be input on Timer Oscillator 1 (TOSC1) pin instead of a 32 kHz crystal. Writing to EXCLK should be done before asynchronous operation is selected. Note that the crystal Oscillator will only run when this bit is zero.

## - Bit 5 - AS2: Asynchronous Timer/Counter2

When AS2 is written to zero, Timer/Counter2 is clocked from the I/O clock, clk ${ }_{1 / 0}$. When AS2 is written to one, Timer/Counter2 is clocked from a crystal Oscillator connected to the Timer Oscillator 1 (TOSC1) pin. When the value of AS2 is changed, the contents of TCNT2, OCR2A, OCR2B, TCCR2A and TCCR2B might be corrupted.

## - Bit 4 - TCN2UB: Timer/Counter2 Update Busy

When Timer/Counter2 operates asynchronously and TCNT2 is written, this bit becomes set. When TCNT2 has been updated from the temporary storage register, this bit is cleared by hardware. A logical zero in this bit indicates that TCNT2 is ready to be updated with a new value.

## - Bit 3 - OCR2AUB: Output Compare Register2 Update Busy

When Timer/Counter2 operates asynchronously and OCR2A is written, this bit becomes set. When OCR2A has been updated from the temporary storage register, this bit is cleared by hardware. A logical zero in this bit indicates that OCR2A is ready to be updated with a new value.

- Bit 2 - OCR2BUB: Output Compare Register2 Update Busy

When Timer/Counter2 operates asynchronously and OCR2B is written, this bit becomes set. When OCR2B has been updated from the temporary storage register, this bit is cleared by hardware. A logical zero in this bit indicates that OCR2B is ready to be updated with a new value.

- Bit 1 - TCR2AUB: Timer/Counter Control Register2 Update Busy

When Timer/Counter2 operates asynchronously and TCCR2A is written, this bit becomes set. When TCCR2A has been updated from the temporary storage register, this bit is cleared by hardware. A logical zero in this bit indicates that TCCR2A is ready to be updated with a new value.

## - Bit 0 - TCR2BUB: Timer/Counter Control Register2 Update Busy

When Timer/Counter2 operates asynchronously and TCCR2B is written, this bit becomes set. When TCCR2B has been updated from the temporary storage register, this bit is cleared by hardware. A logical zero in this bit indicates that TCCR2B is ready to be updated with a new value.

If a write is performed to any of the five Timer/Counter2 Registers while its update busy flag is set, the updated value might get corrupted and cause an unintentional interrupt to occur.

The mechanisms for reading TCNT2, OCR2A, OCR2B, TCCR2A and TCCR2B are different. When reading TCNT2, the actual timer value is read. When reading OCR2A, OCR2B, TCCR2A and TCCR2B the value in the temporary storage register is read.

### 20.10.7 TIMSK2 - Timer/Counter2 Interrupt Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TIMSK2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x70) | - | - | - | - | - | OCIE2B | OCIE2A | TOIE2 |  |
| Read/Write | R | R | R | R | R | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 2 - OCIE2B: Timer/Counter2 Output Compare Match B Interrupt Enable

When the OCIE2B bit is written to one and the I-bit in the Status Register is set (one), the Timer/Counter2 Compare Match B interrupt is enabled. The corresponding interrupt is executed if a compare match in Timer/Counter2 occurs, that is, when the OCF2B bit is set in the Timer/Counter 2 Interrupt Flag Register - TIFR2.

## - Bit 1 - OCIE2A: Timer/Counter2 Output Compare Match A Interrupt Enable

When the OCIE2A bit is written to one and the l-bit in the Status Register is set (one), the Timer/Counter2 Compare Match A interrupt is enabled. The corresponding interrupt is executed if a compare match in Timer/Counter2 occurs, that is, when the OCF2A bit is set in the Timer/Counter 2 Interrupt Flag Register - TIFR2.

## - Bit 0 - TOIE2: Timer/Counter2 Overflow Interrupt Enable

When the TOIE2 bit is written to one and the I-bit in the Status Register is set (one), the Timer/Counter2 Overflow interrupt is enabled. The corresponding interrupt is executed if an overflow in Timer/Counter2 occurs, that is, when the TOV2 bit is set in the Timer/Counter2 Interrupt Flag Register - TIFR2.

### 20.10.8 TIFR2 - Timer/Counter2 Interrupt Flag Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x17 (0x37) | - | - | - | - | - | OCF2B | OCF2A | TOV2 |
| Read/Write | R | R | R | R | R | R/W | R/W | R/W |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

## - Bit 2 - OCF2B: Output Compare Flag 2 B

The OCF2B bit is set (one) when a compare match occurs between the Timer/Counter2 and the data in OCR2B Output Compare Register2. OCF2B is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, OCF2B is cleared by writing a logic one to the flag. When the I-bit in SREG, OCIE2B (Timer/Counter2 Compare match Interrupt Enable), and OCF2B are set (one), the Timer/Counter2 Compare match Interrupt is executed.

## - Bit 1 - OCF2A: Output Compare Flag 2 A

The OCF2A bit is set (one) when a compare match occurs between the Timer/Counter2 and the data in OCR2A Output Compare Register2. OCF2A is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, OCF2A is cleared by writing a logic one to the flag. When the I-bit in SREG, OCIE2A (Timer/Counter2 Compare match Interrupt Enable), and OCF2A are set (one), the Timer/Counter2 Compare match Interrupt is executed.

## - Bit 0 - TOV2: Timer/Counter2 Overflow Flag

The TOV2 bit is set (one) when an overflow occurs in Timer/Counter2. TOV2 is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, TOV2 is cleared by writing a logic one to the flag. When the SREG I-bit, TOIE2A (Timer/Counter2 Overflow Interrupt Enable), and TOV2 are set (one), the Timer/Counter2 Overflow interrupt is executed. In PWM mode, this bit is set when Timer/Counter2 changes counting direction at $0 \times 00$.

### 20.10.9 GTCCR - General Timer/Counter Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | GTCCR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x23 (0x43) | TSM | - | - | - | - | - | PSRASY | PSRSYNC |  |
| Read/Write | R/W | R | R | R | R | R | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 1 - PSRASY: Prescaler Reset Timer/Counter2

When this bit is one, the Timer/Counter2 prescaler will be reset. This bit is normally cleared immediately by hardware. If the bit is written when Timer/Counter2 is operating in asynchronous mode, the bit will remain one until the prescaler has been reset. The bit will not be cleared by hardware if the TSM bit is set. Refer to the description of the "Bit 7 - TSM: Timer/Counter Synchronization Mode" on page 166 for a description of the Timer/Counter Synchronization mode.

## 21. SPI - Serial Peripheral Interface

The Serial Peripheral Interface (SPI) allows high-speed synchronous data transfer between the ATmega640/1280/1281/2560/2561 and peripheral devices or between several AVR devices. The ATmega640/1280/1281/2560/2561 SPI includes the following features:

- Full-duplex, Three-wire Synchronous Data Transfer
- Master or Slave Operation
- LSB First or MSB First Data Transfer
- Seven Programmable Bit Rates
- End of Transmission Interrupt Flag
- Write Collision Flag Protection
- Wake-up from Idle Mode
- Double Speed (CK/2) Master SPI Mode

USART can also be used in Master SPI mode, see "USART in SPI Mode" on page 227.
The Power Reduction SPI bit, PRSPI, in "PRR0 - Power Reduction Register 0 " on page 55 on page 50 must be written to zero to enable SPI module.

Figure 21-1. SPI Block Diagram ${ }^{(1)}$


Note: 1. Refer to Figure 1-1 on page 2, and Table 13-6 on page 76 for SPI pin placement.

The interconnection between Master and Slave CPUs with SPI is shown in Figure 21-2. The system consists of two shift Registers, and a Master clock generator. The SPI Master initiates the communication cycle when pulling low the Slave Select SS pin of the desired Slave. Master and Slave prepare the data to be sent in their respective shift Registers, and the Master generates the required clock pulses on the SCK line to interchange data. Data is always shifted from Master to Slave on the Master Out - Slave In, MOSI, line, and from Slave to Master on the Master In Slave Out, MISO, line. After each data packet, the Master will synchronize the Slave by pulling high the Slave Select, $\overline{\text { SS, }}$, line.

When configured as a Master, the SPI interface has no automatic control of the $\overline{\mathrm{SS}}$ line. This must be handled by user software before communication can start. When this is done, writing a byte to the SPI Data Register starts the SPI clock generator, and the hardware shifts the eight bits into the Slave. After shifting one byte, the SPI clock generator stops, setting the end of Transmission Flag (SPIF). If the SPI Interrupt Enable bit (SPIE) in the SPCR Register is set, an interrupt is requested. The Master may continue to shift the next byte by writing it into SPDR, or signal the end of packet by pulling high the Slave Select, SS line. The last incoming byte will be kept in the Buffer Register for later use.
When configured as a Slave, the SPI interface will remain sleeping with MISO tri-stated as long as the $\overline{\text { SS }}$ pin is driven high. In this state, software may update the contents of the SPI Data Register, SPDR, but the data will not be shifted out by incoming clock pulses on the SCK pin until the $\overline{\mathrm{SS}}$ pin is driven low. As one byte has been completely shifted, the end of Transmission Flag, SPIF is set. If the SPI Interrupt Enable bit, SPIE, in the SPCR Register is set, an interrupt is requested. The Slave may continue to place new data to be sent into SPDR before reading the incoming data. The last incoming byte will be kept in the Buffer Register for later use.

Figure 21-2. SPI Master-slave Interconnection


The system is single buffered in the transmit direction and double buffered in the receive direction. This means that bytes to be transmitted cannot be written to the SPI Data Register before the entire shift cycle is completed. When receiving data, however, a received character must be read from the SPI Data Register before the next character has been completely shifted in. Otherwise, the first byte is lost.

In SPI Slave mode, the control logic will sample the incoming signal of the SCK pin. To ensure correct sampling of the clock signal, the minimum low and high periods should be:

Low period: Longer than two CPU clock cycles.
High period: Longer than two CPU clock cycles.

When the SPI is enabled, the data direction of the MOSI, MISO, SCK, and $\overline{S S}$ pins is overridden according to Table 21-1. For more details on automatic port overrides, refer to "Alternate Port Functions" on page 72.

Table 21-1. $\quad$ SPI Pin Overrides ${ }^{(1)}$

| Pin | Direction, Master SPI | Direction, Slave SPI |
| :---: | :--- | :--- |
| MOSI | User Defined | Input |
| MISO | Input | User Defined |
| SCK | User Defined | Input |
| $\overline{\text { SS }}$ | User Defined | Input |

Note: 1. See "Alternate Functions of Port B" on page 76 for a detailed description of how to define the direction of the user defined SPI pins.
The following code examples show how to initialize the SPI as a Master and how to perform a simple transmission. DDR_SPI in the examples must be replaced by the actual Data Direction Register controlling the SPI pins. DD_MOSI, DD_MISO and DD_SCK must be replaced by the actual data direction bits for these pins. For example, if MOSI is placed on pin PB5, replace DD_MOSI with DDB5 and DDR_SPI with DDRB.

```
Assembly Code Example \({ }^{(1)}\)
    SPI_MasterInit:
    ; Set MOSI and SCK output, all others input
    ldi \(r 17,(1 \ll\) DD_MOSI \() \mid\left(1 \ll D D \_S C K\right)\)
    out DDR_SPI,r17
    ; Enable SPI, Master, set clock rate fck/16
    ldi \(r 17,(1 \ll S P E)|(1 \ll M S T R)|(1 \ll S P R 0)\)
    out SPCR,r17
    ret
SPI_MasterTransmit:
    ; Start transmission of data (r16)
    out SPDR,r16
Wait_Transmit:
    ; Wait for transmission complete
    sbis SPSR,SPIF
    rjmp Wait_Transmit
    ret
C Code Example \({ }^{(1)}\)
    void SPI_MasterInit(void)
    \{
        /* Set MOSI and SCK output, all others input */
        DDR_SPI \(=\left(1 \ll D D \_M O S I\right) \mid\left(1 \ll D D \_S C K\right) ;\)
        /* Enable SPI, Master, set clock rate fck/16 */
        \(\mathrm{SPCR}=(1 \ll \mathrm{SPE})|(1 \ll \mathrm{MSTR})|(1 \ll \mathrm{SPR} 0)\);
    \}
    void SPI_MasterTransmit(char cData)
    \{
        /* Start transmission */
        SPDR = cData;
        /* Wait for transmission complete */
        while(! (SPSR \& ( \(1 \ll\) SPIF) ))
        ;
    \}
```

Note: 1. See "About Code Examples" on page 10.

The following code examples show how to initialize the SPI as a Slave and how to perform a simple reception.

```
Assembly Code Example \({ }^{(1)}\)
    SPI_SlaveInit:
    ; Set MISO output, all others input
    1di r17,(1<<DD_MISO)
    out DDR_SPI,r17
        ; Enable SPI
        ldi r17,(1<<SPE)
        out SPCR,r17
        ret
    SPI_SlaveReceive:
        ; Wait for reception complete
        sbis SPSR,SPIF
        rjmp SPI_SlaveReceive
        ; Read received data and return
        in r16,SPDR
        ret
C Code Example \({ }^{(1)}\)
    void SPI_SlaveInit(void)
    \{
        /* Set MISO output, all others input */
        DDR_SPI = (1<<DD_MISO);
        /* Enable SPI */
        SPCR = (1<<SPE);
    \}
    char SPI_SlaveReceive(void)
    \{
        /* Wait for reception complete */
        while(!(SPSR \& (1<<SPIF)))
            ;
        /* Return Data Register */
        return SPDR;
    \}
```

Note: 1. See "About Code Examples" on page 10.

### 21.1 SS Pin Functionality

### 21.1.1 Slave Mode

When the SPI is configured as a Slave, the Slave Select $(\overline{\mathrm{SS}})$ pin is always input. When $\overline{\mathrm{SS}}$ is held low, the SPI is activated, and MISO becomes an output if configured so by the user. All other pins are inputs. When $\overline{\mathrm{SS}}$ is driven high, all pins are inputs, and the SPI is passive, which means that it will not receive incoming data. Note that the SPI logic will be reset once the $\overline{\mathrm{SS}}$ pin is driven high.

The $\overline{\mathrm{SS}}$ pin is useful for packet/byte synchronization to keep the slave bit counter synchronous with the master clock generator. When the SS pin is driven high, the SPI slave will immediately reset the send and receive logic, and drop any partially received data in the Shift Register.

### 21.1.2 Master Mode

When the SPI is configured as a Master (MSTR in SPCR is set), the user can determine the direction of the $\overline{\mathrm{SS}}$ pin.
If $\overline{S S}$ is configured as an output, the pin is a general output pin which does not affect the SPI system. Typically, the pin will be driving the $\overline{S S}$ pin of the SPI Slave.
If $\overline{\mathrm{SS}}$ is configured as an input, it must be held high to ensure Master SPI operation. If the $\overline{\mathrm{SS}}$ pin is driven low by peripheral circuitry when the SPI is configured as a Master with the $\overline{\mathrm{SS}}$ pin defined as an input, the SPI system interprets this as another master selecting the SPI as a slave and starting to send data to it. To avoid bus contention, the SPI system takes the following actions:

1. The MSTR bit in SPCR is cleared and the SPI system becomes a Slave. As a result of the SPI becoming a Slave, the MOSI and SCK pins become inputs.
2. The SPIF Flag in SPSR is set, and if the SPI interrupt is enabled, and the I-bit in SREG is set, the interrupt routine will be executed.
Thus, when interrupt-driven SPI transmission is used in Master mode, and there exists a possibility that $\overline{\mathrm{SS}}$ is driven low, the interrupt should always check that the MSTR bit is still set. If the MSTR bit has been cleared by a slave select, it must be set by the user to re-enable SPI Master mode.

### 21.1.3 Data Modes

There are four combinations of SCK phase and polarity with respect to serial data, which are determined by control bits CPHA and CPOL. The SPI data transfer formats are shown in Figure 21-3 on page 196 and Figure 21-4 on page 196. Data bits are shifted out and latched in on opposite edges of the SCK signal, ensuring sufficient time for data signals to stabilize. This is clearly seen by summarizing Table 21-3 on page 197 and Table 21-4 on page 197 in Table 21-2.

Table 21-2. CPOL Functionality

|  | Leading Edge | Trailing Edge | SPI Mode |
| :---: | :---: | :---: | :---: |
| $\mathrm{CPOL}=0, \mathrm{CPHA}=0$ | Sample (Rising) | Setup (Falling) | 0 |
| $\mathrm{CPOL}=0, \mathrm{CPHA}=1$ | Setup (Rising) | Sample (Falling) | 1 |
| $\mathrm{CPOL}=1, \mathrm{CPHA}=0$ | Sample (Falling) | Setup (Rising) | 2 |
| $\mathrm{CPOL}=1, \mathrm{CPHA}=1$ | Setup (Falling) | Sample (Rising) | 3 |

Figure 21-3. SPI Transfer Format with $\mathrm{CPHA}=0$


Figure 21-4. SPI Transfer Format with CPHA = 1


### 21.2 Register Description

### 21.2.1 SPCR - SPI Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | SPCR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \times 2 \mathrm{C}$ (0x4C) | SPIE | SPE | DORD | MSTR | CPOL | CPHA | SPR1 | SPR0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - SPIE: SPI Interrupt Enable

This bit causes the SPI interrupt to be executed if SPIF bit in the SPSR Register is set and the if the Global Interrupt Enable bit in SREG is set.

## - Bit 6 - SPE: SPI Enable

When the SPE bit is written to one, the SPI is enabled. This bit must be set to enable any SPI operations.

## - Bit 5 - DORD: Data Order

When the DORD bit is written to one, the LSB of the data word is transmitted first.
When the DORD bit is written to zero, the MSB of the data word is transmitted first.

## - Bit 4 - MSTR: Master/Slave Select

This bit selects Master SPI mode when written to one, and Slave SPI mode when written logic zero. If $\overline{\text { SS }}$ is configured as an input and is driven low while MSTR is set, MSTR will be cleared, and SPIF in SPSR will become set. The user will then have to set MSTR to re-enable SPI Master mode.

## - Bit 3 - CPOL: Clock Polarity

When this bit is written to one, SCK is high when idle. When CPOL is written to zero, SCK is low when idle. Refer to Figure 21-3 on page 196 and Figure 21-4 on page 196 for an example. The CPOL functionality is summarized in Table 21-3.

Table 21-3. CPOL Functionality

| CPOL | Leading Edge | Trailing Edge |
| :---: | :---: | :---: |
| 0 | Rising | Falling |
| 1 | Falling | Rising |

## - Bit 2 - CPHA: Clock Phase

The settings of the Clock Phase bit (CPHA) determine if data is sampled on the leading (first) or trailing (last) edge of SCK. Refer to Figure 21-3 on page 196 and Figure 21-4 on page 196 for an example. The CPOL functionality is summarized in Table 21-4.

Table 21-4. CPHA Functionality

| CPHA | Leading Edge | Trailing Edge |
| :---: | :---: | :---: |
| 0 | Sample | Setup |
| 1 | Setup | Sample |

## - Bits 1, 0 - SPR1, SPR0: SPI Clock Rate Select 1 and 0

These two bits control the SCK rate of the device configured as a Master. SPR1 and SPR0 have no effect on the Slave. The relationship between SCK and the Oscillator Clock frequency $f_{\text {osc }}$ is shown in Table 21-5.

Table 21-5. Relationship Between SCK and the Oscillator Frequency

| SPI2X | SPR1 | SPR0 | SCK Frequency |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | $\mathrm{f}_{\text {osd }} / 4$ |
| 0 | 0 | 1 | $\mathrm{f}_{\text {osd }} / 16$ |
| 0 | 1 | 0 | $\mathrm{f}_{\text {osd }} / 64$ |
| 0 | 1 | 1 | $\mathrm{f}_{\text {osd }} / 128$ |
| 1 | 0 | 0 | $\mathrm{f}_{\text {osd }} / 2$ |
| 1 | 0 | 1 | $\mathrm{f}_{\text {osd }} / 8$ |
| 1 | 1 | 0 | $\mathrm{f}_{\text {osc }} / 32$ |
| 1 | 1 | 1 | $\mathrm{f}_{\text {osc }} / 64$ |

### 21.2.2 SPSR - SPI Status Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | SPSR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x2D (0x4D) | SPIF | WCOL | - | - | - | - | - | SPI2X |  |
| Read/Write | R | R | R | R | R | R | R | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - SPIF: SPI Interrupt Flag

When a serial transfer is complete, the SPIF Flag is set. An interrupt is generated if SPIE in SPCR is set and global interrupts are enabled. If $\overline{S S}$ is an input and is driven low when the SPI is in Master mode, this will also set the SPIF Flag. SPIF is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, the SPIF bit is cleared by first reading the SPI Status Register with SPIF set, then accessing the SPI Data Register (SPDR).

## - Bit 6 - WCOL: Write COLlision Flag

The WCOL bit is set if the SPI Data Register (SPDR) is written during a data transfer. The WCOL bit (and the SPIF bit) are cleared by first reading the SPI Status Register with WCOL set, and then accessing the SPI Data Register.

## - Bit 5:1 - Res: Reserved Bits

These bits are reserved bits and will always read as zero.

## - Bit 0 - SPI2X: Double SPI Speed Bit

When this bit is written logic one the SPI speed (SCK Frequency) will be doubled when the SPI is in Master mode (see Table 21-5). This means that the minimum SCK period will be two CPU clock periods. When the SPI is configured as Slave, the SPI is only guaranteed to work at $\mathrm{f}_{\text {osc }} / 4$ or lower.

The SPI interface on the ATmega640/1280/1281/2560/2561 is also used for program memory and EEPROM downloading or uploading. See "Serial Downloading" on page 338 for serial programming and verification.

### 21.2.3 SPDR - SPI Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x2E (0x4E) | MSB |  |  |  |  |  |  | LSB | SPDR |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | X | X | X | X | X | X | X | X | Undefined |

The SPI Data Register is a read/write register used for data transfer between the Register File and the SPI Shift Register. Writing to the register initiates data transmission. Reading the register causes the Shift Register Receive buffer to be read.

## 22. USART

### 22.1 Features

- Full Duplex Operation (Independent Serial Receive and Transmit Registers)
- Asynchronous or Synchronous Operation
- Master or Slave Clocked Synchronous Operation
- High Resolution Baud Rate Generator
- Supports Serial Frames with 5, 6, 7, 8, or 9 Data Bits and 1 or 2 Stop Bits
- Odd or Even Parity Generation and Parity Check Supported by Hardware
- Data OverRun Detection
- Framing Error Detection
- Noise Filtering Includes False Start Bit Detection and Digital Low Pass Filter
- Three Separate Interrupts on TX Complete, TX Data Register Empty and RX Complete
- Multi-processor Communication Mode
- Double Speed Asynchronous Communication Mode


### 22.2 Overview

The Universal Synchronous and Asynchronous serial Receiver and Transmitter (USART) is a highly flexible serial communication device.

The ATmega640/1280/2560 has four USART's, USART0, USART1, USART2, and USART3. The functionality for all four USART's is described below. USART0, USART1, USART2, and USART3 have different I/O registers as shown in "Register Summary" on page 399.
A simplified block diagram of the USART Transmitter is shown in Figure 22-1 on page 201. CPU accessible I/O Registers and I/O pins are shown in bold.

The Power Reducion USART0 bit, PRUSART0, in "PRR0 - Power Reduction Register 0" on page 55 must be disabled by writing a logical zero to it.

The Power Reducion USART1 bit, PRUSART1, in "PRR1 - Power Reduction Register 1" on page 56 must be disabled by writing a logical zero to it.

The Power Reducion USART2 bit, PRUSART2, in "PRR1 - Power Reduction Register 1" on page 56 must be disabled by writing a logical zero to it.

The Power Reducion USART3 bit, PRUSART3, in "PRR1 - Power Reduction Register 1" on page 56 must be disabled by writing a logical zero to it.

Figure 22-1. USART Block Diagram ${ }^{(1)}$


Note: 1. See Figure 1-1 on page 2, Figure 1-3 on page 4, Table 13-12 on page 80, Table 13-15 on page 82, Table 13-24 on page 88 and Table 13-27 on page 90 for USART pin placement.
The dashed boxes in the block diagram separate the three main parts of the USART (listed from the top): Clock Generator, Transmitter and Receiver. Control Registers are shared by all units. The Clock Generation logic consists of synchronization logic for external clock input used by synchronous slave operation, and the baud rate generator. The XCKn (Transfer Clock) pin is only used by synchronous transfer mode. The Transmitter consists of a single write buffer, a serial Shift Register, Parity Generator and Control logic for handling different serial frame formats. The write buffer allows a continuous transfer of data without any delay between frames. The Receiver is the most complex part of the USART module due to its clock and data recovery units. The recovery units are used for asynchronous data reception. In addition to the recovery units, the Receiver includes a Parity Checker, Control logic, a Shift Register and a two level receive buffer (UDRn). The Receiver supports the same frame formats as the Transmitter, and can detect Frame Error, Data OverRun and Parity Errors.

### 22.3 Clock Generation

The Clock Generation logic generates the base clock for the Transmitter and Receiver. The USARTn supports four modes of clock operation: Normal asynchronous, Double Speed asynchronous, Master synchronous and Slave synchronous mode. The UMSELn bit in USART Control and Status Register C (UCSRnC) selects between asynchronous and synchronous operation. Double Speed (asynchronous mode only) is controlled by the U2Xn found in the UCSRnA Register. When using synchronous mode (UMSELn = 1), the Data Direction Register for the XCKn pin (DDR_XCKn) controls whether the clock source is internal (Master mode) or external (Slave mode). The XCKn pin is only active when using synchronous mode.

Figure 22-2 shows a block diagram of the clock generation logic.
Figure 22-2. Clock Generation Logic, Block Diagram


Signal description:

$$
\begin{array}{ll}
\text { txclk } & \text { Transmitter clock (Internal Signal). } \\
\text { rxclk } & \text { Receiver base clock (Internal Signal). } \\
\text { xcki } & \text { Input from XCK pin (internal Signal). Used for synchronous slave operation. } \\
\text { xcko } & \text { Clock output to XCK pin (Internal Signal). Used for synchronous master operation. } \\
\mathbf{f}_{\text {Osc }} & \text { XTAL pin frequency (System Clock). }
\end{array}
$$

### 22.3.1 Internal Clock Generation - The Baud Rate Generator

Internal clock generation is used for the asynchronous and the synchronous master modes of operation. The description in this section refers to Figure 22-2.

The USART Baud Rate Register (UBRRn) and the down-counter connected to it function as a programmable prescaler or baud rate generator. The down-counter, running at system clock ( $\mathrm{f}_{\text {osc }}$ ), is loaded with the UBRRn value each time the counter has counted down to zero or when the UBRRLn Register is written. A clock is generated each time the counter reaches zero. This clock is the baud rate generator clock output ( $=f_{\text {osc }} /(U B R R n+1)$ ). The Transmitter divides the baud rate generator clock output by 2,8 , or 16 depending on mode. The baud rate generator output is used directly by the Receiver's clock and data recovery units. However, the recovery units use a state machine that uses 2,8 or 16 states depending on mode set by the state of the UMSELn, U2Xn and DDR_XCKn bits.

Table 22-1 on page 203 contains equations for calculating the baud rate (in bits per second) and for calculating the UBRRn value for each mode of operation using an internally generated clock source.

Table 22-1. Equations for Calculating Baud Rate Register Setting

| Operating Mode | Equation for Calculating Baud Rate ${ }^{(1)}$ | Equation for Calculating UBRR Value |
| :--- | :---: | :---: |
| Asynchronous Normal mode <br> (U2Xn = 0) | BAUD $=\frac{f_{\text {OSC }}}{16(\text { UBRR } n+1)}$ | UBRRn $=\frac{f_{\text {OSC }}}{16 B A U D}-1$ |
| Asynchronous Double Speed <br> mode (U2Xn = 1) | BAUD $=\frac{f_{\text {OSC }}}{8(\text { UBRR } n+1)}$ | UBRRn $=\frac{f_{\text {OSC }}}{8 B A U D}-1$ |
| Synchronous Master mode | BAUD $=\frac{f_{O S C}}{2(\text { UBRR } n+1)}$ | UBRRn $=\frac{f_{O S C}}{2 B A U D}-1$ |

Note: 1. The baud rate is defined to be the transfer rate in bit per second (bps).
BAUD Baud rate (in bits per second, bps).
$f_{\text {osc }} \quad$ System Oscillator clock frequency.
UBRRn Contents of the UBRRHn and UBRRLn Registers, (0-4095).
Some examples of UBRRn values for some system clock frequencies are found in Table 22-9 on page 223.

### 22.3.2 Double Speed Operation (U2Xn)

The transfer rate can be doubled by setting the U2Xn bit in UCSRnA. Setting this bit only has effect for the asynchronous operation. Set this bit to zero when using synchronous operation.

Setting this bit will reduce the divisor of the baud rate divider from 16 to 8 , effectively doubling the transfer rate for asynchronous communication. Note however that the Receiver will in this case only use half the number of samples (reduced from 16 to 8) for data sampling and clock recovery, and therefore a more accurate baud rate setting and system clock are required when this mode is used. For the Transmitter, there are no downsides.

### 22.3.3 External Clock

External clocking is used by the synchronous slave modes of operation. The description in this section refers to Figure 22-2 on page 202 for details.

External clock input from the XCKn pin is sampled by a synchronization register to minimize the chance of metastability. The output from the synchronization register must then pass through an edge detector before it can be used by the Transmitter and Receiver. This process introduces a two CPU clock period delay and therefore the maximum external XCKn clock frequency is limited by the following equation:

$$
f_{X C K}<\frac{f_{O S C}}{4}
$$

Note that $\mathrm{f}_{\text {osc }}$ depends on the stability of the system clock source. It is therefore recommended to add some margin to avoid possible loss of data due to frequency variations.

### 22.3.4 Synchronous Clock Operation

When synchronous mode is used (UMSELn = 1), the XCKn pin will be used as either clock input (Slave) or clock output (Master). The dependency between the clock edges and data sampling or data change is the same. The basic principle is that data input ( $o n \mathrm{RxDn}$ ) is sampled at the opposite XCKn clock edge of the edge the data output ( TxDn ) is changed.

Figure 22-3. Synchronous Mode XCKn Timing.


The UCPOLn bit UCRSC selects which XCKn clock edge is used for data sampling and which is used for data change. As Figure 22-3 shows, when UCPOLn is zero the data will be changed at rising XCKn edge and sampled at falling XCKn edge. If UCPOLn is set, the data will be changed at falling XCKn edge and sampled at rising XCKn edge.

### 22.4 Frame Formats

A serial frame is defined to be one character of data bits with synchronization bits (start and stop bits), and optionally a parity bit for error checking. The USART accepts all 30 combinations of the following as valid frame formats:

- 1 start bit
- $5,6,7,8$, or 9 data bits
- no, even or odd parity bit
- 1 or 2 stop bits

A frame starts with the start bit followed by the least significant data bit. Then the next data bits, up to a total of nine, are succeeding, ending with the most significant bit. If enabled, the parity bit is inserted after the data bits, before the stop bits. When a complete frame is transmitted, it can be directly followed by a new frame, or the communication line can be set to an idle (high) state. Figure 22-4 illustrates the possible combinations of the frame formats. Bits inside brackets are optional.

Figure 22-4. Frame Formats


St Start bit, always low.
(n) Data bits (0 to 8).

P Parity bit. Can be odd or even.
Sp Stop bit, always high.
IDLE No transfers on the communication line (RxDn or TxDn). An IDLE line must be high.

The frame format used by the USART is set by the UCSZn2:0, UPMn1:0 and USBSn bits in UCSRnB and UCSRnC. The Receiver and Transmitter use the same setting. Note that changing the setting of any of these bits will corrupt all ongoing communication for both the Receiver and Transmitter.

The USART Character SiZe (UCSZn2:0) bits select the number of data bits in the frame. The USART Parity mode (UPMn1:0) bits enable and set the type of parity bit. The selection between one or two stop bits is done by the USART Stop Bit Select (USBSn) bit. The Receiver ignores the second stop bit. An FE (Frame Error) will therefore only be detected in the cases where the first stop bit is zero.

### 22.4.1 Parity Bit Calculation

The parity bit is calculated by doing an exclusive-or of all the data bits. If odd parity is used, the result of the exclusive or is inverted. The parity bit is located between the last data bit and first stop bit of a serial frame. The relation between the parity bit and data bits is as follows:

$$
\begin{aligned}
P_{\text {even }} & =d_{n-1} \oplus \ldots \oplus d_{3} \oplus d_{2} \oplus d_{1} \oplus d_{0} \oplus 0 \\
P_{\text {odd }} & =d_{n-1} \oplus \ldots \oplus d_{3} \oplus d_{2} \oplus d_{1} \oplus d_{0} \oplus 1
\end{aligned}
$$

| $\mathbf{P}_{\text {even }}$ | Parity bit using even parity. |
| :--- | :--- |
| $\mathbf{P}^{o d d}$ | Parity bit using odd parity. |
| $\mathbf{d}_{\mathbf{n}}$ | Data bit n of the character. |

### 22.5 USART Initialization

The USART has to be initialized before any communication can take place. The initialization process normally consists of setting the baud rate, setting frame format and enabling the Transmitter or the Receiver depending on the usage. For interrupt driven USART operation, the Global Interrupt Flag should be cleared (and interrupts globally disabled) when doing the initialization.
Before doing a re-initialization with changed baud rate or frame format, be sure that there are no ongoing transmissions during the period the registers are changed. The TXCn Flag can be used to check that the Transmitter has completed all transfers, and the RXC Flag can be used to check that there are no unread data in the receive buffer. Note that the TXCn Flag must be cleared before each transmission (before UDRn is written) if it is used for this purpose.

The following simple USART initialization code examples show one assembly and one C function that are equal in functionality. The examples assume asynchronous operation using polling (no interrupts enabled) and a fixed frame format. The baud rate is given as a function parameter. For the assembly code, the baud rate parameter is assumed to be stored in the r17:r16 Registers.

```
Assembly Code Example }\mp@subsup{}{}{(1)
    USART_Init:
    ; Set baud rate
    sts UBRRnH, r17
    sts UBRRnL, r16
    ldi r16, (1<<U2Xn)
    sts UCRnA, r16
    ; Enable receiver and transmitter
    ldi r16, (1<<RXENn)|(1<<TXENn)
    sts UCSRnB,r16
        ; Set frame format: 8data, 1stop bit
        ldi r16, (2<<UMSELn)|(3<<UCSZn0)
        sts UCSRnC,r16
        ret
C Code Example }\mp@subsup{}{(1)}{
#define FOSC 1843200// Clock Speed
#define BAUD 9600
#define (MYUBRR FOSC/16/BAUD-1)
void main( void )
{ ...
USART_Init ( MYUBRR );
...} // main
void USART_Init( unsigned int ubrr){
/* Set baud rate */
UBRRH = (unsigned char)(ubrr>>8);
UBRRL = (unsigned char)ubrr;
/* Enable receiver and transmitter */
UCSRB = (1<<RXEN)| (1<<TXEN );
/* Set frame format: 8data, 2stop bit */
UCSRC = (1<<USBS) | ( 3<<UCSZ0);
} // USART_Init
```

Note: 1. See "About Code Examples" on page 10.
More advanced initialization routines can be made that include frame format as parameters, disable interrupts and so on. However, many applications use a fixed setting of the baud and control registers, and for these types of applications the initialization code can be placed directly in the main routine, or be combined with initialization code for other I/O modules.

### 22.6 Data Transmission - The USART Transmitter

The USART Transmitter is enabled by setting the Transmit Enable (TXEN) bit in the UCSRnB Register. When the Transmitter is enabled, the normal port operation of the TxDn pin is overridden by the USART and given the function as the Transmitter's serial output. The baud rate, mode of operation and frame format must be set up once before doing any transmissions. If synchronous operation is used, the clock on the XCKn pin will be overridden and used as transmission clock.

### 22.6.1 Sending Frames with 5 to 8 Data Bit

A data transmission is initiated by loading the transmit buffer with the data to be transmitted. The CPU can load the transmit buffer by writing to the UDRn I/O location. The buffered data in the transmit buffer will be moved to the Shift Register when the Shift Register is ready to send a new frame. The Shift Register is loaded with new data if it is in idle state (no ongoing transmission) or immediately after the last stop bit of the previous frame is transmitted. When the Shift Register is loaded with new data, it will transfer one complete frame at the rate given by the Baud Register, U2Xn bit or by XCKn depending on mode of operation.

The following code examples show a simple USART transmit function based on polling of the Data Register Empty (UDREn) Flag. When using frames with less than eight bits, the most significant bits written to the UDRn are ignored. The USART has to be initialized before the function can be used. For the assembly code, the data to be sent is assumed to be stored in Register R16.

```
Assembly Code Example }\mp@subsup{}{}{(1)
    USART_Transmit:
        ; Wait for empty transmit buffer
        1ds r17, UCSRnA
        sbrs r17, UDREn
        rjmp USART_Transmit
        ; Put data (r16) into buffer, sends the data
        sts UDRn,r16
        ret
C Code Example }\mp@subsup{}{}{(1)
    void USART_Transmit( unsigned char data )
    {
        /* Wait for empty transmit buffer */
        while ( !( UCSRnA & (1<<UDREn)) )
            ;
        /* Put data into buffer, sends the data */
        UDRn = data;
    }
```

Note: 1. See "About Code Examples" on page 10.
The function simply waits for the transmit buffer to be empty by checking the UDREn Flag, before loading it with new data to be transmitted. If the Data Register Empty interrupt is utilized, the interrupt routine writes the data into the buffer.

### 22.6.2 Sending Frames with 9 Data Bit

If 9-bit characters are used ( $\mathrm{UCSZn}=7$ ), the ninth bit must be written to the TXB8 bit in UCSRnB before the low byte of the character is written to UDRn. The following code examples show a transmit function that handles 9-bit characters. For the assembly code, the data to be sent is assumed to be stored in registers R17:R16.

```
Assembly Code Example \({ }^{(1)(2)}\)
    USART_Transmit:
    ; Wait for empty transmit buffer
    sbis UCSRnA, UDREn
    rjmp USART_Transmit
    ; Copy 9th bit from r17 to TXB8
    cbi UCSRnB,TXB8
    sbrc r17,0
    sbi UCSRnB,TXB8
    ; Put LSB data (r16) into buffer, sends the data
    sts UDRn,r16
    ret
C Code Example \({ }^{(1)(2)}\)
    void USART_Transmit( unsigned int data )
    \{
    /* Wait for empty transmit buffer */
    while ( ! ( UCSRnA \& (1<<UDREn))) )
    /* Copy 9th bit to TXB8 */
    UCSRnB \(\&=\sim(1 \ll\) TXB 8\() ;\)
    if ( data \& 0x0100 )
            UCSRnB |= ( \(1 \ll\) TXB 8\()\);
    /* Put data into buffer, sends the data */
    UDRn = data;
    \}
```

Notes: 1. These transmit functions are written to be general functions. They can be optimized if the contents of the UCSRnB is static. For example, only the TXB8 bit of the UCSRnB Register is used after initialization.
2. See "About Code Examples" on page 10.

The ninth bit can be used for indicating an address frame when using multi processor communication mode or for other protocol handling as for example synchronization.

### 22.6.3 Transmitter Flags and Interrupts

The USART Transmitter has two flags that indicate its state: USART Data Register Empty (UDREn) and Transmit Complete (TXCn). Both flags can be used for generating interrupts.

The Data Register Empty (UDREn) Flag indicates whether the transmit buffer is ready to receive new data. This bit is set when the transmit buffer is empty, and cleared when the transmit buffer contains data to be transmitted that has not yet been moved into the Shift Register. For compatibility with future devices, always write this bit to zero when writing the UCSRnA Register.

When the Data Register Empty Interrupt Enable (UDRIEn) bit in UCSRnB is written to one, the USART Data Register Empty Interrupt will be executed as long as UDREn is set (provided that global interrupts are enabled).

UDREn is cleared by writing UDRn. When interrupt-driven data transmission is used, the Data Register Empty interrupt routine must either write new data to UDRn in order to clear UDREn or disable the Data Register Empty interrupt, otherwise a new interrupt will occur once the interrupt routine terminates.

The Transmit Complete (TXCn) Flag bit is set one when the entire frame in the Transmit Shift Register has been shifted out and there are no new data currently present in the transmit buffer. The TXCn Flag bit is automatically cleared when a transmit complete interrupt is executed, or it can be cleared by writing a one to its bit location. The TXCn Flag is useful in half-duplex communication interfaces (like the RS-485 standard), where a transmitting application must enter receive mode and free the communication bus immediately after completing the transmission.

When the Transmit Compete Interrupt Enable (TXCIEn) bit in UCSRnB is set, the USART Transmit Complete Interrupt will be executed when the TXCn Flag becomes set (provided that global interrupts are enabled). When the transmit complete interrupt is used, the interrupt handling routine does not have to clear the TXCn Flag, this is done automatically when the interrupt is executed.

### 22.6.4 Parity Generator

The Parity Generator calculates the parity bit for the serial frame data. When parity bit is enabled (UPMn1 = 1), the transmitter control logic inserts the parity bit between the last data bit and the first stop bit of the frame that is sent.

### 22.6.5 Disabling the Transmitter

The disabling of the Transmitter (setting the TXEN to zero) will not become effective until ongoing and pending transmissions are completed, that is, when the Transmit Shift Register and Transmit Buffer Register do not contain data to be transmitted. When disabled, the Transmitter will no longer override the TxDn pin.

### 22.7 Data Reception - The USART Receiver

The USART Receiver is enabled by writing the Receive Enable (RXENn) bit in the UCSRnB Register to one. When the Receiver is enabled, the normal pin operation of the RxDn pin is overridden by the USART and given the function as the Receiver's serial input. The baud rate, mode of operation and frame format must be set up once before any serial reception can be done. If synchronous operation is used, the clock on the XCKn pin will be used as transfer clock.

### 22.7.1 Receiving Frames with 5 to 8 Data Bits

The Receiver starts data reception when it detects a valid start bit. Each bit that follows the start bit will be sampled at the baud rate or XCKn clock, and shifted into the Receive Shift Register until the first stop bit of a frame is received. A second stop bit will be ignored by the Receiver. When the first stop bit is received, that is, a complete serial frame is present in the Receive Shift Register, the contents of the Shift Register will be moved into the receive buffer. The receive buffer can then be read by reading the UDRn I/O location.

The following code example shows a simple USART receive function based on polling of the Receive Complete ( RXCn ) Flag. When using frames with less than eight bits the most significant bits of the data read from the UDRn will be masked to zero. The USART has to be initialized before the function can be used.


Note: 1. See "About Code Examples" on page 10.
The function simply waits for data to be present in the receive buffer by checking the RXCn Flag, before reading the buffer and returning the value.

### 22.7.2 Receiving Frames with 9 Data Bits

If 9-bit characters are used ( $\mathrm{UCSZn=7)}$ the ninth bit must be read from the RXB8n bit in UCSRnB before reading the low bits from the UDRn. This rule applies to the FEn, DORn and UPEn Status Flags as well. Read status from UCSRnA, then data from UDRn. Reading the UDRn I/O location will change the state of the receive buffer FIFO and consequently the TXB8n, FEn, DORn and UPEn bits, which all are stored in the FIFO, will change.

The following code example shows a simple USART receive function that handles both nine bit characters and the status bits.

```
Assembly Code Example }\mp@subsup{}{}{(1)
    USART_Receive:
        ; Wait for data to be received
        lds r17, UCSRnA
        sbrs r17, RXCn
        rjmp USART_Receive
        ; Get status and 9th bit, then data from buffer
        lds r18, UCSRnA
        lds r17, UCSRnB
        lds r16, UDRn
        ; If error, return -1
        andi r18,(1<<FEn)|(1<<DORn)|(1<<UPEn)
        breq USART_ReceiveNoError
        ldi r17, HIGH(-1)
        ldi r16, LOW(-1)
USART_ReceiveNoError:
        ; Filter the 9th bit, then return
        lsr r17
        andi r17, 0x01
        ret
C Code Example }\mp@subsup{}{}{(1)
unsigned int USART_Receive( void )
    {
        unsigned char status, resh, resl;
        /* Wait for data to be received */
        while ( ! (UCSRnA & (1<<RXCn)) )
                ;
        /* Get status and 9th bit, then data */
        /* from buffer */
        status = UCSRnA;
        resh = UCSRnB;
        resl = UDRn;
        /* If error, return -1 */
        if ( status & (1<<FEn)|(1<<DORn)| (1<<UPEn) )
            return -1;
        /* Filter the 9th bit, then return */
        resh = (resh >> 1) & 0x01;
        return ((resh << 8) | resl);
    }
```

Note: 1. See "About Code Examples" on page 10.
The receive function example reads all the I/O Registers into the Register File before any computation is done.
This gives an optimal receive buffer utilization since the buffer location read will be free to accept new data as early as possible.

### 22.7.3 Receive Compete Flag and Interrupt

The USART Receiver has one flag that indicates the Receiver state.
The Receive Complete (RXCn) Flag indicates if there are unread data present in the receive buffer. This flag is one when unread data exist in the receive buffer, and zero when the receive buffer is empty (that is, does not contain any unread data). If the Receiver is disabled ( $\operatorname{RXEN} n=0$ ), the receive buffer will be flushed and consequently the RXCn bit will become zero.

When the Receive Complete Interrupt Enable (RXCIEn) in UCSRnB is set, the USART Receive Complete interrupt will be executed as long as the RXCn Flag is set (provided that global interrupts are enabled). When interruptdriven data reception is used, the receive complete routine must read the received data from UDRn in order to clear the RXCn Flag, otherwise a new interrupt will occur once the interrupt routine terminates.

### 22.7.4 Receiver Error Flags

The USART Receiver has three Error Flags: Frame Error (FEn), Data OverRun (DORn) and Parity Error (UPEn). All can be accessed by reading UCSRnA. Common for the Error Flags is that they are located in the receive buffer together with the frame for which they indicate the error status. Due to the buffering of the Error Flags, the UCSRnA must be read before the receive buffer (UDRn), since reading the UDRn I/O location changes the buffer read location. Another equality for the Error Flags is that they can not be altered by software doing a write to the flag location. However, all flags must be set to zero when the UCSRnA is written for upward compatibility of future USART implementations. None of the Error Flags can generate interrupts.

The Frame Error (FEn) Flag indicates the state of the first stop bit of the next readable frame stored in the receive buffer. The FEn Flag is zero when the stop bit was correctly read (as one), and the FEn Flag will be one when the stop bit was incorrect (zero). This flag can be used for detecting out-of-sync conditions, detecting break conditions and protocol handling. The FEn Flag is not affected by the setting of the USBSn bit in UCSRnC since the Receiver ignores all, except for the first, stop bits. For compatibility with future devices, always set this bit to zero when writing to UCSRnA.

The Data OverRun (DORn) Flag indicates data loss due to a receiver buffer full condition. A Data OverRun occurs when the receive buffer is full (two characters), it is a new character waiting in the Receive Shift Register, and a new start bit is detected. If the DORn Flag is set there was one or more serial frame lost between the frame last read from UDRn, and the next frame read from UDRn. For compatibility with future devices, always write this bit to zero when writing to UCSRnA. The DORn Flag is cleared when the frame received was successfully moved from the Shift Register to the receive buffer.

The Parity Error (UPEn) Flag indicates that the next frame in the receive buffer had a Parity Error when received. If Parity Check is not enabled the UPEn bit will always be read zero. For compatibility with future devices, always set this bit to zero when writing to UCSRnA. For more details see "Parity Bit Calculation" on page 205 and "Parity Checker".

### 22.7.5 Parity Checker

The Parity Checker is active when the high USART Parity mode (UPMn1) bit is set. Type of Parity Check to be performed (odd or even) is selected by the UPMn0 bit. When enabled, the Parity Checker calculates the parity of the data bits in incoming frames and compares the result with the parity bit from the serial frame. The result of the check is stored in the receive buffer together with the received data and stop bits. The Parity Error (UPEn) Flag can then be read by software to check if the frame had a Parity Error.

The UPEn bit is set if the next character that can be read from the receive buffer had a Parity Error when received and the Parity Checking was enabled at that point (UPMn1 = 1). This bit is valid until the receive buffer (UDRn) is read.

### 22.7.6 Disabling the Receiver

In contrast to the Transmitter, disabling of the Receiver will be immediate. Data from ongoing receptions will therefore be lost. When disabled (that is, the RXENn is set to zero) the Receiver will no longer override the normal function of the RxDn port pin. The Receiver buffer FIFO will be flushed when the Receiver is disabled. Remaining data in the buffer will be lost.

### 22.7.7 Flushing the Receive Buffer

The receiver buffer FIFO will be flushed when the Receiver is disabled, that is, the buffer will be emptied of its contents. Unread data will be lost. If the buffer has to be flushed during normal operation, due to for instance an error condition, read the UDRn I/O location until the RXCn Flag is cleared. The following code example shows how to flush the receive buffer.


Note: 1. See "About Code Examples" on page 10.

### 22.8 Asynchronous Data Reception

The USART includes a clock recovery and a data recovery unit for handling asynchronous data reception. The clock recovery logic is used for synchronizing the internally generated baud rate clock to the incoming asynchronous serial frames at the RxDn pin. The data recovery logic samples and low pass filters each incoming bit, thereby improving the noise immunity of the Receiver. The asynchronous reception operational range depends on the accuracy of the internal baud rate clock, the rate of the incoming frames, and the frame size in number of bits.

### 22.8.1 Asynchronous Clock Recovery

The clock recovery logic synchronizes internal clock to the incoming serial frames. Figure 22-5 on page 214 illustrates the sampling process of the start bit of an incoming frame. The sample rate is 16 times the baud rate for Normal mode, and eight times the baud rate for Double Speed mode. The horizontal arrows illustrate the synchronization variation due to the sampling process. Note the larger time variation when using the Double Speed mode ( $\mathrm{U} 2 \mathrm{Xn}=1$ ) of operation. Samples denoted zero are samples done when the RxDn line is idle (that is, no communication activity).

Figure 22-5. Start Bit Sampling


When the clock recovery logic detects a high (idle) to low (start) transition on the RxDn line, the start bit detection sequence is initiated. Let sample 1 denote the first zero-sample as shown in the figure. The clock recovery logic then uses samples 8,9 , and 10 for Normal mode, and samples 4, 5 , and 6 for Double Speed mode (indicated with sample numbers inside boxes on the figure), to decide if a valid start bit is received. If two or more of these three samples have logical high levels (the majority wins), the start bit is rejected as a noise spike and the Receiver starts looking for the next high to low-transition. If however, a valid start bit is detected, the clock recovery logic is synchronized and the data recovery can begin. The synchronization process is repeated for each start bit.

### 22.8.2 Asynchronous Data Recovery

When the receiver clock is synchronized to the start bit, the data recovery can begin. The data recovery unit uses a state machine that has 16 states for each bit in Normal mode and eight states for each bit in Double Speed mode. Figure $22-6$ shows the sampling of the data bits and the parity bit. Each of the samples is given a number that is equal to the state of the recovery unit.

Figure 22-6. Sampling of Data and Parity Bit


The decision of the logic level of the received bit is taken by doing a majority voting of the logic value to the three samples in the center of the received bit. The center samples are emphasized on the figure by having the sample number inside boxes. The majority voting process is done as follows: If two or all three samples have high levels, the received bit is registered to be a logic 1. If two or all three samples have low levels, the received bit is registered to be a logic 0 . This majority voting process acts as a low pass filter for the incoming signal on the RxDn pin. The recovery process is then repeated until a complete frame is received. Including the first stop bit. Note that the Receiver only uses the first stop bit of a frame.

Figure 22-7 shows the sampling of the stop bit and the earliest possible beginning of the start bit of the next frame.
Figure 22-7. Stop Bit Sampling and Next Start Bit Sampling


The same majority voting is done to the stop bit as done for the other bits in the frame. If the stop bit is registered to have a logic 0 value, the Frame Error (FEn) Flag will be set.

A new high to low transition indicating the start bit of a new frame can come right after the last of the bits used for majority voting. For Normal Speed mode, the first low level sample can be at point marked (A) in Figure 22-7 on page 214. For Double Speed mode the first low level must be delayed to (B). (C) marks a stop bit of full length. The early start bit detection influences the operational range of the Receiver.

### 22.8.3 Asynchronous Operational Range

The operational range of the Receiver is dependent on the mismatch between the received bit rate and the internally generated baud rate. If the Transmitter is sending frames at too fast or too slow bit rates, or the internally generated baud rate of the Receiver does not have a similar (see Table 22-2) base frequency, the Receiver will not be able to synchronize the frames to the start bit.

The following equations can be used to calculate the ratio of the incoming data rate and internal receiver baud rate.

$$
R_{\text {slow }}=\frac{(D+1) S}{S-1+D \cdot S+S_{F}} \quad R_{\text {fast }}=\frac{(D+2) S}{(D+1) S+S_{M}}
$$

D $\quad$ Sum of character size and parity size ( $\mathrm{D}=5$ to 10 bit).
S Samples per bit. $S=16$ for Normal Speed mode and $S=8$ for Double Speed mode.
$S_{F} \quad$ First sample number used for majority voting. $S_{F}=8$ for normal speed and $S_{F}=4$ for Double Speed mode.
$S_{M} \quad$ Middle sample number used for majority voting. $S_{M}=9$ for normal speed and $S_{M}=5$ for Double Speed mode.
$\mathbf{R}_{\text {slow }} \quad$ is the ratio of the slowest incoming data rate that can be accepted in relation to the receiver baud rate. $\mathrm{R}_{\text {fast }}$ is the ratio of the fastest incoming data rate that can be accepted in relation to the receiver baud rate.
Table 22-2 and Table 22-3 on page 216 list the maximum receiver baud rate error that can be tolerated. Note that Normal Speed mode has higher toleration of baud rate variations.

Table 22-2. Recommended Maximum Receiver Baud Rate Error for Normal Speed Mode (U2Xn = 0)

| D (Data+Parity Bit) | $\mathbf{R}_{\text {slow }}$ (\%) | $\mathbf{R}_{\text {fast }}$ (\%) | Max. total error (\%) | Recommended max. <br> receiver error (\%) |
| :---: | :---: | :---: | :---: | :---: |
| 5 | 93.20 | 106.67 | $+6.67 /-6.8$ | $\pm 3.0$ |
| 6 | 94.12 | 105.79 | $+5.79 /-5.88$ | $\pm 2.5$ |
| 7 | 94.81 | 105.11 | $+5.11 /-5.19$ | $\pm 2.0$ |
| 8 | 95.36 | 104.58 | $+4.58 /-4.54$ | $\pm 2.0$ |
| 9 | 95.81 | 104.14 | $+4.14 /-4.19$ | $\pm 1.5$ |
| 10 | 96.17 | 103.78 | $+3.78 /-3.83$ | $\pm 1.5$ |

Table 22-3. $\quad$ Recommended Maximum Receiver Baud Rate Error for Double Speed Mode (U2Xn = 1)

| D (Data+Parity Bit) | $\mathbf{R}_{\text {slow }}$ (\%) | $\mathbf{R}_{\text {fast }}(\%)$ | Max. total error (\%) | Recommended max. receiver <br> error (\%) |
| :---: | :---: | :---: | :---: | :---: |
| 5 | 94.12 | 105.66 | $+5.66 /-5.88$ | $\pm 2.5$ |
| 6 | 94.92 | 104.92 | $+4.92 /-5.08$ | $\pm 2.0$ |
| 7 | 95.52 | 104.35 | $+4.35 /-4.48$ | $\pm 1.5$ |
| 8 | 96.00 | 103.90 | $+3.90 /-4.00$ | $\pm 1.5$ |
| 9 | 96.39 | 103.53 | $+3.53 /-3.61$ | $\pm 1.5$ |
| 10 | 96.70 | 103.23 | $+3.23 /-3.30$ | $\pm 1.0$ |

The recommendations of the maximum receiver baud rate error was made under the assumption that the Receiver and Transmitter equally divides the maximum total error.

There are two possible sources for the receivers baud rate error. The Receiver's system clock (XTAL) will always have some minor instability over the supply voltage range and the temperature range. When using a crystal to generate the system clock, this is rarely a problem, but for a resonator the system clock may differ more than $2 \%$ depending of the resonators tolerance. The second source for the error is more controllable. The baud rate generator can not always do an exact division of the system frequency to get the baud rate wanted. In this case an UBRR value that gives an acceptable low error can be used if possible.

### 22.9 Multi-processor Communication Mode

Setting the Multi-processor Communication mode (MPCMn) bit in UCSRnA enables a filtering function of incoming frames received by the USART Receiver. Frames that do not contain address information will be ignored and not put into the receive buffer. This effectively reduces the number of incoming frames that has to be handled by the CPU, in a system with multiple MCUs that communicate via the same serial bus. The Transmitter is unaffected by the MPCMn setting, but has to be used differently when it is a part of a system utilizing the Multi-processor Communication mode.
If the Receiver is set up to receive frames that contain 5 to 8 data bits, then the first stop bit indicates if the frame contains data or address information. If the Receiver is set up for frames with nine data bits, then the ninth bit (RXB8n) is used for identifying address and data frames. When the frame type bit (the first stop or the ninth bit) is one, the frame contains an address. When the frame type bit is zero the frame is a data frame.

The Multi-processor Communication mode enables several slave MCUs to receive data from a master MCU. This is done by first decoding an address frame to find out which MCU has been addressed. If a particular slave MCU has been addressed, it will receive the following data frames as normal, while the other slave MCUs will ignore the received frames until another address frame is received.

### 22.9.1 Using MPCMn

For an MCU to act as a master MCU, it can use a 9-bit character frame format (UCSZn = 7). The ninth bit (TXB8n) must be set when an address frame ( $\mathrm{TXB8} 8=1$ ) or cleared when a data frame ( $\mathrm{TXB}=0$ ) is being transmitted. The slave MCUs must in this case be set to use a 9-bit character frame format.

The following procedure should be used to exchange data in Multi-processor Communication mode:

1. All Slave MCUs are in Multi-processor Communication mode (MPCMn in UCSRnA is set).
2. The Master MCU sends an address frame, and all slaves receive and read this frame. In the Slave MCUs, the RXCn Flag in UCSRnA will be set as normal.
3. Each Slave MCU reads the UDRn Register and determines if it has been selected. If so, it clears the MPCMn bit in UCSRnA, otherwise it waits for the next address byte and keeps the MPCMn setting.
4. The addressed MCU will receive all data frames until a new address frame is received. The other Slave MCUs, which still have the MPCMn bit set, will ignore the data frames.
5. When the last data frame is received by the addressed MCU, the addressed MCU sets the MPCMn bit and waits for a new address frame from master. The process then repeats from 2.
Using any of the 5 -bit to 8 -bit character frame formats is possible, but impractical since the Receiver must change between using n and $\mathrm{n}+1$ character frame formats. This makes full-duplex operation difficult since the Transmitter and Receiver uses the same character size setting. If 5 -bit to 8 -bit character frames are used, the Transmitter must be set to use two stop bit ( $\mathrm{USBSn}=1$ ) since the first stop bit is used for indicating the frame type.

Do not use Read-Modify-Write instructions (SBI and CBI) to set or clear the MPCMn bit. The MPCMn bit shares the same I/O location as the TXCn Flag and this might accidentally be cleared when using SBI or CBI instructions.

### 22.10 Register Description

The following section describes the USART's registers.

### 22.10.1 UDRn - USART I/O Data Register n



The USART Transmit Data Buffer Register and USART Receive Data Buffer Registers share the same I/O address referred to as USART Data Register or UDRn. The Transmit Data Buffer Register (TXB) will be the destination for data written to the UDRn Register location. Reading the UDRn Register location will return the contents of the Receive Data Buffer Register (RXB).

For 5-bit, 6-bit, or 7-bit characters the upper unused bits will be ignored by the Transmitter and set to zero by the Receiver.

The transmit buffer can only be written when the UDREn Flag in the UCSRnA Register is set. Data written to UDRn when the UDREn Flag is not set, will be ignored by the USART Transmitter. When data is written to the transmit buffer, and the Transmitter is enabled, the Transmitter will load the data into the Transmit Shift Register when the Shift Register is empty. Then the data will be serially transmitted on the TxDn pin.
The receive buffer consists of a two level FIFO. The FIFO will change its state whenever the receive buffer is accessed. Due to this behavior of the receive buffer, do not use Read-Modify-Write instructions (SBI and CBI) on this location. Be careful when using bit test instructions (SBIC and SBIS), since these also will change the state of the FIFO.

### 22.10.2 UCSRnA - USART Control and Status Register A



## - Bit 7 - RXCn: USART Receive Complete

This flag bit is set when there are unread data in the receive buffer and cleared when the receive buffer is empty (that is, does not contain any unread data). If the Receiver is disabled, the receive buffer will be flushed and consequently the RXCn bit will become zero. The RXCn Flag can be used to generate a Receive Complete interrupt (see description of the RXCIEn bit).

## - Bit 6 - TXCn: USART Transmit Complete

This flag bit is set when the entire frame in the Transmit Shift Register has been shifted out and there are no new data currently present in the transmit buffer (UDRn). The TXCn Flag bit is automatically cleared when a transmit complete interrupt is executed, or it can be cleared by writing a one to its bit location. The TXCn Flag can generate a Transmit Complete interrupt (see description of the TXCIEn bit).

## - Bit 5 - UDREn: USART Data Register Empty

The UDREn Flag indicates if the transmit buffer (UDRn) is ready to receive new data. If UDREn is one, the buffer is empty, and therefore ready to be written. The UDREn Flag can generate a Data Register Empty interrupt (see description of the UDRIEn bit).

UDREn is set after a reset to indicate that the Transmitter is ready.

## - Bit 4 - FEn: Frame Error

This bit is set if the next character in the receive buffer had a Frame Error when received, that is, when the first stop bit of the next character in the receive buffer is zero. This bit is valid until the receive buffer (UDRn) is read. The FEn bit is zero when the stop bit of received data is one. Always set this bit to zero when writing to UCSRnA.

## - Bit 3 - DORn: Data OverRun

This bit is set if a Data OverRun condition is detected. A Data OverRun occurs when the receive buffer is full (two characters), it is a new character waiting in the Receive Shift Register, and a new start bit is detected. This bit is valid until the receive buffer (UDRn) is read. Always set this bit to zero when writing to UCSRnA.

## - Bit 2 - UPEn: USART Parity Error

This bit is set if the next character in the receive buffer had a Parity Error when received and the Parity Checking was enabled at that point (UPMn1 = 1). This bit is valid until the receive buffer (UDRn) is read. Always set this bit to zero when writing to UCSRnA.

## - Bit 1 - U2Xn: Double the USART Transmission Speed

This bit only has effect for the asynchronous operation. Write this bit to zero when using synchronous operation.
Writing this bit to one will reduce the divisor of the baud rate divider from 16 to 8 effectively doubling the transfer rate for asynchronous communication.

## - Bit 0 - MPCMn: Multi-processor Communication Mode

This bit enables the Multi-processor Communication mode. When the MPCMn bit is written to one, all the incoming frames received by the USART Receiver that do not contain address information will be ignored. The Transmitter is unaffected by the MPCMn setting. For more detailed information see "Multi-processor Communication Mode" on page 216.

### 22.10.3 UCSRnB - USART Control and Status Register n B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | RXCIEn | TXCIEn | UDRIEn | RXENn | TXENn | UCSZn2 | RXB8n | TXB8n | UCSRnB |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - RXCIEn: RX Complete Interrupt Enable n

Writing this bit to one enables interrupt on the RXCn Flag. A USART Receive Complete interrupt will be generated only if the RXCIEn bit is written to one, the Global Interrupt Flag in SREG is written to one and the RXCn bit in UCSRnA is set.

## - Bit 6 - TXCIEn: TX Complete Interrupt Enable n

Writing this bit to one enables interrupt on the TXCn Flag. A USART Transmit Complete interrupt will be generated only if the TXCIEn bit is written to one, the Global Interrupt Flag in SREG is written to one and the TXCn bit in UCSRnA is set.

- Bit 5 - UDRIEn: USART Data Register Empty Interrupt Enable n

Writing this bit to one enables interrupt on the UDREn Flag. A Data Register Empty interrupt will be generated only if the UDRIEn bit is written to one, the Global Interrupt Flag in SREG is written to one and the UDREn bit in UCSRnA is set.

## - Bit 4 - RXENn: Receiver Enable n

Writing this bit to one enables the USART Receiver. The Receiver will override normal port operation for the RxDn pin when enabled. Disabling the Receiver will flush the receive buffer invalidating the FEn, DORn, and UPEn Flags.

## - Bit 3 - TXENn: Transmitter Enable n

Writing this bit to one enables the USART Transmitter. The Transmitter will override normal port operation for the TxDn pin when enabled. The disabling of the Transmitter (writing TXENn to zero) will not become effective until ongoing and pending transmissions are completed, that is, when the Transmit Shift Register and Transmit Buffer Register do not contain data to be transmitted. When disabled, the Transmitter will no longer override the TxDn port.

## - Bit 2 - UCSZn2: Character Size n

The UCSZn2 bits combined with the UCSZn1:0 bit in UCSRnC sets the number of data bits (Character SiZe) in a frame the Receiver and Transmitter use.

## - Bit 1 - RXB8n: Receive Data Bit 8 n

RXB8n is the ninth data bit of the received character when operating with serial frames with nine data bits. Must be read before reading the low bits from UDRn.

- Bit 0 - TXB8n: Transmit Data Bit 8 n

TXB8n is the ninth data bit in the character to be transmitted when operating with serial frames with nine data bits. Must be written before writing the low bits to UDRn.

### 22.10.4 UCSRnC - USART Control and Status Register n C



## - Bits 7:6 - UMSELn1:0 USART Mode Select

These bits select the mode of operation of the USARTn as shown in Table 22-4.
Table 22-4. UMSELn Bits Settings

| UMSELn1 | UMSELn0 | Mode |
| :---: | :---: | :---: |
| 0 | 0 | Asynchronous USART |
| 0 | 1 | Synchronous USART |
| 1 | 0 | (Reserved) |
| 1 | 1 | Master SPI (MSPIM) ${ }^{(1)}$ |

Note: 1. See "USART in SPI Mode" on page 227 for full description of the Master SPI Mode (MSPIM) operation.

## - Bits 5:4 - UPMn1:0: Parity Mode

These bits enable and set type of parity generation and check. If enabled, the Transmitter will automatically generate and send the parity of the transmitted data bits within each frame. The Receiver will generate a parity value for the incoming data and compare it to the UPMn setting. If a mismatch is detected, the UPEn Flag in UCSRnA will be set.

Table 22-5. UPMn Bits Settings

| UPMn1 | UPMn0 | Parity Mode |
| :---: | :---: | :---: |
| 0 | 0 | Disabled |
| 0 | 1 | Reserved |
| 1 | 0 | Enabled, Even Parity |
| 1 | 1 | Enabled, Odd Parity |

## - Bit 3 - USBSn: Stop Bit Select

This bit selects the number of stop bits to be inserted by the Transmitter. The Receiver ignores this setting.
Table 22-6. USBS Bit Settings

| USBSn | Stop Bit(s) |
| :---: | :---: |
| 0 | 1-bit |
| 1 | 2 -bit |

## - Bit 2:1 - UCSZn1:0: Character Size

The UCSZn1:0 bits combined with the UCSZn2 bit in UCSRnB sets the number of data bits (Character SiZe) in a frame the Receiver and Transmitter use.

Table 22-7. UCSZn Bits Settings

| UCSZn2 | UCSZn1 | UCSZn0 | Character Size |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 5-bit |
| 0 | 0 | 1 | 6-bit |
| 0 | 1 | 0 | 7-bit |
| 0 | 1 | 1 | 8-bit |
| 1 | 0 | 0 | Reserved |
| 1 | 0 | 1 | Reserved |
| 1 | 1 | 0 | Reserved |
| 1 | 1 | 1 | $9-b i t$ |

## - Bit 0 - UCPOLn: Clock Polarity

This bit is used for synchronous mode only. Write this bit to zero when asynchronous mode is used. The UCPOLn bit sets the relationship between data output change and data input sample, and the synchronous clock (XCKn).

Table 22-8. UCPOLn Bit Settings

| UCPOLn | Transmitted Data Changed <br> (Output of TxDn Pin) | Received Data Sampled <br> (Input on RxDn Pin) |
| :---: | :---: | :---: |
| 0 | Rising XCKn Edge | Falling XCKn Edge |
| 1 | Falling XCKn Edge | Rising XCKn Edge |

### 22.10.5 UBRRnL and UBRRnH - USART Baud Rate Registers

| Bit | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | UBRRHn UBRRLn |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | - | - | - | - |  | UBRR[11:8] |  |  |  |
|  | UBRR[7:0] |  |  |  |  |  |  |  |  |
|  | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| Read/Write | R | R | R | R | R/W | R/W | R/W | R/W |  |
|  | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 15:12 - Reserved Bits

These bits are reserved for future use. For compatibility with future devices, these bit must be written to zero when UBRRH is written.

## - Bit 11:0 - UBRR11:0: USART Baud Rate Register

This is a 12-bit register which contains the USART baud rate. The UBRRH contains the four most significant bits, and the UBRRL contains the eight least significant bits of the USART baud rate. Ongoing transmissions by the Transmitter and Receiver will be corrupted if the baud rate is changed. Writing UBRRL will trigger an immediate update of the baud rate prescaler.

### 22.11 Examples of Baud Rate Setting

For standard crystal and resonator frequencies, the most commonly used baud rates for asynchronous operation can be generated by using the UBRR settings in Table 22-9 to Table 22-12 on page 226. UBRR values which yield an actual baud rate differing less than $0.5 \%$ from the target baud rate, are bold in the table. Higher error ratings are acceptable, but the Receiver will have less noise resistance when the error ratings are high, especially for large serial frames (see "Asynchronous Operational Range" on page 215). The error values are calculated using the following equation:

$$
\text { Error }[\%]=\left(\frac{\text { BaudRate }_{\text {closest Match }}}{\text { BaudRate }}-1\right) \cdot 100 \%
$$

Table 22-9. Examples of UBRRn Settings for Commonly Used Oscillator Frequencies

| Baud Rate [bps] | $\mathrm{f}_{\text {osc }}=1.0000 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=1.8432 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=2.0000 \mathrm{MHz}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  |
|  | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error |
| 2400 | 25 | 0.2\% | 51 | 0.2\% | 47 | 0.0\% | 95 | 0.0\% | 51 | 0.2\% | 103 | 0.2\% |
| 4800 | 12 | 0.2\% | 25 | 0.2\% | 23 | 0.0\% | 47 | 0.0\% | 25 | 0.2\% | 51 | 0.2\% |
| 9600 | 6 | -7.0\% | 12 | 0.2\% | 11 | 0.0\% | 23 | 0.0\% | 12 | 0.2\% | 25 | 0.2\% |
| 14.4K | 3 | 8.5\% | 8 | -3.5\% | 7 | 0.0\% | 15 | 0.0\% | 8 | -3.5\% | 16 | 2.1\% |
| 19.2K | 2 | 8.5\% | 6 | -7.0\% | 5 | 0.0\% | 11 | 0.0\% | 6 | -7.0\% | 12 | 0.2\% |
| 28.8K | 1 | 8.5\% | 3 | 8.5\% | 3 | 0.0\% | 7 | 0.0\% | 3 | 8.5\% | 8 | -3.5\% |
| 38.4K | 1 | -18.6\% | 2 | 8.5\% | 2 | 0.0\% | 5 | 0.0\% | 2 | 8.5\% | 6 | -7.0\% |
| 57.6K | 0 | 8.5\% | 1 | 8.5\% | 1 | 0.0\% | 3 | 0.0\% | 1 | 8.5\% | 3 | 8.5\% |
| 76.8K | - | - | 1 | -18.6\% | 1 | -25.0\% | 2 | 0.0\% | 1 | -18.6\% | 2 | 8.5\% |
| 115.2K | - | - | 0 | 8.5\% | 0 | 0.0\% | 1 | 0.0\% | 0 | 8.5\% | 1 | 8.5\% |
| 230.4K | - | - | - | - | - | - | 0 | 0.0\% | - | - | - | - |
| 250K | - | - | - | - | - | - | - | - | - | - | 0 | 0.0\% |
| Max. ${ }^{(1)}$ | 62.5Kbps |  | 125Kbps |  | 115.2Kbps |  | 230.4Kbps |  | 125Kbps |  | 250Kbps |  |

Note: 1. $\mathrm{UBRR}=0$, Error $=0.0 \%$

Table 22-10. Examples of UBRRn Settings for Commonly Used Oscillator Frequencies

| Baud Rate [bps] | $\mathrm{f}_{\text {osc }}=3.6864 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=4.0000 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=7.3728 \mathrm{MHz}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | U2Xn = 0 |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  |
|  | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error |
| 2400 | 95 | 0.0\% | 191 | 0.0\% | 103 | 0.2\% | 207 | 0.2\% | 191 | 0.0\% | 383 | 0.0\% |
| 4800 | 47 | 0.0\% | 95 | 0.0\% | 51 | 0.2\% | 103 | 0.2\% | 95 | 0.0\% | 191 | 0.0\% |
| 9600 | 23 | 0.0\% | 47 | 0.0\% | 25 | 0.2\% | 51 | 0.2\% | 47 | 0.0\% | 95 | 0.0\% |
| 14.4K | 15 | 0.0\% | 31 | 0.0\% | 16 | 2.1\% | 34 | -0.8\% | 31 | 0.0\% | 63 | 0.0\% |
| 19.2K | 11 | 0.0\% | 23 | 0.0\% | 12 | 0.2\% | 25 | 0.2\% | 23 | 0.0\% | 47 | 0.0\% |
| 28.8K | 7 | 0.0\% | 15 | 0.0\% | 8 | -3.5\% | 16 | 2.1\% | 15 | 0.0\% | 31 | 0.0\% |
| 38.4K | 5 | 0.0\% | 11 | 0.0\% | 6 | -7.0\% | 12 | 0.2\% | 11 | 0.0\% | 23 | 0.0\% |
| 57.6K | 3 | 0.0\% | 7 | 0.0\% | 3 | 8.5\% | 8 | -3.5\% | 7 | 0.0\% | 15 | 0.0\% |
| 76.8K | 2 | 0.0\% | 5 | 0.0\% | 2 | 8.5\% | 6 | -7.0\% | 5 | 0.0\% | 11 | 0.0\% |
| 115.2K | 1 | 0.0\% | 3 | 0.0\% | 1 | 8.5\% | 3 | 8.5\% | 3 | 0.0\% | 7 | 0.0\% |
| 230.4K | 0 | 0.0\% | 1 | 0.0\% | 0 | 8.5\% | 1 | 8.5\% | 1 | 0.0\% | 3 | 0.0\% |
| 250K | 0 | -7.8\% | 1 | -7.8\% | 0 | 0.0\% | 1 | 0.0\% | 1 | -7.8\% | 3 | -7.8\% |
| 0.5M | - | - | 0 | -7.8\% | - | - | 0 | 0.0\% | 0 | -7.8\% | 1 | -7.8\% |
| 1M | - | - | - | - | - | - | - | - | - | - | 0 | -7.8\% |
| Max. ${ }^{(1)}$ | 230. | bps | 460. | bps | 250 | ps | 0.5 |  | 460. | Spps | 921. | Kpps |

Note: 1. $\operatorname{UBRR}=0$, Error $=0.0 \%$

Table 22-11. Examples of UBRRn Settings for Commonly Used Oscillator Frequencies

| Baud Rate [bps] | $\mathrm{f}_{\text {osc }}=8.0000 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=11.0592 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=14.7456 \mathrm{MHz}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  |
|  | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error |
| 2400 | 207 | 0.2\% | 416 | -0.1\% | 287 | 0.0\% | 575 | 0.0\% | 383 | 0.0\% | 767 | 0.0\% |
| 4800 | 103 | 0.2\% | 207 | 0.2\% | 143 | 0.0\% | 287 | 0.0\% | 191 | 0.0\% | 383 | 0.0\% |
| 9600 | 51 | 0.2\% | 103 | 0.2\% | 71 | 0.0\% | 143 | 0.0\% | 95 | 0.0\% | 191 | 0.0\% |
| 14.4K | 34 | -0.8\% | 68 | 0.6\% | 47 | 0.0\% | 95 | 0.0\% | 63 | 0.0\% | 127 | 0.0\% |
| 19.2K | 25 | 0.2\% | 51 | 0.2\% | 35 | 0.0\% | 71 | 0.0\% | 47 | 0.0\% | 95 | 0.0\% |
| 28.8K | 16 | 2.1\% | 34 | -0.8\% | 23 | 0.0\% | 47 | 0.0\% | 31 | 0.0\% | 63 | 0.0\% |
| 38.4K | 12 | 0.2\% | 25 | 0.2\% | 17 | 0.0\% | 35 | 0.0\% | 23 | 0.0\% | 47 | 0.0\% |
| 57.6K | 8 | -3.5\% | 16 | 2.1\% | 11 | 0.0\% | 23 | 0.0\% | 15 | 0.0\% | 31 | 0.0\% |
| 76.8K | 6 | -7.0\% | 12 | 0.2\% | 8 | 0.0\% | 17 | 0.0\% | 11 | 0.0\% | 23 | 0.0\% |
| 115.2K | 3 | 8.5\% | 8 | -3.5\% | 5 | 0.0\% | 11 | 0.0\% | 7 | 0.0\% | 15 | 0.0\% |
| 230.4K | 1 | 8.5\% | 3 | 8.5\% | 2 | 0.0\% | 5 | 0.0\% | 3 | 0.0\% | 7 | 0.0\% |
| 250K | 1 | 0.0\% | 3 | 0.0\% | 2 | -7.8\% | 5 | -7.8\% | 3 | -7.8\% | 6 | 5.3\% |
| 0.5M | 0 | 0.0\% | 1 | 0.0\% | - | - | 2 | -7.8\% | 1 | -7.8\% | 3 | -7.8\% |
| 1M | - | - | 0 | 0.0\% | - | - | - | - | 0 | -7.8\% | 1 | -7.8\% |
| Max. ${ }^{(1)}$ | 0.5 Mbps |  | 1Mbps |  | 691.2Kbps |  | 1.3824 Mbps |  | $921.6 \mathrm{Kbps}$ |  | 1.8432 Mbps |  |

Note: 1. $\operatorname{UBRR}=0$, Error $=0.0 \%$

Table 22-12. Examples of UBRRn Settings for Commonly Used Oscillator Frequencies

| Baud <br> Rate <br> [bps] | $f_{\text {osc }}=16.0000 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=18.4320 \mathrm{MHz}$ |  |  |  | $\mathrm{f}_{\text {osc }}=20.0000 \mathrm{MHz}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | U2Xn = 0 |  | U2Xn = 1 |  | U2Xn = 0 |  | U2Xn = 1 |  | $\mathrm{U} 2 \mathrm{Xn}=0$ |  | $\mathrm{U} 2 \mathrm{Xn}=1$ |  |
|  | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error | UBRR | Error |
| 2400 | 416 | -0.1\% | 832 | 0.0\% | 479 | 0.0\% | 959 | 0.0\% | 520 | 0.0\% | 1041 | 0.0\% |
| 4800 | 207 | 0.2\% | 416 | -0.1\% | 239 | 0.0\% | 479 | 0.0\% | 259 | 0.2\% | 520 | 0.0\% |
| 9600 | 103 | 0.2\% | 207 | 0.2\% | 119 | 0.0\% | 239 | 0.0\% | 129 | 0.2\% | 259 | 0.2\% |
| 14.4K | 68 | 0.6\% | 138 | -0.1\% | 79 | 0.0\% | 159 | 0.0\% | 86 | -0.2\% | 173 | -0.2\% |
| 19.2K | 51 | 0.2\% | 103 | 0.2\% | 59 | 0.0\% | 119 | 0.0\% | 64 | 0.2\% | 129 | 0.2\% |
| 28.8K | 34 | -0.8\% | 68 | 0.6\% | 39 | 0.0\% | 79 | 0.0\% | 42 | 0.9\% | 86 | -0.2\% |
| 38.4K | 25 | 0.2\% | 51 | 0.2\% | 29 | 0.0\% | 59 | 0.0\% | 32 | -1.4\% | 64 | 0.2\% |
| 57.6K | 16 | 2.1\% | 34 | -0.8\% | 19 | 0.0\% | 39 | 0.0\% | 21 | -1.4\% | 42 | 0.9\% |
| 76.8K | 12 | 0.2\% | 25 | 0.2\% | 14 | 0.0\% | 29 | 0.0\% | 15 | 1.7\% | 32 | -1.4\% |
| 115.2K | 8 | -3.5\% | 16 | 2.1\% | 9 | 0.0\% | 19 | 0.0\% | 10 | -1.4\% | 21 | -1.4\% |
| 230.4K | 3 | 8.5\% | 8 | -3.5\% | 4 | 0.0\% | 9 | 0.0\% | 4 | 8.5\% | 10 | -1.4\% |
| 250K | 3 | 0.0\% | 7 | 0.0\% | 4 | -7.8\% | 8 | 2.4\% | 4 | 0.0\% | 9 | 0.0\% |
| 0.5M | 1 | 0.0\% | 3 | 0.0\% | - | - | 4 | -7.8\% | - | - | 4 | 0.0\% |
| 1M | 0 | 0.0\% | 1 | 0.0\% | - | - | - | - | - | - | - | - |
| Max. ${ }^{(1)}$ | 1Mbps |  | 2Mbps |  | 1.152 Mbps |  | 2.304 Mbps |  | 1.25Mbps |  | 2.5 Mbps |  |

Note: 1. $\operatorname{UBRR}=0$, Error $=0.0 \%$

## 23. USART in SPI Mode

The Universal Synchronous and Asynchronous serial Receiver and Transmitter (USART) can be set to a master SPI compliant mode of operation. The Master SPI Mode (MSPIM) has the following features:

- Full Duplex, Three-wire Synchronous Data Transfer
- Master Operation
- Supports all four SPI Modes of Operation (Mode 0, 1, 2, and 3)
- LSB First or MSB First Data Transfer (Configurable Data Order)
- Queued Operation (Double Buffered)
- High Resolution Baud Rate Generator
- High Speed Operation (fXCKmax = fCK/2)
- Flexible Interrupt Generation


### 23.1 Overview

Setting both UMSELn1:0 bits to one enables the USART in MSPIM logic. In this mode of operation the SPI master control logic takes direct control over the USART resources. These resources include the transmitter and receiver shift register and buffers, and the baud rate generator. The parity generator and checker, the data and clock recovery logic, and the RX and TX control logic is disabled. The USART RX and TX control logic is replaced by a common SPI transfer control logic. However, the pin control logic and interrupt generation logic is identical in both modes of operation.

The I/O register locations are the same in both modes. However, some of the functionality of the control registers changes when using MSPIM.

### 23.2 USART MSPIM vs. SPI

The AVR USART in MSPIM mode is fully compatible with the AVR SPI regarding:

- Master mode timing diagram
- The UCPOLn bit functionality is identical to the SPI CPOL bit
- The UCPHAn bit functionality is identical to the SPI CPHA bit
- The UDORDn bit functionality is identical to the SPI DORD bit

However, since the USART in MSPIM mode reuses the USART resources, the use of the USART in MSPIM mode is somewhat different compared to the SPI. In addition to differences of the control register bits, and that only master operation is supported by the USART in MSPIM mode, the following features differ between the two modules:

- The USART in MSPIM mode includes (double) buffering of the transmitter. The SPI has no buffer
- The USART in MSPIM mode receiver includes an additional buffer level
- The SPI WCOL (Write Collision) bit is not included in USART in MSPIM mode
- The SPI double speed mode (SPI2X) bit is not included. However, the same effect is achieved by setting UBRRn accordingly
- Interrupt timing is not compatible
- Pin control differs due to the master only operation of the USART in MSPIM mode

A comparison of the USART in MSPIM mode and the SPI pins is shown in Table 23-4 on page 235.

### 23.2.1 Clock Generation

The Clock Generation logic generates the base clock for the Transmitter and Receiver. For USART MSPIM mode of operation only internal clock generation (that is, master operation) is supported. The Data Direction Register for the XCKn pin (DDR_XCKn) must therefore be set to one (that is, as output) for the USART in MSPIM to operate
correctly. Preferably the DDR_XCKn should be set up before the USART in MSPIM is enabled (that is, TXENn and RXENn bit set to one).

The internal clock generation used in MSPIM mode is identical to the USART synchronous master mode. The baud rate or UBRRn setting can therefore be calculated using the same equations, see Table 23-1.

Table 23-1. Equations for Calculating Baud Rate Register Setting

| Operating Mode | Equation for Calculating Baud Rate ${ }^{(1)}$ | Equation for Calculating UBRRn Value |
| :--- | :---: | :---: |
| Synchronous Master mode | $B A U D=\frac{f_{O S C}}{2(U B R R n+1)}$ | UBRRn $=\frac{f_{O S C}}{2 B A U D}-1$ |

Note: 1. The baud rate is defined to be the transfer rate in bit per second (bps).
BAUD Baud rate (in bits per second, bps).
$f_{\text {Osc }} \quad$ System Oscillator clock frequency.
UBRRn Contents of the UBRRnH and UBRRnL Registers, (0-4095).

### 23.3 SPI Data Modes and Timing

There are four combinations of XCKn (SCK) phase and polarity with respect to serial data, which are determined by control bits UCPHAn and UCPOLn. The data transfer timing diagrams are shown in Figure 23-1. Data bits are shifted out and latched in on opposite edges of the XCKn signal, ensuring sufficient time for data signals to stabilize. The UCPOLn and UCPHAn functionality is summarized in Table 23-2. Note that changing the setting of any of these bits will corrupt all ongoing communication for both the Receiver and Transmitter.

Table 23-2. UCPOLn and UCPHAn Functionality.

| UCPOLn | UCPHAn | SPI Mode | Leading Edge | Trailing Edge |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | Sample (Rising) | Setup (Falling) |
| 0 | 1 | 1 | Setup (Rising) | Sample (Falling) |
| 1 | 0 | 2 | Sample (Falling) | Setup (Rising) |
| 1 | 1 | 3 | Setup (Falling) | Sample (Rising) |

Figure 23-1. UCPHAn and UCPOLn data transfer timing diagrams.


### 23.4 Frame Formats

A serial frame for the MSPIM is defined to be one character of 8 data bits. The USART in MSPIM mode has two valid frame formats:

- 8-bit data with MSB first
- 8-bit data with LSB first

A frame starts with the least or most significant data bit. Then the next data bits, up to a total of eight, are succeeding, ending with the most or least significant bit accordingly. When a complete frame is transmitted, a new frame can directly follow it, or the communication line can be set to an idle (high) state.
The UDORDn bit in UCSRnC sets the frame format used by the USART in MSPIM mode. The Receiver and Transmitter use the same setting. Note that changing the setting of any of these bits will corrupt all ongoing communication for both the Receiver and Transmitter.

16-bit data transfer can be achieved by writing two data bytes to UDRn. A UART transmit complete interrupt will then signal that the 16-bit value has been shifted out.

### 23.4.1 USART MSPIM Initialization

The USART in MSPIM mode has to be initialized before any communication can take place. The initialization process normally consists of setting the baud rate, setting master mode of operation (by setting DDR_XCKn to one), setting frame format and enabling the Transmitter and the Receiver. Only the transmitter can operate independently. For interrupt driven USART operation, the Global Interrupt Flag should be cleared (and thus interrupts globally disabled) when doing the initialization.

Note: To ensure immediate initialization of the XCKn output the baud-rate register (UBRRn) must be zero at the time the transmitter is enabled. Contrary to the normal mode USART operation the UBRRn must then be written to the desired value after the transmitter is enabled, but before the first transmission is started. Setting UBRRn to zero before enabling the transmitter is not necessary if the initialization is done immediately after a reset since UBRRn is reset to zero.
Before doing a re-initialization with changed baud rate, data mode, or frame format, be sure that there is no ongoing transmissions during the period the registers are changed. The TXCn Flag can be used to check that the Transmitter has completed all transfers, and the RXCn Flag can be used to check that there are no unread data in the receive buffer. Note that the TXCn Flag must be cleared before each transmission (before UDRn is written) if it is used for this purpose.

The following simple USART initialization code examples show one assembly and one C function that are equal in functionality. The examples assume polling (no interrupts enabled). The baud rate is given as a function parameter. For the assembly code, the baud rate parameter is assumed to be stored in the r17:r16 registers.

```
Assembly Code Example (1)
    USART_Init:
        clr r18
        out UBRRnH,r18
        out UBRRnL,r18
        ; Setting the XCKn port pin as output, enables master mode.
        sbi XCKn_DDR, XCKn
        ; Set MSPI mode of operation and SPI data mode 0.
        ldi r18, (1<<UMSELn1)|(1<<UMSELn0)|(0<<UCPHAn)|(0<<UCPOLn)
        out UCSRnC,r18
        ; Enable receiver and transmitter.
        1di r18, (1<<RXENn)|(1<<TXENn)
        out UCSRnB,r18
        ; Set baud rate.
        ; IMPORTANT: The Baud Rate must be set after the transmitter is
enabled!
        out UBRRnH, r17
        out UBRRnL, r18
        ret
C Code Example }\mp@subsup{}{}{(1)
    void USART_Init( unsigned int baud )
    {
        UBRRn = 0;
        /* Setting the XCKn port pin as output, enables master mode. */
        XCKn_DDR |= (1<<XCKn);
        /* Set MSPI mode of operation and SPI data mode 0. */
        UCSRnC = (1<<UMSELn1)|(1<<UMSELn0)|(0<<UCPHAn)|(0<<UCPOLn);
        /* Enable receiver and transmitter. */
        UCSRnB = (1<<RXENn) | (1<<TXENn);
        /* Set baud rate. */
        /* IMPORTANT: The Baud Rate must be set after the transmitter is
    enabled */
        UBRRn = baud;
    }
```

Note: 1. See "About Code Examples" on page 10.

### 23.5 Data Transfer

Using the USART in MSPI mode requires the Transmitter to be enabled, that is, the TXENn bit in the UCSRnB register is set to one. When the Transmitter is enabled, the normal port operation of the TxDn pin is overridden and given the function as the Transmitter's serial output. Enabling the receiver is optional and is done by setting the RXENn bit in the UCSRnB register to one. When the receiver is enabled, the normal pin operation of the RxDn pin is overridden and given the function as the Receiver's serial input. The XCKn will in both cases be used as the transfer clock.

After initialization the USART is ready for doing data transfers. A data transfer is initiated by writing to the UDRn I/O location. This is the case for both sending and receiving data since the transmitter controls the transfer clock. The data written to UDRn is moved from the transmit buffer to the shift register when the shift register is ready to send a new frame.

Note: To keep the input buffer in sync with the number of data bytes transmitted, the UDRn register must be read once for each byte transmitted. The input buffer operation is identical to normal USART mode, that is, if an overflow occurs the character last received will be lost, not the first data in the buffer. This means that if four bytes are transferred, byte 1 first, then byte 2,3 , and 4 , and the UDRn is not read before all transfers are completed, then byte 3 to be received will be lost, and not byte 1 .
The following code examples show a simple USART in MSPIM mode transfer function based on polling of the Data Register Empty (UDREn) Flag and the Receive Complete (RXCn) Flag. The USART has to be initialized before the function can be used. For the assembly code, the data to be sent is assumed to be stored in Register R16 and the data received will be available in the same register (R16) after the function returns.

The function simply waits for the transmit buffer to be empty by checking the UDREn Flag, before loading it with new data to be transmitted. The function then waits for data to be present in the receive buffer by checking the RXCn Flag, before reading the buffer and returning the value.

```
Assembly Code Example(1)
    USART_MSPIM_Transfer:
    ; Wait for empty transmit buffer
    sbis UCSRnA, UDREn
    rjmp USART_MSPIM_Transfer
    ; Put data (r16) into buffer, sends the data
    out UDRn,r16
    ; Wait for data to be received
    USART_MSPIM_Wait_RXCn:
    sbis UCSRnA, RXCn
    rjmp USART_MSPIM_Wait_RXCn
    ; Get and return received data from buffer
    in r16, UDRn
    ret
C Code Example }\mp@subsup{}{}{(1)
    unsigned char USART_Receive( void )
    {
    /* Wait for empty transmit buffer */
    while ( !( UCSRnA & (1<<UDREn)) );
    /* Put data into buffer, sends the data */
    UDRn = data;
    /* Wait for data to be received */
    while ( !(UCSRnA & (1<<RXCn)) );
    /* Get and return received data from buffer */
    return UDRn;
}
```

Note: 1. See "About Code Examples" on page 10.

### 23.5.1 Transmitter and Receiver Flags and Interrupts

The RXCn, TXCn, and UDREn flags and corresponding interrupts in USART in MSPIM mode are identical in function to the normal USART operation. However, the receiver error status flags (FE, DOR, and PE) are not in use and is always read as zero.

### 23.5.2 Disabling the Transmitter or Receiver

The disabling of the transmitter or receiver in USART in MSPIM mode is identical in function to the normal USART operation.

### 23.6 USART MSPIM Register Description

The following section describes the registers used for SPI operation using the USART.

### 23.6.1 UDRn - USART MSPIM I/O Data Register

The function and bit description of the USART data register (UDRn) in MSPI mode is identical to normal USART operation. See "UDRn - USART I/O Data Register n" on page 218.

### 23.6.2 UCSRnA - USART MSPIM Control and Status Register n A

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | UCSRnA |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | RXCn | TXCn | UDREn | - | - | - | - | - |  |
| Read/Write | R/W | R/W | R/W | R | R | R | R | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |  |

- Bit 7 - RXCn: USART Receive Complete

This flag bit is set when there are unread data in the receive buffer and cleared when the receive buffer is empty (that is, does not contain any unread data). If the Receiver is disabled, the receive buffer will be flushed and consequently the RXCn bit will become zero. The RXCn Flag can be used to generate a Receive Complete interrupt (see description of the RXCIEn bit).

## - Bit 6 - TXCn: USART Transmit Complete

This flag bit is set when the entire frame in the Transmit Shift Register has been shifted out and there are no new data currently present in the transmit buffer (UDRn). The TXCn Flag bit is automatically cleared when a transmit complete interrupt is executed, or it can be cleared by writing a one to its bit location. The TXCn Flag can generate a Transmit Complete interrupt (see description of the TXCIEn bit).

## - Bit 5 - UDREn: USART Data Register Empty

The UDREn Flag indicates if the transmit buffer (UDRn) is ready to receive new data. If UDREn is one, the buffer is empty, and therefore ready to be written. The UDREn Flag can generate a Data Register Empty interrupt (see description of the UDRIE bit). UDREn is set after a reset to indicate that the Transmitter is ready.

## - Bit 4:0-Reserved Bits in MSPI mode

When in MSPI mode, these bits are reserved for future use. For compatibility with future devices, these bits must be written to zero when UCSRnA is written.

### 23.6.3 UCSRnB - USART MSPIM Control and Status Register n B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | UCSRnB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | RXCIEn | TXCIEn | UDRIE | RXENn | TXENn | - | - | - |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R | R | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |  |

## - Bit 7 - RXCIEn: RX Complete Interrupt Enable

Writing this bit to one enables interrupt on the RXCn Flag. A USART Receive Complete interrupt will be generated only if the RXCIEn bit is written to one, the Global Interrupt Flag in SREG is written to one and the RXCn bit in UCSRnA is set.

## - Bit 6 - TXCIEn: TX Complete Interrupt Enable

Writing this bit to one enables interrupt on the TXCn Flag. A USART Transmit Complete interrupt will be generated only if the TXCIEn bit is written to one, the Global Interrupt Flag in SREG is written to one and the TXCn bit in UCSRnA is set.

- Bit 5 - UDRIE: USART Data Register Empty Interrupt Enable

Writing this bit to one enables interrupt on the UDREn Flag. A Data Register Empty interrupt will be generated only if the UDRIE bit is written to one, the Global Interrupt Flag in SREG is written to one and the UDREn bit in UCSRnA is set.

## - Bit 4 - RXENn: Receiver Enable

Writing this bit to one enables the USART Receiver in MSPIM mode. The Receiver will override normal port operation for the RxDn pin when enabled. Disabling the Receiver will flush the receive buffer. Only enabling the receiver in MSPI mode (that is, setting RXENn=1 and TXENn=0) has no meaning since it is the transmitter that controls the transfer clock and since only master mode is supported.

## - Bit 3 - TXENn: Transmitter Enable

Writing this bit to one enables the USART Transmitter. The Transmitter will override normal port operation for the TxDn pin when enabled. The disabling of the Transmitter (writing TXENn to zero) will not become effective until ongoing and pending transmissions are completed, that is, when the Transmit Shift Register and Transmit Buffer Register do not contain data to be transmitted. When disabled, the Transmitter will no longer override the TxDn port.

## - Bit 2:0-Reserved Bits in MSPI mode

When in MSPI mode, these bits are reserved for future use. For compatibility with future devices, these bits must be written to zero when UCSRnB is written.

### 23.6.4 UCSRnC - USART MSPIM Control and Status Register n C

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | UCSRnC |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | UMSELn1 | UMSELn0 | - | - | - | UDORDn | UCPHAn | UCPOLn |  |
| Read/Write | R/W | R/W | R | R | R | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |  |

## - Bit 7:6-UMSELn1:0: USART Mode Select

These bits select the mode of operation of the USART as shown in Table 23-3. See "UCSRnC - USART Control and Status Register $n$ C" on page 221 for full description of the normal USART operation. The MSPIM is enabled when both UMSELn bits are set to one. The UDORDn, UCPHAn, and UCPOLn can be set in the same write operation where the MSPIM is enabled.

Table 23-3. UMSELn Bits Settings

| UMSELn1 | UMSELn0 | Mode |
| :---: | :---: | :---: |
| 0 | 0 | Asynchronous USART |
| 0 | 1 | Synchronous USART |
| 1 | 0 | (Reserved) |
| 1 | 1 | Master SPI (MSPIM) |

## - Bit 5:3-Reserved Bits in MSPI mode

When in MSPI mode, these bits are reserved for future use. For compatibility with future devices, these bits must be written to zero when UCSRnC is written.

## - Bit 2 - UDORDn: Data Order

When set to one the LSB of the data word is transmitted first. When set to zero the MSB of the data word is transmitted first. Refer to "SPI Data Modes and Timing" on page 228 for details.

## - Bit 1 - UCPHAn: Clock Phase

The UCPHAn bit setting determine if data is sampled on the leasing edge (first) or tailing (last) edge of XCKn. Refer to "SPI Data Modes and Timing" on page 228 for details.

## - Bit 0 - UCPOLn: Clock Polarity

The UCPOLn bit sets the polarity of the XCKn clock. The combination of the UCPOLn and UCPHAn bit settings determine the timing of the data transfer. Refer to "SPI Data Modes and Timing" on page 228 for details.

### 23.6.5 UBRRnL and UBRRnH - USART MSPIM Baud Rate Registers

The function and bit description of the baud rate registers in MSPI mode is identical to normal USART operation. See "UBRRnL and UBRRnH - USART Baud Rate Registers" on page 222.

Table 23-4. Comparison of USART in MSPIM mode and SPI pins.

| USART_MSPIM | SPI | Comment |
| :---: | :---: | :---: |
| TxDn | MOSI | Master Out only |
| RxDn | MISO | Master In only |
| XCKn | SCK | (Functionally identical) |
| $(\mathrm{N} / \mathrm{A})$ | $\overline{S S}$ | Not supported by USART in MSPIM |

## 24. 2-wire Serial Interface

### 24.1 Features

- Simple yet Powerful and Flexible Communication Interface, only two Bus Lines needed
- Both Master and Slave Operation Supported
- Device can Operate as Transmitter or Receiver
- 7-bit Address Space Allows up to 128 Different Slave Addresses
- Multi-master Arbitration Support
- Up to 400kHz Data Transfer Speed
- Slew-rate Limited Output Drivers
- Noise Suppression Circuitry Rejects Spikes on Bus Lines
- Fully Programmable Slave Address with General Call Support
- Address Recognition Causes Wake-up When AVR is in Sleep Mode


### 24.2 2-wire Serial Interface Bus Definition

The 2-wire Serial Interface (TWI) is ideally suited for typical microcontroller applications. The TWI protocol allows the systems designer to interconnect up to 128 different devices using only two bi-directional bus lines, one for clock (SCL) and one for data (SDA). The only external hardware needed to implement the bus is a single pull-up resistor for each of the TWI bus lines. All devices connected to the bus have individual addresses, and mechanisms for resolving bus contention are inherent in the TWI protocol.

Figure 24-1. TWI Bus Interconnection


### 24.2.1 TWI Terminology

The following definitions are frequently encountered in this section.
Table 24-1. TWI Terminology

| Term | Description |
| :--- | :--- |
| Master | The device that initiates and terminates a transmission. The Master also generates the SCL clock |
| Slave | The device addressed by a Master |
| Transmitter | The device placing data on the bus |
| Receiver | The device reading data from the bus |

The Power Reduction TWI bit, PRTWI bit in "PRR0 - Power Reduction Register 0" on page 55 must be written to zero to enable the 2-wire Serial Interface.

### 24.2.2 Electrical Interconnection

As depicted in Figure 24-1 on page 236, both bus lines are connected to the positive supply voltage through pull-up resistors. The bus drivers of all TWI-compliant devices are open-drain or open-collector. This implements a wiredAND function which is essential to the operation of the interface. A low level on a TWI bus line is generated when one or more TWI devices output a zero. A high level is output when all TWI devices trim-state their outputs, allowing the pull-up resistors to pull the line high. Note that all AVR devices connected to the TWI bus must be powered in order to allow any bus operation.

The number of devices that can be connected to the bus is only limited by the bus capacitance limit of 400pF and the 7-bit slave address space. A detailed specification of the electrical characteristics of the TWI is given in "SPI Timing Characteristics" on page 363. Two different sets of specifications are presented there, one relevant for bus speeds below 100 kHz , and one valid for bus speeds up to 400 kHz .

### 24.3 Data Transfer and Frame Format

### 24.3.1 Transferring Bits

Each data bit transferred on the TWI bus is accompanied by a pulse on the clock line. The level of the data line must be stable when the clock line is high. The only exception to this rule is for generating start and stop conditions.

Figure 24-2. Data Validity


Data Change

### 24.3.2 START and STOP Conditions

The Master initiates and terminates a data transmission. The transmission is initiated when the Master issues a START condition on the bus, and it is terminated when the Master issues a STOP condition. Between a START and a STOP condition, the bus is considered busy, and no other master should try to seize control of the bus. A special case occurs when a new START condition is issued between a START and STOP condition. This is referred to as a REPEATED START condition, and is used when the Master wishes to initiate a new transfer without relinquishing control of the bus. After a REPEATED START, the bus is considered busy until the next STOP. This is identical to the START behavior, and therefore START is used to describe both START and REPEATED START for the remainder of this datasheet, unless otherwise noted. As depicted below, START and STOP conditions are signalled by changing the level of the SDA line when the SCL line is high.

Figure 24-3. START, REPEATED START and STOP conditions


### 24.3.3 Address Packet Format

All address packets transmitted on the TWI bus are 9 bits long, consisting of 7 address bits, one READ/WRITE control bit and an acknowledge bit. If the READ/WRITE bit is set, a read operation is to be performed, otherwise a write operation should be performed. When a Slave recognizes that it is being addressed, it should acknowledge by pulling SDA low in the ninth SCL (ACK) cycle. If the addressed Slave is busy, or for some other reason can not service the Master's request, the SDA line should be left high in the ACK clock cycle. The Master can then transmit a STOP condition, or a REPEATED START condition to initiate a new transmission. An address packet consisting of a slave address and a READ or a WRITE bit is called SLA + R or SLA+W, respectively.
The MSB of the address byte is transmitted first. Slave addresses can freely be allocated by the designer, but the address 0000000 is reserved for a general call.
When a general call is issued, all slaves should respond by pulling the SDA line low in the ACK cycle. A general call is used when a Master wishes to transmit the same message to several slaves in the system. When the general call address followed by a Write bit is transmitted on the bus, all slaves set up to acknowledge the general call will pull the SDA line low in the ack cycle. The following data packets will then be received by all the slaves that acknowledged the general call. Note that transmitting the general call address followed by a Read bit is meaningless, as this would cause contention if several slaves started transmitting different data.

All addresses of the format 1111 xxx should be reserved for future purposes.
Figure 24-4. Address Packet Format


### 24.3.4 Data Packet Format

All data packets transmitted on the TWI bus are nine bits long, consisting of one data byte and an acknowledge bit. During a data transfer, the Master generates the clock and the START and STOP conditions, while the Receiver is responsible for acknowledging the reception. An Acknowledge (ACK) is signalled by the Receiver pulling the SDA line low during the ninth SCL cycle. If the Receiver leaves the SDA line high, a NACK is signalled. When the Receiver has received the last byte, or for some reason cannot receive any more bytes, it should inform the Transmitter by sending a NACK after the final byte. The MSB of the data byte is transmitted first.

Figure 24-5. Data Packet Format


### 24.3.5 Combining Address and Data Packets into a Transmission

A transmission basically consists of a START condition, a SLA+R/W, one or more data packets and a STOP condition. An empty message, consisting of a START followed by a STOP condition, is illegal. Note that the WiredANDing of the SCL line can be used to implement handshaking between the Master and the Slave. The Slave can extend the SCL low period by pulling the SCL line low. This is useful if the clock speed set up by the Master is too fast for the Slave, or the Slave needs extra time for processing between the data transmissions. The Slave extending the SCL low period will not affect the SCL high period, which is determined by the Master. As a consequence, the Slave can reduce the TWI data transfer speed by prolonging the SCL duty cycle.
Figure 24-6 shows a typical data transmission. Note that several data bytes can be transmitted between the SLA+R/W and the STOP condition, depending on the software protocol implemented by the application software.

Figure 24-6. Typical Data Transmission


### 24.4 Multi-master Bus Systems, Arbitration, and Synchronization

The TWI protocol allows bus systems with several masters. Special concerns have been taken in order to ensure that transmissions will proceed as normal, even if two or more masters initiate a transmission at the same time. Two problems arise in multi-master systems:

- An algorithm must be implemented allowing only one of the masters to complete the transmission. All other masters should cease transmission when they discover that they have lost the selection process. This selection process is called arbitration. When a contending master discovers that it has lost the arbitration process, it should immediately switch to Slave mode to check whether it is being addressed by the winning master. The fact that multiple masters have started transmission at the same time should not be detectable to the slaves, that is, the data being transferred on the bus must not be corrupted.
- Different masters may use different SCL frequencies. A scheme must be devised to synchronize the serial clocks from all masters, in order to let the transmission proceed in a lockstep fashion. This will facilitate the arbitration process.

The wired-ANDing of the bus lines is used to solve both these problems. The serial clocks from all masters will be wired-ANDed, yielding a combined clock with a high period equal to the one from the Master with the shortest high period. The low period of the combined clock is equal to the low period of the Master with the longest low period. Note that all masters listen to the SCL line, effectively starting to count their SCL high and low time-out periods when the combined SCL line goes high or low, respectively.

Figure 24-7. SCL Synchronization Between Multiple Masters


Arbitration is carried out by all masters continuously monitoring the SDA line after outputting data. If the value read from the SDA line does not match the value the Master had output, it has lost the arbitration. Note that a Master can only lose arbitration when it outputs a high SDA value while another Master outputs a low value. The losing Master should immediately go to Slave mode, checking if it is being addressed by the winning Master. The SDA line should be left high, but losing masters are allowed to generate a clock signal until the end of the current data or address packet. Arbitration will continue until only one Master remains, and this may take many bits. If several masters are trying to address the same Slave, arbitration will continue into the data packet.

Figure 24-8. Arbitration Between Two Masters


Note that arbitration is not allowed between:

- A REPEATED START condition and a data bit
- A STOP condition and a data bit
- A REPEATED START and a STOP condition

It is the user software's responsibility to ensure that these illegal arbitration conditions never occur. This implies that in multi-master systems, all data transfers must use the same composition of SLA+R/W and data packets. In other words: All transmissions must contain the same number of data packets, otherwise the result of the arbitration is undefined.

### 24.5 Overview of the TWI Module

The TWI module is comprised of several sub-modules, as shown in Figure 24-9 on page 242. All registers drawn in a thick line are accessible through the AVR data bus.

Figure 24-9. Overview of the TWI Module


### 24.5.1 SCL and SDA Pins

These pins interface the AVR TWI with the rest of the MCU system. The output drivers contain a slew-rate limiter in order to conform to the TWI specification. The input stages contain a spike suppression unit removing spikes shorter than 50 ns . Note that the internal pull-ups in the AVR pads can be enabled by setting the PORT bits corresponding to the SCL and SDA pins, as explained in the I/O Port section. The internal pull-ups can in some systems eliminate the need for external ones.

### 24.5.2 Bit Rate Generator Unit

This unit controls the period of SCL when operating in a Master mode. The SCL period is controlled by settings in the TWI Bit Rate Register (TWBR) and the Prescaler bits in the TWI Status Register (TWSR). Slave operation does not depend on Bit Rate or Prescaler settings, but the CPU clock frequency in the Slave must be at least 16 times higher than the SCL frequency. Note that slaves may prolong the SCL low period, thereby reducing the average TWI bus clock period.

The SCL frequency is generated according to the following equation:

$$
\text { SCL frequency }=\frac{\text { CPU Clock frequency }}{16+2(\text { TWBR }) \cdot 4^{T W P S}}
$$

- $\quad$ TWBR $=$ Value of the TWI Bit Rate Register
- TWPS = Value of the prescaler bits in the TWI Status Register

Note: Pull-up resistor values should be selected according to the SCL frequency and the capacitive bus line load. See " 2 -wire Serial Interface Characteristics" on page 361 for value of pull-up resistor.

### 24.5.3 Bus Interface Unit

This unit contains the Data and Address Shift Register (TWDR), a START/STOP Controller and Arbitration detection hardware. The TWDR contains the address or data bytes to be transmitted, or the address or data bytes received. In addition to the 8-bit TWDR, the Bus Interface Unit also contains a register containing the (N)ACK bit to be transmitted or received. This (N)ACK Register is not directly accessible by the application software. However, when receiving, it can be set or cleared by manipulating the TWI Control Register (TWCR). When in Transmitter mode, the value of the received ( N )ACK bit can be determined by the value in the TWSR.

The START/STOP Controller is responsible for generation and detection of START, REPEATED START, and STOP conditions. The START/STOP controller is able to detect START and STOP conditions even when the AVR MCU is in one of the sleep modes, enabling the MCU to wake up if addressed by a Master.

If the TWI has initiated a transmission as Master, the Arbitration Detection hardware continuously monitors the transmission trying to determine if arbitration is in process. If the TWI has lost an arbitration, the Control Unit is informed. Correct action can then be taken and appropriate status codes generated.

### 24.5.4 Address Match Unit

The Address Match unit checks if received address bytes match the seven-bit address in the TWI Address Register (TWAR). If the TWI General Call Recognition Enable (TWGCE) bit in the TWAR is written to one, all incoming address bits will also be compared against the General Call address. Upon an address match, the Control Unit is informed, allowing correct action to be taken. The TWI may or may not acknowledge its address, depending on settings in the TWCR. The Address Match unit is able to compare addresses even when the AVR MCU is in sleep mode, enabling the MCU to wake up if addressed by a Master. If another interrupt (for example, INTO) occurs during TWI Power-down address match and wakes up the CPU, the TWI aborts operation and return to it's idle state. If this cause any problems, ensure that TWI Address Match is the only enabled interrupt when entering Powerdown.

### 24.5.5 Control Unit

The Control unit monitors the TWI bus and generates responses corresponding to settings in the TWI Control Register (TWCR). When an event requiring the attention of the application occurs on the TWI bus, the TWI Interrupt Flag (TWINT) is asserted. In the next clock cycle, the TWI Status Register (TWSR) is updated with a status code identifying the event. The TWSR only contains relevant status information when the TWI Interrupt Flag is asserted. At all other times, the TWSR contains a special status code indicating that no relevant status information is available. As long as the TWINT Flag is set, the SCL line is held low. This allows the application software to complete its tasks before allowing the TWI transmission to continue.

The TWINT Flag is set in the following situations:

- After the TWI has transmitted a START/REPEATED START condition
- After the TWI has transmitted SLA+R/W
- After the TWI has transmitted an address byte
- After the TWI has lost arbitration
- After the TWI has been addressed by own slave address or general call
- After the TWI has received a data byte
- After a STOP or REPEATED START has been received while still addressed as a Slave
- When a bus error has occurred due to an illegal START or STOP condition


### 24.6 Using the TWI

The AVR TWI is byte-oriented and interrupt based. Interrupts are issued after all bus events, like reception of a byte or transmission of a START condition. Because the TWI is interrupt-based, the application software is free to carry on other operations during a TWI byte transfer. Note that the TWI Interrupt Enable (TWIE) bit in TWCR together with the Global Interrupt Enable bit in SREG allow the application to decide whether or not assertion of the TWINT Flag should generate an interrupt request. If the TWIE bit is cleared, the application must poll the TWINT Flag in order to detect actions on the TWI bus.

When the TWINT Flag is asserted, the TWI has finished an operation and awaits application response. In this case, the TWI Status Register (TWSR) contains a value indicating the current state of the TWI bus. The application software can then decide how the TWI should behave in the next TWI bus cycle by manipulating the TWCR and TWDR Registers.

Figure 24-10 is a simple example of how the application can interface to the TWI hardware. In this example, a Master wishes to transmit a single data byte to a Slave. This description is quite abstract, a more detailed explanation follows later in this section. A simple code example implementing the desired behavior is also presented.

Figure 24-10. Interfacing the Application to the TWI in a Typical Transmission


1. The first step in a TWI transmission is to transmit a START condition. This is done by writing a specific value into TWCR, instructing the TWI hardware to transmit a START condition. Which value to write is described later on. However, it is important that the TWINT bit is set in the value written. Writing a one to TWINT clears the flag. The TWI will not start any operation as long as the TWINT bit in TWCR is set. Immediately after the application has cleared TWINT, the TWI will initiate transmission of the START condition.
2. When the START condition has been transmitted, the TWINT Flag in TWCR is set, and TWSR is updated with a status code indicating that the START condition has successfully been sent.
3. The application software should now examine the value of TWSR, to make sure that the START condition was successfully transmitted. If TWSR indicates otherwise, the application software might take some special action, like calling an error routine. Assuming that the status code is as expected, the application must load SLA+W into TWDR. Remember that TWDR is used both for address and data. After TWDR has been loaded with the desired SLA+W, a specific value must be written to TWCR, instructing the TWI hardware to transmit the SLA+W present in TWDR. Which value to write is described later on. However, it is important that the TWINT bit is set in the value written. Writing a one to TWINT clears the flag. The TWI will not start any operation as long as the TWINT bit in TWCR is set. Immediately after the application has cleared TWINT, the TWI will initiate transmission of the address packet.
4. When the address packet has been transmitted, the TWINT Flag in TWCR is set, and TWSR is updated with a status code indicating that the address packet has successfully been sent. The status code will also reflect whether a Slave acknowledged the packet or not.
5. The application software should now examine the value of TWSR, to make sure that the address packet was successfully transmitted, and that the value of the ACK bit was as expected. If TWSR indicates otherwise, the application software might take some special action, like calling an error routine. Assuming that the status code is as expected, the application must load a data packet into TWDR. Subsequently, a specific value must be written to TWCR, instructing the TWI hardware to transmit the data packet present in TWDR. Which value to write is described later on. However, it is important that the TWINT bit is set in the value written. Writing a one to TWINT clears the flag. The TWI will not start any operation as long as the TWINT bit in TWCR is set. Immediately after the application has cleared TWINT, the TWI will initiate transmission of the data packet.
6. When the data packet has been transmitted, the TWINT Flag in TWCR is set, and TWSR is updated with a status code indicating that the data packet has successfully been sent. The status code will also reflect whether a Slave acknowledged the packet or not.
7. The application software should now examine the value of TWSR, to make sure that the data packet was successfully transmitted, and that the value of the ACK bit was as expected. If TWSR indicates otherwise, the application software might take some special action, like calling an error routine. Assuming that the status code is as expected, the application must write a specific value to TWCR, instructing the TWI hardware to transmit a STOP condition. Which value to write is described later on. However, it is important that the TWINT bit is set in the value written. Writing a one to TWINT clears the flag. The TWI will not start any operation as long as the TWINT bit in TWCR is set. Immediately after the application has cleared TWINT, the TWI will initiate transmission of the STOP condition. Note that TWINT is NOT set after a STOP condition has been sent.
Even though this example is simple, it shows the principles involved in all TWI transmissions. These can be summarized as follows:

- When the TWI has finished an operation and expects application response, the TWINT Flag is set. The SCL line is pulled low until TWINT is cleared.
- When the TWINT Flag is set, the user must update all TWI Registers with the value relevant for the next TWI bus cycle. As an example, TWDR must be loaded with the value to be transmitted in the next bus cycle.
- After all TWI Register updates and other pending application software tasks have been completed, TWCR is written. When writing TWCR, the TWINT bit should be set. Writing a one to TWINT clears the flag. The TWI will then commence executing whatever operation was specified by the TWCR setting.
In the following an assembly and C implementation of the example is given. Note that the code below assumes that several definitions have been made, for example by using include-files.

|  | Assembly Code Example | C Example | Comments |
| :---: | :---: | :---: | :---: |
| 1 | ```ldi r16, (1<<TWINT)\|(1<<TWSTA)| (1<<TWEN) out TWCR, r16``` | $\begin{aligned} & \text { TWCR }=(1 \ll \text { TWINT }) \mid(1 \ll \text { TWSTA }) \mid \\ & (1 \ll \text { TWEN }) \end{aligned}$ | Send START condition |
| 2 | ```wait1: in r16,TWCR sbrs r16,TWINT rjmp wait1``` | while (!(TWCR \& (1<<TWINT))) ; | Wait for TWINT Flag set. This indicates that the START condition has been transmitted |
| 3 | ```in r16,TWSR andi r16, 0xF8 cpi r16, START brne ERROR``` | if ((TWSR \& 0xF8) != START) ERROR(); | Check value of TWI Status Register. Mask prescaler bits. If status different from START go to ERROR |
|  | ```ldi r16, SLA_W out TWDR, r16 ldi r16, (1<<TWINT) \| (1<<TWEN) out TWCR, r16``` | ```TWDR = SLA_W; TWCR = (1<<TWINT) (1<<TWEN);``` | Load SLA_W into TWDR Register. Clear TWINT bit in TWCR to start transmission of address |
| 4 | ```wait2: in r16,TWCR sbrs r16,TWINT rjmp wait2``` | while (!(TWCR \& (1<<TWINT))) | Wait for TWINT Flag set. This indicates that the SLA+W has been transmitted, and ACK/NACK has been received. |
| 5 | ```in r16,TWSR andi r16, 0xF8 cpi r16, MT_SLA_ACK brne ERROR``` | ```if ((TWSR & 0xF8) != MT_SLA_ACK) ERROR();``` | Check value of TWI Status Register. Mask prescaler bits. If status different from MT_SLA_ACK go to ERROR |
|  | ```1di r16, DATA out TWDR, r16 ldi r16, (1<<TWINT) \| (1<<TWEN) out TWCR, r16``` | ```TWDR = DATA; TWCR = (1<<TWINT) \| (1<<TWEN);``` | Load DATA into TWDR Register. Clear TWINT bit in TWCR to start transmission of data |
| 6 | wait3: <br> in r16,TWCR <br> sbrs r16, TWINT <br> rjmp wait3 | while (!(TWCR \& (1<<TWINT))) | Wait for TWINT Flag set. This indicates that the DATA has been transmitted, and ACK/NACK has been received. |
| 7 | ```in r16,TWSR andi r16, 0xF8 cpi r16, MT_DATA_ACK brne ERROR``` | ```if ((TWSR & 0xF8) != MT_DATA_ACK) ERROR();``` | Check value of TWI Status <br> Register. Mask prescaler bits. If status different from <br> MT_DATA_ACK go to ERROR |
|  | ```ldi r16, (1<<TWINT)\|(1<<TWEN)| (1<<TWSTO) out TWCR, r16``` | ```TWCR = (1<<TWINT) \| (1<<TWEN ) | (1<<TWSTO);``` | Transmit STOP condition |

### 24.7 Transmission Modes

The TWI can operate in one of four major modes. These are named Master Transmitter (MT), Master Receiver (MR), Slave Transmitter (ST) and Slave Receiver (SR). Several of these modes can be used in the same application. As an example, the TWI can use MT mode to write data into a TWI EEPROM, MR mode to read the data back from the EEPROM. If other masters are present in the system, some of these might transmit data to the TWI, and then SR mode would be used. It is the application software that decides which modes are legal.
The following sections describe each of these modes. Possible status codes are described along with figures detailing data transmission in each of the modes. These figures contain the following abbreviations:
S: START condition
Rs: REPEATED START condition
R: Read bit (high level at SDA)
W: Write bit (low level at SDA)
A: Acknowledge bit (low level at SDA)
$\overline{\text { A }} \quad$ Not acknowledge bit (high level at SDA)
Data: 8-bit data byte
P: STOP condition

## SLA: Slave Address

In Figure 24-12 on page 250 to Figure 24-18 on page 258, circles are used to indicate that the TWINT Flag is set. The numbers in the circles show the status code held in TWSR, with the prescaler bits masked to zero. At these points, actions must be taken by the application to continue or complete the TWI transfer. The TWI transfer is suspended until the TWINT Flag is cleared by software.
When the TWINT Flag is set, the status code in TWSR is used to determine the appropriate software action. For each status code, the required software action and details of the following serial transfer are given in Table 24-2 on page 249 to Table 24-5 on page 257. Note that the prescaler bits are masked to zero in these tables.

### 24.7.1 Master Transmitter Mode

In the Master Transmitter mode, a number of data bytes are transmitted to a Slave Receiver (see Figure 24-11). In order to enter a Master mode, a START condition must be transmitted. The format of the following address packet determines whether Master Transmitter or Master Receiver mode is to be entered. If SLA+W is transmitted, MT mode is entered, if SLA+R is transmitted, MR mode is entered. All the status codes mentioned in this section assume that the prescaler bits are zero or are masked to zero.

Figure 24-11. Data Transfer in Master Transmitter Mode


A START condition is sent by writing the following value to TWCR:

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 1 | X | 1 | 0 | X | 1 | 0 | X |

TWEN must be set to enable the 2-wire Serial Interface, TWSTA must be written to one to transmit a START condition and TWINT must be written to one to clear the TWINT Flag. The TWI will then test the 2-wire Serial Bus and generate a START condition as soon as the bus becomes free. After a START condition has been transmitted, the TWINT Flag is set by hardware, and the status code in TWSR will be $0 x 08$ (see Table 24-2 on page 249). In order to enter MT mode, SLA+W must be transmitted. This is done by writing SLA+W to TWDR. Thereafter the TWINT bit should be cleared (by writing it to one) to continue the transfer. This is accomplished by writing the following value to TWCR:

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 1 | X | 0 | 0 | X | 1 | 0 | X |

When SLA+W have been transmitted and an acknowledgement bit has been received, TWINT is set again and a number of status codes in TWSR are possible. Possible status codes in Master mode are $0 \times 18,0 \times 20$, or $0 \times 38$. The appropriate action to be taken for each of these status codes is detailed in Table 24-2 on page 249.

When SLA+W has been successfully transmitted, a data packet should be transmitted. This is done by writing the data byte to TWDR. TWDR must only be written when TWINT is high. If not, the access will be discarded, and the Write Collision bit (TWWC) will be set in the TWCR Register. After updating TWDR, the TWINT bit should be cleared (by writing it to one) to continue the transfer. This is accomplished by writing the following value to TWCR:

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 1 | X | 0 | 0 | X | 1 | 0 | X |

This scheme is repeated until the last byte has been sent and the transfer is ended by generating a STOP condition or a repeated START condition. A STOP condition is generated by writing the following value to TWCR:

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 1 | X | 0 | 1 | X | 1 | 0 | X |

A REPEATED START condition is generated by writing the following value to TWCR:
TWCR
value

| TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | X | 1 | 0 | X | 1 | 0 | X |

After a repeated START condition (state $0 \times 10$ ) the 2 -wire Serial Interface can access the same Slave again, or a new Slave without transmitting a STOP condition. Repeated START enables the Master to switch between Slaves, Master Transmitter mode and Master Receiver mode without losing control of the bus.

Table 24-2. Status codes for Master Transmitter Mode

| Status Code (TWSR) Prescaler Bits are 0 | Status of the 2-wire Serial Bus and 2-wire Serial Interface Hardware | Application Software Response |  |  |  |  | Next Action Taken by TWI Hardware |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | To/from TWDR | To TWCR |  |  |  |  |
|  |  |  | STA | STO | $\begin{gathered} \text { TWIN } \\ \mathrm{T} \end{gathered}$ | $\begin{gathered} \text { TWE } \\ \text { A } \\ \hline \end{gathered}$ |  |
| 0x08 | A START condition has been transmitted | Load SLA+W | 0 | 0 | 1 | X | SLA +W will be transmitted; ACK or NOT ACK will be received |
| 0x10 | A repeated START condition has been transmitted | Load SLA+W or Load SLA+R | $\begin{aligned} & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & x \\ & x \end{aligned}$ | SLA+W will be transmitted; ACK or NOT ACK will be received SLA+R will be transmitted; Logic will switch to Master Receiver mode |
| 0x18 | SLA+W has been transmitted; ACK has been received | Load data byte or <br> No TWDR action or No TWDR action or <br> No TWDR action | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | Data byte will be transmitted and ACK or NOT ACK will be received <br> Repeated START will be transmitted <br> STOP condition will be transmitted and <br> TWSTO Flag will be reset <br> STOP condition followed by a START condition will be transmitted and TWSTO Flag will be reset |
| 0x20 | SLA+W has been transmitted; NOT ACK has been received | Load data byte or <br> No TWDR action or No TWDR action or <br> No TWDR action | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | Data byte will be transmitted and ACK or NOT ACK will be received <br> Repeated START will be transmitted <br> STOP condition will be transmitted and TWSTO Flag will be reset <br> STOP condition followed by a START condition will be transmitted and TWSTO Flag will be reset |
| 0x28 | Data byte has been transmitted; <br> ACK has been received | Load data byte or <br> No TWDR action or No TWDR action or <br> No TWDR action | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | Data byte will be transmitted and ACK or NOT ACK will be received <br> Repeated START will be transmitted <br> STOP condition will be transmitted and <br> TWSTO Flag will be reset <br> STOP condition followed by a START condition will be transmitted and TWSTO Flag will be reset |
| 0x30 | Data byte has been transmitted; <br> NOT ACK has been received | Load data byte or <br> No TWDR action or No TWDR action or <br> No TWDR action | $\begin{aligned} & 0 \\ & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & x \\ & x \\ & x \\ & x \end{aligned}$ | Data byte will be transmitted and ACK or NOT ACK will be received <br> Repeated START will be transmitted <br> STOP condition will be transmitted and <br> TWSTO Flag will be reset <br> STOP condition followed by a START condition will be transmitted and TWSTO Flag will be reset |
| 0x38 | Arbitration lost in SLA+W or data bytes | No TWDR action or No TWDR action | 0 1 | 0 0 | 1 1 | X $\times$ | 2-wire Serial Bus will be released and not addressed Slave mode entered <br> A START condition will be transmitted when the bus becomes free |

Figure 24-12. Formats and States in the Master Transmitter Mode


### 24.7.2 Master Receiver Mode

In the Master Receiver mode, a number of data bytes are received from a Slave Transmitter (see Figure 24-13 on page 251). In order to enter a Master mode, a START condition must be transmitted. The format of the following address packet determines whether Master Transmitter or Master Receiver mode is to be entered. If SLA+W is transmitted, MT mode is entered, if SLA+R is transmitted, MR mode is entered. All the status codes mentioned in this section assume that the prescaler bits are zero or are masked to zero.

Figure 24-13. Data Transfer in Master Receiver Mode


A START condition is sent by writing the following value to TWCR:

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 1 | X | 1 | 0 | X | 1 | 0 | X |

TWEN must be written to one to enable the 2-wire Serial Interface, TWSTA must be written to one to transmit a START condition and TWINT must be set to clear the TWINT Flag. The TWI will then test the 2 -wire Serial Bus and generate a START condition as soon as the bus becomes free. After a START condition has been transmitted, the TWINT Flag is set by hardware, and the status code in TWSR will be 0x08 (see Table 24-2 on page 249). In order to enter MR mode, SLA+R must be transmitted. This is done by writing SLA+R to TWDR. Thereafter the TWINT bit should be cleared (by writing it to one) to continue the transfer. This is accomplished by writing the following value to TWCR:

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 1 | X | 0 | 0 | X | 1 | 0 | X |

When SLA+R have been transmitted and an acknowledgement bit has been received, TWINT is set again and a number of status codes in TWSR are possible. Possible status codes in Master mode are $0 \times 38,0 \times 40$, or $0 \times 48$. The appropriate action to be taken for each of these status codes is detailed in Table 24-3 on page 252. Received data can be read from the TWDR Register when the TWINT Flag is set high by hardware. This scheme is repeated until the last byte has been received. After the last byte has been received, the MR should inform the ST by sending a NACK after the last received data byte. The transfer is ended by generating a STOP condition or a repeated START condition. A STOP condition is generated by writing the following value to TWCR:

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 1 | X | 0 | 1 | X | 1 | 0 | X |

A REPEATED START condition is generated by writing the following value to TWCR:
TWCR
value

| TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | X | 1 | 0 | x | 1 | 0 | x |

After a repeated START condition (state 0x10) the 2-wire Serial Interface can access the same Slave again, or a new Slave without transmitting a STOP condition. Repeated START enables the Master to switch between Slaves, Master Transmitter mode and Master Receiver mode without losing control over the bus.

Table 24-3. Status codes for Master Receiver Mode

| Status Code (TWSR) Prescaler Bits are 0 | Status of the 2-wire Serial Bus and 2-wire Serial Interface Hardware | Application Software Response |  |  |  |  | Next Action Taken by TWI Hardware |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | To/from TWDR | To TWCR |  |  |  |  |
|  |  |  | STA | STO | $\begin{gathered} \text { TWIN } \\ \mathrm{T} \\ \hline \end{gathered}$ | $\begin{gathered} \text { TWE } \\ \text { A } \\ \hline \end{gathered}$ |  |
| 0x08 | A START condition has been transmitted | Load SLA+R | 0 | 0 | 1 | X | SLA+R will be transmitted ACK or NOT ACK will be received |
| 0x10 | A repeated START condition has been transmitted | Load SLA+R or Load SLA+W | $\begin{aligned} & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | SLA+R will be transmitted ACK or NOT ACK will be received SLA+W will be transmitted Logic will switch to Master Transmitter mode |
| $0 \times 38$ | Arbitration lost in SLA+R or NOT ACK bit | No TWDR action or No TWDR action | $0$ $1$ | $\begin{aligned} & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | 2-wire Serial Bus will be released and not addressed Slave mode will be entered A START condition will be transmitted when the bus becomes free |
| 0x40 | SLA+R has been transmitted; ACK has been received | No TWDR action or No TWDR action | 0 <br> 0 | $0$ $0$ | 1 <br> 1 | $0$ $1$ | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| $0 \times 48$ | SLA+R has been transmitted; NOT ACK has been received | No TWDR action or No TWDR action or <br> No TWDR action | $\begin{aligned} & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | Repeated START will be transmitted STOP condition will be transmitted and TWSTO Flag will be reset <br> STOP condition followed by a START condition will be transmitted and TWSTO Flag will be reset |
| 0x50 | Data byte has been received; ACK has been returned | Read data byte or Read data byte | $\begin{aligned} & 0 \\ & 0 \\ & \hline \end{aligned}$ | $0$ | $\begin{aligned} & 1 \\ & 1 \\ & \hline \end{aligned}$ | $0$ <br> 1 | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| $0 \times 58$ | Data byte has been received; NOT ACK has been returned | Read data byte or Read data byte or <br> Read data byte | $\begin{aligned} & 1 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | Repeated START will be transmitted STOP condition will be transmitted and TWSTO Flag will be reset <br> STOP condition followed by a START condition will be transmitted and TWSTO Flag will be reset |

Figure 24-14. Formats and States in the Master Receiver Mode


### 24.7.3 Slave Receiver Mode

In the Slave Receiver mode, a number of data bytes are received from a Master Transmitter (see Figure 24-15). All the status codes mentioned in this section assume that the prescaler bits are zero or are masked to zero.

Figure 24-15. Data transfer in Slave Receiver mode


To initiate the Slave Receiver mode, TWAR and TWCR must be initialized as follows:
TWAR
value

| TWA6 | TWA5 | TWA4 | TWA3 | TWA2 | TWA1 | TWA0 | TWGCE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Device's Own Slave Address |  |  |  |  |  |  |  |

The upper seven bits are the address to which the 2-wire Serial Interface will respond when addressed by a Master. If the LSB is set, the TWI will respond to the general call address ( $0 \times 00$ ), otherwise it will ignore the general call address.

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 0 | 1 | 0 | 0 | 0 | 1 | 0 | X |

TWEN must be written to one to enable the TWI. The TWEA bit must be written to one to enable the acknowledgement of the device's own slave address or the general call address. TWSTA and TWSTO must be written to zero.

When TWAR and TWCR have been initialized, the TWI waits until it is addressed by its own slave address (or the general call address if enabled) followed by the data direction bit. If the direction bit is " 0 " (write), the TWI will operate in SR mode, otherwise ST mode is entered. After its own slave address and the write bit have been received, the TWINT Flag is set and a valid status code can be read from TWSR. The status code is used to determine the appropriate software action. The appropriate action to be taken for each status code is detailed in Table 24-4 on page 255. The Slave Receiver mode may also be entered if arbitration is lost while the TWI is in the Master mode (see states $0 \times 68$ and $0 \times 78$ ).

If the TWEA bit is reset during a transfer, the TWI will return a "Not Acknowledge" ("1") to SDA after the next received data byte. This can be used to indicate that the Slave is not able to receive any more bytes. While TWEA is zero, the TWI does not acknowledge its own slave address. However, the 2 -wire Serial Bus is still monitored and address recognition may resume at any time by setting TWEA. This implies that the TWEA bit may be used to temporarily isolate the TWI from the 2-wire Serial Bus.

In all sleep modes other than Idle mode, the clock system to the TWI is turned off. If the TWEA bit is set, the interface can still acknowledge its own slave address or the general call address by using the 2-wire Serial Bus clock as a clock source. The part will then wake up from sleep and the TWI will hold the SCL clock low during the wake up and until the TWINT Flag is cleared (by writing it to one). Further data reception will be carried out as normal, with the AVR clocks running as normal. Observe that if the AVR is set up with a long start-up time, the SCL line may be held low for a long time, blocking other data transmissions.

Note that the 2-wire Serial Interface Data Register - TWDR does not reflect the last byte present on the bus when waking up from these Sleep modes.

Table 24-4. Status Codes for Slave Receiver Mode

| Status Code (TWSR) Prescaler Bits are 0 | Status of the 2-wire Serial Bus and 2-wire Serial Interface Hardware | Application Software Response |  |  |  |  | Next Action Taken by TWI Hardware |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | To/from TWDR | To TWCR |  |  |  |  |
|  |  |  | STA | STO | $\begin{gathered} \text { TWIN } \\ \mathrm{T} \\ \hline \end{gathered}$ | $\begin{gathered} \text { TWE } \\ \text { A } \\ \hline \end{gathered}$ |  |
| 0x60 | Own SLA+W has been received; ACK has been returned | No TWDR action or <br> No TWDR action | $\begin{aligned} & x \\ & x \\ & \hline \end{aligned}$ | $0$ $0$ | $\begin{aligned} & 1 \\ & 1 \\ & \hline \end{aligned}$ | 0 $1$ | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| 0x68 | Arbitration lost in SLA+R/W as Master; own SLA+W has been received; ACK has been returned | No TWDR action or No TWDR action | $\begin{gathered} x \\ x \\ \hline \end{gathered}$ | $0$ $0$ | $\begin{aligned} & 1 \\ & 1 \end{aligned}$ | $0$ $1$ | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| 0x70 | General call address has been received; ACK has been returned | No TWDR action or No TWDR action | $\begin{aligned} & \text { X } \\ & \text { x } \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \end{aligned}$ | $1$ $1$ | $0$ $1$ | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| 0x78 | Arbitration lost in SLA+R/W as Master; General call address has been received; ACK has been returned | No TWDR action or No TWDR action | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | $0$ $0$ | $\begin{aligned} & 1 \\ & 1 \end{aligned}$ | $0$ $1$ | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| 0x80 | Previously addressed with own SLA+W; data has been received; ACK has been returned | Read data byte or <br> Read data byte | $\begin{aligned} & x \\ & x \\ & \hline \end{aligned}$ | $0$ $0$ | $\begin{aligned} & 1 \\ & 1 \\ & \hline \end{aligned}$ | 0 1 | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| 0x88 | Previously addressed with own SLA+W; data has been received; NOT ACK has been returned | Read data byte or Read data byte or <br> Read data byte or <br> Read data byte | 0 <br> 0 <br> 1 <br> 1 | 0 <br> 0 <br> 0 <br> 0 | 1 <br> 1 <br> 1 <br> 1 | 0 <br> 1 <br> 0 <br> 1 | Switched to the not addressed Slave mode; no recognition of own SLA or GCA Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1" Switched to the not addressed Slave mode; no recognition of own SLA or GCA; a START condition will be transmitted when the bus becomes free <br> Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = " 1 "; a START condition will be transmitted when the bus becomes free |
| 0x90 | Previously addressed with general call; data has been received; ACK has been returned | Read data byte or <br> Read data byte | $\begin{aligned} & x \\ & x \\ & \hline \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \\ & \hline \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \\ & \hline \end{aligned}$ | $0$ $1$ | Data byte will be received and NOT ACK will be returned <br> Data byte will be received and ACK will be returned |
| 0x98 | Previously addressed with general call; data has been received; NOT ACK has been returned | Read data byte or <br> Read data byte or <br> Read data byte or <br> Read data byte | 0 0 1 1 1 | 0 0 0 | 1 1 1 1 1 | 1 <br> 0 <br> 1 | Switched to the not addressed Slave mode; no recognition of own SLA or GCA Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1" Switched to the not addressed Slave mode; no recognition of own SLA or GCA; a START condition will be transmitted when the bus becomes free <br> Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = " 1 "; a START condition will be transmitted when the bus becomes free |
| 0xA0 | A STOP condition or repeated START condition has been received while still addressed as Slave | No action | 0 <br> 0 <br> 1 <br> 1 | 0 <br> 0 <br> 0 <br> 0 | 1 <br> 1 | 1 <br> 0 <br> 1 | Switched to the not addressed Slave mode; no recognition of own SLA or GCA Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1" Switched to the not addressed Slave mode; no recognition of own SLA or GCA; a START condition will be transmitted when the bus becomes free <br> Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = " 1 "; a START condition will be transmitted when the bus becomes free |

Figure 24-16. Formats and States in the Slave Receiver Mode

$\square$ From master to slave


Any number of data bytes and their associated acknowledge bits

From slave to master


This number (contained in TWSR) corresponds to a defined state of the Two-Wire Serial Bus. The prescaler bits are zero or masked to zero

### 24.7.4 Slave Transmitter Mode

In the Slave Transmitter mode, a number of data bytes are transmitted to a Master Receiver (see Figure 24-17). All the status codes mentioned in this section assume that the prescaler bits are zero or are masked to zero.

Figure 24-17. Data Transfer in Slave Transmitter Mode


To initiate the Slave Transmitter mode, TWAR and TWCR must be initialized as follows:

| TWAR <br> value | TWA6 | TWA5 | TWA4 | TWA3 | TWA2 | TWA1 | TWA0 | TWGCE |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Device's Own Slave Address |  |  |  |  |  |  |  |  |

The upper seven bits are the address to which the 2 -wire Serial Interface will respond when addressed by a Master. If the LSB is set, the TWI will respond to the general call address ( $0 \times 00$ ), otherwise it will ignore the general call address.

| TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| value | 0 | 1 | 0 | 0 | 0 | 1 | 0 | X |

TWEN must be written to one to enable the TWI. The TWEA bit must be written to one to enable the acknowledgement of the device's own slave address or the general call address. TWSTA and TWSTO must be written to zero.
When TWAR and TWCR have been initialized, the TWI waits until it is addressed by its own slave address (or the general call address if enabled) followed by the data direction bit. If the direction bit is "1" (read), the TWI will operate in ST mode, otherwise SR mode is entered. After its own slave address and the write bit have been received, the TWINT Flag is set and a valid status code can be read from TWSR. The status code is used to determine the appropriate software action. The appropriate action to be taken for each status code is detailed in Table 24-5. The Slave Transmitter mode may also be entered if arbitration is lost while the TWI is in the Master mode (see state $0 \times B 0$ ).

If the TWEA bit is written to zero during a transfer, the TWI will transmit the last byte of the transfer. State 0xC0 or state 0xC8 will be entered, depending on whether the Master Receiver transmits a NACK or ACK after the final byte. The TWI is switched to the not addressed Slave mode, and will ignore the Master if it continues the transfer. Thus the Master Receiver receives all " 1 " as serial data. State $0 \times C 8$ is entered if the Master demands additional data bytes (by transmitting ACK), even though the Slave has transmitted the last byte (TWEA zero and expecting NACK from the Master).
While TWEA is zero, the TWI does not respond to its own slave address. However, the 2-wire Serial Bus is still monitored and address recognition may resume at any time by setting TWEA. This implies that the TWEA bit may be used to temporarily isolate the TWI from the 2-wire Serial Bus.

In all sleep modes other than Idle mode, the clock system to the TWI is turned off. If the TWEA bit is set, the interface can still acknowledge its own slave address or the general call address by using the 2 -wire Serial Bus clock as a clock source. The part will then wake up from sleep and the TWI will hold the SCL clock will low during the wake up and until the TWINT Flag is cleared (by writing it to one). Further data transmission will be carried out as normal, with the AVR clocks running as normal. Observe that if the AVR is set up with a long start-up time, the SCL line may be held low for a long time, blocking other data transmissions.

Note that the 2-wire Serial Interface Data Register - TWDR does not reflect the last byte present on the bus when waking up from these sleep modes.

Table 24-5. Status Codes for Slave Transmitter Mode

| Status Code <br> (TWSR) <br> Prescaler <br> Bits <br> are 0 | Status of the 2-wire Serial Bus and 2-wire Serial Interface Hardware | Application Software Response |  |  |  |  | Next Action Taken by TWI Hardware |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | To/from TWDR | To TWCR |  |  |  |  |
|  |  |  | STA | STO | $\begin{gathered} \text { TWIN } \\ \mathrm{T} \end{gathered}$ | $\begin{gathered} \text { TWE } \\ \text { A } \end{gathered}$ |  |
| 0xA8 | Own SLA+R has been received; ACK has been returned | Load data byte or Load data byte | $\begin{aligned} & \mathrm{x} \\ & \mathrm{x} \end{aligned}$ | $\begin{aligned} & 0 \\ & 0 \end{aligned}$ | $\begin{aligned} & 1 \\ & 1 \end{aligned}$ | 0 <br> 1 | Last data byte will be transmitted and NOT ACK should be received Data byte will be transmitted and ACK should be received |
| 0xB0 | Arbitration lost in SLA+R/W as Master; own SLA+R has been received; ACK has been returned | Load data byte or Load data byte | $\begin{aligned} & x \\ & x \end{aligned}$ | 0 0 | 1 1 | 0 1 | Last data byte will be transmitted and NOT ACK should be received <br> Data byte will be transmitted and ACK should be received |

Table 24-5. $\quad$ Status Codes for Slave Transmitter Mode (Continued)

| 0xB8 | Data byte in TWDR has been transmitted; ACK has been received | Load data byte or Load data byte | X $\times$ | 0 0 | 1 1 | 0 1 | Last data byte will be transmitted and NOT ACK should be received Data byte will be transmitted and ACK should be received |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0xC0 | Data byte in TWDR has been transmitted; NOT ACK has been received | No TWDR action or No TWDR action or <br> No TWDR action or <br> No TWDR action | 0 0 1 1 | 0 0 0 | 1 1 1 1 | 0 1 0 | Switched to the not addressed Slave mode; no recognition of own SLA or GCA <br> Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1" Switched to the not addressed Slave mode; no recognition of own SLA or GCA; <br> a START condition will be transmitted when the bus becomes free <br> Switched to the not addressed Slave mode; own SLA will be recognized; <br> GCA will be recognized if TWGCE = " 1 "; <br> a START condition will be transmitted when the bus becomes free |
| 0xC8 | Last data byte in TWDR has been transmitted (TWEA = "0"); ACK has been received | No TWDR action or No TWDR action or <br> No TWDR action or <br> No TWDR action | 0 0 1 | 0 0 0 | 1 1 1 | 0 1 0 | Switched to the not addressed Slave mode; no recognition of own SLA or GCA Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = " 1 " Switched to the not addressed Slave mode; no recognition of own SLA or GCA; a START condition will be transmitted when the bus becomes free <br> Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = " 1 "; a START condition will be transmitted when the bus becomes free |

Figure 24-18. Formats and States in the Slave Transmitter Mode


### 24.7.5 Miscellaneous States

There are two status codes that do not correspond to a defined TWI state, see Table 24-6 on page 259.
Status 0xF8 indicates that no relevant information is available because the TWINT Flag is not set. This occurs between other states, and when the TWI is not involved in a serial transfer.

Status 0x00 indicates that a bus error has occurred during a 2 -wire Serial Bus transfer. A bus error occurs when a START or STOP condition occurs at an illegal position in the format frame. Examples of such illegal positions are during the serial transfer of an address byte, a data byte, or an acknowledge bit. When a bus error occurs, TWINT is set. To recover from a bus error, the TWSTO Flag must set and TWINT must be cleared by writing a logic one to
it. This causes the TWI to enter the not addressed Slave mode and to clear the TWSTO Flag (no other bits in TWCR are affected). The SDA and SCL lines are released, and no STOP condition is transmitted.

Table 24-6. Miscellaneous States

| Status Code (TWSR) Prescaler Bits are 0 | Status of the 2-wire Serial Bus and 2-wire Serial Interface Hardware | Application Software Response |  |  |  |  | Next Action Taken by TWI Hardware |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | To/from TWDR | To TWCR |  |  |  |  |
|  |  |  | STA | STO | $\begin{gathered} \text { TWIN } \\ \mathrm{T} \\ \hline \end{gathered}$ | TWE |  |
| 0xF8 | No relevant state information available; TWINT = "0" | No TWDR action |  | No TW | R action |  | Wait or proceed current transfer |
| 0x00 | Bus error due to an illegal START or STOP condition | No TWDR action | 0 | 1 | 1 | X | Only the internal hardware is affected, no STOP condition is sent on the bus. In all cases, the bus is released and TWSTO is cleared. |

### 24.7.6 Combining Several TWI Modes

In some cases, several TWI modes must be combined in order to complete the desired action. Consider for example reading data from a serial EEPROM. Typically, such a transfer involves the following steps:

1. The transfer must be initiated.
2. The EEPROM must be instructed what location should be read.
3. The reading must be performed.
4. The transfer must be finished.

Note that data is transmitted both from Master to Slave and vice versa. The Master must instruct the Slave what location it wants to read, requiring the use of the MT mode. Subsequently, data must be read from the Slave, implying the use of the MR mode. Thus, the transfer direction must be changed. The Master must keep control of the bus during all these steps, and the steps should be carried out as an atomical operation. If this principle is violated in a multimaster system, another Master can alter the data pointer in the EEPROM between steps 2 and 3, and the Master will read the wrong data location. Such a change in transfer direction is accomplished by transmitting a REPEATED START between the transmission of the address byte and reception of the data. After a REPEATED START, the Master keeps ownership of the bus. The following figure shows the flow in this transfer.

Figure 24-19. Combining Several TWI Modes to Access a Serial EEPROM


### 24.8 Multi-master Systems and Arbitration

If multiple masters are connected to the same bus, transmissions may be initiated simultaneously by one or more of them. The TWI standard ensures that such situations are handled in such a way that one of the masters will be allowed to proceed with the transfer, and that no data will be lost in the process. An example of an arbitration situation is depicted below, where two masters are trying to transmit data to a Slave Receiver.

Figure 24-20. An Arbitration Example


Several different scenarios may arise during arbitration, as described below:

- Two or more masters are performing identical communication with the same Slave. In this case, neither the Slave nor any of the masters will know about the bus contention.
- Two or more masters are accessing the same Slave with different data or direction bit. In this case, arbitration will occur, either in the READ/WRITE bit or in the data bits. The masters trying to output a one on SDA while another Master outputs a zero will lose the arbitration. Losing masters will switch to not addressed Slave mode or wait until the bus is free and transmit a new START condition, depending on application software action.
- Two or more masters are accessing different slaves. In this case, arbitration will occur in the SLA bits. Masters trying to output a one on SDA while another Master outputs a zero will lose the arbitration. Masters losing arbitration in SLA will switch to Slave mode to check if they are being addressed by the winning Master. If addressed, they will switch to SR or ST mode, depending on the value of the READ/WRITE bit. If they are not being addressed, they will switch to not addressed Slave mode or wait until the bus is free and transmit a new START condition, depending on application software action.
This is summarized in Figure 24-21. Possible status values are given in circles.
Figure 24-21. Possible Status Codes Caused by Arbitration



### 24.9 Register Description

### 24.9.1 TWBR - TWI Bit Rate Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TWBR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xB8) | TWBR7 | TWBR6 | TWBR5 | TWBR4 | TWBR3 | TWBR2 | TWBR1 | TWBR0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7:0 - TWI Bit Rate Register

TWBR selects the division factor for the bit rate generator. The bit rate generator is a frequency divider which generates the SCL clock frequency in the Master modes. See "Bit Rate Generator Unit" on page 242 for calculating bit rates.

### 24.9.2 TWCR - TWI Control Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TWCR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xBC) | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE |  |
| Read/Write | R/W | R/W | R/W | R/W | R | R/W | R | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

The TWCR is used to control the operation of the TWI. It is used to enable the TWI, to initiate a Master access by applying a START condition to the bus, to generate a Receiver acknowledge, to generate a stop condition, and to control halting of the bus while the data to be written to the bus are written to the TWDR. It also indicates a write collision if data is attempted written to TWDR while the register is inaccessible.

## - Bit 7 - TWINT: TWI Interrupt Flag

This bit is set by hardware when the TWI has finished its current job and expects application software response. If the I-bit in SREG and TWIE in TWCR are set, the MCU will jump to the TWI Interrupt Vector. While the TWINT Flag is set, the SCL low period is stretched. The TWINT Flag must be cleared by software by writing a logic one to it. Note that this flag is not automatically cleared by hardware when executing the interrupt routine. Also note that clearing this flag starts the operation of the TWI, so all accesses to the TWI Address Register (TWAR), TWI Status Register (TWSR), and TWI Data Register (TWDR) must be complete before clearing this flag.

## - Bit 6 - TWEA: TWI Enable Acknowledge Bit

The TWEA bit controls the generation of the acknowledge pulse. If the TWEA bit is written to one, the ACK pulse is generated on the TWI bus if the following conditions are met:

1. The device's own slave address has been received.
2. A general call has been received, while the TWGCE bit in the TWAR is set.
3. A data byte has been received in Master Receiver or Slave Receiver mode.

By writing the TWEA bit to zero, the device can be virtually disconnected from the 2 -wire Serial Bus temporarily. Address recognition can then be resumed by writing the TWEA bit to one again.

## - Bit 5 - TWSTA: TWI START Condition Bit

The application writes the TWSTA bit to one when it desires to become a Master on the 2-wire Serial Bus. The TWI hardware checks if the bus is available, and generates a START condition on the bus if it is free. However, if the bus is not free, the TWI waits until a STOP condition is detected, and then generates a new START condition to claim the bus Master status. TWSTA must be cleared by software when the START condition has been transmitted.

## - Bit 4 - TWSTO: TWI STOP Condition Bit

Writing the TWSTO bit to one in Master mode will generate a STOP condition on the 2-wire Serial Bus. When the STOP condition is executed on the bus, the TWSTO bit is cleared automatically. In Slave mode, setting the TWSTO bit can be used to recover from an error condition. This will not generate a STOP condition, but the TWI returns to a well-defined unaddressed Slave mode and releases the SCL and SDA lines to a high impedance state.

## - Bit 3 - TWWC: TWI Write Collision Flag

The TWWC bit is set when attempting to write to the TWI Data Register - TWDR when TWINT is low. This flag is cleared by writing the TWDR Register when TWINT is high.

## - Bit 2 - TWEN: TWI Enable Bit

The TWEN bit enables TWI operation and activates the TWI interface. When TWEN is written to one, the TWI takes control over the I/O pins connected to the SCL and SDA pins, enabling the slew-rate limiters and spike filters. If this bit is written to zero, the TWI is switched off and all TWI transmissions are terminated, regardless of any ongoing operation.

## - Bit 1 - Res: Reserved Bit

This bit is a reserved bit and will always read as zero.

## - Bit 0 - TWIE: TWI Interrupt Enable

When this bit is written to one, and the I-bit in SREG is set, the TWI interrupt request will be activated for as long as the TWINT Flag is high.

### 24.9.3 TWSR - TWI Status Register



## - Bits 7:3 - TWS: TWI Status

These five bits reflect the status of the TWI logic and the 2 -wire Serial Bus. The different status codes are described later in this section. Note that the value read from TWSR contains both the 5-bit status value and the 2bit prescaler value. The application designer should mask the prescaler bits to zero when checking the Status bits. This makes status checking independent of prescaler setting. This approach is used in this datasheet, unless otherwise noted.

- Bit 2 - Res: Reserved Bit

This bit is reserved and will always read as zero.

- Bits 1:0-TWPS: TWI Prescaler Bits

These bits can be read and written, and control the bit rate prescaler.
Table 24-7. TWI Bit Rate Prescaler

| TWPS1 | TWPS0 | Prescaler Value |
| :--- | :--- | :--- |
| 0 | 0 | 1 |
| 0 | 1 | 4 |
| 1 | 0 | 16 |
| 1 | 1 | 64 |

To calculate bit rates, see "Bit Rate Generator Unit" on page 242. The value of TWPS1:0 is used in the equation.

### 24.9.4 TWDR - TWI Data Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TWDR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xBB) | TWD7 | TWD6 | TWD5 | TWD4 | TWD3 | TWD2 | TWD1 | TWD0 |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |  |

In Transmit mode, TWDR contains the next byte to be transmitted. In Receive mode, the TWDR contains the last byte received. It is writable while the TWI is not in the process of shifting a byte. This occurs when the TWI Interrupt Flag (TWINT) is set by hardware. Note that the Data Register cannot be initialized by the user before the first interrupt occurs. The data in TWDR remains stable as long as TWINT is set. While data is shifted out, data on the bus is simultaneously shifted in. TWDR always contains the last byte present on the bus, except after a wake up from a sleep mode by the TWI interrupt. In this case, the contents of TWDR is undefined. In the case of a lost bus arbitration, no data is lost in the transition from Master to Slave. Handling of the ACK bit is controlled automatically by the TWI logic, the CPU cannot access the ACK bit directly.

## - Bits 7:0 - TWD: TWI Data Register

These eight bits constitute the next data byte to be transmitted, or the latest data byte received on the 2 -wire Serial Bus.

### 24.9.5 TWAR - TWI (Slave) Address Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TWAR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xBA) | TWA6 | TWA5 | TWA4 | TWA3 | TWA2 | TWA1 | TWA0 | TWGCE |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 |  |

The TWAR should be loaded with the 7-bit Slave address (in the seven most significant bits of TWAR) to which the TWI will respond when programmed as a Slave Transmitter or Receiver, and not needed in the Master modes. In multimaster systems, TWAR must be set in masters which can be addressed as Slaves by other Masters.

The LSB of TWAR is used to enable recognition of the general call address ( $0 \times 00$ ). There is an associated address comparator that looks for the slave address (or general call address if enabled) in the received serial address. If a match is found, an interrupt request is generated.

## - Bits 7:1 - TWA: TWI (Slave) Address Register

These seven bits constitute the slave address of the TWI unit.

- Bit 0 - TWGCE: TWI General Call Recognition Enable Bit

If set, this bit enables the recognition of a General Call given over the 2-wire Serial Bus.

### 24.9.6 TWAMR - TWI (Slave) Address Mask Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | TWAMR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xBD) | TWAM[6:0] |  |  |  |  |  |  | - |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

- Bits 7:1 - TWAM: TWI Address Mask

The TWAMR can be loaded with a 7 -bit Slave Address mask. Each of the bits in TWAMR can mask (disable) the corresponding address bit in the TWI Address Register (TWAR). If the mask bit is set to one then the address match logic ignores the compare between the incoming address bit and the corresponding bit in TWAR. Figure 2422 shows the address match logic in detail.

Figure 24-22. TWI Address Match Logic, Block Diagram


## - Bit 0 - Res: Reserved Bit

This bit is reserved and will always read as zero.

## 25. AC - Analog Comparator

The Analog Comparator compares the input values on the positive pin AINO and negative pin AIN1. When the voltage on the positive pin AINO is higher than the voltage on the negative pin AIN1, the Analog Comparator output, ACO, is set. The comparator's output can be set to trigger the Timer/Counter1 Input Capture function. In addition, the comparator can trigger a separate interrupt, exclusive to the Analog Comparator. The user can select Interrupt triggering on comparator output rise, fall or toggle. A block diagram of the comparator and its surrounding logic is shown in Figure 25-1.
The Power Reduction ADC bit, PRADC, in "PRRO - Power Reduction Register 0" on page 55 must be disabled by writing a logical zero to be able to use the ADC input MUX.

Figure 25-1. Analog Comparator Block Diagram ${ }^{(2)}$


Note: 1. See Table 25-1.
2. Refer to Figure 1-1 on page 2 and Table 13-5 on page 76 for Analog Comparator pin placement.

### 25.1 Analog Comparator Multiplexed Input

It is possible to select any of the ADC15:0 pins to replace the negative input to the Analog Comparator. The ADC multiplexer is used to select this input, and consequently, the ADC must be switched off to utilize this feature. If the Analog Comparator Multiplexer Enable bit (ACME in ADCSRB) is set and the ADC is switched off (ADEN in ADCSRA is zero), MUX5 and MUX2:0 in ADMUX select the input pin to replace the negative input to the Analog Comparator, as shown in Table 25-1. If ACME is cleared or ADEN is set, AIN1 is applied to the negative input to the Analog Comparator.

Table 25-1. Analog Comparator Mulitiplexed Input

| ACME | ADEN | MUX5 | MUX2:0 | Analog Comparator Negative Input |
| :---: | :---: | :---: | :---: | :---: |
| 0 | x | x | xxx | AIN1 |
| 1 | 1 | x | xxx | AIN1 |
| 1 | 0 | 0 | 000 | ADC0 |
| 1 | 0 | 0 | 001 | ADC1 |
| 1 | 0 | 0 | 010 | ADC2 |
| 1 | 0 | 0 | 011 | ADC3 |

Table 25-1. Analog Comparator Mulitiplexed Input (Continued)

| ACME | ADEN | MUX5 | MUX2:0 | Analog Comparator Negative Input |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | 0 | 100 | ADC4 |
| 1 | 0 | 0 | 101 | ADC5 |
| 1 | 0 | 0 | 110 | ADC6 |
| 1 | 0 | 0 | 111 | ADC7 |
| 1 | 0 | 1 | 000 | ADC8 |
| 1 | 0 | 1 | 001 | ADC9 |
| 1 | 0 | 1 | 010 | ADC10 |
| 1 | 0 | 1 | 011 | ADC111 |
| 1 | 0 | 1 | 100 | ADC12 |
| 1 | 0 | 1 | 101 | ADC13 |
| 1 | 0 | 1 | 110 | ADC14 |
| 1 | 0 | 1 | 111 | ADC15 |

### 25.2 Register Description

### 25.2.1 ADCSRB - ADC Control and Status Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | ADCSRB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x7B) | - | ACME | - | - | MUX5 | ADTS2 | ADTS1 | ADTS0 |  |
| Read/Write | R | R/W | R | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 6 - ACME: Analog Comparator Multiplexer Enable

When this bit is written logic one and the ADC is switched off (ADEN in ADCSRA is zero), the ADC multiplexer selects the negative input to the Analog Comparator. When this bit is written logic zero, AIN1 is applied to the negative input of the Analog Comparator. For a detailed description of this bit, see "Analog Comparator Multiplexed Input" on page 265.

### 25.2.2 ACSR - Analog Comparator Control and Status Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | ACSR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x30 (0x50) | ACD | ACBG | ACO | ACI | ACIE | ACIC | ACIS1 | ACISO |  |
| Read/Write | R/W | R/W | R | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | N/A | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - ACD: Analog Comparator Disable

When this bit is written logic one, the power to the Analog Comparator is switched off. This bit can be set at any time to turn off the Analog Comparator. This will reduce power consumption in Active and Idle mode. When changing the ACD bit, the Analog Comparator Interrupt must be disabled by clearing the ACIE bit in ACSR. Otherwise an interrupt can occur when the bit is changed.

## - Bit 6 - ACBG: Analog Comparator Bandgap Select

When this bit is set, a fixed bandgap reference voltage replaces the positive input to the Analog Comparator. When this bit is cleared, AINO is applied to the positive input of the Analog Comparator. When the bandgap reference is used as input to the Analog Comparator, it will take a certain time for the voltage to stabilize. If not stabilized, the first conversion may give a wrong value. See "Internal Voltage Reference" on page 60.

## - Bit 5 - ACO: Analog Comparator Output

The output of the Analog Comparator is synchronized and then directly connected to ACO. The synchronization introduces a delay of 1-2 clock cycles.

## - Bit 4 - ACI: Analog Comparator Interrupt Flag

This bit is set by hardware when a comparator output event triggers the interrupt mode defined by ACIS1 and ACISO. The Analog Comparator interrupt routine is executed if the ACIE bit is set and the I-bit in SREG is set. ACI is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, ACl is cleared by writing a logic one to the flag.

## - Bit 3 - ACIE: Analog Comparator Interrupt Enable

When the ACIE bit is written logic one and the I-bit in the Status Register is set, the Analog Comparator interrupt is activated. When written logic zero, the interrupt is disabled.

## - Bit 2 - ACIC: Analog Comparator Input Capture Enable

When written logic one, this bit enables the input capture function in Timer/Counter1 to be triggered by the Analog Comparator. The comparator output is in this case directly connected to the input capture front-end logic, making the comparator utilize the noise canceler and edge select features of the Timer/Counter1 Input Capture interrupt. When written logic zero, no connection between the Analog Comparator and the input capture function exists. To make the comparator trigger the Timer/Counter1 Input Capture interrupt, the ICIE1 bit in the Timer Interrupt Mask Register (TIMSK1) must be set.

- Bits 1, 0 - ACIS1, ACIS0: Analog Comparator Interrupt Mode Select

These bits determine which comparator events that trigger the Analog Comparator interrupt. The different settings are shown in Table 25-2.

Table 25-2. ACIS1/ACISO Settings

| ACIS1 | ACISO | Interrupt Mode |
| :---: | :---: | :--- |
| 0 | 0 | Comparator Interrupt on Output Toggle |
| 0 | 1 | Reserved |
| 1 | 0 | Comparator Interrupt on Falling Output Edge |
| 1 | 1 | Comparator Interrupt on Rising Output Edge |

When changing the ACIS1/ACISO bits, the Analog Comparator Interrupt must be disabled by clearing its Interrupt Enable bit in the ACSR Register. Otherwise an interrupt can occur when the bits are changed.

### 25.2.3 DIDR1 - Digital Input Disable Register 1

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | DIDR1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x7F) | - | - | - | - | - | - | AIN1D | AINOD |  |
| Read/Write | R | R | R | R | R | R | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 1, 0 - AIN1D, AINOD: AIN1, AINO Digital Input Disable

When this bit is written logic one, the digital input buffer on the AIN1/0 pin is disabled. The corresponding PIN Register bit will always read as zero when this bit is set. When an analog signal is applied to the AIN $1 / 0$ pin and the digital input from this pin is not needed, this bit should be written logic one to reduce power consumption in the digital input buffer.

## 26. ADC - Analog to Digital Converter

### 26.1 Features

- 10-bit Resolution
- 1 LSB Integral Non-linearity
- $\pm 2$ LSB Absolute Accuracy
- $13 \mu \mathrm{~s}-260 \mu \mathrm{~s}$ Conversion Time
- Up to 76.9kSPS (Up to 15kSPS at Maximum Resolution)
- 16 Multiplexed Single Ended Input Channels
- 14 Differential input channels
- 4 Differential Input Channels with Optional Gain of 10x and 200x
- Optional Left Adjustment for ADC Result Readout
- OV - $\mathrm{V}_{\mathrm{cc}}$ ADC Input Voltage Range
- 2.7V - V ${ }_{\text {cc }}$ Differential ADC Voltage Range
- Selectable 2.56V or 1.1V ADC Reference Voltage
- Free Running or Single Conversion Mode
- Interrupt on ADC Conversion Complete
- Sleep Mode Noise Canceler

The ATmega640/1280/1281/2560/2561 features a 10-bit successive approximation ADC. The ADC is connected to an 8/16-channel Analog Multiplexer which allows eight/sixteen single-ended voltage inputs constructed from the pins of Port F and Port K. The single-ended voltage inputs refer to OV (GND).
The device also supports 16/32 differential voltage input combinations. Four of the differential inputs (ADC1 \& ADC0, ADC3 \& ADC2, ADC9 \& ADC8 and ADC11 \& ADC10) are equipped with a programmable gain stage, providing amplification steps of $0 \mathrm{~dB}(1 \times), 20 \mathrm{~dB}(10 \times)$ or $46 \mathrm{~dB}(200 x)$ on the differential input voltage before the ADC conversion. The 16 channels are split in two sections of 8 channels where in each section seven differential analog input channels share a common negative terminal (ADC1/ADC9), while any other ADC input in that section can be selected as the positive input terminal. If $1 \times$ or $10 \times$ gain is used, 8 bit resolution can be expected. If $200 \times$ gain is used, 7 bit resolution can be expected.

The ADC contains a Sample and Hold circuit which ensures that the input voltage to the ADC is held at a constant level during conversion. A block diagram of the ADC is shown in Figure 26-1 on page 269.

The ADC has a separate analog supply voltage pin, AVCC. $\mathrm{AV}_{\mathrm{cc}}$ must not differ more than $\pm 0.3 \mathrm{~V}$ from $\mathrm{V}_{\mathrm{cc}}$. See the paragraph "ADC Noise Canceler" on page 275 on how to connect this pin.

Internal reference voltages of nominally $1.1 \mathrm{~V}, 2.56 \mathrm{~V}$ or AVCC are provided On-chip. The voltage reference may be externally decoupled at the AREF pin by a capacitor for better noise performance.
The Power Reduction ADC bit, PRADC, in "PRR0 - Power Reduction Register 0" on page 55 must be disabled by writing a logical zero to enable the ADC.

Figure 26-1. Analog to Digital Converter Block Schematic


### 26.2 Operation

The ADC converts an analog input voltage to a 10 -bit digital value through successive approximation. The minimum value represents GND and the maximum value represents the voltage on the AREF pin minus 1 LSB. Optionally, AVCC or an internal 1.1 V or 2.56 V reference voltage may be connected to the AREF pin by writing to the REFSn bits in the ADMUX Register. The internal voltage reference may thus be decoupled by an external capacitor at the AREF pin to improve noise immunity.

The analog input channel is selected by writing to the MUX bits in ADMUX and ADCSRB. Any of the ADC input pins, as well as GND and a fixed bandgap voltage reference, can be selected as single ended inputs to the ADC. A selection of ADC input pins can be selected as positive and negative inputs to the differential amplifier.

If differential channels are selected, the voltage difference between the selected input channel pair then becomes the analog input to the ADC. If single ended channels are used, the amplifier is bypassed altogether.

The ADC is enabled by setting the ADC Enable bit, ADEN in ADCSRA. Voltage reference and input channel selections will not go into effect until ADEN is set. The ADC does not consume power when ADEN is cleared, so it is recommended to switch off the ADC before entering power saving sleep modes.

The ADC generates a 10-bit result which is presented in the ADC Data Registers, ADCH and ADCL. By default, the result is presented right adjusted, but can optionally be presented left adjusted by setting the ADLAR bit in ADMUX.

If the result is left adjusted and no more than 8-bit precision is required, it is sufficient to read ADCH. Otherwise, ADCL must be read first, then ADCH, to ensure that the content of the Data Registers belongs to the same conversion. Once ADCL is read, ADC access to Data Registers is blocked. This means that if ADCL has been read, and a conversion completes before ADCH is read, neither register is updated and the result from the conversion is lost. When ADCH is read, ADC access to the ADCH and ADCL Registers is re-enabled.

The ADC has its own interrupt which can be triggered when a conversion completes. When ADC access to the Data Registers is prohibited between reading of ADCH and ADCL, the interrupt will trigger even if the result is lost.

### 26.3 Starting a Conversion

A single conversion is started by writing a logical one to the ADC Start Conversion bit, ADSC. This bit stays high as long as the conversion is in progress and will be cleared by hardware when the conversion is completed. If a different data channel is selected while a conversion is in progress, the ADC will finish the current conversion before performing the channel change.
Alternatively, a conversion can be triggered automatically by various sources. Auto Triggering is enabled by setting the ADC Auto Trigger Enable bit, ADATE in ADCSRA. The trigger source is selected by setting the ADC Trigger Select bits, ADTS in ADCSRB (see description of the ADTS bits for a list of the trigger sources). When a positive edge occurs on the selected trigger signal, the ADC prescaler is reset and a conversion is started. This provides a method of starting conversions at fixed intervals. If the trigger signal still is set when the conversion completes, a new conversion will not be started. If another positive edge occurs on the trigger signal during conversion, the edge will be ignored. Note that an Interrupt Flag will be set even if the specific interrupt is disabled or the Global Interrupt Enable bit in SREG is cleared. A conversion can thus be triggered without causing an interrupt. However, the Interrupt Flag must be cleared in order to trigger a new conversion at the next interrupt event.

Figure 26-2. ADC Auto Trigger Logic


Using the ADC Interrupt Flag as a trigger source makes the ADC start a new conversion as soon as the ongoing conversion has finished. The ADC then operates in Free Running mode, constantly sampling and updating the ADC Data Register. The first conversion must be started by writing a logical one to the ADSC bit in ADCSRA. In
this mode the ADC will perform successive conversions independently of whether the ADC Interrupt Flag, ADIF is cleared or not.

If Auto Triggering is enabled, single conversions can be started by writing ADSC in ADCSRA to one. ADSC can also be used to determine if a conversion is in progress. The ADSC bit will be read as one during a conversion, independently of how the conversion was started.

### 26.4 Prescaling and Conversion Timing

Figure 26-3. ADC Prescaler


By default, the successive approximation circuitry requires an input clock frequency between 50 kHz and 200 kHz . If a lower resolution than 10 bits is needed, the input clock frequency to the ADC can be as high as 1000 kHz to get a higher sample rate.

The ADC module contains a prescaler, which generates an acceptable ADC clock frequency from any CPU frequency above 100 kHz . The prescaling is set by the ADPS bits in ADCSRA. The prescaler starts counting from the moment the ADC is switched on by setting the ADEN bit in ADCSRA. The prescaler keeps running for as long as the ADEN bit is set, and is continuously reset when ADEN is low.

When initiating a single ended conversion by setting the ADSC bit in ADCSRA, the conversion starts at the following rising edge of the ADC clock cycle.
A normal conversion takes 13 ADC clock cycles. The first conversion after the ADC is switched on (ADEN in ADCSRA is set) takes 25 ADC clock cycles in order to initialize the analog circuitry.
When the bandgap reference voltage is used as input to the ADC, it will take a certain time for the voltage to stabilize. If not stabilized, the first value read after the first conversion may be wrong.
The actual sample-and-hold takes place 1.5 ADC clock cycles after the start of a normal conversion and 13.5 ADC clock cycles after the start of an first conversion. When a conversion is complete, the result is written to the ADC Data Registers, and ADIF is set. In Single Conversion mode, ADSC is cleared simultaneously. The software may then set ADSC again, and a new conversion will be initiated on the first rising ADC clock edge.

When Auto Triggering is used, the prescaler is reset when the trigger event occurs. This assures a fixed delay from the trigger event to the start of conversion. In this mode, the sample-and-hold takes place two ADC clock cycles after the rising edge on the trigger source signal. Three additional CPU clock cycles are used for synchronization logic.

In Free Running mode, a new conversion will be started immediately after the conversion completes, while ADSC remains high. For a summary of conversion times, see Table 26-1 on page 273.

Figure 26-4. ADC Timing Diagram, First Conversion (Single Conversion Mode)


Figure 26-5. ADC Timing Diagram, Single Conversion


Figure 26-6. ADC Timing Diagram, Auto Triggered Conversion


Figure 26-7. ADC Timing Diagram, Free Running Conversion


Table 26-1. ADC Conversion Time

| Condition | Sample \& Hold <br> (Cycles from Start of Conversion) | Conversion Time <br> (Cycles) |
| :--- | :---: | :---: |
| First conversion | 13.5 | 25 |
| Normal conversions, single ended | 1.5 | 13 |
| Auto Triggered conversions | 2 | 13.5 |
| Normal conversions, differential | $1.5 / 2.5$ | $13 / 14$ |

### 26.4.1 Differential Channels

When using differential channels, certain aspects of the conversion need to be taken into consideration.
Differential conversions are synchronized to the internal clock $\mathrm{CK}_{\mathrm{ADC} 2}$ equal to half the ADC clock. This synchronization is done automatically by the ADC interface in such a way that the sample-and-hold occurs at a specific phase of $\mathrm{CK}_{\mathrm{ADC} 2}$. A conversion initiated by the user (that is, all single conversions, and the first free running conversion) when $\mathrm{CK}_{\mathrm{ADC} 2}$ is low will take the same amount of time as a single ended conversion (13 ADC clock cycles from the next prescaled clock cycle). A conversion initiated by the user when $\mathrm{CK}_{\mathrm{ADC}}$ is high will take 14 ADC clock cycles due to the synchronization mechanism. In Free Running mode, a new conversion is initiated immediately after the previous conversion completes, and since $\mathrm{CK}_{\mathrm{ADC} 2}$ is high at this time, all automatically started (that is, all but the first) Free Running conversions will take 14 ADC clock cycles.

If differential channels are used and conversions are started by Auto Triggering, the ADC must be switched off between conversions. When Auto Triggering is used, the ADC prescaler is reset before the conversion is started. Since the stage is dependent of a stable ADC clock prior to the conversion, this conversion will not be valid. By disabling and then re-enabling the ADC between each conversion (writing ADEN in ADCSRA to " 0 " then to " 1 "), only extended conversions are performed. The result from the extended conversions will be valid. See "Prescaling and Conversion Timing" on page 271 for timing details.

### 26.5 Changing Channel or Reference Selection

The MUXn and REFS1:0 bits in the ADMUX Register are single buffered through a temporary register to which the CPU has random access. This ensures that the channels and reference selection only takes place at a safe point during the conversion. The channel and reference selection is continuously updated until a conversion is started. Once the conversion starts, the channel and reference selection is locked to ensure a sufficient sampling time for the ADC. Continuous updating resumes in the last ADC clock cycle before the conversion completes (ADIF in ADCSRA is set). Note that the conversion starts on the following rising ADC clock edge after ADSC is written. The user is thus advised not to write new channel or reference selection values to ADMUX until one ADC clock cycle after ADSC is written.

If Auto Triggering is used, the exact time of the triggering event can be indeterministic. Special care must be taken when updating the ADMUX Register, in order to control which conversion will be affected by the new settings.

If both ADATE and ADEN is written to one, an interrupt event can occur at any time. If the ADMUX Register is changed in this period, the user cannot tell if the next conversion is based on the old or the new settings. ADMUX can be safely updated in the following ways:

1. When ADATE or ADEN is cleared.
2. During conversion, minimum one ADC clock cycle after the trigger event.
3. After a conversion, before the Interrupt Flag used as trigger source is cleared.

When updating ADMUX in one of these conditions, the new settings will affect the next ADC conversion.
Special care should be taken when changing differential channels. Once a differential channel has been selected, the stage may take as much as $125 \mu \mathrm{~s}$ to stabilize to the new value. Thus conversions should not be started within the first $125 \mu$ s after selecting a new differential channel. Alternatively, conversion results obtained within this period should be discarded.

The same settling time should be observed for the first differential conversion after changing ADC reference (by changing the REFS1:0 bits in ADMUX).

### 26.5.1 ADC Input Channels

When changing channel selections, the user should observe the following guidelines to ensure that the correct channel is selected:

In Single Conversion mode, always select the channel before starting the conversion. The channel selection may be changed one ADC clock cycle after writing one to ADSC. However, the simplest method is to wait for the conversion to complete before changing the channel selection.

In Free Running mode, always select the channel before starting the first conversion. The channel selection may be changed one ADC clock cycle after writing one to ADSC. However, the simplest method is to wait for the first conversion to complete, and then change the channel selection. Since the next conversion has already started automatically, the next result will reflect the previous channel selection. Subsequent conversions will reflect the new channel selection.

When switching to a differential gain channel, the first conversion result may have a poor accuracy due to the required settling time for the automatic offset cancellation circuitry. The user should preferably disregard the first conversion result.

### 26.5.2 ADC Voltage Reference

The reference voltage for the ADC ( $\mathrm{V}_{\text {REF }}$ ) indicates the conversion range for the ADC. Single ended channels that exceed $\mathrm{V}_{\text {REF }}$ will result in codes close to $0 \times 3 F F$. $\mathrm{V}_{\text {REF }}$ can be selected as either AVCC, internal 1.1V reference, internal 2.56 V reference or external AREF pin.

AVCC is connected to the ADC through a passive switch. The internal 1.1 V reference is generated from the internal bandgap reference (VBG) through an internal amplifier. In either case, the external AREF pin is directly connected to the ADC, and the reference voltage can be made more immune to noise by connecting a capacitor between the AREF pin and ground. $\mathrm{V}_{\text {REF }}$ can also be measured at the AREF pin with a high impedant voltmeter. Note that $\mathrm{V}_{\text {REF }}$ is a high impedant source, and only a capacitive load should be connected in a system. The Internal 2.56 V reference is generated from the 1.1 V reference.

If the user has a fixed voltage source connected to the AREF pin, the user may not use the other reference voltage options in the application, as they will be shorted to the external voltage. If no external voltage is applied to the AREF pin, the user may switch between AVCC, 1.1V and 2.56 V as reference selection. The first ADC conversion result after switching reference voltage source may be inaccurate, and the user is advised to discard this result.

If differential channels are used, the selected reference should not be closer to AVCC than indicated in "ADC Characteristics - Preliminary Data" on page 365.

### 26.6 ADC Noise Canceler

The ADC features a noise canceler that enables conversion during sleep mode to reduce noise induced from the CPU core and other I/O peripherals. The noise canceler can be used with ADC Noise Reduction and Idle mode. To make use of this feature, the following procedure should be used:

1. Make sure that the ADC is enabled and is not busy converting. Single Conversion mode must be selected and the ADC conversion complete interrupt must be enabled.
2. Enter ADC Noise Reduction mode (or Idle mode). The ADC will start a conversion once the CPU has been halted.
3. If no other interrupts occur before the ADC conversion completes, the ADC interrupt will wake up the CPU and execute the ADC Conversion Complete interrupt routine. If another interrupt wakes up the CPU before the ADC conversion is complete, that interrupt will be executed, and an ADC Conversion Complete interrupt request will be generated when the ADC conversion completes. The CPU will remain in active mode until a new sleep command is executed.
Note that the ADC will not be automatically turned off when entering other sleep modes than Idle mode and ADC Noise Reduction mode. The user is advised to write zero to ADEN before entering such sleep modes to avoid excessive power consumption.

If the ADC is enabled in such sleep modes and the user wants to perform differential conversions, the user is advised to switch the ADC off and on after waking up from sleep to prompt an extended conversion to get a valid result.

### 26.6.1 Analog Input Circuitry

The analog input circuitry for single ended channels is illustrated in Figure 26-8 on page 276 An analog source applied to ADCn is subjected to the pin capacitance and input leakage of that pin, regardless of whether that channel is selected as input for the ADC. When the channel is selected, the source must drive the S/H capacitor through the series resistance (combined resistance in the input path).

The ADC is optimized for analog signals with an output impedance of approximately $10 \mathrm{k} \Omega$ or less. If such a source is used, the sampling time will be negligible. If a source with higher impedance is used, the sampling time will depend on how long time the source needs to charge the S/H capacitor, which can vary widely. The user is recommended to only use low impedant sources with slowly varying signals, since this minimizes the required charge transfer to the S/H capacitor.

Signal components higher than the Nyquist frequency ( $f_{A D C} / 2$ ) should not be present for either kind of channels, to avoid distortion from unpredictable signal convolution. The user is advised to remove high frequency components with a low-pass filter before applying the signals as inputs to the ADC.

Figure 26-8. Analog Input Circuitry


### 26.6.2 Analog Noise Canceling Techniques

Digital circuitry inside and outside the device generates EMI which might affect the accuracy of analog measurements. If conversion accuracy is critical, the noise level can be reduced by applying the following techniques:

1. Keep analog signal paths as short as possible. Make sure analog tracks run over the ground plane, and keep them well away from high-speed switching digital tracks.
2. The AVCC pin on the device should be connected to the digital $\mathrm{V}_{\mathrm{CC}}$ supply voltage via an LC network as shown in Figure 26-9.
3. Use the ADC noise canceler function to reduce induced noise from the CPU.
4. If any ADC port pins are used as digital outputs, it is essential that these do not switch while a conversion is in progress.

Figure 26-9. ADC Power Connections, ATmega1281/2561.


Figure 26-10. ADC Power Connections, ATmega640/1280/2560


### 26.6.3 Offset Compensation Schemes

The stage has a built-in offset cancellation circuitry that nulls the offset of differential measurements as much as possible. The remaining offset in the analog path can be measured directly by selecting the same channel for both differential inputs. This offset residue can be then subtracted in software from the measurement results. Using this kind of software based offset correction, offset on any channel can be reduced below one LSB.

### 26.6.4 ADC Accuracy Definitions

An $n$-bit single-ended ADC converts a voltage linearly between GND and $V_{\text {REF }}$ in $2^{n}$ steps (LSBs). The lowest code is read as 0 , and the highest code is read as $2^{n}-1$.

Several parameters describe the deviation from the ideal behavior:

- Offset: The deviation of the first transition ( $0 \times 000$ to $0 \times 001$ ) compared to the ideal transition (at 0.5 LSB ). Ideal value: 0 LSB.

Figure 26-11. Offset Error


- Gain Error: After adjusting for offset, the Gain Error is found as the deviation of the last transition (0x3FE to $0 \times 3 F F$ ) compared to the ideal transition (at 1.5 LSB below maximum). Ideal value: 0 LSB.

Figure 26-12. Gain Error


- Integral Non-linearity (INL): After adjusting for offset and gain error, the INL is the maximum deviation of an actual transition compared to an ideal transition for any code. Ideal value: 0 LSB.

Figure 26-13. Integral Non-linearity (INL)


- Differential Non-linearity (DNL): The maximum deviation of the actual code width (the interval between two adjacent transitions) from the ideal code width (1 LSB). Ideal value: 0 LSB.

Figure 26-14. Differential Non-linearity (DNL)


- Quantization Error: Due to the quantization of the input voltage into a finite number of codes, a range of input voltages ( 1 LSB wide) will code to the same value. Always $\pm 0.5 \mathrm{LSB}$.
- Absolute Accuracy: The maximum deviation of an actual (unadjusted) transition compared to an ideal transition for any code. This is the compound effect of offset, gain error, differential error, non-linearity, and quantization error. Ideal value: $\pm 0.5$ LSB.


### 26.7 ADC Conversion Result

After the conversion is complete (ADIF is high), the conversion result can be found in the ADC Result Registers (ADCL, ADCH).

For single ended conversion, the result is

$$
A D C=\frac{V_{I N} \cdot 1024}{V_{\text {REF }}}
$$

where $\mathrm{V}_{\text {IN }}$ is the voltage on the selected input pin and $\mathrm{V}_{\text {REF }}$ the selected voltage reference (see Table 26-3 on page 281 and Table $26-4$ on page 282). $0 \times 000$ represents analog ground, and $0 \times 3 F F$ represents the selected reference voltage minus one LSB.

If differential channels are used, the result is

$$
A D C=\frac{\left(V_{P O S}-V_{N E G}\right) \cdot 512}{V_{R E F}}
$$

where $\mathrm{V}_{\text {POS }}$ is the voltage on the positive input pin, $\mathrm{V}_{\text {NEG }}$ the voltage on the negative input pin, and $\mathrm{V}_{\text {REF }}$ the selected voltage reference. The result is presented in two's complement form, from $0 \times 200(-512 d)$ through 0x1FF ( +511 d ). Note that if the user wants to perform a quick polarity check of the result, it is sufficient to read the MSB of the result (ADC9 in ADCH). If the bit is one, the result is negative, and if this bit is zero, the result is positive. Figure $26-15$ shows the decoding of the differential input range.

Table 26-2 on page 281 shows the resulting output codes if the differential input channel pair (ADCn - ADCm) is selected with a gain of GAIN and a reference voltage of $\mathrm{V}_{\text {REF }}$.

Figure 26-15. Differential Measurement Range


Table 26-2. Correlation Between Input Voltage and Output Codes

| $\mathrm{V}_{\text {ADCn }}$ | Read Code | Corresponding Decimal Value |
| :---: | :---: | :---: |
| $\mathrm{V}_{\text {ADCm }}+\mathrm{V}_{\text {REF }} / \mathrm{GAIN}$ | 0x1FF | 511 |
| $\mathrm{V}_{\text {ADCm }}+0.999 \mathrm{~V}_{\text {REF }} / \mathrm{GAIN}$ | 0x1FF | 511 |
| $\mathrm{V}_{\text {ADCm }}+0.998 \mathrm{~V}_{\text {REF }} / \mathrm{GAIN}$ | 0x1FE | 510 |
| ... | ... | $\ldots$ |
| $\mathrm{V}_{\text {ADCm }}+0.001 \mathrm{~V}_{\text {REF }} / \mathrm{GAIN}$ | 0x001 | 1 |
| $\mathrm{V}_{\text {ADCm }}$ | 0x000 | 0 |
| $\mathrm{V}_{\text {ADCm }}-0.001 \mathrm{~V}_{\text {REF }} / \mathrm{GAIN}$ | 0x3FF | -1 |
| ... | ... | $\ldots$ |
| $\mathrm{V}_{\text {ADCm }}-0.999 \mathrm{~V}_{\text {REF }} / \mathrm{GAIN}$ | 0x201 | -511 |
| $\mathrm{V}_{\text {ADCm }}-\mathrm{V}_{\text {REF }} / \mathrm{GAIN}$ | 0x200 | -512 |

## Example:

ADMUX $=0 \times F B$ (ADC3 - ADC2, $10 \times$ gain, 2.56 V reference, left adjusted result).
Voltage on ADC3 is 300 mV , voltage on ADC2 is 500 mV .
ADCR $=512 \times 10 \times(300-500) / 2560=-400=0 \times 270$.
ADCL will thus read $0 \times 00$, and $A D C H$ will read $0 \times 9 C$. Writing zero to $A D L A R$ right adjusts the result: $A D C L=0 \times 70$, $\mathrm{ADCH}=0 \times 02$.

### 26.8 Register Description

### 26.8.1 ADMUX - ADC Multiplexer Selection Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | ADMUX |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x7C) | REFS1 | REFS0 | ADLAR | MUX4 | MUX3 | MUX2 | MUX1 | MUXO |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7:6 - REFS1:0: Reference Selection Bits

These bits select the voltage reference for the ADC, as shown in Table 26-3. If these bits are changed during a conversion, the change will not go in effect until this conversion is complete (ADIF in ADCSRA is set). The internal voltage reference options may not be used if an external reference voltage is being applied to the AREF pin.

Table 26-3. Voltage Reference Selections for ADC

| REFS1 | REFS0 | Voltage Reference Selection $^{(1)}$ |
| :---: | :---: | :--- |
| 0 | 0 | AREF, Internal V $_{\text {REF }}$ turned off |
| 0 | 1 | AVCC with external capacitor at AREF pin |
| 1 | 0 | Internal 1.1V Voltage Reference with external capacitor at AREF pin |
| 1 | 1 | Internal 2.56V Voltage Reference with external capacitor at AREF pin |

Note: 1. If $10 x$ or $200 x$ gain is selected, only 2.56 V should be used as Internal Voltage Reference. For differential conversion, only 1.1 V cannot be used as internal voltage reference.

## - Bit 5 - ADLAR: ADC Left Adjust Result

The ADLAR bit affects the presentation of the ADC conversion result in the ADC Data Register. Write one to ADLAR to left adjust the result. Otherwise, the result is right adjusted. Changing the ADLAR bit will affect the ADC Data Register immediately, regardless of any ongoing conversions. For a complete description of this bit, see "ADCL and ADCH - The ADC Data Register" on page 286.

## - Bits 4:0 - MUX4:0: Analog Channel and Gain Selection Bits

The value of these bits selects which combination of analog inputs are connected to the ADC. See Table 26-4 for details. If these bits are changed during a conversion, the change will not go in effect until this conversion is complete (ADIF in ADCSRA is set).

### 26.8.2 ADCSRB - ADC Control and Status Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | ADCSRB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x7B) | - | ACME | - | - | MUX5 | ADTS2 | ADTS1 | ADTS0 |  |
| Read/Write | R | R/W | R | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 3 - MUX5: Analog Channel and Gain Selection Bit

This bit is used together with MUX4:0 in ADMUX to select which combination in of analog inputs are connected to the ADC. See Table 26-4 for details. If this bit is changed during a conversion, the change will not go in effect until this conversion is complete.

This bit is not valid for ATmega1281/2561.
Table 26-4. Input Channel Selections

| MUX5:0 | Single Ended Input | Positive Differential Input | Negative Differential Input | Gain |
| :---: | :---: | :---: | :---: | :---: |
| 000000 | ADC0 |  |  |  |
| 000001 | ADC1 |  |  |  |
| 000010 | ADC2 |  | N/A |  |
| 000011 | ADC3 |  |  |  |
| 000100 | ADC4 |  |  |  |
| 000101 | ADC5 |  |  |  |
| 000110 | ADC6 |  |  |  |
| 000111 | ADC7 |  |  |  |

Table 26-4. Input Channel Selections (Continued)

| MUX5:0 | Single Ended Input | Positive Differential Input | Negative Differential Input | Gain |
| :---: | :---: | :---: | :---: | :---: |
| 001000 ${ }^{(1)}$ | N/A | ADC0 | ADC0 | 10x |
| $001001{ }^{(1)}$ |  | ADC1 | ADC0 | 10x |
| 001010 ${ }^{(1)}$ |  | ADC0 | ADC0 | 200x |
| $001011^{(1)}$ |  | ADC1 | ADC0 | 200x |
| 001100 ${ }^{(1)}$ |  | ADC2 | ADC2 | 10x |
| $001101{ }^{(1)}$ |  | ADC3 | ADC2 | 10x |
| 001110 ${ }^{(1)}$ |  | ADC2 | ADC2 | 200x |
| $001111^{(1)}$ |  | ADC3 | ADC2 | 200x |
| 010000 |  | ADC0 | ADC1 | $1 \times$ |
| 010001 |  | ADC1 | ADC1 | 1× |
| 010010 |  | ADC2 | ADC1 | $1 \times$ |
| 010011 |  | ADC3 | ADC1 | $1 \times$ |
| 010100 |  | ADC4 | ADC1 | $1 \times$ |
| 010101 |  | ADC5 | ADC1 | 1× |
| 010110 |  | ADC6 | ADC1 | 1× |
| 010111 |  | ADC7 | ADC1 | 1× |
| 011000 |  | ADC0 | ADC2 | 1× |
| 011001 |  | ADC1 | ADC2 | $1 \times$ |
| 011010 | N/A | ADC2 | ADC2 | 1× |
| 011011 |  | ADC3 | ADC2 | $1 \times$ |
| 011100 |  | ADC4 | ADC2 | 1× |
| 011101 |  | ADC5 | ADC2 | $1 \times$ |
| 011110 | $1.1 \mathrm{~V}\left(\mathrm{~V}_{\mathrm{BG}}\right)$ | N/A |  |  |
| 011111 | OV (GND) |  |  |  |
| 100000 | ADC8 | N/A |  |  |
| 100001 | ADC9 |  |  |  |
| 100010 | ADC10 |  |  |  |
| 100011 | ADC11 |  |  |  |
| 100100 | ADC12 |  |  |  |
| 100101 | ADC13 |  |  |  |
| 100110 | ADC14 |  |  |  |
| 100111 | ADC15 |  |  |  |

Table 26-4. Input Channel Selections (Continued)

| MUX5:0 | Single Ended Input | Positive Differential Input | Negative Differential Input | Gain |
| :---: | :---: | :---: | :---: | :---: |
| $101000{ }^{(1)}$ | N/A | ADC8 | ADC8 | 10x |
| $101001{ }^{(1)}$ |  | ADC9 | ADC8 | $10 \times$ |
| $101010^{(1)}$ |  | ADC8 | ADC8 | 200x |
| $101011^{(1)}$ |  | ADC9 | ADC8 | 200x |
| $101100{ }^{(1)}$ |  | ADC10 | ADC10 | 10x |
| $101101{ }^{(1)}$ |  | ADC11 | ADC10 | $10 \times$ |
| 101110 ${ }^{(1)}$ |  | ADC10 | ADC10 | 200x |
| $101111^{(1)}$ |  | ADC11 | ADC10 | 200x |
| 110000 |  | ADC8 | ADC9 | $1 \times$ |
| 110001 |  | ADC9 | ADC9 | 1× |
| 110010 |  | ADC10 | ADC9 | $1 \times$ |
| 110011 |  | ADC11 | ADC9 | $1 \times$ |
| 110100 |  | ADC12 | ADC9 | $1 \times$ |
| 110101 |  | ADC13 | ADC9 | $1 \times$ |
| 110110 |  | ADC14 | ADC9 | $1 \times$ |
| 110111 |  | ADC15 | ADC9 | $1 \times$ |
| 111000 |  | ADC8 | ADC10 | $1 \times$ |
| 111001 |  | ADC9 | ADC10 | $1 \times$ |
| 111010 |  | ADC10 | ADC10 | $1 \times$ |
| 111011 |  | ADC11 | ADC10 | $1 \times$ |
| 111100 |  | ADC12 | ADC10 | $1 \times$ |
| 111101 | N/A | ADC13 | ADC10 | $1 \times$ |
| 111110 | Reserved | N/A |  |  |
| 111111 | Reserved | N/A |  |  |

Note: 1. To reach the given accuracy, $10 \times$ or $200 \times$ Gain should not be used for operating voltage below 2.7 V .

### 26.8.3 ADCSRA - ADC Control and Status Register A

Bit
(0x7A)
Read/Write
Initial Value

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ADEN | ADSC | ADATE | ADIF | ADIE | ADPS2 | ADPS1 | ADPS0 |
| R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

ADCSRA

## - Bit 7 - ADEN: ADC Enable

Writing this bit to one enables the ADC. By writing it to zero, the ADC is turned off. Turning the ADC off while a conversion is in progress, will terminate this conversion.

## - Bit 6 - ADSC: ADC Start Conversion

In Single Conversion mode, write this bit to one to start each conversion. In Free Running mode, write this bit to one to start the first conversion. The first conversion after ADSC has been written after the ADC has been enabled, or if ADSC is written at the same time as the ADC is enabled, will take 25 ADC clock cycles instead of the normal 13. This first conversion performs initialization of the ADC.

ADSC will read as one as long as a conversion is in progress. When the conversion is complete, it returns to zero. Writing zero to this bit has no effect.

## - Bit 5 - ADATE: ADC Auto Trigger Enable

When this bit is written to one, Auto Triggering of the ADC is enabled. The ADC will start a conversion on a positive edge of the selected trigger signal. The trigger source is selected by setting the ADC Trigger Select bits, ADTS in ADCSRB.

## - Bit 4 - ADIF: ADC Interrupt Flag

This bit is set when an ADC conversion completes and the Data Registers are updated. The ADC Conversion Complete Interrupt is executed if the ADIE bit and the I-bit in SREG are set. ADIF is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, ADIF is cleared by writing a logical one to the flag. Beware that if doing a Read-Modify-Write on ADCSRA, a pending interrupt can be disabled. This also applies if the SBI and CBI instructions are used.

## - Bit 3 - ADIE: ADC Interrupt Enable

When this bit is written to one and the I-bit in SREG is set, the ADC Conversion Complete Interrupt is activated.

- Bits 2:0 - ADPS2:0: ADC Prescaler Select Bits

These bits determine the division factor between the XTAL frequency and the input clock to the ADC.
Table 26-5. ADC Prescaler Selections

| ADPS2 | ADPS1 | ADPS0 | Division Factor |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 2 |
| 0 | 0 | 1 | 2 |
| 0 | 1 | 0 | 4 |
| 0 | 1 | 1 | 8 |
| 1 | 0 | 0 | 16 |
| 1 | 0 | 1 | 32 |
| 1 | 1 | 0 | 64 |
| 1 | 1 | 1 | 128 |

### 26.8.4 ADCL and ADCH - The ADC Data Register

### 26.8.4.1 $\quad A D L A R=0$

| Bit | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | $\begin{aligned} & \text { ADCH } \\ & \text { ADCL } \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x79) | - | - | - | - | - | - | ADC9 | ADC8 |  |
| (0x78) | ADC7 | ADC6 | ADC5 | ADC4 | ADC3 | ADC2 | ADC1 | ADC0 |  |
|  | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| Read/Write | R | R | R | R | R | R | R | R |  |
|  | R | R | R | R | R | R | R | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

### 26.8.4.2 $A D L A R=1$

| Bit | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x79) | ADC9 | ADC8 | ADC7 | ADC6 | ADC5 | ADC4 | ADC3 | ADC2 | ADCH |
| (0x78) | ADC1 | ADC0 | - | - | - | - | - | - | ADCL |
|  | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |  |
| Read/Write | R | R | R | R | R | R | R | R |  |
|  | R | R | R | R | R | R | R | R |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

When an ADC conversion is complete, the result is found in these two registers. If differential channels are used, the result is presented in two's complement form.

When ADCL is read, the ADC Data Register is not updated until ADCH is read. Consequently, if the result is left adjusted and no more than 8 -bit precision ( 7 bit + sign bit for differential input channels) is required, it is sufficient to read ADCH. Otherwise, ADCL must be read first, then ADCH.

The ADLAR bit in ADMUX, and the MUXn bits in ADMUX affect the way the result is read from the registers. If ADLAR is set, the result is left adjusted. If ADLAR is cleared (default), the result is right adjusted.

## - ADC9:0: ADC Conversion Result

These bits represent the result from the conversion, as detailed in "ADC Conversion Result" on page 280.

### 26.8.5 ADCSRB - ADC Control and Status Register B

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | ADCSRB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x7B) | - | ACME | - | - | MUX5 | ADTS2 | ADTS1 | ADTS0 |  |
| Read/Write | R | R/W | R | R | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - Res: Reserved Bit

This bit is reserved for future use. To ensure compatibility with future devices, this bit must be written to zero when ADCSRB is written.

## - Bit 2:0 - ADTS2:0: ADC Auto Trigger Source

If ADATE in ADCSRA is written to one, the value of these bits selects which source will trigger an ADC conversion. If ADATE is cleared, the ADTS2:0 settings will have no effect. A conversion will be triggered by the rising edge of the selected Interrupt Flag. Note that switching from a trigger source that is cleared to a trigger source that is set, will generate a positive edge on the trigger signal. If ADEN in ADCSRA is set, this will start a conversion. Switching to Free Running mode (ADTS[2:0]=0) will not cause a trigger event, even if the ADC Interrupt Flag is set.

Table 26-6. ADC Auto Trigger Source Selections

| ADTS2 | ADTS1 | ADTS0 | Trigger Source |
| :---: | :---: | :---: | :--- |
| 0 | 0 | 0 | Free Running mode |
| 0 | 0 | 1 | Analog Comparator |
| 0 | 1 | 0 | External Interrupt Request 0 |
| 0 | 1 | 1 | Timer/Counter0 Compare Match A |
| 1 | 0 | 0 | Timer/Counter0 Overflow |
| 1 | 0 | 1 | Timer/Counter1 Compare Match B |
| 1 | 1 | 0 | Timer/Counter1 Overflow |
| 1 | 1 | 1 | Timer/Counter1 Capture Event |

Note: Free running mode cannot be used for differential channels (see chapter "Differential Channels" on page 273).

### 26.8.6 DIDR0 - Digital Input Disable Register 0

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | DIDR0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x7E) | ADC7D | ADC6D | ADC5D | ADC4D | ADC3D | ADC2D | ADC1D | ADCOD |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7:0 - ADC7D:ADC0D: ADC7:0 Digital Input Disable

When this bit is written logic one, the digital input buffer on the corresponding ADC pin is disabled. The corresponding PIN Register bit will always read as zero when this bit is set. When an analog signal is applied to the ADC7:0 pin and the digital input from this pin is not needed, this bit should be written logic one to reduce power consumption in the digital input buffer.

### 26.8.7 DIDR2 - Digital Input Disable Register 2

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | DIDR2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x7D) | ADC15D | ADC14D | ADC13D | ADC12D | ADC11D | ADC10D | ADC9D | ADC8D |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7:0 - ADC15D:ADC8D: ADC15:8 Digital Input Disable

When this bit is written logic one, the digital input buffer on the corresponding ADC pin is disabled. The corresponding PIN Register bit will always read as zero when this bit is set. When an analog signal is applied to the ADC15:8 pin and the digital input from this pin is not needed, this bit should be written logic one to reduce power consumption in the digital input buffer.

## 27. JTAG Interface and On-chip Debug System

### 27.1 Features

- JTAG (IEEE std. 1149.1 Compliant) Interface
- Boundary-scan Capabilities According to the IEEE std. 1149.1 (JTAG) Standard
- Debugger Access to:
- All Internal Peripheral Units
- Internal and External RAM
- The Internal Register File
- Program Counter
- EEPROM and Flash Memories
- Extensive On-chip Debug Support for Break Conditions, Including
- AVR Break Instruction
- Break on Change of Program Memory Flow
- Single Step Break
- Program Memory Break Points on Single Address or Address Range
- Data Memory Break Points on Single Address or Address Range
- Programming of Flash, EEPROM, Fuses, and Lock Bits through the JTAG Interface
- On-chip Debugging Supported by AVR Studio ${ }^{\circledR}$


### 27.2 Overview

The AVR IEEE std. 1149.1 compliant JTAG interface can be used for

- Testing PCBs by using the JTAG Boundary-scan capability
- Programming the non-volatile memories, Fuses and Lock bits
- On-chip debugging

A brief description is given in the following sections. Detailed descriptions for Programming via the JTAG interface, and using the Boundary-scan Chain can be found in the sections "Programming via the JTAG Interface" on page 342 and "IEEE 1149.1 (JTAG) Boundary-scan" on page 295, respectively. The On-chip Debug support is considered being private JTAG instructions, and distributed within Atmel and to selected third party vendors only.

Figure 27-1 on page 290 shows a block diagram of the JTAG interface and the On-chip Debug system. The TAP Controller is a state machine controlled by the TCK and TMS signals. The TAP Controller selects either the JTAG Instruction Register or one of several Data Registers as the scan chain (Shift Register) between the TDI - input and TDO - output. The Instruction Register holds JTAG instructions controlling the behavior of a Data Register.

The ID-Register, Bypass Register, and the Boundary-scan Chain are the Data Registers used for board-level testing. The JTAG Programming Interface (actually consisting of several physical and virtual Data Registers) is used for serial programming via the JTAG interface. The Internal Scan Chain and Break Point Scan Chain are used for On-chip debugging only.

Figure 27-1. Block Diagram


### 27.3 TAP - Test Access Port

The JTAG interface is accessed through four of the AVR's pins. In JTAG terminology, these pins constitute the Test Access Port - TAP. These pins are:

- TMS: Test mode select. This pin is used for navigating through the TAP-controller state machine
- TCK: Test Clock. JTAG operation is synchronous to TCK
- TDI: Test Data In. Serial input data to be shifted in to the Instruction Register or Data Register (Scan Chains)
- TDO: Test Data Out. Serial output data from Instruction Register or Data Register

The IEEE std. 1149.1 also specifies an optional TAP signal; TRST - Test ReSeT - which is not provided.
When the JTAGEN Fuse is unprogrammed, these four TAP pins are normal port pins, and the TAP controller is in reset. When programmed, the input TAP signals are internally pulled high and the JTAG is enabled for Boundaryscan and programming. The device is shipped with this fuse programmed.

For the On-chip Debug system, in addition to the JTAG interface pins, the RESET pin is monitored by the debugger to be able to detect external reset sources. The debugger can also pull the RESET pin low to reset the whole system, assuming only open collectors on the reset line are used in the application.

Figure 27-2. TAP Controller State Diagram


### 27.3.1 TAP Controller

The TAP controller is a 16 -state finite state machine that controls the operation of the Boundary-scan circuitry, JTAG programming circuitry, or On-chip Debug system. The state transitions depicted in Figure 27-2 depend on the signal present on TMS (shown adjacent to each state transition) at the time of the rising edge at TCK. The initial state after a Power-on Reset is Test-Logic-Reset.

As a definition in this document, the LSB is shifted in and out first for all Shift Registers.
Assuming Run-Test/Idle is the present state, a typical scenario for using the JTAG interface is:

- At the TMS input, apply the sequence 1, 1, 0, 0 at the rising edges of TCK to enter the Shift Instruction Register - Shift-IR state. While in this state, shift the four bits of the JTAG instructions into the JTAG Instruction Register from the TDI input at the rising edge of TCK. The TMS input must be held low during input of the 3 LSBs in order to remain in the Shift-IR state. The MSB of the instruction is shifted in when this state is left by setting TMS high. While the instruction is shifted in from the TDI pin, the captured IR-state $0 \times 01$ is shifted out on the TDO pin. The JTAG Instruction selects a particular Data Register as path between TDI and TDO and controls the circuitry surrounding the selected Data Register.
- Apply the TMS sequence $1,1,0$ to re-enter the Run-Test/Idle state. The instruction is latched onto the parallel output from the Shift Register path in the Update-IR state. The Exit-IR, Pause-IR, and Exit2-IR states are only used for navigating the state machine.
- At the TMS input, apply the sequence 1, 0, 0 at the rising edges of TCK to enter the Shift Data Register - ShiftDR state. While in this state, upload the selected Data Register (selected by the present JTAG instruction in the JTAG Instruction Register) from the TDI input at the rising edge of TCK. In order to remain in the Shift-DR state, the TMS input must be held low during input of all bits except the MSB. The MSB of the data is shifted in when this state is left by setting TMS high. While the Data Register is shifted in from the TDI pin, the parallel inputs to the Data Register captured in the Capture-DR state is shifted out on the TDO pin.
- Apply the TMS sequence 1, 1,0 to re-enter the Run-Test/Idle state. If the selected Data Register has a latched parallel-output, the latching takes place in the Update-DR state. The Exit-DR, Pause-DR, and Exit2-DR states are only used for navigating the state machine.
As shown in the state diagram, the Run-Test/Idle state need not be entered between selecting JTAG instruction and using Data Registers, and some JTAG instructions may select certain functions to be performed in the RunTest/Idle, making it unsuitable as an Idle state.

Note: Independent of the initial state of the TAP Controller, the Test-Logic-Reset state can always be entered by holding TMS high for five TCK clock periods.
For detailed information on the JTAG specification, refer to the literature listed in "Bibliography" on page 294.

### 27.4 Using the Boundary-scan Chain

A complete description of the Boundary-scan capabilities are given in the section "IEEE 1149.1 (JTAG) Boundaryscan" on page 295.

### 27.5 Using the On-chip Debug System

As shown in Figure 27-1 on page 290, the hardware support for On-chip Debugging consists mainly of:

- A scan chain on the interface between the internal AVR CPU and the internal peripheral units
- Break Point unit
- Communication interface between the CPU and JTAG system

All read or modify/write operations needed for implementing the Debugger are done by applying AVR instructions via the internal AVR CPU Scan Chain. The CPU sends the result to an I/O memory mapped location which is part of the communication interface between the CPU and the JTAG system.
The Break Point Unit implements Break on Change of Program Flow, Single Step Break, two Program Memory Break Points, and two combined Break Points. Together, the four Break Points can be configured as either:

- 4 single Program Memory Break Points
- 3 Single Program Memory Break Point + 1 single Data Memory Break Point
- 2 single Program Memory Break Points +2 single Data Memory Break Points
- 2 single Program Memory Break Points + 1 Program Memory Break Point with mask ("range Break Point")
- 2 single Program Memory Break Points + 1 Data Memory Break Point with mask ("range Break Point")

A debugger, like the AVR Studio, may however use one or more of these resources for its internal purpose, leaving less flexibility to the end-user.
A list of the On-chip Debug specific JTAG instructions is given in "On-chip Debug Specific JTAG Instructions" on page 293.

The JTAGEN Fuse must be programmed to enable the JTAG Test Access Port. In addition, the OCDEN Fuse must be programmed and no Lock bits must be set for the On-chip debug system to work. As a security feature, the Onchip debug system is disabled when either of the LB1 or LB2 Lock bits are set. Otherwise, the On-chip debug system would have provided a back-door into a secured device.
The AVR Studio ${ }^{\circledR}$ enables the user to fully control execution of programs on an AVR device with On-chip Debug capability, AVR In-Circuit Emulator, or the built-in AVR Instruction Set Simulator. AVR Studio supports source level
execution of Assembly programs assembled with Atmel Corporation's AVR Assembler and C programs compiled with third party vendors' compilers.
AVR Studio runs under Microsoft ${ }^{\circledR}$ Windows ${ }^{\circledR} 95 / 98 / 2000$ and Microsoft Windows NT.
For a full description of the AVR Studio, refer to the AVR Studio User Guide. Only highlights are presented in this document.

All necessary execution commands are available in AVR Studio, both on source level and on disassembly level. The user can execute the program, single step through the code either by tracing into or stepping over functions, step out of functions, place the cursor on a statement and execute until the statement is reached, stop the execution, and reset the execution target. In addition, the user can have an unlimited number of code Break Points (using the BREAK instruction) and up to two data memory Break Points, alternatively combined as a mask (range) Break Point.

### 27.6 On-chip Debug Specific JTAG Instructions

The On-chip debug support is considered being private JTAG instructions, and distributed within ATMEL and to selected third party vendors only. Instruction opcodes are listed for reference.

### 27.6.1 PRIVATEO; 0x8

Private JTAG instruction for accessing On-chip debug system.

### 27.6.2 PRIVATE1; 0x9

Private JTAG instruction for accessing On-chip debug system.

### 27.6.3 PRIVATE2; 0xA

Private JTAG instruction for accessing On-chip debug system.

### 27.6.4 PRIVATE3; 0xB

Private JTAG instruction for accessing On-chip debug system.

### 27.7 Using the JTAG Programming Capabilities

Programming of AVR parts via JTAG is performed via the 4-pin JTAG port, TCK, TMS, TDI, and TDO. These are the only pins that need to be controlled/observed to perform JTAG programming (in addition to power pins). It is not required to apply 12 V externally. The JTAGEN Fuse must be programmed and the JTD bit in the MCUCR Register must be cleared to enable the JTAG Test Access Port.

The JTAG programming capability supports:

- Flash programming and verifying
- EEPROM programming and verifying
- Fuse programming and verifying
- Lock bit programming and verifying

The Lock bit security is exactly as in parallel programming mode. If the Lock bits LB1 or LB2 are programmed, the OCDEN Fuse cannot be programmed unless first doing a chip erase. This is a security feature that ensures no back-door exists for reading out the content of a secured device.

The details on programming through the JTAG interface and programming specific JTAG instructions are given in the section "Programming via the JTAG Interface" on page 342.

### 27.8 Bibliography

For more information about general Boundary-scan, the following literature can be consulted:

- IEEE: IEEE Std. 1149.1-1990. IEEE Standard Test Access Port and Boundary-scan Architecture, IEEE, 1993
- Colin Maunder: The Board Designers Guide to Testable Logic Circuits, Addison-Wesley, 1992


### 27.9 On-chip Debug Related Register in I/O Memory

### 27.9.1 OCDR - On-chip Debug Register

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | OCDR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x31 (0x51) | MSB/IDRD |  |  |  |  |  |  | LSB |  |
| Read/Write | R/W | R/W | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

The OCDR Register provides a communication channel from the running program in the microcontroller to the debugger. The CPU can transfer a byte to the debugger by writing to this location. At the same time, an internal flag; I/O Debug Register Dirty - IDRD - is set to indicate to the debugger that the register has been written. When the CPU reads the OCDR Register the 7 LSB will be from the OCDR Register, while the MSB is the IDRD bit. The debugger clears the IDRD bit when it has read the information.

In some AVR devices, this register is shared with a standard I/O location. In this case, the OCDR Register can only be accessed if the OCDEN Fuse is programmed, and the debugger enables access to the OCDR Register. In all other cases, the standard I/O location is accessed.

Refer to the debugger documentation for further information on how to use this register.

## 28. IEEE 1149.1 (JTAG) Boundary-scan

### 28.1 Features

- JTAG (IEEE std. 1149.1 compliant) Interface
- Boundary-scan Capabilities According to the JTAG Standard
- Full Scan of all Port Functions as well as Analog Circuitry having Off-chip Connections
- Supports the Optional IDCODE Instruction
- Additional Public AVR_RESET Instruction to Reset the AVR


### 28.2 System Overview

The Boundary-scan chain has the capability of driving and observing the logic levels on the digital I/O pins, as well as the boundary between digital and analog logic for analog circuitry having off-chip connections. At system level, all ICs having JTAG capabilities are connected serially by the TDI/TDO signals to form a long Shift Register. An external controller sets up the devices to drive values at their output pins, and observe the input values received from other devices. The controller compares the received data with the expected result. In this way, Boundary-scan provides a mechanism for testing interconnections and integrity of components on Printed Circuits Boards by using the four TAP signals only.

The four IEEE 1149.1 defined mandatory JTAG instructions IDCODE, BYPASS, SAMPLE/PRELOAD, and EXTEST, as well as the AVR specific public JTAG instruction AVR_RESET can be used for testing the Printed Circuit Board. Initial scanning of the Data Register path will show the ID-Code of the device, since IDCODE is the default JTAG instruction. It may be desirable to have the AVR device in reset during test mode. If not reset, inputs to the device may be determined by the scan operations, and the internal software may be in an undetermined state when exiting the test mode. Entering reset, the outputs of any port pin will instantly enter the high impedance state, making the HIGHZ instruction redundant. If needed, the BYPASS instruction can be issued to make the shortest possible scan chain through the device. The device can be set in the reset state either by pulling the external $\overline{\text { RESET }}$ pin low, or issuing the AVR_RESET instruction with appropriate setting of the Reset Data Register.

The EXTEST instruction is used for sampling external pins and loading output pins with data. The data from the output latch will be driven out on the pins as soon as the EXTEST instruction is loaded into the JTAG IR-Register. Therefore, the SAMPLE/PRELOAD should also be used for setting initial values to the scan ring, to avoid damaging the board when issuing the EXTEST instruction for the first time. SAMPLE/PRELOAD can also be used for taking a snapshot of the external pins during normal operation of the part.

The JTAGEN Fuse must be programmed and the JTD bit in the I/O Register MCUCR must be cleared to enable the JTAG Test Access Port.

When using the JTAG interface for Boundary-scan, using a JTAG TCK clock frequency higher than the internal chip frequency is possible. The chip clock is not required to run.

### 28.3 Data Registers

The Data Registers relevant for Boundary-scan operations are:

- Bypass Register
- Device Identification Register
- Reset Register
- Boundary-scan Chain


### 28.3.1 Bypass Register

The Bypass Register consists of a single Shift Register stage. When the Bypass Register is selected as path between TDI and TDO, the register is reset to 0 when leaving the Capture-DR controller state. The Bypass Register can be used to shorten the scan chain on a system when the other devices are to be tested.

### 28.3.2 Device Identification Register

Figure 28-1 shows the structure of the Device Identification Register.
Figure 28-1. The Format of the Device Identification Register

|  | MSB |  |  |  |  |  | LSB |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Bit | 31 | 28 | 27 | 12 | 11 | 1 | 0 |
| Device ID | Version |  | Part Number |  | Manufacturer ID |  | 1 |
|  | 4 bits |  | 16 bits |  | 11 bits |  | 1-bit |

### 28.3.2.1 Version

Version is a 4-bit number identifying the revision of the component. The JTAG version number follows the revision of the device. Revision A is $0 \times 0$, revision $B$ is $0 \times 1$ and so on.

### 28.3.2.2 Part Number

The part number is a 16-bit code identifying the component. The JTAG Part Number for ATmega640/1280/1281/2560/2561 is listed in Table 30-6 on page 328.

### 28.3.2.3 Manufacturer ID

The Manufacturer ID is a 11 -bit code identifying the manufacturer. The JTAG manufacturer ID for ATMEL is listed in Table 30-6 on page 328.

### 28.3.3 Reset Register

The Reset Register is a test Data Register used to reset the part. Since the AVR tri-states Port Pins when reset, the Reset Register can also replace the function of the un-implemented optional JTAG instruction HIGHZ.

A high value in the Reset Register corresponds to pulling the external Reset low. The part is reset as long as there is a high value present in the Reset Register. Depending on the fuse settings for the clock options, the part will remain reset for a reset time-out period (see "Clock Sources" on page 40) after releasing the Reset Register. The output from this Data Register is not latched, so the reset will take place immediately, as shown in Figure 28-2 on page 297.

Figure 28-2. Reset Register


### 28.3.4 Boundary-scan Chain

The Boundary-scan Chain has the capability of driving and observing the logic levels on the digital I/O pins, as well as the boundary between digital and analog logic for analog circuitry having off-chip connections.
See "Boundary-scan Chain" on page 298 for a complete description.

### 28.4 Boundary-scan Specific JTAG Instructions

The Instruction Register is 4-bit wide, supporting up to 16 instructions. Listed below are the JTAG instructions useful for Boundary-scan operation. Note that the optional HIGHZ instruction is not implemented, but all outputs with tri-state capability can be set in high-impedant state by using the AVR_RESET instruction, since the initial state for all port pins is tri-state.

As a definition in this datasheet, the LSB is shifted in and out first for all Shift Registers.
The OPCODE for each instruction is shown behind the instruction name in hex format. The text describes which Data Register is selected as path between TDI and TDO for each instruction.

### 28.4.1 EXTEST; 0x0

Mandatory JTAG instruction for selecting the Boundary-scan Chain as Data Register for testing circuitry external to the AVR package. For port-pins, Pull-up Disable, Output Control, Output Data, and Input Data are all accessible in the scan chain. For Analog circuits having off-chip connections, the interface between the analog and the digital logic is in the scan chain. The contents of the latched outputs of the Boundary-scan chain is driven out as soon as the JTAG IR-Register is loaded with the EXTEST instruction.

The active states are:

- Capture-DR: Data on the external pins are sampled into the Boundary-scan Chain
- Shift-DR: The Internal Scan Chain is shifted by the TCK input
- Update-DR: Data from the scan chain is applied to output pins


### 28.4.2 IDCODE; 0x1

Optional JTAG instruction selecting the 32-bit ID-Register as Data Register. The ID-Register consists of a version number, a device number and the manufacturer code chosen by JEDEC. This is the default instruction after powerup.
The active states are:

- Capture-DR: Data in the IDCODE Register is sampled into the Boundary-scan Chain
- Shift-DR: The IDCODE scan chain is shifted by the TCK input


### 28.4.3 SAMPLE_PRELOAD; 0x2

Mandatory JTAG instruction for pre-loading the output latches and taking a snap-shot of the input/output pins without affecting the system operation. However, the output latches are not connected to the pins. The Boundary-scan Chain is selected as Data Register.

The active states are:

- Capture-DR: Data on the external pins are sampled into the Boundary-scan Chain
- Shift-DR: The Boundary-scan Chain is shifted by the TCK input
- Update-DR: Data from the Boundary-scan chain is applied to the output latches. However, the output latches are not connected to the pins


### 28.4.4 AVR_RESET; 0xC

The AVR specific public JTAG instruction for forcing the AVR device into the Reset mode or releasing the JTAG reset source. The TAP controller is not reset by this instruction. The one bit Reset Register is selected as Data Register. Note that the reset will be active as long as there is a logic "one" in the Reset Chain. The output from this chain is not latched.

The active states are:

- Shift-DR: The Reset Register is shifted by the TCK input


### 28.4.5 BYPASS; 0xF

Mandatory JTAG instruction selecting the Bypass Register for Data Register.
The active states are:

- Capture-DR: Loads a logic "0" into the Bypass Register
- Shift-DR: The Bypass Register cell between TDI and TDO is shifted


### 28.5 Boundary-scan Chain

The Boundary-scan chain has the capability of driving and observing the logic levels on the digital I/O pins, as well as the boundary between digital and analog logic for analog circuitry having off-chip connection.

### 28.5.1 Scanning the Digital Port Pins

Figure $28-3$ on page 299 shows the Boundary-scan Cell for a bi-directional port pin. The pull-up function is disabled during Boundary-scan when the JTAG IC contains EXTEST or SAMPLE_PRELOAD. The cell consists of a bidirectional pin cell that combines the three signals Output Control - OCxn, Output Data - ODxn, and Input Data IDxn, into only a two-stage Shift Register. The port and pin indexes are not used in the following description.

The Boundary-scan logic is not included in the figures in the datasheet. Figure $28-4$ on page 300 shows a simple digital port pin as described in the section "I/O-Ports" on page 67. The Boundary-scan details from Figure 28-3 on page 299 replaces the dashed box in Figure 28-4 on page 300.

When no alternate port function is present, the Input Data - ID - corresponds to the PINxn Register value (but ID has no synchronizer), Output Data corresponds to the PORT Register, Output Control corresponds to the Data Direction - DD Register, and the Pull-up Enable - PUExn - corresponds to logic expression PUD • DDxn • PORTxn.
Digital alternate port functions are connected outside the dotted box in Figure $28-4$ on page 300 to make the scan chain read the actual pin value. For analog function, there is a direct connection from the external pin to the analog circuit. There is no scan chain on the interface between the digital and the analog circuitry, but some digital control signal to analog circuitry are turned off to avoid driving contention on the pads.

When JTAG IR contains EXTEST or SAMPLE_PRELOAD the clock is not sent out on the port pins even if the CKOUT fuse is programmed. Even though the clock is output when the JTAG IR contains SAMPLE_PRELOAD, the clock is not sampled by the boundary scan.

Figure 28-3. Boundary-scan Cell for Bi-directional Port Pin with Pull-up Function.


Figure 28-4. General Port Pin Schematic Diagram
See Boundary-scan
Description for Details!


### 28.5.2 Scanning the RESET Pin

The RESET pin accepts 5 V active low logic for standard reset operation, and 12 V active high logic for High Voltage Parallel programming. An observe-only cell as shown in Figure 28-5 is inserted for the 5 V reset signal.

Figure 28-5. Observe-only Cell


### 28.6 Boundary-scan Related Register in I/O Memory

### 28.6.1 MCUCR - MCU Control Register

The MCU Control Register contains control bits for general MCU functions.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | MCUCR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x35 (0x55) | JTD | - | - | PUD | - | - | IVSEL | IVCE |  |
| Read/Write | R/W | R | R | R/W | R | R | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bits 7 - JTD: JTAG Interface Disable

When this bit is zero, the JTAG interface is enabled if the JTAGEN Fuse is programmed. If this bit is one, the JTAG interface is disabled. In order to avoid unintentional disabling or enabling of the JTAG interface, a timed sequence must be followed when changing this bit: The application software must write this bit to the desired value twice within four cycles to change its value. Note that this bit must not be altered when using the On-chip Debug system.

### 28.6.2 MCUSR - MCU Status Register

The MCU Status Register provides information on which reset source caused an MCU reset.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | MCUSR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x34 (0x54) | - | - | - | JTRF | WDRF | BORF | EXTRF | PORF |  |
| Read/Write | R | R | R | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 |  |  | Bit Descrip |  |  |  |

## - Bit 4 - JTRF: JTAG Reset Flag

This bit is set if a reset is being caused by a logic one in the JTAG Reset Register selected by the JTAG instruction AVR_RESET. This bit is reset by a Power-on Reset, or by writing a logic zero to the flag.

### 28.7 ATmega640/1280/1281/2560/2561 Boundary-scan Order

Table 28-1 on page 302 shows the Scan order between TDI and TDO when the Boundary-scan chain is selected as data path. Bit 0 is the LSB; the first bit scanned in, and the first bit scanned out. The scan order follows the pinout order as far as possible. Therefore, the bits of Port A and Port K is scanned in the opposite bit order of the other ports. Exceptions from the rules are the Scan chains for the analog circuits, which constitute the most significant bits of the scan chain regardless of which physical pin they are connected to. In Figure 28-3 on page 299, PXn. Data corresponds to FFO, PXn. Control corresponds to FF1, PXn. Bit 4, bit 5, bit 6 and bit 7 of Port F is not in the scan chain, since these pins constitute the TAP pins when the JTAG is enabled.

### 28.8 Boundary-scan Description Language Files

Boundary-scan Description Language (BSDL) files describe Boundary-scan capable devices in a standard format used by automated test-generation software. The order and function of bits in the Boundary-scan Data Register are included in this description. BSDL files are available for ATmega1281/2561 and ATmega640/1280/2560.

Table 28-1. ATmega640/1280/2560 Boundary-scan Order


Table 28-1. ATmega640/1280/2560 Boundary-scan Order (Continued)

| Bit Number | Signal Name | Module |
| :---: | :---: | :---: |
| 132 | PB0.Data | Port B |
| 131 | PB0.Control |  |
| 130 | PB1.Data |  |
| 129 | PB1.Control |  |
| 128 | PB2.Data |  |
| 127 | PB2.Control |  |
| 126 | PB3.Data |  |
| 125 | PB3.Control |  |
| 124 | PB4.Data |  |
| 123 | PB4.Control |  |
| 122 | PB5.Data |  |
| 121 | PB5.Control |  |
| 120 | PB6.Data |  |
| 119 | PB6.Control |  |
| 118 | PB7.Data |  |
| 117 | PB7.Control |  |
| 116 | PH7.Data | Port H |
| 115 | PH7.Control |  |
| 114 | PG3.Data | Port G |
| 113 | PG3.Control |  |
| 112 | PG4.Data |  |
| 111 | PG4.Control |  |
| 110 | RSTT | Reset Logic (Observe Only) |
| 109 | PLo.Data | Port L |
| 108 | PLO.Control |  |
| 107 | PL1.Data |  |
| 106 | PL1.Control |  |
| 105 | PL2.Data |  |

Table 28-1. ATmega640/1280/2560 Boundary-scan Order (Continued)


Table 28-1. ATmega640/1280/2560 Boundary-scan Order (Continued)


Table 28-1. ATmega640/1280/2560 Boundary-scan Order (Continued)

| Bit Number | Signal Name | Module |  |
| :---: | :---: | :---: | :---: |
| 32 | PA3.Control |  |  |
| 31 | PA2.Data |  |  |
| 30 | PA2.Control |  |  |
| 29 | PA1.Data |  |  |
| 28 | PA1.Control |  |  |
| 27 | PA0.Data |  |  |
| 26 | PA0.Control |  |  |
| 25 | PJ7.Data |  | Port J |
| 24 | PJ7.Control |  |  |
| 23 | PK7.Data |  | Port K |
| 22 | PK7.Control |  |  |
| 21 | PK6.Data |  |  |
| 20 | PK6.Control |  |  |
| 19 | PK5.Data |  |  |
| 18 | PK5.Control |  |  |
| 17 | PK4.Data |  |  |
| 16 | PK4.Control |  |  |
| 15 | PK3.Data |  |  |
| 14 | PK3.Control |  |  |
| 13 | PK2.Data |  |  |
| 12 | PK2.Control |  |  |
| 11 | PK1.Data |  |  |
| 10 | PK1.Control |  |  |
| 9 | PK0.Data |  |  |
| 8 | PKO.Control |  |  |
| 7 | PF3.Data |  | Port F |
| 6 | PF3.Control |  |  |
| 5 | PF2.Data |  |  |
| 4 | PF2.Control |  |  |
| 3 | PF1.Data |  |  |
| 2 | PF1.Control |  |  |
| 1 | PFO.Data |  |  |
| 0 | PFO.Control |  |  |

Table 28-2. ATmega1281/2561 Boundary-scan Order

| Bit Number | Signal Name | Module |
| :---: | :---: | :---: |
| 100 | PG5.Data | Port G |
| 99 | PG5.Control |  |
| 98 | PE0.Data | Port E |
| 97 | PE0.Control |  |
| 96 | PE1.Data |  |
| 95 | PE1.Control |  |
| 94 | PE2.Data |  |
| 93 | PE2.Control |  |
| 92 | PE3.Data |  |
| 91 | PE3.Control |  |
| 90 | PE4.Data |  |
| 89 | PE4.Control |  |
| 88 | PE5.Data |  |
| 87 | PE5.Control |  |
| 86 | PE6.Data |  |
| 85 | PE6.Control |  |
| 84 | PE7.Data |  |
| 83 | PE7.Control |  |
| 82 | PB0.Data | Port B |
| 81 | PB0.Control |  |
| 80 | PB1.Data |  |
| 79 | PB1.Control |  |
| 78 | PB2.Data |  |
| 77 | PB2.Control |  |
| 76 | PB3.Data |  |
| 75 | PB3.Control |  |
| 74 | PB4.Data |  |
| 73 | PB4.Control |  |
| 72 | PB5.Data |  |
| 71 | PB5.Control |  |
| 70 | PB6.Data |  |
| 69 | PB6.Control |  |
| 68 | PB7.Data |  |
| 67 | PB7.Control |  |
| 66 | PG3.Data | Port G |

Table 28-2. ATmega1281/2561 Boundary-scan Order (Continued)

| Bit Number | Signal Name | Module |
| :---: | :---: | :---: |
| 65 | PG3.Control |  |
| 64 | PG4.Data |  |
| 63 | PG4.Control |  |
| 62 | RSTT | Reset Logic (Observe Only) |
| 61 | PD0.Data | Port D |
| 60 | PD0.Control |  |
| 59 | PD1.Data |  |
| 58 | PD1.Control |  |
| 57 | PD2.Data |  |
| 56 | PD2.Control |  |
| 55 | PD3.Data |  |
| 54 | PD3.Control |  |
| 53 | PD4.Data |  |
| 52 | PD4.Control |  |
| 51 | PD5.Data |  |
| 50 | PD5.Control |  |
| 49 | PD6.Data |  |
| 48 | PD6.Control |  |
| 47 | PD7.Data |  |
| 46 | PD7.Control |  |
| 45 | PGo.Data | Port G |
| 44 | PG0.Control |  |
| 43 | PG1.Data |  |
| 42 | PG1.Control |  |
| 41 | PC0.Data | Port C |
| 40 | PC0.Control |  |
| 39 | PC1.Data |  |
| 38 | PC1.Control |  |
| 37 | PC2.Data |  |
| 36 | PC2.Control |  |
| 35 | PC3.Data |  |
| 34 | PC3.Control |  |
| 33 | PC4.Data |  |
| 32 | PC4.Control |  |
| 31 | PC5.Data |  |
| 30 | PC5.Control |  |

Table 28-2. ATmega1281/2561 Boundary-scan Order (Continued)

| Bit Number | Signal Name | Module |  |
| :---: | :---: | :---: | :---: |
| 29 | PC6.Data |  |  |
| 28 | PC6.Control |  |  |
| 27 | PC7.Data |  |  |
| 26 | PC7.Control |  |  |
| 25 | PG2.Data | Port G |  |
| 24 | PG2.Control |  |  |
| 23 | PA7.Data |  | Port A |
| 22 | PA7.Control |  |  |
| 21 | PA6.Data |  |  |
| 20 | PA6.Control |  |  |
| 19 | PA5.Data |  |  |
| 18 | PA5.Control |  |  |
| 17 | PA4.Data |  |  |
| 16 | PA4.Control |  |  |
| 15 | PA3.Data |  |  |
| 14 | PA3.Control |  |  |
| 13 | PA2.Data |  |  |
| 12 | PA2.Control |  |  |
| 11 | PA1.Data |  |  |
| 10 | PA1.Control |  |  |
| 9 | PA0.Data |  |  |
| 8 | PAO.Control |  |  |
| 7 | PF3.Data |  | Port F |
| 6 | PF3.Control |  |  |
| 5 | PF2.Data |  |  |
| 4 | PF2.Control |  |  |
| 3 | PF1.Data |  |  |
| 2 | PF1.Control |  |  |
| 1 | PFO.Data |  |  |
| 0 | PF0.Control |  |  |

## 29. Boot Loader Support - Read-While-Write Self-Programming

The Boot Loader Support provides a real Read-While-Write Self-Programming mechanism for downloading and uploading program code by the MCU itself. This feature allows flexible application software updates controlled by the MCU using a Flash-resident Boot Loader program. The Boot Loader program can use any available data interface and associated protocol to read code and write (program) that code into the Flash memory, or read the code from the program memory. The program code within the Boot Loader section has the capability to write into the entire Flash, including the Boot Loader memory. The Boot Loader can thus even modify itself, and it can also erase itself from the code if the feature is not needed anymore. The size of the Boot Loader memory is configurable with fuses and the Boot Loader has two separate sets of Boot Lock bits which can be set independently. This gives the user a unique flexibility to select different levels of protection.

### 29.1 Features

- Read-While-Write Self-Programming
- Flexible Boot Memory Size
- High Security (Separate Boot Lock Bits for a Flexible Protection)
- Separate Fuse to Select Reset Vector
- Optimized Page ${ }^{(1)}$ Size
- Code Efficient Algorithm
- Efficient Read-Modify-Write Support

Note: 1. A page is a section in the Flash consisting of several bytes (see Table 30-7 on page 328) used during programming. The page organization does not affect normal operation.

### 29.2 Application and Boot Loader Flash Sections

The Flash memory is organized in two main sections, the Application section and the Boot Loader section (see Figure 29-2 on page 312). The size of the different sections is configured by the BOOTSZ Fuses as shown in Table 29-7 on page 320 and Figure 29-2 on page 312. These two sections can have different level of protection since they have different sets of Lock bits.

### 29.2.1 Application Section

The Application section is the section of the Flash that is used for storing the application code. The protection level for the Application section can be selected by the application Boot Lock bits (Boot Lock bits 0), see Table 29-2 on page 313. The Application section can never store any Boot Loader code since the SPM instruction is disabled when executed from the Application section.

### 29.2.2 BLS - Boot Loader Section

While the Application section is used for storing the application code, the The Boot Loader software must be located in the BLS since the SPM instruction can initiate a programming when executing from the BLS only. The SPM instruction can access the entire Flash, including the BLS itself. The protection level for the Boot Loader section can be selected by the Boot Loader Lock bits (Boot Lock bits 1), see Table 29-3 on page 313.

### 29.3 Read-While-Write and No Read-While-Write Flash Sections

Whether the CPU supports Read-While-Write or if the CPU is halted during a Boot Loader software update is dependent on which address that is being programmed. In addition to the two sections that are configurable by the BOOTSZ Fuses as described above, the Flash is also divided into two fixed sections, the Read-While-Write (RWW) section and the No Read-While-Write (NRWW) section. The limit between the RWW- and NRWW sections is given in Table 29-1 and Figure 29-1 on page 311. The main difference between the two sections is:

- When erasing or writing a page located inside the RWW section, the NRWW section can be read during the operation
- When erasing or writing a page located inside the NRWW section, the CPU is halted during the entire operation
Note that the user software can never read any code that is located inside the RWW section during a Boot Loader software operation. The syntax "Read-While-Write section" refers to which section that is being programmed (erased or written), not which section that actually is being read during a Boot Loader software update.


### 29.3.1 RWW - Read-While-Write Section

If a Boot Loader software update is programming a page inside the RWW section, it is possible to read code from the Flash, but only code that is located in the NRWW section. During an on-going programming, the software must ensure that the RWW section never is being read. If the user software is trying to read code that is located inside the RWW section (that is, by load program memory, call, or jump instructions or an interrupt) during programming, the software might end up in an unknown state. To avoid this, the interrupts should either be disabled or moved to the Boot Loader section. The Boot Loader section is always located in the NRWW section. The RWW Section Busy bit (RWWSB) in the Store Program Memory Control and Status Register (SPMCSR) will be read as logical one as long as the RWW section is blocked for reading. After a programming is completed, the RWWSB must be cleared by software before reading code located in the RWW section. See "SPMCSR - Store Program Memory Control and Status Register" on page 323. for details on how to clear RWWSB.

### 29.3.2 NRWW - No Read-While-Write Section

The code located in the NRWW section can be read when the Boot Loader software is updating a page in the RWW section. When the Boot Loader code updates the NRWW section, the CPU is halted during the entire Page Erase or Page Write operation.

Table 29-1. Read-While-Write Features

| Which Section does the Z-pointer <br> Address during the Programming? | Which Section can be <br> Read during Programming? | CPU Halted? | Read-While-Write <br> Supported? |
| :---: | :---: | :---: | :---: |
| RWW Section | NRWW Section | No | Yes |
| NRWW Section | None | Yes | No |

Figure 29-1. Read-While-Write vs. No Read-While-Write


Figure 29-2. Memory Sections


Note: 1. The parameters in the figure above are given in Table 29-7 on page 320.

### 29.4 Boot Loader Lock Bits

If no Boot Loader capability is needed, the entire Flash is available for application code. The Boot Loader has two separate sets of Boot Lock bits which can be set independently. This gives the user a unique flexibility to select different levels of protection.
The user can select:

- To protect the entire Flash from a software update by the MCU
- To protect only the Boot Loader Flash section from a software update by the MCU
- To protect only the Application Flash section from a software update by the MCU
- Allow software update in the entire Flash

See Table 29-2 on page 313 and Table 29-3 on page 313 for further details. The Boot Lock bits can be set in software and in Serial or Parallel Programming mode, but they can be cleared by a Chip Erase command only. The general Write Lock (Lock Bit mode 2) does not control the programming of the Flash memory by SPM instruction. Similarly, the general Read/Write Lock (Lock Bit mode 1) does not control reading nor writing by (E)LPM/SPM, if it is attempted.

Table 29-2. Boot Lock Bit0 Protection Modes (Application Section) ${ }^{(1)}$

| BLB0 Mode | BLB02 | BLB01 | Protection |
| :---: | :---: | :---: | :--- |
| 1 | 1 | 1 | No restrictions for SPM or (E)LPM accessing the Application section. |
| 2 | 1 | 0 | SPM is not allowed to write to the Application section. |
| 3 | 0 | 0 | SPM is not allowed to write to the Application section, and (E)LPM executing from <br> the Boot Loader section is not allowed to read from the Application section. If <br> Interrupt Vectors are placed in the Boot Loader section, interrupts are disabled while <br> executing from the Application section. |
| 4 | 0 | 1 | (E)LPM executing from the Boot Loader section is not allowed to read from the <br> Application section. If Interrupt Vectors are placed in the Boot Loader section, <br> interrupts are disabled while executing from the Application section. |

Note: 1. "1" means unprogrammed, " 0 " means programmed.
Table 29-3. Boot Lock Bit1 Protection Modes (Boot Loader Section) ${ }^{(1)}$

| BLB1 Mode | BLB12 | BLB11 | Protection |
| :---: | :---: | :---: | :--- |
| 1 | 1 | 1 | No restrictions for SPM or (E)LPM accessing the Boot Loader section. |
| 2 | 1 | 0 | SPM is not allowed to write to the Boot Loader section. |
| 3 | 0 | 0 | SPM is not allowed to write to the Boot Loader section, and (E)LPM executing from <br> the Application section is not allowed to read from the Boot Loader section. If <br> Interrupt Vectors are placed in the Application section, interrupts are disabled while <br> executing from the Boot Loader section. |
| 4 | 0 | 1 | (E)LPM executing from the Application section is not allowed to read from the Boot <br> Loader section. If Interrupt Vectors are placed in the Application section, interrupts <br> are disabled while executing from the Boot Loader section. |

Note: 1. "1" means unprogrammed, "0" means programmed.

### 29.4.1 Entering the Boot Loader Program

Entering the Boot Loader takes place by a jump or call from the application program. This may be initiated by a trigger such as a command received via USART, or SPI interface. Alternatively, the Boot Reset Fuse can be programmed so that the Reset Vector is pointing to the Boot Flash start address after a reset. In this case, the Boot Loader is started after a reset. After the application code is loaded, the program can start executing the application code. Note that the fuses cannot be changed by the MCU itself. This means that once the Boot Reset Fuse is programmed, the Reset Vector will always point to the Boot Loader Reset and the fuse can only be changed through the serial or parallel programming interface.

Table 29-4. Boot Reset Fuse ${ }^{(1)}$

| BOOTRST | Reset Address |
| :---: | :--- |
| 1 | Reset Vector $=$ Application Reset (address 0x0000) |
| 0 | Reset Vector = Boot Loader Reset (see Table 29-7 on page 320) |

Note: 1. "1" means unprogrammed, " 0 " means programmed.

### 29.5 Addressing the Flash During Self-Programming

The Z-pointer is used to address the SPM commands. The Z pointer consists of the Z-registers ZL and ZH in the register file, and RAMPZ in the I/O space. The number of bits actually used is implementation dependent. Note that the RAMPZ register is only implemented when the program space is larger than 64Kbytes.

| Bit | 23 | 22 | 21 | 20 | 19 | 18 | 17 | 16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 |
| RAMPZ | RAMPZ7 | RAMPZ6 | RAMPZ5 | RAMPZ4 | RAMPZ3 | RAMPZ2 | RAMPZ1 | RAMPZ0 |
| ZH (R31) | Z15 | Z14 | Z13 | Z12 | Z11 | Z10 | Z9 | Z8 |
| ZL (R30) | Z7 | Z6 | Z5 | Z4 | Z3 | Z2 | Z1 | Z0 |
|  | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |

Since the Flash is organized in pages (see Table 30-7 on page 328), the Program Counter can be treated as having two different sections. One section, consisting of the least significant bits, is addressing the words within a page, while the most significant bits are addressing the pages. This is shown in Figure 29-3. Note that the Page Erase and Page Write operations are addressed independently. Therefore it is of major importance that the Boot Loader software addresses the same page in both the Page Erase and Page Write operation. Once a programming operation is initiated, the address is latched and the Z-pointer can be used for other operations.

The (E)LPM instruction use the Z-pointer to store the address. Since this instruction addresses the Flash byte-bybyte, also bit ZO of the Z-pointer is used.

Figure 29-3. Addressing the Flash During SPM ${ }^{(1)}$


Note: 1. The different variables used in Figure 29-3 are listed in Table 29-9 on page 320.

### 29.6 Self-Programming the Flash

The program memory is updated in a page by page fashion. Before programming a page with the data stored in the temporary page buffer, the page must be erased. The temporary page buffer is filled one word at a time using SPM and the buffer can be filled either before the Page Erase command or between a Page Erase and a Page Write operation:
Alternative 1, fill the buffer before a Page Erase

- Fill temporary page buffer
- Perform a Page Erase
- Perform a Page Write

Alternative 2, fill the buffer after Page Erase

- Perform a Page Erase
- Fill temporary page buffer
- Perform a Page Write

If only a part of the page needs to be changed, the rest of the page must be stored (for example in the temporary page buffer) before the erase, and then be rewritten. When using alternative 1, the Boot Loader provides an effective Read-Modify-Write feature which allows the user software to first read the page, do the necessary changes, and then write back the modified data. If alternative 2 is used, it is not possible to read the old data while loading since the page is already erased. The temporary page buffer can be accessed in a random sequence. It is essential that the page address used in both the Page Erase and Page Write operation is addressing the same page. See "Simple Assembly Code Example for a Boot Loader" on page 318 for an assembly code example.

### 29.6.1 Performing Page Erase by SPM

To execute Page Erase, set up the address in the Z-pointer, write "X0000011" to SPMCSR and execute SPM within four clock cycles after writing SPMCSR. The data in R1 and R0 is ignored. The page address must be written to PCPAGE in the Z-register. Other bits in the Z-pointer will be ignored during this operation.

- Page Erase to the RWW section: The NRWW section can be read during the Page Erase
- Page Erase to the NRWW section: The CPU is halted during the operation


### 29.6.2 Filling the Temporary Buffer (Page Loading)

To write an instruction word, set up the address in the Z-pointer and data in R1:R0, write "00000001" to SPMCSR and execute SPM within four clock cycles after writing SPMCSR. The content of PCWORD in the Z-register is used to address the data in the temporary buffer. The temporary buffer will auto-erase after a Page Write operation or by writing the RWWSRE bit in SPMCSR. It is also erased after a system reset. Note that it is not possible to write more than one time to each address without erasing the temporary buffer.
If the EEPROM is written in the middle of an SPM Page Load operation, all data loaded is still buffered.

### 29.6.3 Performing a Page Write

To execute Page Write, set up the address in the Z-pointer, write "X0000101" to SPMCSR and execute SPM within four clock cycles after writing SPMCSR. The data in R1 and R0 is ignored. The page address must be written to PCPAGE. Other bits in the Z-pointer must be written to zero during this operation.

- Page Write to the RWW section: The NRWW section can be read during the Page Write
- Page Write to the NRWW section: The CPU is halted during the operation


### 29.6.4 Using the SPM Interrupt

If the SPM interrupt is enabled, the SPM interrupt will generate a constant interrupt when the SPMEN bit in SPMCSR is cleared. This means that the interrupt can be used instead of polling the SPMCSR Register in soft-
ware. When using the SPM interrupt, the Interrupt Vectors should be moved to the BLS section to avoid that an interrupt is accessing the RWW section when it is blocked for reading. How to move the interrupts is described in "Interrupts" on page 101.

### 29.6.5 Consideration While Updating BLS

Special care must be taken if the user allows the Boot Loader section to be updated by leaving Boot Lock bit11 unprogrammed. An accidental write to the Boot Loader itself can corrupt the entire Boot Loader, and further software updates might be impossible. If it is not necessary to change the Boot Loader software itself, it is recommended to program the Boot Lock bit11 to protect the Boot Loader software from any internal software changes.

### 29.6.6 Prevent Reading the RWW Section During Self-Programming

During Self-Programming (either Page Erase or Page Write), the RWW section is always blocked for reading. The user software itself must prevent that this section is addressed during the self programming operation. The RWWSB in the SPMCSR will be set as long as the RWW section is busy. During Self-Programming the Interrupt Vector table should be moved to the BLS as described in "Interrupts" on page 101, or the interrupts must be disabled. Before addressing the RWW section after the programming is completed, the user software must clear the RWWSB by writing the RWWSRE. See "Simple Assembly Code Example for a Boot Loader" on page 318 for an example.

### 29.6.7 Setting the Boot Loader Lock Bits by SPM

To set the Boot Loader Lock bits and general Lock bits, write the desired data to R0, write "X0001001" to SPMCSR and execute SPM within four clock cycles after writing SPMCSR.

Bit
R0

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{1}$ | $\mathbf{1}$ | BLB12 | BLB11 | BLB02 | BLB01 | LB2 | LB1 |

See Table 29-2 on page 313 and Table 29-3 on page 313 for how the different settings of the Boot Loader bits affect the Flash access.

If bits 5:0 in R0 are cleared (zero), the corresponding Lock bit will be programmed if an SPM instruction is executed within four cycles after BLBSET and SPMEN are set in SPMCSR. The Z-pointer is don't care during this operation, but for future compatibility it is recommended to load the Z-pointer with $0 \times 0001$ (same as used for reading the $\mathrm{IO}_{\mathrm{ck}}$ bits). For future compatibility it is also recommended to set bits 7 and 6 in R0 to " 1 " when writing the Lock bits. When programming the Lock bits the entire Flash can be read during the operation.

### 29.6.8 EEPROM Write Prevents Writing to SPMCSR

Note that an EEPROM write operation will block all software programming to Flash. Reading the Fuses and Lock bits from software will also be prevented during the EEPROM write operation. It is recommended that the user checks the status bit (EEPE) in the EECR Register and verifies that the bit is cleared before writing to the SPMCSR Register.

### 29.6.9 Reading the Fuse and Lock Bits from Software

It is possible to read both the Fuse and Lock bits from software. To read the Lock bits, load the Z-pointer with $0 \times 0001$ and set the BLBSET and SPMEN bits in SPMCSR. When an (E)LPM instruction is executed within three CPU cycles after the BLBSET and SPMEN bits are set in SPMCSR, the value of the Lock bits will be loaded in the destination register. The BLBSET and SPMEN bits will auto-clear upon completion of reading the Lock bits or if no (E)LPM instruction is executed within three CPU cycles or no SPM instruction is executed within four CPU cycles. When BLBSET and SPMEN are cleared, (E)LPM will work as described in the Instruction set Manual.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Rd | - | - | BLB12 | BLB11 | BLB02 | BLB01 | LB2 | LB1 |

The algorithm for reading the Fuse Low byte is similar to the one described above for reading the Lock bits. To read the Fuse Low byte, load the Z-pointer with $0 x 0000$ and set the BLBSET and SPMEN bits in SPMCSR. When an (E)LPM instruction is executed within three cycles after the BLBSET and SPMEN bits are set in the SPMCSR, the value of the Fuse Low byte (FLB) will be loaded in the destination register as shown below. Refer to Table 305 on page 327 for a detailed description and mapping of the Fuse Low byte.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Rd | FLB7 | FLB6 | FLB5 | FLB4 | FLB3 | FLB2 | FLB1 | FLB0 |

Similarly, when reading the Fuse High byte, load 0x0003 in the Z-pointer. When an (E)LPM instruction is executed within three cycles after the BLBSET and SPMEN bits are set in the SPMCSR, the value of the Fuse High byte (FHB) will be loaded in the destination register as shown below. Refer to Table 30-4 on page 327 for detailed description and mapping of the Fuse High byte.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Rd | FHB7 | FHB6 | FHB5 | FHB4 | FHB3 | FHB2 | FHB1 | FHB0 |

When reading the Extended Fuse byte, load $0 x 0002$ in the Z-pointer. When an (E)LPM instruction is executed within three cycles after the BLBSET and SPMEN bits are set in the SPMCSR, the value of the Extended Fuse byte (EFB) will be loaded in the destination register as shown below. Refer to Table 30-3 on page 326 for detailed description and mapping of the Extended Fuse byte.


Fuse and Lock bits that are programmed, will be read as zero. Fuse and Lock bits that are unprogrammed, will be read as one.

### 29.6.10 Reading the Signature Row from Software

To read the Signature Row from software, load the Z-pointer with the signature byte address given in Table 29-5 on page 317 and set the SIGRD and SPMEN bits in SPMCSR. When an LPM instruction is executed within three CPU cycles after the SIGRD and SPMEN bits are set in SPMCSR, the signature byte value will be loaded in the destination register. The SIGRD and SPMEN bits will auto-clear upon completion of reading the Signature Row Lock bits or if no LPM instruction is executed within three CPU cycles. When SIGRD and SPMEN are cleared, LPM will work as described in the Instruction set Manual.

Table 29-5. $\quad$ Signature Row Addressing

| Signature Byte | Z-Pointer Address |
| :--- | :---: |
| Device Signature Byte 1 | $0 \times 0000$ |
| Device Signature Byte 2 | $0 \times 0002$ |
| Device Signature Byte 3 | $0 \times 0004$ |
| RC Oscillator Calibration Byte | $0 \times 0001$ |

Note: All other addresses are reserved for future use.

### 29.6.11 Preventing Flash Corruption

During periods of low $\mathrm{V}_{\mathrm{CC}}$, the Flash program can be corrupted because the supply voltage is too low for the CPU and the Flash to operate properly. These issues are the same as for board level systems using the Flash, and the same design solutions should be applied.

A Flash program corruption can be caused by two situations when the voltage is too low. First, a regular write sequence to the Flash requires a minimum voltage to operate correctly. Secondly, the CPU itself can execute instructions incorrectly, if the supply voltage for executing instructions is too low.

Flash corruption can easily be avoided by following these design recommendations (one is sufficient):

1. If there is no need for a Boot Loader update in the system, program the Boot Loader Lock bits to prevent any Boot Loader software updates.
2. Keep the AVR RESET active (low) during periods of insufficient power supply voltage. This can be done by enabling the internal Brown-out Detector (BOD) if the operating voltage matches the detection level. If not, an external low $\mathrm{V}_{\mathrm{cc}}$ reset protection circuit can be used. If a reset occurs while a write operation is in progress, the write operation will be completed provided that the power supply voltage is sufficient.
3. Keep the AVR core in Power-down sleep mode during periods of low $\mathrm{V}_{\mathrm{cc}}$. This will prevent the CPU from attempting to decode and execute instructions, effectively protecting the SPMCSR Register and thus the Flash from unintentional writes.

### 29.6.12 Programming Time for Flash when Using SPM

The calibrated RC Oscillator is used to time Flash accesses. Table 29-6 shows the typical programming time for Flash accesses from the CPU.

Table 29-6. SPM Programming Time

| Symbol | Min Programming Time | Max Programming Time |
| :---: | :---: | :---: |
| Flash write (Page Erase, Page Write, and write Lock bits by SPM) | 3.7 ms | 4.5 ms |

### 29.6.13 Simple Assembly Code Example for a Boot Loader

```
    ;-the routine writes one page of data from RAM to Flash
    ; the first data location in RAM is pointed to by the Y pointer
    ; the first data location in Flash is pointed to by the Z-pointer
    ;-error handling is not included
    ;-the routine must be placed inside the Boot space
    ; (at least the Do_spm sub routine). Only code inside NRWW section can
    ; be read during Self-Programming (Page Erase and Page Write).
    ;-registers used: r0, r1, temp1 (r16), temp2 (r17), looplo (r24),
    ; loophi (r25), spmcrval (r20)
    ; storing and restoring of registers is not included in the routine
    ; register usage can be optimized at the expense of code size
    ;-It is assumed that either the interrupt table is moved to the Boot
    ; loader section or that the interrupts are disabled.
.equ PAGESIZEB = PAGESIZE*2 ;PAGESIZEB is page size in BYTES, not words
.org SMALLBOOTSTART
Write_page:
    ; Page Erase
    ldi spmcrval, (1<<PGERS) | (1<<SPMEN)
    call Do_spm
    ; re-enable the RWW section
    ldi spmcrval, (1<<RWWSRE) | (1<<SPMEN)
    call Do_spm
    ; transfer data from RAM to Flash page buffer
    ldi looplo, low(PAGESIZEB) ;init loop variable
    ldi loophi, high(PAGESIZEB) ; not required for PAGESIZEB<=256
Wrloop:
    ld r0, Y+
    ld r1, Y+
    ldi spmcrval, (1<<SPMEN)
    call Do_spm
    adiw ZH:ZL, 2
    sbiw loophi:looplo, 2 ;use subi for PAGESIZEB<=256
    brne Wrloop
```

```
    ; execute Page Write
    subi ZL, low(PAGESIZEB) ;restore pointer
    sbci ZH, high(PAGESIZEB) ; not required for PAGESIZEB<=256
    ldi spmcrval, (1<<PGWRT) | (1<<SPMEN)
    call Do_spm
    ; re-enable the RWW section
    ldi spmcrval, ( }1<<\mathrm{ RWWSRE) | ( }1<<\mathrm{ SPMEN)
    call Do_spm
    ; read back and check, optional
    ldi looplo, low(PAGESIZEB) ;init loop variable
    ldi loophi, high(PAGESIZEB) ; not required for PAGESIZEB<=256
    subi YL, low(PAGESIZEB) ;restore pointer
    sbci YH, high(PAGESIZEB)
Rdloop:
    elpm r0, Z+
    ld r1, Y+
    cpse r0, r1
    jmp Error
    sbiw loophi:looplo, 1 ;use subi for PAGESIZEB<=256
    brne Rdloop
    ; return to RWW section
    ; verify that RWW section is safe to read
Return:
    in temp1, SPMCSR
    sbrs temp1, RWWSB ; If RWWSB is set, the RWW section is not ready yet
    ret
    ; re-enable the RWW section
    ldi spmcrval, ( }1<<\mathrm{ RWWSRE) | ( }1<<<SPMEN
    call Do_spm
    rjmp Return
Do_spm:
    ; check for previous SPM complete
Wait_spm:
    in temp1, SPMCSR
    sbrc temp1, SPMEN
    rjmp Wait_spm
    ; input: spmcrval determines SPM action
    ; disable interrupts if enabled, store status
    in temp2, SREG
    cli
    ; check that no EEPROM write access is present
Wait_ee:
    sbic EECR, EEPE
    rjmp Wait_ee
    ; SPM timed sequence
    out SPMCSR, spmcrval
    spm
    ; restore SREG (to enable interrupts if originally enabled)
    out SREG, temp2
    ret
```


### 29.6.14 ATmega640 Boot Loader Parameters

In Table 29-7 through Table 29-9 on page 320, the parameters used in the description of the Self-Programming are given.

Table 29-7. Boot Size Configuration, ATmega640 ${ }^{(1)}$

| N N O O 0 | $\begin{aligned} & \text { O} \\ & \text { N } \\ & \text { O } \\ & \text { O } \end{aligned}$ | $\begin{aligned} & \text { O } \\ & \stackrel{N}{0} \\ & \stackrel{0}{0} \\ & \dot{0} \end{aligned}$ | ® \% ¢ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 512 words | 4 | 0x0000-0x7DFF | 0x7E00-0x7FFF | 0x7DFF | 0x7E00 |
| 1 | 0 | 1024 words | 8 | 0x0000-0x7BFF | 0x7C00-0x7FFF | 0x7BFF | 0x7C00 |
| 0 | 1 | 2048 words | 16 | 0x0000-0x77FF | 0x7800-0x7FFF | 0x77FF | 0x7800 |
| 0 | 0 | 4096 words | 32 | 0x0000-0x6FFF | 0x7000-0x7FFF | 0x6FFF | 0x7000 |

Note: 1. The different BOOTSZ Fuse configurations are shown in Figure 29-2 on page 312.

Table 29-8. Read-While-Write Limit, ATmega640

| Section $^{(1)}$ | Pages | Address |
| :--- | :--- | :--- |
| Read-While-Write section (RWW) | 224 | $0 \times 0000-0 \times 6 F F F$ |
| No Read-While-Write section (NRWW) | 32 | $0 \times 7000-0 \times 7 F F F$ |

Note: 1. For details about these two section, see "NRWW - No Read-While-Write Section" on page 311 and "RWW - Read-While-Write Section" on page 311.

Table 29-9. Explanation of different variables used in Figure 29-3 on page 314 and the mapping to the Z-pointer, ATmega640

| Variable |  | Corresponding Z-value ${ }^{(2)}$ | Description ${ }^{(1)}$ |
| :---: | :---: | :---: | :---: |
| PCMSB | 14 |  | Most significant bit in the Program Counter. (The Program Counter is 15 bits PC[14:0]). |
| PAGEMSB | 6 |  | Most significant bit which is used to address the words within one page (128 words in a page requires seven bits PC [6:0]). |
| ZPCMSB |  | Z15 | Bit in Z-pointer that is mapped to PCMSB. Because Z0 is not used, the ZPCMSB equals PCMSB +1 . |
| ZPAGEMSB |  | Z7 | Bit in Z-pointer that is mapped to PCMSB. Because Z0 is not used, the ZPAGEMSB equals PAGEMSB +1 . |
| PCPAGE | PC[14:7] | Z15:Z8 | Program Counter page address: Page select, for Page Erase and Page Write. |
| PCWORD | PC[6:0] | Z7:Z1 | Program Counter word address: Word select, for filling temporary buffer (must be zero during Page Write operation). |

Note: 1. Z0: should be zero for all SPM commands, byte select for the (E)LPM instruction.
2. See "Addressing the Flash During Self-Programming" on page 314 for details about the use of Z-pointer during Self-Programming.

### 29.6.15 ATmega1280/1281 Boot Loader Parameters

In Table 29-10 and Table 29-11, the parameters used in the description of the Self-Programming are given.
Table 29-10. Boot Size Configuration, ATmega1280/1281 ${ }^{(1)}$

| $\begin{aligned} & \bar{N} \\ & \text { N } \\ & \text { O } \\ & 0 \end{aligned}$ | $\begin{aligned} & \text { O} \\ & \text { N } \\ & \text { O } \\ & \mathbf{o} \end{aligned}$ | $\begin{aligned} & \text { N } \\ & \text { N } \\ & \text { + } \\ & \text { O } \end{aligned}$ | ® ® ¢ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 512 words | 4 | 0x0000-0xFDFF | 0xFE00-0xFFFF | 0xFDFF | 0xFE00 |
| 1 | 0 | 1024 words | 8 | 0x0000-0xFBFF | 0xFC00-0xFFFF | 0xFBFF | 0xFC00 |
| 0 | 1 | 2048 words | 16 | 0x0000-0xF7FF | 0xF800-0xFFFFF | 0xF7FF | 0xF800 |
| 0 | 0 | 4096 words | 32 | 0x0000-0xEFFFF | 0xF000-0xFFFFF | 0xEFFF | 0xF000 |

Note: 1. The different BOOTSZ Fuse configurations are shown in Figure 29-2 on page 312.

Table 29-11. Read-While-Write Limit, ATmega1280/1281

| Section $^{(1)}$ | Pages | Address |
| :--- | :---: | :--- |
| Read-While-Write section (RWW) | 480 | 0x0000-0xEFFF |
| No Read-While-Write section (NRWW) | 32 | 0xF000 - 0xFFFF |

Note: 1. For details about these two section, see "NRWW - No Read-While-Write Section" on page 311 and "RWW - Read-While-Write Section" on page 311.

Table 29-12. Explanation of different variables used in Figure 29-3 on page 314 and the mapping to the Z-pointer, ATmega1280/1281

| Variable |  | Corresponding Z-value ${ }^{(2)}$ | Description ${ }^{(1)}$ |
| :---: | :---: | :---: | :---: |
| PCMSB | 15 |  | Most significant bit in the Program Counter. (The Program Counter is 16 bits PC[15:0]) |
| PAGEMSB | 6 |  | Most significant bit which is used to address the words within one page (128 words in a page requires seven bits PC [6:0]). |
| ZPCMSB |  | Z16 ${ }^{(3)}$ | Bit in Z-pointer that is mapped to PCMSB. Because Z0 is not used, the ZPCMSB equals PCMSB +1 . |
| ZPAGEMSB |  | Z7 | Bit in Z-pointer that is mapped to PCMSB. Because Z0 is not used, the ZPAGEMSB equals PAGEMSB +1 . |
| PCPAGE | $\mathrm{PC}[15: 7]$ | Z16 ${ }^{(3)}$ : $\mathrm{Z8}$ | Program Counter page address: Page select, for Page Erase and Page Write |
| PCWORD | $\mathrm{PC}[6: 0]$ | Z7:Z1 | Program Counter word address: Word select, for filling temporary buffer (must be zero during Page Write operation) |

[^1]
### 29.6.16 ATmega2560/2561 Boot Loader Parameters

In Table 29-13 through Table 29-15, the parameters used in the description of the Self-Programming are given.
Table 29-13. Boot Size Configuration, ATmega2560/2561 ${ }^{(1)}$

| $\begin{aligned} & \bar{N} \\ & \text { N } \\ & \mathbf{O} \\ & 0 \end{aligned}$ | $\begin{aligned} & \text { O } \\ & \text { N } \\ & \text { O } \\ & \text { O} \end{aligned}$ | $\begin{aligned} & \text { N } \\ & \stackrel{N}{\mathbf{N}} \\ & \stackrel{0}{\mathbf{O}} \\ & \mathbf{0} \end{aligned}$ | - |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 512 words | 4 | 0x00000-0x1FDFF | 0x1FE00-0x1FFFF | 0x1FDFF | 0x1FE00 |
| 1 | 0 | 1024 words | 8 | 0x00000-0x1FBFF | 0x1FC00-0x1FFFF | 0x1FBFF | 0x1FC00 |
| 0 | 1 | 2048 words | 16 | 0x00000-0x1F7FF | 0x1F800-0x1FFFF | 0x1F7FF | 0x1F800 |
| 0 | 0 | 4096 words | 32 | 0x00000-0x1EFFF | 0x1F000-0x1FFFF | 0x1EFFF | 0x1F000 |

Note: 1. The different BOOTSZ Fuse configurations are shown in Figure 29-2 on page 312.

Table 29-14. Read-While-Write Limit, ATmega2560/2561

| Section $^{(1)}$ | Pages | Address |
| :--- | :---: | :--- |
| Read-While-Write section (RWW) | 992 | $0 \times 00000-0 \times 1 \mathrm{EFFF}$ |
| No Read-While-Write section (NRWW) | 32 | $0 \times 1 F 000-0 \times 1 F F F F$ |

Note: 1. For details about these two section, see "NRWW - No Read-While-Write Section" on page 311 and "RWW - Read-While-Write Section" on page 311.

Table 29-15. Explanation of different variables used in Figure 29-3 on page 314 and the mapping to the Z-pointer, ATmega2560/2561

| Variable |  | Corresponding Z-value ${ }^{(2)}$ | Description ${ }^{(1)}$ |
| :---: | :---: | :---: | :---: |
| PCMSB | 16 |  | Most significant bit in the Program Counter. (The Program Counter is 17 bits PC[16:0]). |
| PAGEMSB | 6 |  | Most significant bit which is used to address the words within one page (128 words in a page requires seven bits PC [6:0]). |
| ZPCMSB |  | Z17:Z16 ${ }^{(3)}$ | Bit in Z-pointer that is mapped to PCMSB. Because Z0 is not used, the ZPCMSB equals PCMSB +1 . |
| ZPAGEMSB |  | Z7 | Bit in Z-pointer that is mapped to PCMSB. Because Z0 is not used, the ZPAGEMSB equals PAGEMSB +1 . |
| PCPAGE | PC[16:7] | Z17 ${ }^{(3)}$ : $\mathrm{Z8}$ | Program Counter page address: Page select, for Page Erase and Page Write. |
| PCWORD | PC[6:0] | Z7:Z1 | Program Counter word address: Word select, for filling temporary buffer (must be zero during Page Write operation). |

[^2]
### 29.7 Register Description

### 29.7.1 SPMCSR - Store Program Memory Control and Status Register

The Store Program Memory Control and Status Register contains the control bits needed to control the Boot Loader operations.

| Bit | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | SPMCSR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x37 (0x57) | SPMIE | RWWSB | SIGRD | RWWSRE | BLBSET | PGWRT | PGERS | SPMEN |  |
| Read/Write | R/W | R | R/W | R/W | R/W | R/W | R/W | R/W |  |
| Initial Value | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |

## - Bit 7 - SPMIE: SPM Interrupt Enable

When the SPMIE bit is written to one, and the I-bit in the Status Register is set (one), the SPM ready interrupt will be enabled. The SPM ready Interrupt will be executed as long as the SPMEN bit in the SPMCSR Register is cleared.

## - Bit 6 - RWWSB: Read-While-Write Section Busy

When a Self-Programming (Page Erase or Page Write) operation to the RWW section is initiated, the RWWSB will be set (one) by hardware. When the RWWSB bit is set, the RWW section cannot be accessed. The RWWSB bit will be cleared if the RWWSRE bit is written to one after a Self-Programming operation is completed. Alternatively the RWWSB bit will automatically be cleared if a page load operation is initiated.

## - Bit 5 - SIGRD: Signature Row Read

If this bit is written to one at the same time as SPMEN, the next LPM instruction within three clock cycles will read a byte from the signature row into the destination register. see "Reading the Signature Row from Software" on page 317 for details. An SPM instruction within four cycles after SIGRD and SPMEN are set will have no effect. This operation is reserved for future use and should not be used.

- Bit 4 - RWWSRE: Read-While-Write Section Read Enable

When programming (Page Erase or Page Write) to the RWW section, the RWW section is blocked for reading (the RWWSB will be set by hardware). To re-enable the RWW section, the user software must wait until the programming is completed (SPMEN will be cleared). Then, if the RWWSRE bit is written to one at the same time as SPMEN, the next SPM instruction within four clock cycles re-enables the RWW section. The RWW section cannot be re-enabled while the Flash is busy with a Page Erase or a Page Write (SPMEN is set). If the RWWSRE bit is written while the Flash is being loaded, the Flash load operation will abort and the data loaded will be lost.

## - Bit 3 - BLBSET: Boot Lock Bit Set

If this bit is written to one at the same time as SPMEN, the next SPM instruction within four clock cycles sets Boot Lock bits, according to the data in R0. The data in R1 and the address in the Z-pointer are ignored. The BLBSET bit will automatically be cleared upon completion of the Lock bit set, or if no SPM instruction is executed within four clock cycles.
An (E)LPM instruction within three cycles after BLBSET and SPMEN are set in the SPMCSR Register, will read either the Lock bits or the Fuse bits (depending on ZO in the Z-pointer) into the destination register. See "Reading the Fuse and Lock Bits from Software" on page 316 for details.

## - Bit 2 - PGWRT: Page Write

If this bit is written to one at the same time as SPMEN, the next SPM instruction within four clock cycles executes Page Write, with the data stored in the temporary buffer. The page address is taken from the high part of the Zpointer. The data in R1 and R0 are ignored. The PGWRT bit will auto-clear upon completion of a Page Write, or if no SPM instruction is executed within four clock cycles. The CPU is halted during the entire Page Write operation if the NRWW section is addressed.

## - Bit 1 - PGERS: Page Erase

If this bit is written to one at the same time as SPMEN, the next SPM instruction within four clock cycles executes Page Erase. The page address is taken from the high part of the Z-pointer. The data in R1 and R0 are ignored. The PGERS bit will auto-clear upon completion of a Page Erase, or if no SPM instruction is executed within four clock cycles. The CPU is halted during the entire Page Write operation if the NRWW section is addressed.

## - Bit 0 - SPMEN: Store Program Memory Enable

This bit enables the SPM instruction for the next four clock cycles. If written to one together with either RWWSRE, BLBSET, PGWRT' or PGERS, the following SPM instruction will have a special meaning, see description above. If only SPMEN is written, the following SPM instruction will store the value in R1:R0 in the temporary page buffer addressed by the Z-pointer. The LSB of the Z-pointer is ignored. The SPMEN bit will auto-clear upon completion of an SPM instruction, or if no SPM instruction is executed within four clock cycles. During Page Erase and Page Write, the SPMEN bit remains high until the operation is completed.

Writing any other combination than "10001", "01001", "00101", "00011" or "00001" in the lower five bits will have no effect.

Note: Only one SPM instruction should be active at any time.

## 30. Memory Programming

### 30.1 Program And Data Memory Lock Bits

The ATmega640/1280/1281/2560/2561 provides six Lock bits which can be left unprogrammed ("1") or can be programmed ("0") to obtain the additional features listed in Table 30-2. The Lock bits can only be erased to " 1 " with the Chip Erase command.

Table 30-1. Lock Bit Byte ${ }^{(1)}$

| Lock Bit Byte | Bit No | Description | Default Value |
| :--- | :---: | :--- | :--- |
|  | 7 | - | 1 (unprogrammed) |
|  | 6 | - | 1 (unprogrammed) |
| BLB12 | 5 | Boot Lock bit | 1 (unprogrammed) |
| BLB11 | 4 | Boot Lock bit | 1 (unprogrammed) |
| BLB02 | 3 | Boot Lock bit | 1 (unprogrammed) |
| BLB01 | 2 | Boot Lock bit | 1 (unprogrammed) |
| LB2 | 1 | Lock bit | 1 (unprogrammed) |
| LB1 | 0 | Lock bit | 1 (unprogrammed) |

Note: 1. " 1 " means unprogrammed, "0" means programmed
Table 30-2. Lock Bit Protection Modes ${ }^{(1)(2)}$

| Memory Lock Bits |  | Protection Type |  |
| :---: | :---: | :---: | :--- |
| LB Mode | LB2 | LB1 |  |
| 1 | 1 | 1 | No memory lock features enabled. |
| 2 | 1 | 0 | Further programming of the Flash and EEPROM is disabled in Parallel and Serial <br> Programming mode. The Fuse bits are locked in both Serial and Parallel <br> Programming mode. |
| (1) |  |  |  |

Table 30-2. Lock Bit Protection Modes ${ }^{(1)(2)}$ (Continued)

| Memory Lock Bits |  |  | Protection Type |
| :---: | :---: | :---: | :--- |
| BLB1 Mode | BLB12 | BLB11 |  |
| 1 | 1 | 1 | No restrictions for SPM or (E)LPM accessing the Boot Loader section. |
| 2 | 1 | 0 | SPM is not allowed to write to the Boot Loader section. |
| 3 | 0 | 0 | SPM is not allowed to write to the Boot Loader section, and (E)LPM executing from <br> the Application section is not allowed to read from the Boot Loader section. If <br> Interrupt Vectors are placed in the Application section, interrupts are disabled while <br> executing from the Boot Loader section. |
| 4 | 0 | 1 | (E)LPM executing from the Application section is not allowed to read from the Boot <br> Loader section. If Interrupt Vectors are placed in the Application section, interrupts <br> are disabled while executing from the Boot Loader section. |

Notes: 1. Program the Fuse bits and Boot Lock bits before programming the LB1 and LB2.
2. "1" means unprogrammed, " 0 " means programmed.

### 30.2 Fuse Bits

The ATmega640/1280/1281/2560/2561 has three Fuse bytes. Table 30-3 through Table 30-5 on page 327 describe briefly the functionality of all the fuses and how they are mapped into the Fuse bytes. Note that the fuses are read as logical zero, "0", if they are programmed.

Table 30-3. Extended Fuse Byte

| Extended Fuse Byte | Bit No | Description | Default Value |
| :---: | :---: | :---: | :---: |
| - | 7 | - | 1 |
| - | 6 | - | 1 |
| - | 5 | - | 1 |
| - | 4 | - | 1 |
| - | 3 | Brown-out Detector trigger level | 1 |
| BODLEVEL2 $^{(1)}$ | 2 | Brown-out Detector trigger level | 1 (unprogrammed) |
| BODLEVEL1 $^{(1)}$ | 1 | Brown-out Detector trigger level | 1 (unprogrammed) |
| BODLEVEL0 $^{(1)}$ | 0 | 1 (unprogrammed) |  |

Note: 1. See "System and Reset Characteristics" on page 360 for BODLEVEL Fuse decoding.

Table 30-4. Fuse High Byte

| Fuse High Byte | Bit No | Description | Default Value |
| :--- | :---: | :--- | :--- |
| OCDEN $^{(4)}$ | 7 | Enable OCD | 1 (unprogrammed, OCD disabled) |
| JTAGEN $^{(1)}$ | 6 | Enable JTAG | 0 (programmed, JTAG enabled) |
| SPIEN $^{(1)}$ | 5 | Enable Serial Program and Data Downloading | 0 (programmed, SPI prog. enabled) |
| WDTON $^{(3)}$ | 4 | Watchdog Timer always on | 1 (unprogrammed) |
| EESAVE | 3 | EEPROM memory is preserved through the Chip <br> Erase | 1 (unprogrammed, EEPROM not <br> preserved) |
| BOOTSZ1 | 2 | Select Boot Size (see Table 30-9 on page 329 for <br> details) | 0 (programmed) ${ }^{(2)}$ |
| BOOTSZ0 | 1 | Select Boot Size (see Table 30-9 on page 329 for <br> details) | 0 (programmed) ${ }^{(2)}$ |
| BOOTRST | 0 | Select Reset Vector | 1 (unprogrammed) |

Notes: 1. The SPIEN Fuse is not accessible in serial programming mode.
2. The default value of BOOTSZ1:0 results in maximum Boot Size. See Table 29-7 on page 320 for details.
3. See "WDTCSR - Watchdog Timer Control Register" on page 65 for details.
4. Never ship a product with the OCDEN Fuse programmed regardless of the setting of Lock bits and JTAGEN Fuse. A programmed OCDEN Fuse enables some parts of the clock system to be running in all sleep modes. This may increase the power consumption.
Table 30-5. Fuse Low Byte

| Fuse Low Byte | Bit No | Description | Default Value |
| :--- | :---: | :--- | :--- |
| CKDIV8 $^{(4)}$ | 7 | Divide clock by 8 | 0 (programmed) |
| CKOUT $^{(3)}$ | 6 | Clock output | 1 (unprogrammed) |
| SUT1 $^{\text {SUT0 }}$ | 5 | Select start-up time | 1 (unprogrammed) ${ }^{(1)}$ |
| CKSEL3 | 4 | Select start-up time | 0 (programmed) $^{(1)}$ |
| CKSEL2 | 3 | Select Clock source | 0 (programmed) $^{(2)}$ |
| CKSEL1 | 2 | Select Clock source | 0 (programmed) $^{(2)}$ |
| CKSEL0 | 0 | Select Clock source | 1 (unprogrammed) $^{(2)}$ |

Notes: 1. The default value of SUT1:0 results in maximum start-up time for the default clock source. See "System and Reset Characteristics" on page 360 for details.
2. The default setting of CKSEL3:0 results in internal RC Oscillator @ 8 MHz . See Table 10-1 on page 40 for details.
3. The CKOUT Fuse allow the system clock to be output on PORTE7. See "Clock Output Buffer" on page 47 for details.
4. See "System Clock Prescaler" on page 47 for details.

The status of the Fuse bits is not affected by Chip Erase. Note that the Fuse bits are locked if Lock bit1 (LB1) is programmed. Program the Fuse bits before programming the Lock bits.

### 30.2.1 Latching of Fuses

The fuse values are latched when the device enters programming mode and changes of the fuse values will have no effect until the part leaves Programming mode. This does not apply to the EESAVE Fuse which will take effect once it is programmed. The fuses are also latched on Power-up in Normal mode.

### 30.3 Signature Bytes

All Atmel microcontrollers have a three-byte signature code which identifies the device. This code can be read in both serial and parallel mode, also when the device is locked. The three bytes reside in a separate address space. For the ATmega640/1280/1281/2560/2561 the signature bytes are given in Table 30-6.

Table 30-6. Device and JTAG ID

| Part | Signature Bytes Address |  | JTAG |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: |
|  | $0 \times 000$ | $0 \times 001$ | $\mathbf{0 x 0 0 2}$ | Part Number | Manufacture ID |
| ATmega640 | $0 \times 1 \mathrm{E}$ | $0 \times 96$ | $0 \times 08$ | 9608 | $0 \times 1 \mathrm{~F}$ |
| ATmega1280 | $0 \times 1 \mathrm{E}$ | $0 \times 97$ | $0 \times 03$ | 9703 | $0 \times 1 \mathrm{~F}$ |
| ATmega1281 | $0 \times 1 \mathrm{E}$ | $0 \times 97$ | $0 \times 04$ | 9704 | $0 \times 1 \mathrm{~F}$ |
| ATmega2560 | $0 \times 1 \mathrm{E}$ | $0 \times 98$ | $0 \times 01$ | 9801 | $0 \times 1 \mathrm{~F}$ |
| ATmega2561 | $0 \times 1 \mathrm{E}$ | $0 \times 98$ | $0 \times 02$ | 9802 | $0 \times 1 \mathrm{~F}$ |

### 30.4 Calibration Byte

The ATmega640/1280/1281/2560/2561 has a byte calibration value for the internal RC Oscillator. This byte resides in the high byte of address $0 x 000$ in the signature address space. During reset, this byte is automatically written into the OSCCAL Register to ensure correct frequency of the calibrated RC Oscillator.

### 30.5 Page Size

Table 30-7. $\quad$ No. of Words in a Page and No. of Pages in the Flash

| Flash Size | Page Size | PCWORD | No. of Pages | PCPAGE | PCMSB |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 128K words (256Kbytes) | 128 words | PC[6:0] | 1024 | PC[16:7] | 16 |

Table 30-8. No. of Words in a Page and No. of Pages in the EEPROM

| EEPROM Size | Page Size | PCWORD | No. of Pages | PCPAGE | EEAMSB |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 4 Kbytes | 8 bytes | EEA[2:0] | 512 | EEA[11:3] | 11 |

### 30.6 Parallel Programming Parameters, Pin Mapping, and Commands

This section describes how to parallel program and verify Flash Program memory, EEPROM Data memory, Memory Lock bits, and Fuse bits in the ATmega640/1280/1281/2560/2561. Pulses are assumed to be at least 250ns unless otherwise noted.

### 30.6.1 Signal Names

In this section, some pins of the ATmega640/1280/1281/2560/2561 are referenced by signal names describing their functionality during parallel programming, see Figure $30-1$ and Table 30-9 on page 329. Pins not described in the following table are referenced by pin names.

The XA1/XA0 pins determine the action executed when the XTAL1 pin is given a positive pulse. The bit coding is shown in Table 30-12 on page 330.

When pulsing $\overline{W R}$ or $\overline{\mathrm{OE}}$, the command loaded determines the action executed. The different commands are shown in Table 30-13 on page 330.

Figure 30-1. Parallel Programming ${ }^{(1)}$


Note: 1. Unused Pins should be left floating.
Table 30-9. Pin Name Mapping

| Signal Name in <br> Programming Mode | Pin Name | I/O | Function |
| :---: | :---: | :---: | :--- |
| RDY/BSY | PD1 | O | 0: Device is busy programming, 1: Device is ready for new command |
| $\overline{\text { OE }}$ | PD2 | I | Output Enable (Active low) |
| $\overline{W R}$ | PD3 | I | Write Pulse (Active low) |
| BS1 | PD4 | I | Byte Select 1 |
| XA0 | PD5 | I | XTAL Action Bit 0 |
| XA1 | PD6 | I | XTAL Action Bit 1 |
| PAGEL | PD7 | I | Program Memory and EEPROM data Page Load |
| BS2 | PA0 | I | Byte Select 2 |
| DATA | PB7-0 | I/O | Bi-directional Data bus (Output when $\overline{\text { OE is low) }}$ |

Table 30-10. BS2 and BS1 Encoding

| BS2 | BS1 | Flash / EEPROM <br> Address | Flash Data Loading / <br> Reading | Fuse Programming | Reading Fuse and Lock <br> Bits |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | Low Byte | Low Byte | Low Byte | Fuse Low Byte |
| 0 | 1 | High Byte | High Byte | High Byte | Lockbits |
| 1 | 0 | Extended High <br> Byte | Reserved | Extended Byte | Extended Fuse Byte |
| 1 | 1 | Reserved | Reserved | Reserved | Fuse High Byte |

Table 30-11. Pin Values Used to Enter Programming Mode

| Pin | Symbol | Value |
| :---: | :---: | :---: |
| PAGEL | Prog_enable[3] | 0 |
| XA1 | Prog_enable[2] | 0 |
| XA0 | Prog_enable[1] | 0 |
| BS1 | Prog_enable[0] | 0 |

Table 30-12. XA1 and XAO Enoding

| XA1 | XA0 | Action when XTAL1 is Pulsed |
| :---: | :---: | :--- |
| 0 | 0 | Load Flash or EEPROM Address (High or low address byte determined by BS2 and BS1) |
| 0 | 1 | Load Data (High or Low data byte for Flash determined by BS1) |
| 1 | 0 | Load Command |
| 1 | 1 | No Action, Idle |

Table 30-13. Command Byte Bit Encoding

| Command Byte | Command Executed |
| :---: | :--- |
| 10000000 | Chip Erase |
| 01000000 | Write Fuse bits |
| 00100000 | Write Lock bits |
| 00010000 | Write Flash |
| 00010001 | Write EEPROM |
| 00001000 | Read Signature Bytes and Calibration byte |
| 00000100 | Read Fuse and Lock bits |
| 00000010 | Read Flash |
| 00000011 | Read EEPROM |

### 30.7 Parallel Programming

### 30.7.1 Enter Programming Mode

The following algorithm puts the device in parallel programming mode:

1. Apply $4.5 \mathrm{~V}-5.5 \mathrm{~V}$ between $\mathrm{V}_{\mathrm{CC}}$ and GND .
2. Set RESET to "0" and toggle XTAL1 at least six times.
3. Set the Prog_enable pins listed in Table 30-11 to "0000" and wait at least 100ns.
4. Apply $11.5 \mathrm{~V}-12.5 \mathrm{~V}$ to RESET. Any activity on Prog_enable pins within 100 ns after +12 V has been applied to RESET, will cause the device to fail entering programming mode.
5. Wait at least $50 \mu$ s before sending a new command.

### 30.7.2 Considerations for Efficient Programming

The loaded command and address are retained in the device during programming. For efficient programming, the following should be considered.

- The command needs only be loaded once when writing or reading multiple memory locations
- Skip writing the data value 0xFF, that is the contents of the entire EEPROM (unless the EESAVE Fuse is programmed) and Flash after a Chip Erase
- Address high byte needs only be loaded before programming or reading a new 256 word window in Flash or 256 byte EEPROM. This consideration also applies to Signature bytes reading


### 30.7.3 Chip Erase

The Chip Erase will erase the Flash and EEPROM ${ }^{(1)}$ memories plus Lock bits. The Lock bits are not reset until the program memory has been completely erased. The Fuse bits are not changed. A Chip Erase must be performed before the Flash and/or EEPROM are reprogrammed.

Note: 1. The EEPRPOM memory is preserved during Chip Erase if the EESAVE Fuse is programmed.
Load Command "Chip Erase"

1. Set XA1, XA0 to " 10 ". This enables command loading.
2. Set BS1 to "0".
3. Set DATA to " 10000000 ". This is the command for Chip Erase.
4. Give XTAL1 a positive pulse. This loads the command.
5. Give $\overline{W R}$ a negative pulse. This starts the Chip Erase. RDY/ $\overline{B S Y}$ goes low.
6. Wait until RDY/ $\overline{\mathrm{BSY}}$ goes high before loading a new command.

### 30.7.4 Programming the Flash

The Flash is organized in pages, see Table 30-7 on page 328. When programming the Flash, the program data is latched into a page buffer. This allows one page of program data to be programmed simultaneously. The following procedure describes how to program the entire Flash memory:
A. Load Command "Write Flash"

1. Set XA1, XA0 to " 10 ". This enables command loading.
2. Set BS1 to "0".
3. Set DATA to "0001 0000". This is the command for Write Flash.
4. Give XTAL1 a positive pulse. This loads the command.
B. Load Address Low byte (Address bits 7:0)
5. Set XA1, XA0 to " 00 ". This enables address loading.
6. Set BS2, BS1 to " 00 ". This selects the address low byte.
7. Set DATA = Address low byte ( $0 \times 00-0 x F F$ ).
8. Give XTAL1 a positive pulse. This loads the address low byte.
C. Load Data Low Byte
9. Set XA1, XA0 to " 01 ". This enables data loading.
10. Set DATA = Data low byte ( $0 \times 00-0 x F F$ ).
11. Give XTAL1 a positive pulse. This loads the data byte.

## D. Load Data High Byte

1. Set BS1 to "1". This selects high data byte.
2. Set XA1, XA0 to " 01 ". This enables data loading.
3. Set DATA $=$ Data high byte $(0 \times 00-0 x F F)$.
4. Give XTAL1 a positive pulse. This loads the data byte.
E. Latch Data
5. Set BS1 to "1". This selects high data byte.
6. Give PAGEL a positive pulse. This latches the data bytes. See Figure $30-3$ on page 333 for signal waveforms.
F. Repeat $B$ through $E$ until the entire buffer is filled or until all data within the page is loaded

While the lower bits in the address are mapped to words within the page, the higher bits address the pages within the FLASH. This is illustrated in Figure 30-2 on page 333. Note that if less than eight bits are required to address words in the page (pagesize < 256), the most significant bit(s) in the address low byte are used to address the page when performing a Page Write.
G. Load Address High byte (Address bits15:8)

1. Set XA1, XAO to " 00 ". This enables address loading.
2. Set BS2, BS1 to " 01 ". This selects the address high byte.
3. Set DATA = Address high byte ( $0 \times 00-0 x F F$ ).
4. Give XTAL1 a positive pulse. This loads the address high byte.
H. Load Address Extended High byte (Address bits 23:16)
5. Set XA1, XA0 to " 00 ". This enables address loading.
6. Set BS2, BS1 to "10". This selects the address extended high byte.
7. Set DATA = Address extended high byte ( $0 \times 00-0 x F F$ ).
8. Give XTAL1 a positive pulse. This loads the address high byte.

## I. Program Page

1. Set BS2, BS1 to " 00 ".
2. Give $\overline{W R}$ a negative pulse. This starts programming of the entire page of data. RDY/ $\overline{B S Y}$ goes low.
3. Wait until RDY/ $\overline{\mathrm{BSY}}$ goes high (see Figure 30-3 on page 333 for signal waveforms).
J. Repeat B through I until the entire Flash is programmed or until all data has been programmed
K. End Page Programming
4. 5. Set XA1, XA0 to " 10 ". This enables command loading.
1. Set DATA to " 00000000 ". This is the command for No Operation.
2. Give XTAL1 a positive pulse. This loads the command, and the internal write signals are reset.

Figure 30-2. Addressing the Flash Which is Organized in Pages ${ }^{(1)}$


Note: 1. PCPAGE and PCWORD are listed in Table 30-7 on page 328.
Figure 30-3. Programming the Flash Waveforms ${ }^{(1)}$


Note: 1. " $X X$ " is don't care. The letters refer to the programming description above.

### 30.7.5 Programming the EEPROM

The EEPROM is organized in pages, see Table $30-8$ on page 328 . When programming the EEPROM, the program data is latched into a page buffer. This allows one page of data to be programmed simultaneously. The programming algorithm for the EEPROM data memory is as follows (refer to "Programming the Flash" on page 331 for details on Command, Address and Data loading):

1. A: Load Command "0001 0001".
2. G: Load Address High Byte ( $0 \times 00-0 x F F$ ).
3. B: Load Address Low Byte (0x00-0xFF).
4. C: Load Data (0x00-0xFF).
5. E: Latch data (give PAGEL a positive pulse).

K : Repeat 3 through 5 until the entire buffer is filled.
L: Program EEPROM page.

1. Set BS2, BS1 to " 00 ".
2. Give $\overline{\mathrm{WR}}$ a negative pulse. This starts programming of the EEPROM page. RDY/ $\overline{\mathrm{BSY}}$ goes low.
3. Wait until to RDY/ $\overline{B S Y}$ goes high before programming the next page (see Figure $30-4$ for signal waveforms).

Figure 30-4. Programming the EEPROM Waveforms


### 30.7.6 Reading the Flash

The algorithm for reading the Flash memory is as follows (refer to "Programming the Flash" on page 331 for details on Command and Address loading):

1. A: Load Command "0000 0010".
2. H: Load Address Extended Byte ( $0 \times 00-0 x F F$ ).
3. G: Load Address High Byte ( $0 \times 00-0 x F F$ ).
4. B: Load Address Low Byte ( $0 \times 00-0 x F F$ ).
5. Set $\overline{\mathrm{OE}}$ to " 0 ", and BS1 to "0". The Flash word low byte can now be read at DATA.
6. Set BS to "1". The Flash word high byte can now be read at DATA.
7. Set $\overline{O E}$ to " 1 ".

### 30.7.7 Reading the EEPROM

The algorithm for reading the EEPROM memory is as follows (refer to "Programming the Flash" on page 331 for details on Command and Address loading):

1. A: Load Command "0000 0011".
2. G: Load Address High Byte ( $0 \times 00-0 \times F F$ ).
3. B: Load Address Low Byte ( $0 \times 00-0 \times F F$ ).
4. Set $\overline{O E}$ to " "0", and BS1 to "0". The EEPROM Data byte can now be read at DATA.
5. Set $\overline{O E}$ to " 1 ".

### 30.7.8 Programming the Fuse Low Bits

The algorithm for programming the Fuse Low bits is as follows (refer to "Programming the Flash" on page 331 for details on Command and Data loading):

1. A: Load Command "0100 0000".
2. C: Load Data Low Byte. Bit $\mathrm{n}=$ " 0 " programs and bit $\mathrm{n}=$ " 1 " erases the Fuse bit.
3. Give $\overline{\mathrm{WR}}$ a negative pulse and wait for RDY/ $\overline{\mathrm{BSY}}$ to go high.

### 30.7.9 Programming the Fuse High Bits

The algorithm for programming the Fuse High bits is as follows (refer to "Programming the Flash" on page 331 for details on Command and Data loading):

1. A: Load Command "0100 0000".
2. C: Load Data Low Byte. Bit $\mathrm{n}=$ " 0 " programs and bit $\mathrm{n}=$ " 1 " erases the Fuse bit.
3. Set BS2, BS1 to " 01 ". This selects high data byte.
4. Give $\overline{W R}$ a negative pulse and wait for RDY/ $\overline{B S Y}$ to go high.
5. Set BS2, BS1 to "00". This selects low data byte.

### 30.7.10 Programming the Extended Fuse Bits

The algorithm for programming the Extended Fuse bits is as follows (refer to "Programming the Flash" on page 331 for details on Command and Data loading):

1. 2. A: Load Command "0100 0000".
1. 2. C: Load Data Low Byte. Bit $\mathrm{n}=$ " 0 " programs and bit $\mathrm{n}=$ " 1 " erases the Fuse bit.
1. 3. Set BS2, BS1 to " 10 ". This selects extended data byte.
1. 4. Give $\overline{W R}$ a negative pulse and wait for RDY/ $\overline{B S Y}$ to go high.
1. 5. Set BS2, BS1 to "00". This selects low data byte.

Figure 30-5. Programming the FUSES Waveforms


### 30.7.11 Programming the Lock Bits

The algorithm for programming the Lock bits is as follows (refer to "Programming the Flash" on page 331 for details on Command and Data loading):

1. A: Load Command "0010 0000".
2. C: Load Data Low Byte. Bit $\mathrm{n}=$ " 0 " programs the Lock bit. If LB mode 3 is programmed (LB1 and LB2 is programmed), it is not possible to program the Boot Lock bits by any External Programming mode.
3. Give $\overline{W R}$ a negative pulse and wait for RDY/BSY to go high.

The Lock bits can only be cleared by executing Chip Erase.

### 30.7.12 Reading the Fuse and Lock Bits

The algorithm for reading the Fuse and Lock bits is as follows (refer to "Programming the Flash" on page 331 for details on Command loading):

1. A: Load Command "0000 0100".
2. Set $\overline{O E}$ to " 0 ", and $B S 2, B S 1$ to " 00 ". The status of the Fuse Low bits can now be read at DATA (" 0 " means programmed).
3. Set $\overline{O E}$ to " 0 ", and BS2, BS1 to " 11 ". The status of the Fuse High bits can now be read at DATA (" 0 " means programmed).
4. Set OE to " 0 ", and BS2, BS1 to " 10 ". The status of the Extended Fuse bits can now be read at DATA (" 0 " means programmed).
5. Set $\overline{O E}$ to " 0 ", and BS2, BS1 to " 01 ". The status of the Lock bits can now be read at DATA ("0" means programmed).
6. Set $\overline{O E}$ to " 1 ".

Figure 30-6. Mapping Between BS1, BS2 and the Fuse and Lock Bits During Read


### 30.7.13 Reading the Signature Bytes

The algorithm for reading the Signature bytes is as follows (refer to "Programming the Flash" on page 331 for details on Command and Address loading):

1. A: Load Command "0000 1000".
2. B: Load Address Low Byte (0x00-0x02).
3. Set $\overline{O E}$ to " 0 ", and $B S$ to " 0 ". The selected Signature byte can now be read at DATA.
4. Set $\overline{O E}$ to " 1 ".

### 30.7.14 Reading the Calibration Byte

The algorithm for reading the Calibration byte is as follows (refer to "Programming the Flash" on page 331 for details on Command and Address loading):

1. A: Load Command "0000 1000".
2. B: Load Address Low Byte, $0 \times 00$.
3. Set $\overline{O E}$ to " ", and BS1 to " 1 ". The Calibration byte can now be read at DATA.
4. Set $\overline{O E}$ to "1".

### 30.7.15 Parallel Programming Characteristics

Figure 30-7. Parallel Programming Timing, Including some General Timing Requirements


Figure 30-8. Parallel Programming Timing, Loading Sequence with Timing Requirements ${ }^{(1)}$


Note: 1. The timing requirements shown in Figure $30-7$ (that is, $\mathrm{t}_{\mathrm{DVXH}}, \mathrm{t}_{\mathrm{XHXL}}$, and $\mathrm{t}_{\mathrm{XLDX}}$ ) also apply to loading operation.
Figure 30-9. Parallel Programming Timing, Reading Sequence (within the Same Page) with Timing Requirements ${ }^{(1)}$


Note: 1. The timing requirements shown in Figure $30-7$ (that is, $\mathrm{t}_{\mathrm{DVXH}}, \mathrm{t}_{\mathrm{XHXL}}$, and $\mathrm{t}_{\mathrm{XLDX}}$ ) also apply to reading operation.

Table 30-14. Parallel Programming Characteristics, $\mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V} \pm 10 \%$

| Symbol | Parameter | Min | Typ | Max | Units |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{V}_{\text {PP }}$ | Programming Enable Voltage | 11.5 |  | 12.5 | V |
| $\mathrm{I}_{\mathrm{PP}}$ | Programming Enable Current |  |  | 250 | $\mu \mathrm{A}$ |
| $\mathrm{t}_{\text {DVXH }}$ | Data and Control Valid before XTAL1 High | 67 |  |  | ns |
| $\mathrm{t}_{\text {XLXH }}$ | XTAL1 Low to XTAL1 High | 200 |  |  |  |
| ${ }^{\text {XHXL }}$ | XTAL1 Pulse Width High | 150 |  |  |  |
| $t_{\text {XLDX }}$ | Data and Control Hold after XTAL1 Low | 67 |  |  |  |
| ${ }^{\text {XLWL }}$ | XTAL1 Low to $\overline{W R}$ Low | 0 |  |  |  |
| $\mathrm{t}_{\text {XLPH }}$ | XTAL1 Low to PAGEL high | 0 |  |  |  |
| $\mathrm{t}_{\text {PLXH }}$ | PAGEL low to XTAL1 high | 150 |  |  |  |
| $\mathrm{t}_{\text {BVPH }}$ | BS1 Valid before PAGEL High | 67 |  |  |  |
| $\mathrm{t}_{\text {PHPL }}$ | PAGEL Pulse Width High | 150 |  |  |  |
| $t_{\text {PLBX }}$ | BS1 Hold after PAGEL Low | 67 |  |  |  |
| $t_{\text {WLBX }}$ | BS2/1 Hold after $\overline{W R}$ Low | 67 |  |  |  |
| $t_{\text {PLWL }}$ | PAGEL Low to $\overline{W R}$ Low | 67 |  |  |  |
| $\mathrm{t}_{\text {BVWL }}$ | BS2/1 Valid to WR Low | 67 |  |  |  |
| $\mathrm{t}_{\text {WLWH }}$ | $\overline{\text { WR Pulse Width Low }}$ | 150 |  |  |  |
| $t_{\text {WLRL }}$ |  | 0 |  | 1 | $\mu \mathrm{s}$ |
| $t_{\text {WLRH }}$ | $\overline{\text { WR Low to RDY/BSY High }}{ }^{(1)}$ | 3.7 |  | 4.5 | ms |
| ${ }^{\text {W WLRH_CE }}$ | $\overline{\text { WR Low to RDY/BSY High for Chip Erase }}{ }^{(2)}$ | 7.5 |  | 9 |  |
| ${ }^{\text {XLOL }}$ | XTAL1 Low to $\overline{\mathrm{OE}}$ Low | 0 |  |  | ns |
| $\mathrm{t}_{\text {BVDV }}$ | BS1 Valid to DATA valid | 0 |  | 250 |  |
| toldv | $\overline{\text { OE Low to DATA Valid }}$ |  |  | 250 |  |
| $\mathrm{t}_{\text {OHDZ }}$ | $\overline{\text { OE High to DATA Tri-stated }}$ |  |  | 250 |  |

Notes: 1. $t_{\text {WLRH }}$ is valid for the Write Flash, Write EEPROM, Write Fuse bits and Write Lock bits commands.
2. $t_{\text {WLRH_CE }}$ is valid for the Chip Erase command.

### 30.8 Serial Downloading

Both the Flash and EEPROM memory arrays can be programmed using a serial programming bus while RESET is pulled to GND. The serial programming interface consists of pins SCK, PDI (input) and PDO (output). After RESET is set low, the Programming Enable instruction needs to be executed first before program/erase operations can be executed. NOTE, in Table 30-15 on page 339, the pin mapping for serial programming is listed. Not all packages use the SPI pins dedicated for the internal Serial Peripheral Interface - SPI.

### 30.8.1 Serial Programming Pin Mapping

Table 30-15. Pin Mapping Serial Programming

| Symbol | Pins (TQFP-100) | Pins (TQFP-64) | I/O | Description |
| :---: | :---: | :---: | :---: | :---: |
| PDI | PB2 | PE0 | I | Serial Data in |
| PDO | PB3 | PE1 | O | Serial Data out |
| SCK | PB1 | PB1 | I | Serial Clock |

Figure 30-10. Serial Programming and Verify ${ }^{(1)}$


Notes: 1. If the device is clocked by the internal Oscillator, it is no need to connect a clock source to the XTAL1 pin.
2. $\mathrm{V}_{C C}-0.3 \mathrm{~V}<\mathrm{AVCC}<\mathrm{V}_{C C}+0.3 \mathrm{~V}$, however, AVCC should always be within $1.8 \mathrm{~V}-5.5 \mathrm{~V}$. When programming the EEPROM, an auto-erase cycle is built into the self-timed programming operation (in the Serial mode ONLY) and there is no need to first execute the Chip Erase instruction. The Chip Erase operation turns the content of every memory location in both the Program and EEPROM arrays into 0xFF.
Depending on CKSEL Fuses, a valid clock must be present. The minimum low and high periods for the serial clock (SCK) input are defined as follows:

Low: > 2 CPU clock cycles for $\mathrm{f}_{\mathrm{ck}}<12 \mathrm{MHz}, 3 \mathrm{CPU}$ clock cycles for $\mathrm{f}_{\mathrm{ck}}>=12 \mathrm{MHz}$
High: > 2 CPU clock cycles for $f_{c k}<12 \mathrm{MHz}, 3 \mathrm{CPU}$ clock cycles for $\mathrm{f}_{\mathrm{ck}}>=12 \mathrm{MHz}$

### 30.8.2 Serial Programming Algorithm

When writing serial data to the ATmega640/1280/1281/2560/2561, data is clocked on the rising edge of SCK.
When reading data from the ATmega640/1280/1281/2560/2561, data is clocked on the falling edge of SCK. See Figure $30-12$ on page 342 for timing details.

To program and verify the ATmega640/1280/1281/2560/2561 in the serial programming mode, the following sequence is recommended (see four byte instruction formats in Table 30-17 on page 340):

1. Power-up sequence:

Apply power between $\mathrm{V}_{\mathrm{CC}}$ and GND while RESET and SCK are set to " 0 ". In some systems, the programmer can not guarantee that SCK is held low during power-up. In this case, $\overline{\text { RESET }}$ must be given a positive pulse of at least two CPU clock cycles duration after SCK has been set to "0".
2. Wait for at least 20 ms and enable serial programming by sending the Programming Enable serial instruction to pin PDI.
3. The serial programming instructions will not work if the communication is out of synchronization. When in sync. the second byte ( $0 \times 53$ ), will echo back when issuing the third byte of the Programming Enable instruction. Whether the echo is correct or not, all four bytes of the instruction must be transmitted. If the $0 \times 53$ did not echo back, give $\overline{\text { RESET }}$ a positive pulse and issue a new Programming Enable command.
4. The Flash is programmed one page at a time. The memory page is loaded one byte at a time by supplying the 7 LSB of the address and data together with the Load Program Memory Page instruction. To ensure correct loading of the page, the data low byte must be loaded before data high byte is applied for a given address. The Program Memory Page is stored by loading the Write Program Memory Page instruction with the address lines $15: 8$. Before issuing this command, make sure the instruction Load Extended Address Byte has been used to define the MSB of the address. The extended address byte is stored until the command is re-issued, that is, the command needs only be issued for the first page, and when crossing the 64 KWord boundary. If polling (RDY/ $\overline{\mathrm{BSY}}$ ) is not used, the user must wait at least $\mathrm{t}_{\text {WD_FLASH }}$ before issuing the next page (see Table 30-16). Accessing the serial programming interface before the Flash write operation completes can result in incorrect programming.
5. The EEPROM array is programmed one byte at a time by supplying the address and data together with the appropriate Write instruction. An EEPROM memory location is first automatically erased before new data is written. If polling is not used, the user must wait at least $t_{\text {wd_eeprom }}$ before issuing the next byte (see Table $30-16$ ). In a chip erased device, no 0xFFs in the data file(s) need to be programmed.
6. Any memory location can be verified by using the Read instruction which returns the content at the selected address at serial output PDO. When reading the Flash memory, use the instruction Load Extended Address Byte to define the upper address byte, which is not included in the Read Program Memory instruction. The extended address byte is stored until the command is re-issued, that is, the command needs only be issued for the first page, and when crossing the 64KWord boundary.
7. At the end of the programming session, $\overline{\text { RESET }}$ can be set high to commence normal operation.
8. Power-off sequence (if needed):

Set RESET to " 1 ".
Turn $\mathrm{V}_{\mathrm{CC}}$ power off.

Table 30-16. Minimum Wait Delay Before Writing the Next Flash or EEPROM Location

| Symbol | Minimum Wait Delay |
| :--- | :---: |
| $t_{\text {WD_FLASH }}$ | 4.5 ms |
| $t_{\text {WD_EEPROM }}$ | 3.6 ms |
| $t_{\text {WD_ERASE }}$ | 9.0 ms |

### 30.8.3 Serial Programming Instruction set

Table 30-17 and Figure 30-11 on page 342 describes the Instruction set.
Table 30-17. Serial Programming Instruction Set

| Instruction/Operation | Instruction Format |  |  |  |
| :--- | :---: | :---: | :---: | :---: |
|  | Byte 1 | Byte 2 | Byte 3 | Byte 4 |
| Programming Enable | $\$ A C$ | $\$ 53$ | $\$ 00$ | $\$ 00$ |
| Chip Erase (Program Memory/EEPROM) | $\$ A C$ | $\$ 80$ | $\$ 00$ | $\$ 00$ |
| Poll RDY/BSY | $\$ F 0$ | $\$ 00$ | $\$ 00$ | data byte out |
| Load Instructions | $\$ 4 D$ | $\$ 00$ | Extended adr | $\$ 00$ |
| Load Extended Address byte ${ }^{(1)}$ | $\$ 48$ | $\$ 00$ | adr LSB | high data byte in |
| Load Program Memory Page, High byte | $\$ 40$ | $\$ 00$ | adr LSB | low data byte in |
| Load Program Memory Page, Low byte |  |  |  |  |

Table 30-17. Serial Programming Instruction Set (Continued)

| Instruction/Operation | Instruction Format |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: |
|  | Byte 1 | Byte 2 | Byte 3 | Byte 4 |
| Load EEPROM Memory Page (page access) | $\$ C 1$ | $\$ 00$ | $0000000 a \mathrm{a}$ | data byte in |
| Read Instructions |  |  |  |  |
| Read Program Memory, High byte | $\$ 28$ | adr MSB | adr LSB | high data byte <br> out |
| Read Program Memory, Low byte | $\$ 20$ | adr MSB | adr LSB | low data byte out |
| Read EEPROM Memory | $\$ A 0$ | 0000 aaaa | aaaa aaaa | data byte out |
| Read Lock bits | $\$ 58$ | $\$ 00$ | $\$ 00$ | data byte out |
| Read Signature Byte | $\$ 30$ | $\$ 00$ | $000000 a a$ | data byte out |
| Read Fuse bits | $\$ 50$ | $\$ 00$ | $\$ 00$ | data byte out |
| Read Fuse High bits | $\$ 58$ | $\$ 08$ | $\$ 00$ | data byte out |
| Read Extended Fuse Bits | $\$ 50$ | $\$ 08$ | $\$ 00$ | data byte out |
| Read Calibration Byte | $\$ 38$ | $\$ 00$ | $\$ 00$ | data byte out |
| Write Instructions |  |  |  |  |
| Write Program Memory Page | $\$ 4 C$ | adr MSB | adr LSB | \$00 |
| Write EEPROM Memory | $\$ C 0$ | 0000 aaaa | aaaa aaaa | data byte in |
| Write EEPROM Memory Page (page access) | $\$ C 2$ | 0000 aaaa | aaaa 00 | $\$ \$ 00$ |
| Write Lock bits | $\$ A C$ | $\$ E 0$ | $\$ 00$ | data byte in |
| Write Fuse bits | $\$ A C$ | $\$ A 0$ | $\$ 00$ | data byte in |
| Write Fuse High bits | $\$ A C$ | $\$ A 8$ | $\$ 00$ | data byte in |
| Write Extended Fuse Bits | $\$ A C$ | $\$ A 4$ | $\$ 00$ | data byte in |

Notes: 1. Not all instructions are applicable for all parts.
2. $\mathrm{a}=$ address.
3. Bits are programmed ' 0 ', unprogrammed ' 1 '.
4. To ensure future compatibility, unused Fuses and Lock bits should be unprogrammed ('1').
5. Refer to the corresponding section for Fuse and Lock bits, Calibration and Signature bytes and Page size.
6. See http://www.atmel.com/avr for Application Notes regarding programming and programmers.

If the LSB in RDY/BSY data byte out is ' 1 ', a programming operation is still pending. Wait until this bit returns ' 0 ' before the next instruction is carried out.

Within the same page, the low data byte must be loaded prior to the high data byte.
After data is loaded to the page buffer, program the EEPROM page, see Figure 30-11 on page 342.

Figure 30-11. Serial Programming Instruction example
Serial Programming Instruction
$\begin{array}{ll}\text { Load Program Memory Page (High/Low Byte)/ } & \text { Write Program Memory Page/ } \\ \text { Load EEPROM Memory Page (page access) } & \text { Write EEPROM Memory Page }\end{array}$


### 30.8.4 Serial Programming Characteristics

For characteristics of the Serial Programming module, see "SPI Timing Characteristics" on page 363.
Figure 30-12. Serial Programming Waveforms


### 30.9 Programming via the JTAG Interface

Programming through the JTAG interface requires control of the four JTAG specific pins: TCK, TMS, TDI, and TDO. Control of the reset and clock pins is not required.

To be able to use the JTAG interface, the JTAGEN Fuse must be programmed. The device is default shipped with the fuse programmed. In addition, the JTD bit in MCUCR must be cleared. Alternatively, if the JTD bit is set, the external reset can be forced low. Then, the JTD bit will be cleared after two chip clocks, and the JTAG pins are available for programming. This provides a means of using the JTAG pins as normal port pins in Running mode
while still allowing In-System Programming via the JTAG interface. Note that this technique can not be used when using the JTAG pins for Boundary-scan or On-chip Debug. In these cases the JTAG pins must be dedicated for this purpose.

During programming the clock frequency of the TCK Input must be less than the maximum frequency of the chip. The System Clock Prescaler can not be used to divide the TCK Clock Input into a sufficiently low frequency.

As a definition in this datasheet, the LSB is shifted in and out first of all Shift Registers.

### 30.9.1 Programming Specific JTAG Instructions

The Instruction Register is 4-bit wide, supporting up to 16 instructions. The JTAG instructions useful for programming are listed below.

The OPCODE for each instruction is shown behind the instruction name in hex format. The text describes which Data Register is selected as path between TDI and TDO for each instruction.

The Run-Test/Idle state of the TAP controller is used to generate internal clocks. It can also be used as an idle state between JTAG sequences. The state machine sequence for changing the instruction word is shown in Figure 30-13.

Figure 30-13. State Machine Sequence for Changing the Instruction Word


### 30.9.2 AVR_RESET (0xC)

The AVR specific public JTAG instruction for setting the AVR device in the Reset mode or taking the device out from the Reset mode. The TAP controller is not reset by this instruction. The one bit Reset Register is selected as Data Register. Note that the reset will be active as long as there is a logic "one" in the Reset Chain. The output from this chain is not latched.

The active states are:

- Shift-DR: The Reset Register is shifted by the TCK input
30.9.3 PROG_ENABLE (0x4)

The AVR specific public JTAG instruction for enabling programming via the JTAG port. The 16 -bit Programming Enable Register is selected as Data Register. The active states are the following:

- Shift-DR: The programming enable signature is shifted into the Data Register
- Update-DR: The programming enable signature is compared to the correct value, and Programming mode is entered if the signature is valid


### 30.9.4 PROG_COMMANDS ( $0 \times 5$ )

The AVR specific public JTAG instruction for entering programming commands via the JTAG port. The 15-bit Programming Command Register is selected as Data Register. The active states are the following:

- Capture-DR: The result of the previous command is loaded into the Data Register
- Shift-DR: The Data Register is shifted by the TCK input, shifting out the result of the previous command and shifting in the new command
- Update-DR: The programming command is applied to the Flash inputs
- Run-Test/Idle: One clock cycle is generated, executing the applied command


### 30.9.5 PROG_PAGELOAD (0x6)

The AVR specific public JTAG instruction to directly load the Flash data page via the JTAG port. An 8-bit Flash Data Byte Register is selected as the Data Register. This is physically the 8 LSBs of the Programming Command Register. The active states are the following:

- Shift-DR: The Flash Data Byte Register is shifted by the TCK input.
- Update-DR: The content of the Flash Data Byte Register is copied into a temporary register. A write sequence is initiated that within 11 TCK cycles loads the content of the temporary register into the Flash page buffer. The AVR automatically alternates between writing the low and the high byte for each new Update-DR state, starting with the low byte for the first Update-DR encountered after entering the PROG_PAGELOAD command. The Program Counter is pre-incremented before writing the low byte, except for the first written byte. This ensures that the first data is written to the address set up by PROG_COMMANDS, and loading the last location in the page buffer does not make the program counter increment into the next page.


### 30.9.6 PROG_PAGEREAD (0x7)

The AVR specific public JTAG instruction to directly capture the Flash content via the JTAG port. An 8-bit Flash Data Byte Register is selected as the Data Register. This is physically the 8 LSBs of the Programming Command Register. The active states are the following:

- Capture-DR: The content of the selected Flash byte is captured into the Flash Data Byte Register. The AVR automatically alternates between reading the low and the high byte for each new Capture-DR state, starting with the low byte for the first Capture-DR encountered after entering the PROG_PAGEREAD command. The Program Counter is post-incremented after reading each high byte, including the first read byte. This ensures that the first data is captured from the first address set up by PROG_COMMANDS, and reading the last location in the page makes the program counter increment into the next page.
- Shift-DR: The Flash Data Byte Register is shifted by the TCK input.


### 30.9.7 Data Registers

The Data Registers are selected by the JTAG instruction registers described in section "Programming Specific JTAG Instructions" on page 343. The Data Registers relevant for programming operations are:

- Reset Register
- Programming Enable Register
- Programming Command Register
- Flash Data Byte Register


### 30.9.8 Reset Register

The Reset Register is a Test Data Register used to reset the part during programming. It is required to reset the part before entering Programming mode.
A high value in the Reset Register corresponds to pulling the external reset low. The part is reset as long as there is a high value present in the Reset Register. Depending on the Fuse settings for the clock options, the part will remain reset for a Reset Time-out period (refer to "Clock Sources" on page 40) after releasing the Reset Register. The output from this Data Register is not latched, so the reset will take place immediately, as shown in Figure 28-2 on page 297.

### 30.9.9 Programming Enable Register

The Programming Enable Register is a 16 -bit register. The contents of this register is compared to the programming enable signature, binary code 0b1010_0011_0111_0000. When the contents of the register is equal to the programming enable signature, programming via the JTAG port is enabled. The register is reset to 0 on Power-on Reset, and should always be reset when leaving Programming mode.

Figure 30-14. Programming Enable Register


### 30.9.10 Programming Command Register

The Programming Command Register is a 15 -bit register. This register is used to serially shift in programming commands, and to serially shift out the result of the previous command, if any. The JTAG Programming Instruction Set is shown in Table 30-18 on page 347. The state sequence when shifting in the programming commands is illustrated in Figure 30-16 on page 350.

Figure 30-15. Programming Command Register


Table 30-18. JTAG Programming Instruction
Set $\mathbf{a}=$ address high bits, $\mathbf{b}=$ address low bits, $\mathbf{c}=$ address extended bits, $\mathbf{H}=0$ - Low byte, $1-$ High Byte, $\mathbf{0}=$ data out, i = data in, $\mathrm{x}=$ don't care

| Instruction | TDI Sequence | TDO Sequence | Notes |
| :---: | :---: | :---: | :---: |
| 1a. Chip Erase | $\begin{aligned} & 0100011 \_10000000 \\ & 0110001 \_10000000 \\ & 0110011 \_10000000 \\ & 0110011 \_10000000 \end{aligned}$ | xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx |  |
| 1b. Poll for Chip Erase Complete | 0110011_10000000 | xxxxxox_xxxxxxxx | (2) |
| 2a. Enter Flash Write | 0100011_00010000 | xxxxxxx_xxxxxxxx |  |
| 2b. Load Address Extended High Byte | 0001011_cccccccc | xxxxxxx_xxxxxxxx | (10) |
| 2c. Load Address High Byte | 0000111_aaaaaaaa |  |  |
| 2d. Load Address Low Byte | 0000011_bbbbbbbb | xxxxxxx_xxxxxxxx |  |
| 2e. Load Data Low Byte | 0010011_iiiiiiii |  |  |
| 2f. Load Data High Byte | 0010111_iiiiiiii |  |  |
| 2g. Latch Data | $\begin{aligned} & \text { 0110111_00000000 } \\ & \text { 1110111_00000000 } \\ & 0110111 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx <br> XXXXXXX_XXXXXXXX <br> xxxxxxx_xxxxxxxx | (1) |
| 2h. Write Flash Page | 0110111_00000000 <br> 0110101_00000000 <br> 0110111_00000000 <br> 0110111_00000000 | xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx | (1) |
| 2i. Poll for Page Write Complete | 0110111_00000000 | xxxxxox_xxxxxxxx | (2) |
| 3a. Enter Flash Read | 0100011_00000010 | xxxxxxx_xxxxxxxx |  |
| 3b. Load Address Extended High Byte | 0001011_cccccccc | xxxxxxx_xxxxxxxx | (10) |
| 3c. Load Address High Byte | 0000111_aaaaaaaa |  |  |
| 3d. Load Address Low Byte | 0000011_bbbbbbbb | xxxxxxx_xxxxxxxx |  |
| 3e. Read Data Low and High Byte | $\begin{aligned} & \text { 0110010_00000000 } \\ & 0110110 \_00000000 \\ & 0110111 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx <br> xxxxxxx_00000000 <br> xxxxxxx_00000000 | Low byte High byte |
| 4a. Enter EEPROM Write | 0100011_00010001 | xxxxxxx_xxxxxxxx |  |
| 4b. Load Address High Byte | 0000111_aaaaaaaa | xxxxxxx_xxxxxxxx | (10) |
| 4c. Load Address Low Byte | 0000011_bbbbbbbb | xxxxxxx_xxxxxxxx |  |
| 4d. Load Data Byte | 0010011_iiiiiiii | xxxxxxx_xxxxxxxx |  |
| 4e. Latch Data | $\begin{aligned} & \text { 0110111_000000000 } \\ & \text { 1110111_000000000 } \\ & 0110111 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx <br> XXXXXXX_XXXXXXXX <br> xxxxxxx_xxxxxxxx | (1) |
| 4f. Write EEPROM Page | $\begin{aligned} & 0110011 \_00000000 \\ & 0110001 \_00000000 \\ & 0110011 \_00000000 \\ & 0110011 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx | (1) |
| 4g. Poll for Page Write Complete | 0110011_00000000 | xxxxxox_xxxxxxxx | (2) |

Table 30-18. JTAG Programming Instruction (Continued)
Set (Continued) $\mathbf{a}=$ address high bits, $\mathbf{b}=$ address low bits, $\mathbf{c}=$ address extended bits, $\mathbf{H}=0$ - Low byte, 1 - High Byte, $\mathbf{o}=$ data out, $\mathbf{i}=$ data $\mathrm{in}, \mathrm{x}=$ don't care

| Instruction | TDI Sequence | TDO Sequence | Notes |
| :---: | :---: | :---: | :---: |
| 5a. Enter EEPROM Read | 0100011_00000011 | xxxxxxx_xxxxxxxx |  |
| 5b. Load Address High Byte | 0000111_aaaaaaaa | xxxxxxx_xxxxxxxx | (10) |
| 5c. Load Address Low Byte | 0000011_bbbbbbbbb | xxxxxxx_xxxxxxxx |  |
| 5d. Read Data Byte | $\begin{aligned} & \text { 0110011_bbbbbbbb } \\ & 0110010 \_00000000 \\ & 0110011 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx <br> xxXXXXX_XXXXXXXX <br> xxxxxxx_00000000 |  |
| 6a. Enter Fuse Write | 0100011_01000000 | xxxxxxx_xxxxxxxx |  |
| 6b. Load Data Low Byte ${ }^{(6)}$ | 0010011_iiiiiiii | xxxxxxx_xxxxxxxx | (3) |
| 6c. Write Fuse Extended Byte | 0111011_00000000 0111001_00000000 0111011_00000000 0111011_00000000 | xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx | (1) |
| 6d. Poll for Fuse Write Complete | 0110111_00000000 | xxxxxox_xxxxxxxx | (2) |
| 6e. Load Data Low Byte ${ }^{(7)}$ | 0010011_iiiiiiii | xxxxxxx_xxxxxxxx | (3) |
| 6f. Write Fuse High Byte | $\begin{aligned} & 0110111 \_00000000 \\ & 0110101 \_00000000 \\ & 0110111 \_00000000 \\ & 0110111 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx | (1) |
| 6g. Poll for Fuse Write Complete | 0110111_00000000 | xxxxx0x_xxxxxxxx | (2) |
| 6h. Load Data Low Byte ${ }^{(7)}$ | 0010011_iiiiiiii | xxxxxxx_xxxxxxxx | (3) |
| 6i. Write Fuse Low Byte | 0110011_00000000 <br> 0110001_00000000 <br> 0110011_00000000 <br> 0110011_00000000 | xxxxxxx_xxxxxxxx XXXXXXX_XXXXXXXX xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx | (1) |
| 6j. Poll for Fuse Write Complete | 0110011_00000000 | xxxxx0x_xxxxxxxx | (2) |
| 7a. Enter Lock Bit Write | 0100011_00100000 | xxxxxxx_xxxxxxxx |  |
| 7b. Load Data Byte ${ }^{(9)}$ | 0010011_11iiiiii | xxxxxxx_xxxxxxxx | (4) |
| 7c. Write Lock Bits | 0110011_00000000 0110001_00000000 <br> 0110011_00000000 <br> 0110011_00000000 | xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxXXXXX_XXXXXXXX xxxxxxx_xxxxxxxx | (1) |
| 7d. Poll for Lock Bit Write complete | 0110011_00000000 | xxxxxox_xxxxxxxx | (2) |
| 8a. Enter Fuse/Lock Bit Read | 0100011_00000100 | xxxxxxx_xxxxxxxx |  |
| 8b. Read Extended Fuse Byte ${ }^{(6)}$ | $\begin{aligned} & 0111010 \_00000000 \\ & 0111011 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx <br> xxxxxxx_00000000 |  |
| 8c. Read Fuse High Byte ${ }^{(7)}$ | $\begin{aligned} & 0111110 \_00000000 \\ & 0111111 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx xxxxxxx_00000000 |  |
| 8d. Read Fuse Low Byte ${ }^{(8)}$ | $\begin{aligned} & 0110010 \_00000000 \\ & 0110011 \_00000000 \end{aligned}$ | xxxxxxx_xxxxxxxx xxxxxxx_00000000 |  |

Table 30-18. JTAG Programming Instruction (Continued)
Set (Continued) $\mathbf{a}=$ address high bits, $\mathbf{b}=$ address low bits, $\mathbf{c}=$ address extended bits, $\mathbf{H}=0$ - Low byte, 1 - High Byte, $\mathbf{0}=$ data out, $\mathbf{i}=$ data in, $\mathbf{x}=$ don't care

| Instruction | TDI Sequence | TDO Sequence | Notes |
| :--- | :--- | :--- | :--- |
| 8e. Read Lock Bits ${ }^{(9)}$ | $0110110 \_00000000$ <br> $0110111 \_00000000$ | xxxxxxx_xxxxxxxx <br> xxxxxxx_xxoooooo | (5) |
|  | $0111010 \_00000000$ <br> $0111110 \_00000000$ <br> $0110010 \_00000000$ <br> $0110110 \_00000000$ <br> $0110111 \_00000000$ | xxxxxxx_xxxxxxxx <br> xxxxxx_00000000 <br> xxxxxx_00000000 <br> xxxxxxx_00000000 <br> xxxxxx_00000000 | (5) <br> Fuse Ext. byte <br> Fuse High byte <br> Fuse Low byte <br> Lock bits |
| 9a. Read Fuses and Lock Bits | $0100011 \_00001000$ | xxxxxxx_xxxxxxxx |  |
| 9b. Load Address Byte | $0000011 \_b b b b b b b b$ | xxxxxxx_xxxxxxxx |  |
| 9c. Read Signature Byte | $0110010 \_00000000$ <br> $0110011 \_00000000$ | xxxxxxx_xxxxxxxx <br> xxxxxxx_oooooooo |  |
| 10a. Enter Calibration Byte Read | $0100011 \_00001000$ | xxxxxxx_xxxxxxxx |  |
| 10b. Load Address Byte | $0000011 \_b b b b b b b b$ | xxxxxxx_xxxxxxxx |  |
| 10c. Read Calibration Byte | $0110110 \_00000000$ <br> $0110111 \_00000000$ | xxxxxxx_xxxxxxxx <br> xxxxxx_000ooooo |  |
| 11a. Load No Operation Command | $0100011 \_00000000$ <br> $0110011 \_00000000$ | xxxxxxx_xxxxxxxx <br> xxxxxx_xxxxxxxx |  |

Notes: 1. This command sequence is not required if the seven MSB are correctly set by the previous command sequence (which is normally the case).
2. Repeat until $\mathbf{0}=$ " 1 ".
3. Set bits to " 0 " to program the corresponding Fuse, " 1 " to unprogram the Fuse.
4. Set bits to " 0 " to program the corresponding Lock bit, " 1 " to leave the Lock bit unchanged.
5. " 0 " = programmed, " 1 " = unprogrammed.
6. The bit mapping for Fuses Extended byte is listed in Table 30-3 on page 326.
7. The bit mapping for Fuses High byte is listed in Table 30-4 on page 327.
8. The bit mapping for Fuses Low byte is listed in Table 30-5 on page 327.
9. The bit mapping for Lock bits byte is listed in Table 30-1 on page 325.
10. Address bits exceeding PCMSB and EEAMSB (Table 30-7 and Table 30-8 on page 328) are don't care.
11. All TDI and TDO sequences are represented by binary digits (0b...).

Figure 30-16. State Machine Sequence for Changing/Reading the Data Word


### 30.9.11 Flash Data Byte Register

The Flash Data Byte Register provides an efficient way to load the entire Flash page buffer before executing Page Write, or to read out/verify the content of the Flash. A state machine sets up the control signals to the Flash and senses the strobe signals from the Flash, thus only the data words need to be shifted in/out.
The Flash Data Byte Register actually consists of the 8-bit scan chain and a 8-bit temporary register. During page load, the Update-DR state copies the content of the scan chain over to the temporary register and initiates a write sequence that within 11 TCK cycles loads the content of the temporary register into the Flash page buffer. The AVR automatically alternates between writing the low and the high byte for each new Update-DR state, starting with the low byte for the first Update-DR encountered after entering the PROG_PAGELOAD command. The Program Counter is pre-incremented before writing the low byte, except for the first written byte. This ensures that the first data is written to the address set up by PROG_COMMANDS, and loading the last location in the page buffer does not make the Program Counter increment into the next page.
During Page Read, the content of the selected Flash byte is captured into the Flash Data Byte Register during the Capture-DR state. The AVR automatically alternates between reading the low and the high byte for each new Cap-ture-DR state, starting with the low byte for the first Capture-DR encountered after entering the PROG_PAGEREAD command. The Program Counter is post-incremented after reading each high byte, including the first read byte. This ensures that the first data is captured from the first address set up by PROG_COMMANDS, and reading the last location in the page makes the program counter increment into the next page.

Figure 30-17. Flash Data Byte Register


The state machine controlling the Flash Data Byte Register is clocked by TCK. During normal operation in which eight bits are shifted for each Flash byte, the clock cycles needed to navigate through the TAP controller automatically feeds the state machine for the Flash Data Byte Register with sufficient number of clock pulses to complete its operation transparently for the user. However, if too few bits are shifted between each Update-DR state during page load, the TAP controller should stay in the Run-Test/Idle state for some TCK cycles to ensure that there are at least 11 TCK cycles between each Update-DR state.

### 30.9.12 Programming Algorithm

All references below of type "1a", "1b", and so on, refer to Table 30-18 on page 347.

### 30.9.13 Entering Programming Mode

1. Enter JTAG instruction AVR_RESET and shift 1 in the Reset Register.
2. Enter instruction PROG_ENABLE and shift Ob1010_0011_0111_0000 in the Programming Enable Register.

### 30.9.14 Leaving Programming Mode

1. Enter JTAG instruction PROG_COMMANDS.
2. Disable all programming instructions by using no operation instruction 11a.
3. Enter instruction PROG_ENABLE and shift 0b0000_0000_0000_0000 in the programming Enable Register.
4. Enter JTAG instruction AVR_RESET and shift 0 in the Reset Register.

### 30.9.15 Performing Chip Erase

1. Enter JTAG instruction PROG_COMMANDS.
2. Start Chip Erase using programming instruction 1a.
3. Poll for Chip Erase complete using programming instruction 1b, or wait for $\mathrm{t}_{\text {WLRH_CE }}$ (refer to Table 30-14 on page 338).

### 30.9.16 Programming the Flash

Before programming the Flash a Chip Erase must be performed, see "Performing Chip Erase" on page 351.

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Flash write using programming instruction 2a.
3. Load address Extended High byte using programming instruction 2 b .
4. Load address High byte using programming instruction 2c.
5. Load address Low byte using programming instruction 2 d .
6. Load data using programming instructions $2 \mathrm{e}, 2 \mathrm{ff}$ and 2 g .
7. Repeat steps 5 and 6 for all instruction words in the page.
8. Write the page using programming instruction 2 h .
9. Poll for Flash write complete using programming instruction 2i, or wait for $\mathrm{t}_{\text {wLRH }}$ (refer to Table 30-14 on page 338).
10. Repeat steps 3 to 9 until all data have been programmed.

A more efficient data transfer can be achieved using the PROG_PAGELOAD instruction:

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Flash write using programming instruction $2 a$.
3. Load the page address using programming instructions 2 b , 2 c and 2d. PCWORD (refer to Table 30-7 on page 328) is used to address within one page and must be written as 0 .
4. Enter JTAG instruction PROG_PAGELOAD.
5. Load the entire page by shifting in all instruction words in the page byte-by-byte, starting with the LSB of the first instruction in the page and ending with the MSB of the last instruction in the page. Use Update-DR to copy the contents of the Flash Data Byte Register into the Flash page location and to auto-increment the Program Counter before each new word.
6. Enter JTAG instruction PROG_COMMANDS.
7. Write the page using programming instruction 2 h .
8. Poll for Flash write complete using programming instruction 2 i , or wait for $\mathrm{t}_{\text {wLRH }}$ (refer to Table $30-14$ on page 338).
9. Repeat steps 3 to 8 until all data have been programmed.

### 30.9.17 Reading the Flash

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Flash read using programming instruction 3 a.
3. Load address using programming instructions 3b, 3c and 3d.
4. Read data using programming instruction $3 e$.
5. Repeat steps 3 and 4 until all data have been read.

A more efficient data transfer can be achieved using the PROG_PAGEREAD instruction:

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Flash read using programming instruction 3 a .
3. Load the page address using programming instructions 3b, 3c and 3d. PCWORD (refer to Table 30-7 on page 328) is used to address within one page and must be written as 0 .
4. Enter JTAG instruction PROG_PAGEREAD.
5. Read the entire page (or Flash) by shifting out all instruction words in the page (or Flash), starting with the LSB of the first instruction in the page (Flash) and ending with the MSB of the last instruction in the page (Flash). The Capture-DR state both captures the data from the Flash, and also auto-increments the program counter after each word is read. Note that Capture-DR comes before the shift-DR state. Hence, the first byte which is shifted out contains valid data.
6. Enter JTAG instruction PROG_COMMANDS.
7. Repeat steps 3 to 6 until all data have been read.

### 30.9.18 Programming the EEPROM

Before programming the EEPROM a Chip Erase must be performed, see "Performing Chip Erase" on page 351.

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable EEPROM write using programming instruction 4 a .
3. Load address High byte using programming instruction 4b.
4. Load address Low byte using programming instruction 4c.
5. Load data using programming instructions 4d and 4 e .
6. Repeat steps 4 and 5 for all data bytes in the page.
7. Write the data using programming instruction 4 f .
8. Poll for EEPROM write complete using programming instruction 4 g , or wait for $\mathrm{t}_{\text {WLRH }}$ (refer to Table $30-14$ on page 338).
9. Repeat steps 3 to 8 until all data have been programmed.

Note that the PROG_PAGELOAD instruction can not be used when programming the EEPROM.

### 30.9.19 Reading the EEPROM

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable EEPROM read using programming instruction $5 a$.
3. Load address using programming instructions $5 b$ and $5 c$.
4. Read data using programming instruction 5d.
5. Repeat steps 3 and 4 until all data have been read.

Note that the PROG_PAGEREAD instruction can not be used when reading the EEPROM.

### 30.9.20 Programming the Fuses

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Fuse write using programming instruction 6a.
3. Load data high byte using programming instructions 6 b . A bit value of " 0 " will program the corresponding fuse, a "1" will unprogram the fuse.
4. Write Fuse High byte using programming instruction 6c.
5. Poll for Fuse write complete using programming instruction 6 d , or wait for $\mathrm{t}_{\text {WLRH }}$ (refer to Table 30-14 on page 338).
6. Load data low byte using programming instructions 6 e . A " 0 " will program the fuse, a " 1 " will unprogram the fuse.
7. Write Fuse low byte using programming instruction $6 f$.
8. Poll for Fuse write complete using programming instruction 6 g , or wait for $\mathrm{t}_{\mathrm{WLRH}}$ (refer to Table 30-14 on page 338).

### 30.9.21 Programming the Lock Bits

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Lock bit write using programming instruction 7a.
3. Load data using programming instructions 7 b . A bit value of " 0 " will program the corresponding lock bit, a "1" will leave the lock bit unchanged.
4. Write Lock bits using programming instruction 7c.
5. Poll for Lock bit write complete using programming instruction 7d, or wait for $\mathrm{t}_{\mathrm{WLRH}}$ (refer to Table 30-14 on page 338).

### 30.9.22 Reading the Fuses and Lock Bits

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Fuse/Lock bit read using programming instruction 8a.
3. To read all Fuses and Lock bits, use programming instruction 8 e . To only read Fuse High byte, use programming instruction 8b. To only read Fuse Low byte, use programming instruction 8c. To only read Lock bits, use programming instruction 8d.

### 30.9.23 Reading the Signature Bytes

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Signature byte read using programming instruction 9a.
3. Load address $0 \times 00$ using programming instruction 9 b .
4. Read first signature byte using programming instruction 9 c .
5. Repeat steps 3 and 4 with address $0 \times 01$ and address $0 \times 02$ to read the second and third signature bytes, respectively.

### 30.9.24 Reading the Calibration Byte

1. Enter JTAG instruction PROG_COMMANDS.
2. Enable Calibration byte read using programming instruction 10a.
3. Load address $0 \times 00$ using programming instruction 10 b .
4. Read the calibration byte using programming instruction 10c.

## 31. Electrical Characteristics

## Absolute Maximum Ratings*

|  | Operating Temperature ............................. $-55^{\circ} \mathrm{C}$ to $+125^{\circ} \mathrm{C}$ |
| :---: | :---: |
|  | Storage Temperature ................................ $65^{\circ} \mathrm{C}$ to $+150^{\circ} \mathrm{C}$ |
|  | Voltage on any Pin except RESET with respect to Ground $\qquad$ -0.5 V to $\mathrm{V}_{\mathrm{CC}}+0.5 \mathrm{~V}$ |
|  | Voltage on RESET with respect to Ground......-0.5V to +13.0 V |
|  | Maximum Operating Voltage ........................................ 6.0V |
|  | DC Current per I/O Pin ....................................... 40.0 mA |
|  | DC Current $\mathrm{V}_{\text {CC }}$ and GND Pins............................ 200.0 mA |

*NOTICE: Stresses beyond those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or other conditions beyond those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

### 31.1 DC Characteristics

$\mathrm{T}_{\mathrm{A}}=-40^{\circ} \mathrm{C}$ to $85^{\circ} \mathrm{C}, \mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}$ to 5.5 V (unless otherwise noted)

| Symbol | Parameter | Condition | Min. | Typ. | Max. | Units |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $V_{\text {IL }}$ | Input Low Voltage, Except XTAL1 and Reset pin | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}-2.4 \mathrm{~V} \\ & \mathrm{~V}_{\mathrm{CC}}=2.4 \mathrm{~V}-5.5 \mathrm{~V} \end{aligned}$ | $\begin{aligned} & \hline-0.5 \\ & -0.5 \end{aligned}$ |  | $\begin{aligned} & 0.2 V_{c c}^{(1)} \\ & 0.3 V_{c c}^{(1)} \end{aligned}$ | V |
| $V_{\text {IL1 }}$ | Input Low Voltage, XTAL1 pin | $\mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}-5.5 \mathrm{~V}$ | -0.5 |  | $0.1 \mathrm{~V}_{\mathrm{CC}}{ }^{(1)}$ |  |
| $\mathrm{V}_{\text {IL2 }}$ | Input Low Voltage, RESET pin | $\mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}-5.5 \mathrm{~V}$ | -0.5 |  | $0.1 \mathrm{~V}_{\mathrm{CC}}{ }^{(1)}$ |  |
| $\mathrm{V}_{\mathrm{IH}}$ | Input High Voltage, Except XTAL1 and RESET pins | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}-2.4 \mathrm{~V} \\ & \mathrm{~V}_{\mathrm{cc}}=2.4 \mathrm{~V}-5.5 \mathrm{~V} \end{aligned}$ | $\begin{aligned} & 0.7 \mathrm{~V}_{\mathrm{cc}}{ }^{(2)} \\ & 0.6 \mathrm{~V}_{\mathrm{Cc}}{ }^{(2)} \end{aligned}$ |  | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}+0.5 \\ & \mathrm{~V}_{\mathrm{CC}}+0.5 \end{aligned}$ |  |
| $\mathrm{V}_{\mathrm{IH} 1}$ | Input High Voltage, XTAL1 pin | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}-2.4 \mathrm{~V} \\ & \mathrm{~V}_{\mathrm{CC}}=2.4 \mathrm{~V}-5.5 \mathrm{~V} \end{aligned}$ | $\begin{aligned} & 0.8 \mathrm{~V}_{\mathrm{cc}}{ }^{(2)} \\ & 0.7 \mathrm{~V}_{\mathrm{Cc}}{ }^{(2)} \end{aligned}$ |  | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}+0.5 \\ & \mathrm{~V}_{\mathrm{CC}}+0.5 \end{aligned}$ |  |
| $\mathrm{V}_{\mathrm{IH} 2}$ | Input High Voltage, RESET pin | $\mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}-5.5 \mathrm{~V}$ | $0.9 \mathrm{~V}_{\mathrm{CC}}{ }^{(2)}$ |  | $\mathrm{V}_{\mathrm{CC}}+0.5$ |  |
| $\mathrm{V}_{\mathrm{OL}}$ | Output Low Voltage ${ }^{(3)}$, <br> Except RESET pin | $\begin{aligned} & \mathrm{I}_{\mathrm{OL}}=20 \mathrm{~mA}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \mathrm{I}_{\mathrm{OL}}=10 \mathrm{~mA}, \mathrm{~V}_{\mathrm{CC}}=3 \mathrm{~V} \end{aligned}$ |  |  | $\begin{aligned} & 0.9 \\ & 0.6 \end{aligned}$ |  |
| $\mathrm{V}_{\mathrm{OH}}$ | Output High Voltage ${ }^{(4)}$, Except RESET pin | $\begin{aligned} & \mathrm{I}_{\mathrm{OH}}=-20 \mathrm{~mA}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \mathrm{I}_{\mathrm{OH}}=-10 \mathrm{~mA}, \mathrm{~V}_{\mathrm{CC}}=3 \mathrm{~V} \end{aligned}$ | $\begin{aligned} & 4.2 \\ & 2.3 \end{aligned}$ |  |  |  |
| $I_{\text {IL }}$ | Input Leakage Current I/O Pin | $\mathrm{V}_{\mathrm{CC}}=5.5 \mathrm{~V}$, pin low (absolute value) |  |  | 1 | $\mu \mathrm{A}$ |
| $\mathrm{I}_{\mathrm{IH}}$ | Input Leakage Current I/O Pin | $\mathrm{V}_{\mathrm{CC}}=5.5 \mathrm{~V}$, pin high (absolute value) |  |  | 1 |  |
| $\mathrm{R}_{\text {RST }}$ | Reset Pull-up Resistor |  | 30 |  | 60 | $\mathrm{k} \Omega$ |
| $\mathrm{R}_{\mathrm{PU}}$ | I/O Pin Pull-up Resistor |  | 20 |  | 50 |  |

$\mathrm{T}_{\mathrm{A}}=-40^{\circ} \mathrm{C}$ to $85^{\circ} \mathrm{C}, \mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}$ to 5.5 V (unless otherwise noted) (Continued)

| Symbol | Parameter | Condition | Min. | Typ. | Max. | Units |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{I}_{\mathrm{CC}}$ | Power Supply Current ${ }^{(5)}$ | Active $1 \mathrm{MHz}, \mathrm{V}_{\mathrm{CC}}=2 \mathrm{~V}$ <br> (ATmega640/1280/2560/1V) |  | 0.5 | 0.8 | mA |
|  |  | Active $4 \mathrm{MHz}, \mathrm{V}_{\mathrm{CC}}=3 \mathrm{~V}$ <br> (ATmega640/1280/2560/1L) |  | 3.2 | 5 |  |
|  |  | Active $8 \mathrm{MHz}, \mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V}$ <br> (ATmega640/1280/1281/2560/2561) |  | 10 | 14 |  |
|  |  | Idle $1 \mathrm{MHz}, \mathrm{V}_{\mathrm{CC}}=2 \mathrm{~V}$ <br> (ATmega640/1280/2560/1V) |  | 0.14 | 0.22 |  |
|  |  | Idle $4 \mathrm{MHz}, \mathrm{V}_{\mathrm{CC}}=3 \mathrm{~V}$ <br> (ATmega640/1280/2560/1L) |  | 0.7 | 1.1 |  |
|  |  | Idle $8 \mathrm{MHz}, \mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V}$ <br> (ATmega640/1280/1281/2560/2561) |  | 2.7 | 4 |  |
|  | Power-down mode | WDT enabled, $\mathrm{V}_{\mathrm{CC}}=3 \mathrm{~V}$ |  | <5 | 15 | $\mu \mathrm{A}$ |
|  |  | WDT disabled, $\mathrm{V}_{\mathrm{CC}}=3 \mathrm{~V}$ |  | <1 | 7.5 |  |
| $\mathrm{V}_{\text {ACIO }}$ | Analog Comparator Input Offset Voltage | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \mathrm{~V}_{\mathrm{in}}=\mathrm{V}_{\mathrm{cc}} / 2 \end{aligned}$ |  | <10 | 40 | mV |
| $\mathrm{I}_{\text {ACLK }}$ | Analog Comparator Input Leakage Current | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \mathrm{~V}_{\mathrm{in}}=\mathrm{V}_{\mathrm{CC}} / 2 \end{aligned}$ | -50 |  | 50 | nA |
| $\mathrm{t}_{\text {ACID }}$ | Analog Comparator Propagation Delay | $\begin{aligned} & \mathrm{V}_{\mathrm{CC}}=2.7 \mathrm{~V} \\ & \mathrm{~V}_{\mathrm{CC}}=4.0 \mathrm{~V} \end{aligned}$ |  | $\begin{aligned} & 750 \\ & 500 \end{aligned}$ |  | ns |

Notes: 1. "Max" means the highest value where the pin is guaranteed to be read as low.
2. "Min" means the lowest value where the pin is guaranteed to be read as high.
3. Although each I/O port can sink more than the test conditions ( 20 mA at $\mathrm{VCC}=5 \mathrm{~V}, 10 \mathrm{~mA}$ at $\mathrm{VCC}=3 \mathrm{~V}$ ) under steady state conditions (non-transient), the following must be observed:
ATmega1281/2561:
1.) The sum of all IOL, for ports A0-A7, G2, C4-C7 should not exceed 100 mA .
2.) The sum of all IOL, for ports C0-C3, G0-G1, D0-D7 should not exceed 100 mA .
3.) The sum of all IOL, for ports G3-G5, B0-B7, E0-E7 should not exceed 100 mA .
4.) The sum of all IOL, for ports F0-F7 should not exceed 100 mA .

ATmega640/1280/2560:
1.) The sum of all IOL, for ports J0-J7, A0-A7, G2 should not exceed 200 mA .
2.) The sum of all IOL, for ports C0-C7, G0-G1, D0-D7, L0-L7 should not exceed 200 mA .
3.) The sum of all IOL, for ports G3-G4, B0-B7, H0-B7 should not exceed 200 mA .
4.) The sum of all IOL, for ports E0-E7, G5 should not exceed 100 mA .
5.) The sum of all IOL, for ports F0-F7, K0-K7 should not exceed 100 mA .

If IOL exceeds the test condition, VOL may exceed the related specification. Pins are not guaranteed to sink current greater than the listed test condition.
4. Although each $\mathrm{I} / \mathrm{O}$ port can source more than the test conditions ( 20 mA at $\mathrm{VCC}=5 \mathrm{~V}, 10 \mathrm{~mA}$ at $\mathrm{VCC}=3 \mathrm{~V}$ ) under steady state conditions (non-transient), the following must be observed:
ATmega1281/2561:
1)The sum of all IOH, for ports A0-A7, G2, C4-C7 should not exceed 100 mA .
2)The sum of all IOH, for ports C0-C3, G0-G1, D0-D7 should not exceed 100 mA .
3)The sum of all IOH, for ports G3-G5, B0-B7, E0-E7 should not exceed 100 mA .
4) The sum of all IOH, for ports F0-F7 should not exceed 100 mA .

ATmega640/1280/2560:
1)The sum of all IOH , for ports $\mathrm{J} 0-\mathrm{J} 7$, G2, A0-A7 should not exceed 200 mA .
2) The sum of all IOH, for ports C0-C7, G0-G1, D0-D7, L0-L7 should not exceed 200 mA .
3)The sum of all IOH, for ports G3-G4, B0-B7, H0-H7 should not exceed 200 mA .
4) The sum of all IOH, for ports E0-E7, G5 should not exceed 100 mA .
5) The sum of all IOH, for ports F0-F7, K0-K7 should not exceed 100 mA .

If IOH exceeds the test condition, VOH may exceed the related specification. Pins are not guaranteed to source current greater than the listed test condition.
5. Values with "PRR1 - Power Reduction Register 1" on page 56 enabled (0xFF).

### 31.2 Speed Grades

Maximum frequency is depending on $V_{c c}$. As shown in Figure 31-1 trough Figure 31-4 on page 358, the Maximum Frequency vs. $\mathrm{V}_{\mathrm{CC}}$ curve is linear between $1.8 \mathrm{~V}<\mathrm{V}_{\mathrm{Cc}}<2.7 \mathrm{~V}$ and between $2.7 \mathrm{~V}<\mathrm{V}_{\mathrm{Cc}}<4.5 \mathrm{~V}$.

### 31.2.1 8MHz

Figure 31-1. Maximum Frequency vs. $\mathrm{V}_{\mathrm{CC}}$, ATmega640V/1280V/1281V/2560V/2561V


Figure 31-2. Maximum Frequency vs. $\mathrm{V}_{\mathrm{Cc}}$ when also No-Read-While-Write Section ${ }^{(1)}$, ATmega2560V/ATmega2561V, is used


Note: 1. When only using the Read-While-Write Section of the program memory, a higher speed can be achieved at low voltage, see "Read-While-Write and No Read-While-Write Flash Sections" on page 310 for addresses.

### 31.2.2 16MHz

Figure 31-3. Maximum Frequency vs. $\mathrm{V}_{\mathrm{CC}}$, ATmega640/ATmega1280/ATmega1281


Figure 31-4. Maximum Frequency vs. $\mathrm{V}_{\mathrm{CC}}$, ATmega2560/ATmega2561


### 31.3 Clock Characteristics

### 31.3.1 Calibrated Internal RC Oscillator Accuracy

Table 31-1. Calibration Accuracy of Internal RC Oscillator

|  | Frequency | $\mathbf{V}_{\mathbf{C C}}$ | Temperature | Calibration Accuracy |
| :---: | :---: | :---: | :---: | :---: |
| Factory Calibration | 8.0 MHz | 3 V | $25^{\circ} \mathrm{C}$ | $\pm 10 \%$ |
| User Calibration | $7.3 \mathrm{MHz}-8.1 \mathrm{MHz}$ | $1.8 \mathrm{~V}-5.5 \mathrm{~V}^{(1)}$ <br> $2.7 \mathrm{~V}-5.5 \mathrm{~V}^{(2)}$ | $-40^{\circ} \mathrm{C}-85^{\circ} \mathrm{C}$ | $\pm 1 \%$ |

Notes: 1. Voltage range for ATmega640V/1281V/1280V/2561V/2560V.
2. Voltage range for ATmega640/1281/1280/2561/2560.

### 31.3.2 External Clock Drive Waveforms

Figure 31-5. External Clock Drive Waveforms


### 31.4 External Clock Drive

Table 31-2. External Clock Drive

| Symbol | Parameter | $\mathrm{V}_{\mathrm{cc}}=1.8 \mathrm{~V}-5.5 \mathrm{~V}$ |  | $\mathrm{V}_{\mathrm{cc}}=2.7 \mathrm{~V}-5.5 \mathrm{~V}$ |  | $\mathrm{V}_{\mathrm{CC}}=4.5 \mathrm{~V}-5.5 \mathrm{~V}$ |  | Units |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Min. | Max. | Min. | Max. | Min. | Max. |  |
| $1 / \mathrm{t}_{\text {CLCL }}$ | Oscillator Frequency | 0 | 2 | 0 | 8 | 0 | 16 | MHz |
| $\mathrm{t}_{\mathrm{CLCL}}$ | Clock Period | 500 |  | 125 |  | 62.5 |  |  |
| ${ }^{\text {CHCX }}$ | High Time | 200 |  | 50 |  | 25 |  | ns |
| $\mathrm{t}_{\text {CLCX }}$ | Low Time | 200 |  | 50 |  | 25 |  |  |
| $\mathrm{t}_{\mathrm{CLCH}}$ | Rise Time |  | 2.0 |  | 1.6 |  | 0.5 |  |
| $\mathrm{t}_{\mathrm{CHCL}}$ | Fall Time |  | 2.0 |  | 1.6 |  | 0.5 | $\mu \mathrm{S}$ |
| $\Delta \mathrm{t}_{\mathrm{CLCL}}$ | Change in period from one clock cycle to the next |  | 2 |  | 2 |  | 2 | \% |

### 31.5 System and Reset Characteristics

Table 31-3. Reset, Brown-out and Internal voltage CharacteristicsCharacteristics

| Symbol | Parameter | Condition | Min | Typ | Max | Units |
| :---: | :--- | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{V}_{\mathrm{RST}}$ | $\overline{R E S E T}$ Pin Threshold Voltage |  | $0.2 \mathrm{~V}_{\mathrm{CC}}$ |  | $0.9 \mathrm{~V}_{\mathrm{CC}}$ | V |
| $\mathrm{t}_{\mathrm{RST}}$ | Minimum pulse width on $\overline{R E S E T}$ Pin |  |  |  | 2.5 | $\mu \mathrm{~s}$ |
| $\mathrm{~V}_{\mathrm{HYST}}$ | Brown-out Detector Hysteresis |  |  | 50 |  | mV |
| $\mathrm{t}_{\mathrm{BOD}}$ | Min Pulse Width on Brown-out Reset |  |  | 2 |  | $\mu \mathrm{~s}$ |
| $\mathrm{~V}_{\mathrm{BG}}$ | Bandgap reference voltage | $\mathrm{V}_{\mathrm{CC}}=2.7 \mathrm{~V}, \mathrm{~T}_{\mathrm{A}}=25^{\circ} \mathrm{C}$ | 1.0 | 1.1 | 1.2 | V |
| $\mathrm{t}_{\mathrm{BG}}$ | Bandgap reference start-up time | $\mathrm{V}_{\mathrm{CC}}=2.7 \mathrm{~V}, \mathrm{~T}_{\mathrm{A}}=25^{\circ} \mathrm{C}$ |  | 40 | 70 | $\mu \mathrm{~s}$ |
| $\mathrm{I}_{\mathrm{BG}}$ | Bandgap reference current consumption | $\mathrm{V}_{\mathrm{CC}}=2.7 \mathrm{~V}, \mathrm{~T}_{\mathrm{A}}=25^{\circ} \mathrm{C}$ |  | 10 |  | $\mu \mathrm{~A}$ |

Note: 1. The Power-on Reset will not work unless the supply voltage has been below $\mathrm{V}_{\mathrm{POT}}$ (falling).

### 31.5.1 Standard Power-On Reset

This implementation of power-on reset existed in early versions of ATmega640/1280/1281/2560/2561. The table below describes the characteristics of this power-on reset and it is valid for the following devices only:

- ATmega640: revision A
- ATmega1280: revision A
- ATmega1281: revision A
- ATmega2560: revision A to E
- ATmega2561: revision A to E

Table 31-4. Characteristics of Standard Power-On Reset. $\mathrm{T}_{\mathrm{A}}=-40$ to $+85^{\circ} \mathrm{C}$.

| Symbol | Parameter | Min. $^{(1)}$ | Typ. $^{(1)}$ | Max. $^{(1)}$ | Units |
| :--- | :--- | :---: | :---: | :---: | :---: |
| $\mathrm{V}_{\text {POT }}$ | Power-on Reset Threshold Voltage (rising) ${ }^{(2)}$ | 0.7 | 1.0 | 1.4 | V |
|  | Power-on Reset Threshold Voltage (falling) ${ }^{(3)}$ | 0.05 | 0.9 | 1.3 | V |
| $\mathrm{~V}_{\text {PSR }}$ | Power-on slope rate | 0.01 |  | 4.5 | $\mathrm{~V} / \mathrm{ms}$ |

Notes: 1. Values are guidelines only.
2. Threshold where device is released from reset when voltage is rising.
3. The power-on reset threshold voltage (falling) will not work unless the supply voltage has been below $\mathrm{V}_{\text {POT }}$.

### 31.5.2 Enhanced Power-On Reset

This implementation of power-on reset exists in newer versions of ATmega640/1280/1281/2560/2561. The table below describes the characteristics of this power-on reset and it is valid for the following devices only:

- ATmega640: revision B and newer
- ATmega1280: revision B and newer
- ATmega1281: revision B and newer
- ATmega2560: revision $F$ and newer
- ATmega2561: revision F and newer

Table 31-5. Characteristics of Enhanced Power-On Reset. $\mathrm{T}_{\mathrm{A}}=-40$ to $+85^{\circ} \mathrm{C}$.

| Symbol | Parameter | Min. $^{(1)}$ | Typ. ${ }^{(1)}$ | Max. $^{(1)}$ | Units |
| :--- | :--- | :---: | :---: | :---: | :---: |
| $\mathrm{V}_{\text {POT }}$ | Power-on Reset Threshold Voltage (rising) ${ }^{(2)}$ | 1.1 | 1.4 | 1.6 | V |
|  | Power-on Reset Threshold Voltage (falling) ${ }^{(3)}$ | 0.6 | 1.3 | 1.6 | V |
| $\mathrm{~V}_{\text {PSR }}$ | Power-On Slope Rate | 0.01 |  |  | $\mathrm{~V} / \mathrm{ms}$ |

Notes: 1. Values are guidelines only.
2. Threshold where device is released from reset when voltage is rising.
3. The power-on reset threshold voltage (falling) will not work unless the supply voltage has been below $\mathrm{V}_{\text {POT }}$.

Table 31-6. BODLEVEL Fuse Coding ${ }^{(1)}$


Note: 1. $V_{\text {вот }}$ may be below nominal minimum operating voltage for some devices. For devices where this is the case, the device is tested down to $\mathrm{V}_{\mathrm{CC}}=\mathrm{V}_{\text {BOT }}$ during the production test. This guarantees that a Brown-Out Reset will occur before $\mathrm{V}_{\mathrm{CC}}$ drops to a voltage where correct operation of the microcontroller is no longer guaranteed. The test is performed using BODLEVEL $=110$ for 4 MHz operation of ATmega640V/1280V/1281V/2560V/2561V, BODLEVEL $=101$ for 8 MHz operation of ATmega640V/1280V/1281V/2560V/2561V and ATmega640/1280/1281, and BODLEVEL $=100$ for 16 MHz operation of ATmega640/1280/1281/2560/2561.

### 31.6 2-wire Serial Interface Characteristics

Table 31-7 on page 362 describes the requirements for devices connected to the 2 -wire Serial Bus. The ATmega640/1280/1281/2560/2561 2-wire Serial Interface meets or exceeds these requirements under the noted conditions.

Timing symbols refer to Figure 31-6 on page 363.

Table 31-7. 2 -wire Serial Bus Requirements

| Symbol | Parameter | Condition | Min. | Max. | Units |
| :---: | :---: | :---: | :---: | :---: | :---: |
| VIL | Input Low-voltage |  | -0.5 | $0.3 \mathrm{~V}_{\mathrm{CC}}$ | V |
| VIH | Input High-voltage |  | $0.7 \mathrm{~V}_{\mathrm{CC}}$ | $\mathrm{V}_{\mathrm{CC}}+0.5$ |  |
| Vhys ${ }^{(1)}$ | Hysteresis of Schmitt Trigger Inputs |  | $0.05 \mathrm{~V}_{\mathrm{CC}}{ }^{(2)}$ | - |  |
| $\mathrm{voL}^{(1)}$ | Output Low-voltage | 3mA sink current | 0 | 0.4 |  |
| $\mathrm{tr}^{(1)}$ | Rise Time for both SDA and SCL |  | $\begin{gathered} 20+ \\ 0.1 \mathrm{C}_{\mathrm{b}}^{(3)(2)} \end{gathered}$ | 300 | ns |
| tof ${ }^{(1)}$ | Output Fall Time from $\mathrm{V}_{\text {IHmin }}$ to $\mathrm{V}_{\text {ILmax }}$ | $10 \mathrm{pF}<\mathrm{C}_{\mathrm{b}}<400 \mathrm{pF}^{(3)}$ | $\begin{gathered} 20+ \\ 0.1 \mathrm{C}_{\mathrm{b}}{ }^{(3)(2)} \end{gathered}$ | 250 |  |
| tSP ${ }^{(1)}$ | Spikes Suppressed by Input Filter |  | 0 | $50^{(2)}$ |  |
| $\mathrm{I}_{\mathrm{i}}$ | Input Current each I/O Pin | $0.1 \mathrm{~V}_{\mathrm{CC}}<\mathrm{V}_{\mathrm{i}}<0.9 \mathrm{~V}_{\mathrm{CC}}$ | -10 | 10 | $\mu \mathrm{A}$ |
| $\mathrm{Ci}^{(1)}$ | Capacitance for each I/O Pin |  | - | 10 | pF |
| $\mathrm{f}_{\text {SCL }}$ | SCL Clock Frequency | $\begin{gathered} \mathrm{f}_{\mathrm{CK}}{ }^{(4)}>\max \left(16 \mathrm{f}_{\mathrm{SCL}},\right. \\ 250 \mathrm{kHz})^{(5)} \end{gathered}$ | 0 | 400 | kHz |
| Rp | Value of Pull-up resistor | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | $\frac{V_{C C}-0.4 \mathrm{~V}}{3 \mathrm{~mA}}$ | $\frac{1000 \mathrm{~ns}}{C_{b}}$ | $\Omega$ |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | $\frac{V_{C C}-0.4 \mathrm{~V}}{3 \mathrm{~mA}}$ | $\frac{300 \mathrm{~ns}}{C_{b}}$ |  |
| $\mathrm{t}_{\text {HD; STA }}$ | Hold Time (repeated) START Condition | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | 4.0 | - | $\mu \mathrm{s}$ |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | 0.6 | - |  |
| $\mathrm{t}_{\text {Low }}$ | Low Period of the SCL Clock | $\mathrm{f}_{\mathrm{SCL}} \leq 100 \mathrm{kHz}{ }^{(6)}$ | 4.7 | - |  |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}^{(7)}$ | 1.3 | - |  |
| $\mathrm{t}_{\mathrm{HIGH}}$ | High period of the SCL clock | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | 4.0 | - |  |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | 0.6 | - |  |
| $\mathrm{t}_{\text {SU; }}$ STA | Set-up time for a repeated START condition | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | 4.7 | - |  |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | 0.6 | - |  |
| $\mathrm{t}_{\mathrm{HD} ; \mathrm{DAT}}$ | Data hold time | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | 0 | 3.45 |  |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | 0 | 0.9 |  |
| $\mathrm{t}_{\text {SU; DAT }}$ | Data setup time | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | 250 | - |  |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | 100 | - |  |
| $\mathrm{t}_{\text {SU; }}$ | Setup time for STOP condition | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | 4.0 | - |  |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | 0.6 | - |  |
| $\mathrm{t}_{\text {BUF }}$ | Bus free time between a STOP and START condition | $\mathrm{f}_{\text {SCL }} \leq 100 \mathrm{kHz}$ | 4.7 | - |  |
|  |  | $\mathrm{f}_{\text {SCL }}>100 \mathrm{kHz}$ | 1.3 | - |  |

Note: 1. In ATmega640/1280/1281/2560/2561, this parameter is characterized and not $100 \%$ tested.
2. Required only for $f_{S C L}>100 \mathrm{kHz}$.
3. $\mathrm{C}_{\mathrm{b}}=$ capacitance of one bus line in pF .
4. $\mathrm{f}_{\mathrm{CK}}=\mathrm{CPU}$ clock frequency.
5. This requirement applies to all ATmega640/1280/1281/2560/2561 2-wire Serial Interface operation. Other devices connected to the 2 -wire Serial Bus need only obey the general $f_{S C L}$ requirement.
6. The actual low period generated by the ATmega640/1280/1281/2560/2561 2 -wire Serial Interface is $\left(1 / f_{S C L}-2 / f_{C K}\right)$, thus $\mathrm{f}_{\mathrm{CK}}$ must be greater than 6 MHz for the low time requirement to be strictly met at $\mathrm{f}_{\mathrm{SCL}}=100 \mathrm{kHz}$.
7. The actual low period generated by the ATmega640/1280/1281/2560/2561 2 -wire Serial Interface is $\left(1 / f_{S C L}-2 / f_{C K}\right)$, thus the low time requirement will not be strictly met for $\mathrm{f}_{\mathrm{SCL}}>308 \mathrm{kHz}$ when $\mathrm{f}_{\mathrm{CK}}=8 \mathrm{MHz}$. Still, ATmega640/1280/1281/2560/2561 devices connected to the bus may communicate at full speed ( 400 kHz ) with other ATmega640/1280/1281/2560/2561 devices, as well as any other device with a proper $t_{\text {Low }}$ acceptance margin.

Figure 31-6. 2-wire Serial Bus Timing


### 31.7 SPI Timing Characteristics

See Figure 31-7 and Figure 31-8 on page 364 for details.
Table 31-8. SPI Timing Parameters

|  | Description | Mode | Min. | Typ. | Max. |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | SCK period | Master |  | See Table 21-5 on page 198 |  |
| 2 | SCK high/low | Master |  | $50 \%$ duty cycle |  |
| 3 | Rise/Fall time | Master |  | 3.6 | 10 |
| 4 | Setup | Master |  | 10 |  |
| 5 | Hold | Master |  | $0.5 \cdot \mathrm{t}_{\mathrm{sck}}$ | 10 |
| 6 | Out to SCK | Master |  | 10 |  |
| 7 | SCK to out | Master |  | 15 |  |
| 8 | SCK to out high | Master |  |  |  |
| 9 | $\overline{\text { SS low to out }}$ | Slave |  |  |  |
| 10 | SCK period | Slave | $4 \bullet t_{\text {ck }}$ |  |  |
| 11 | SCK high/low ${ }^{(1)}$ | Slave | $2 \bullet \mathrm{t}_{\mathrm{ck}}$ |  |  |
| 12 | Rise/Fall time | Slave |  |  |  |
| 13 | Setup | Slave | 10 |  |  |
| 14 | Hold | Slave | $\mathrm{t}_{\mathrm{ck}}$ |  |  |
| 15 | SCK to out | Slave |  |  |  |
| 16 | SCK to $\overline{\text { SS high }}$ | Slave | 20 |  |  |
| 17 | $\overline{S S}$ high to tri-state | Slave |  |  |  |
| 18 | SS low to SCK | Slave | 20 |  |  |

Note: 1. In SPI Programming mode the minimum SCK high/low period is:
$-2 \mathrm{t}_{\mathrm{CLCL}}$ for $\mathrm{f}_{\mathrm{CK}}<12 \mathrm{MHz}$
$-3 \mathrm{t}_{\mathrm{CLCL}}$ for $\mathrm{f}_{\mathrm{CK}}>12 \mathrm{MHz}$

Figure 31-7. SPI Interface Timing Requirements (Master Mode)


Figure 31-8. SPI Interface Timing Requirements (Slave Mode)


### 31.8 ADC Characteristics - Preliminary Data

Table 31-9. ADC Characteristics, Singel Ended Channels

| Symbol | Parameter | Condition | Min. ${ }^{(1)}$ | Typ. ${ }^{(1)}$ | Max. ${ }^{(1)}$ | Units |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Resolution | Single Ended Conversion |  | 10 |  | Bits |
|  | Absolute accuracy (Including INL, DNL, quantization error, gain and offset error) | Single Ended Conversion $\begin{aligned} & \mathrm{V}_{\mathrm{REF}}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V}, \\ & \mathrm{CLK}_{\mathrm{ADC}}=200 \mathrm{kHz} \end{aligned}$ |  | 2.25 | 2.5 | LSB |
|  |  | Single Ended Conversion $\mathrm{V}_{\mathrm{REF}}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V}$ $\mathrm{CLK}_{\mathrm{ADC}}=1 \mathrm{MHz}$ |  | 3 |  |  |
|  |  | Single Ended Conversion $V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V},$ <br> $\mathrm{CLK}_{\text {ADC }}=200 \mathrm{kHz}$ <br> Noise Reduction Mode |  | 2 |  |  |
|  |  | Single Ended Conversion $V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V},$ <br> $C L K_{\text {ADC }}=1 \mathrm{MHz}$ <br> Noise Reduction Mode |  | 3 |  |  |
|  | Integral Non-Linearity (INL) | Single Ended Conversion $\begin{aligned} & \mathrm{V}_{\mathrm{REF}}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V}, \\ & \mathrm{CLK}_{\mathrm{ADC}}=200 \mathrm{kHz} \end{aligned}$ |  | 1.25 |  |  |
|  | Differential Non-Linearity (DNL) | Single Ended Conversion $\begin{aligned} & \mathrm{V}_{\mathrm{REF}}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V}, \\ & \mathrm{CLK}_{\mathrm{ADC}}=200 \mathrm{kHz} \end{aligned}$ |  | 0.5 |  |  |
|  | Gain Error | Single Ended Conversion $\begin{aligned} & \mathrm{V}_{\mathrm{REF}}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V}, \\ & \mathrm{CLK}_{\mathrm{ADC}}=200 \mathrm{kHz} \end{aligned}$ |  | 2 |  |  |
|  | Offset Error | Single Ended Conversion $\begin{aligned} & \mathrm{V}_{\mathrm{REF}}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=4 \mathrm{~V}, \\ & \mathrm{CLK}_{\mathrm{ADC}}=200 \mathrm{kHz} \end{aligned}$ |  | -2 |  |  |
|  | Conversion Time | Free Running Conversion | 13 |  | 260 | $\mu \mathrm{s}$ |
|  | Clock Frequency | Single Ended Conversion | 50 |  | 1000 | kHz |
| AVCC | Analog Supply Voltage |  | $\mathrm{V}_{C C}-0.3$ |  | $\mathrm{V}_{\mathrm{CC}}+0.3$ | V |
| $\mathrm{V}_{\text {REF }}$ | Reference Voltage |  | 1.0 |  | AVCC |  |
| $\mathrm{V}_{\text {IN }}$ | Input Voltage |  | GND |  | $V_{\text {REF }}$ |  |
|  | Input Bandwidth |  |  | 38,5 |  | kHz |
| $\mathrm{V}_{\text {INT1 }}$ | Internal Voltage Reference | 1.1 V | 1.0 | 1.1 | 1.2 | V |
| $\mathrm{V}_{\text {INT2 }}$ | Internal Voltage Reference | 2.56 V | 2.4 | 2.56 | 2.8 |  |
| $\mathrm{R}_{\text {REF }}$ | Reference Input Resistance |  |  | 32 |  | $\mathrm{k} \Omega$ |
| $\mathrm{R}_{\text {AIN }}$ | Analog Input Resistance |  |  | 100 |  | $\mathrm{M} \Omega$ |

Note: 1. Values are guidelines only.

Table 31-10. ADC Characteristics, Differential Channels

| Symbol | Parameter | Condition | Min. ${ }^{(1)}$ | Typ. ${ }^{(1)}$ | Max. ${ }^{(1)}$ | $\begin{gathered} \hline \text { Unit } \\ \mathrm{s} \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Resolution | Gain $=1 \times$ |  | 8 |  | Bits |
|  |  | Gain $=10 x$ |  | 8 |  |  |
|  |  | Gain $=200 \times$ |  | 7 |  |  |
|  | Absolute Accuracy(Including INL, DNL, Quantization Error, Gain and Offset Error) | $\begin{aligned} & \text { Gain }=1 \times \\ & V_{\text {REF }}=4 V, V_{C C}=5 \mathrm{~V} \\ & C_{\text {CLI }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 18 |  | LSB |
|  |  | $\begin{aligned} & \text { Gain }=10 x \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 17 |  |  |
|  |  | $\begin{aligned} & \text { Gain }=200 x \\ & V_{\text {REE }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 9 |  |  |
|  | Integral Non-Linearity (INL) | $\begin{aligned} & \text { Gain }=1 \times \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \\ & \hline \end{aligned}$ |  | 2.5 |  |  |
|  |  | $\begin{aligned} & \text { Gain }=10 x \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 5 |  |  |
|  |  | $\begin{aligned} & \text { Gain }=200 x \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 9 |  |  |
|  | Differential Non-Linearity (DNL) | $\begin{aligned} & \text { Gain }=1 \times \\ & V_{\text {REF }}=4 \mathrm{~V}, V_{C C}=5 \mathrm{~V} \\ & C L K_{A D C}=50-200 \mathrm{kHz} \end{aligned}$ |  | 0.75 |  |  |
|  |  | $\begin{aligned} & \text { Gain }=10 x \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \mathrm{CLK}_{\mathrm{ADC}}=50-200 \mathrm{kHz} \end{aligned}$ |  | 1.5 |  |  |
|  |  | $\begin{aligned} & \text { Gain }=200 x \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 10 |  |  |
|  | Gain Error | Gain $=1 \times$ |  | 1.7 |  | \% |
|  |  | Gain = 10x |  | 1.7 |  |  |
|  |  | Gain $=200 \times$ |  | 0.5 |  |  |
|  | Offset Error | $\begin{aligned} & \text { Gain }=1 \times \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \mathrm{CLK}_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 2 |  | LSB |
|  |  | $\begin{aligned} & \text { Gain }=10 x \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 2 |  |  |
|  |  | $\begin{aligned} & \text { Gain }=200 \times \\ & V_{\text {REF }}=4 \mathrm{~V}, \mathrm{~V}_{\mathrm{CC}}=5 \mathrm{~V} \\ & \text { CLK }_{\text {ADC }}=50-200 \mathrm{kHz} \end{aligned}$ |  | 3 |  |  |
|  | Clock Frequency |  | 50 |  | 200 | kHz |

Table 31-10. ADC Characteristics, Differential Channels (Continued)

| Symbol | Parameter | Condition | Min. ${ }^{(1)}$ | Typ. ${ }^{(1)}$ | Max. ${ }^{(1)}$ | Unit s |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Conversion Time |  | 65 |  | 260 | $\mu \mathrm{s}$ |
| AVCC | Analog Supply Voltage |  | $\mathrm{V}_{C C}-0.3$ |  | $V_{C C}+0.3$ | V |
| $\mathrm{V}_{\text {REF }}$ | Reference Voltage |  | 2.7 |  | AVCC - 0.5 |  |
| $\mathrm{V}_{\text {IN }}$ | Input Voltage |  | GND |  | $\mathrm{V}_{\mathrm{CC}}$ |  |
| $\mathrm{V}_{\text {DIFF }}$ | Input Differential Voltage |  | $-\mathrm{V}_{\text {REF }} /$ Gain |  | $\mathrm{V}_{\text {REF }} /$ Gain |  |
|  | ADC Conversion Output |  | -511 |  | 511 | LSB |
|  | Input Bandwidth |  |  | 4 |  | kHz |
| $\mathrm{V}_{\text {INT }}$ | Internal Voltage Reference |  | 2.3 | 2.56 | 2.8 | V |
| $\mathrm{R}_{\text {REF }}$ | Reference Input Resistance |  |  | 32 |  | k $\Omega$ |
| $\mathrm{R}_{\text {AIN }}$ | Analog Input Resistance |  |  | 100 |  | $\mathrm{M} \Omega$ |

Note: Values are guidelines only.

### 31.9 External Data Memory Timing

Table 31-11. External Data Memory Characteristics, 4.5 to 5.5 Volts, No Wait-state

|  | Symbol | Parameter | 8MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | $1 / \mathrm{t}_{\text {CLCL }}$ | Oscillator Frequency |  |  | 0.0 | 16 | MHz |
| 1 | $\mathrm{t}_{\text {LHLL }}$ | ALE Pulse Width | 115 |  | $1.0 \mathrm{t}_{\text {CLCL }}-10$ |  | ns |
| 2 | $\mathrm{t}_{\text {AVLL }}$ | Address Valid A to ALE Low | 57.5 |  | $0.5 \mathrm{t}_{\text {CLCL }}-5^{(1)}$ |  |  |
| 3a | $t_{\text {LLAX_St }}$ | Address Hold After ALE Low, write access | 5 |  | 5 |  |  |
| 3b | $t_{\text {LLAX_LD }}$ | Address Hold after ALE Low, read access | 5 |  | 5 |  |  |
| 4 | $\mathrm{t}_{\text {AVLLC }}$ | Address Valid C to ALE Low | 57.5 |  | $0.5 \mathrm{t}_{\mathrm{CLCL}}-5^{(1)}$ |  |  |
| 5 | $\mathrm{t}_{\text {AVRL }}$ | Address Valid to RD Low | 115 |  | $1.0 \mathrm{t}_{\text {CLCL }}-10$ |  |  |
| 6 | $\mathrm{t}_{\text {AVWL }}$ | Address Valid to WR Low | 115 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}-10$ |  |  |
| 7 | $t_{\text {LLWL }}$ | ALE Low to WR Low | 47.5 | 67.5 | $0.5 \mathrm{t}_{\text {CLCL }}-15^{(2)}$ | $0.5 \mathrm{t}_{\mathrm{CLCL}}+5^{(2)}$ |  |
| 8 | tLLRL | ALE Low to RD Low | 47.5 | 67.5 | $0.5 \mathrm{t}_{\text {CLCL }}-15^{(2)}$ | $0.5 \mathrm{t}_{\mathrm{CLCL}}+5^{(2)}$ |  |
| 9 | $\mathrm{t}_{\text {DVRH }}$ | Data Setup to RD High | 40 |  | 40 |  |  |
| 10 | $t_{\text {RLDV }}$ | Read Low to Data Valid |  | 75 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}-50$ |  |
| 11 | $\mathrm{t}_{\text {RHDX }}$ | Data Hold After RD High | 0 |  | 0 |  |  |
| 12 | $\mathrm{t}_{\text {RLRH }}$ | RD Pulse Width | 115 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}-10$ |  |  |
| 13 | $\mathrm{t}_{\text {DVWL }}$ | Data Setup to WR Low | 42.5 |  | $0.5 \mathrm{t}_{\text {CLCL }}-20^{(1)}$ |  |  |
| 14 | $\mathrm{t}_{\text {WHDX }}$ | Data Hold After WR High | 115 |  | $1.0 \mathrm{t}_{\text {CLCL }}-10$ |  |  |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 125 |  | $1.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 115 |  | $1.0 \mathrm{t}_{\text {CLCL }}-10$ |  |  |

Notes: 1. This assumes $50 \%$ clock duty cycle. The half period is actually the high time of the external clock, XTAL1.
2. This assumes $50 \%$ clock duty cycle. The half period is actually the low time of the external clock, XTAL1.

Table 31-12. External Data Memory Characteristics, 4.5 to 5.5 Volts, 1 Cycle Wait-state

|  | Symbol | Parameter | 8MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | 1/t ${ }_{\text {CLCL }}$ | Oscillator Frequency |  |  | 0.0 | 16 | MHz |
| 10 | $\mathrm{t}_{\text {RLDV }}$ | Read Low to Data Valid |  | 200 |  | $2.0 \mathrm{t}_{\text {CLCL }}-50$ | ns |
| 12 | $\mathrm{t}_{\text {RLRH }}$ | RD Pulse Width | 240 |  | 2.0t ${ }_{\text {CLCL }}-10$ |  |  |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 240 |  | $2.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 240 |  | $2.0 \mathrm{t}_{\text {CLCL }}-10$ |  |  |

Table 31-13. External Data Memory Characteristics, 4.5 to 5.5 Volts, $\mathrm{SRWn} 1=1, \mathrm{SRWn0}=0$

|  | Symbol | Parameter | 4MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | 1/t ${ }_{\text {CLCL }}$ | Oscillator Frequency |  |  | 0.0 | 16 | MHz |
| 10 | $\mathrm{t}_{\text {RLDV }}$ | Read Low to Data Valid |  | 325 |  | $3.0 \mathrm{t}_{\mathrm{CLCL}}-50$ |  |
| 12 | $\mathrm{t}_{\text {RLRH }}$ | RD Pulse Width | 365 |  | 3.04 $\mathrm{CLCL}^{-10}$ |  | ns |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 375 |  | $3.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 365 |  | $3.0 \mathrm{t}_{\mathrm{CLCL}}-10$ |  |  |

Table 31-14. External Data Memory Characteristics, 4.5 to 5.5 Volts, $\mathrm{SRWn} 1=1, \mathrm{SRWnO}=1$

|  | Symbol | Parameter | 4MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | 1/t ${ }_{\text {cLCL }}$ | Oscillator Frequency |  |  | 0.0 | 16 | MHz |
| 10 | $\mathrm{t}_{\text {RLDV }}$ | Read Low to Data Valid |  | 325 |  | $3.0 \mathrm{t}_{\mathrm{CLCL}}-50$ | ns |
| 12 | $\mathrm{t}_{\text {RLRH }}$ | RD Pulse Width | 365 |  | 3.0t $\mathrm{CLCL}^{-10}$ |  |  |
| 14 | $t_{\text {WHDX }}$ | Data Hold After WR High | 240 |  | $2.0 \mathrm{t}_{\text {CLCL }}-10$ |  |  |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 375 |  | $3.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 365 |  | $3.0 \mathrm{t}_{\text {CLCL }}-10$ |  |  |

Table 31-15. External Data Memory Characteristics, 2.7 to 5.5 Volts, No Wait-state

|  | Symbol | Parameter | 4MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | 1/t ${ }_{\text {CLCL }}$ | Oscillator Frequency |  |  | 0.0 | 8 | MHz |
| 1 | $t_{\text {LHLL }}$ | ALE Pulse Width | 235 |  | $\mathrm{t}_{\mathrm{CLCL}}{ }^{-15}$ |  | ns |
| 2 | $\mathrm{t}_{\text {AVLL }}$ | Address Valid A to ALE Low | 115 |  | $0.5 \mathrm{t}_{\text {CLCL }}-10^{(1)}$ |  |  |
| 3а | tLLAX_ST | Address Hold After ALE Low, write access | 5 |  | 5 |  |  |
| 3b | tLLAX_LD | Address Hold after ALE Low, read access | 5 |  | 5 |  |  |
| 4 | $\mathrm{t}_{\text {AVLLC }}$ | Address Valid C to ALE Low | 115 |  | $0.5 \mathrm{t}_{\text {CLCL }}-10^{(1)}$ |  |  |
| 5 | $\mathrm{t}_{\text {AVRL }}$ | Address Valid to RD Low | 235 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}{ }^{-15}$ |  |  |
| 6 | $\mathrm{t}_{\text {AVWL }}$ | Address Valid to WR Low | 235 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}{ }^{-15}$ |  |  |
| 7 | tLLWL | ALE Low to WR Low | 115 | 130 | $0.5 \mathrm{t}_{\text {CLCL }}-10^{(2)}$ | $0.5 \mathrm{t}_{\mathrm{CLCL}}+5^{(2)}$ |  |
| 8 | $\mathrm{t}_{\text {LLRL }}$ | ALE Low to RD Low | 115 | 130 | $0.5 \mathrm{t}_{\text {cLCL }}-10^{(2)}$ | $0.5 \mathrm{t}_{\mathrm{CLCL}}+5^{(2)}$ |  |
| 9 | $\mathrm{t}_{\text {DVRH }}$ | Data Setup to RD High | 45 |  | 45 |  |  |
| 10 | $t_{\text {RLDV }}$ | Read Low to Data Valid |  | 190 |  | $1.0 \mathrm{t}_{\text {CLCL- }} 60$ |  |
| 11 | $\mathrm{t}_{\text {RHDX }}$ | Data Hold After RD High | 0 |  | 0 |  |  |
| 12 | $t_{\text {RLRH }}$ | RD Pulse Width | 235 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}{ }^{-15}$ |  |  |
| 13 | $\mathrm{t}_{\text {DVWL }}$ | Data Setup to WR Low | 105 |  | $0.5 \mathrm{t}_{\text {cLCL }}-20^{(1)}$ |  |  |
| 14 | $\mathrm{t}_{\text {WHDX }}$ | Data Hold After WR High | 235 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}{ }^{-15}$ |  |  |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 250 |  | $1.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 235 |  | $1.0 \mathrm{t}_{\mathrm{CLCL}}{ }^{-15}$ |  |  |

Notes: 1. This assumes $50 \%$ clock duty cycle. The half period is actually the high time of the external clock, XTAL1.
2. This assumes $50 \%$ clock duty cycle. The half period is actually the low time of the external clock, XTAL1.

Table 31-16. External Data Memory Characteristics, 2.7 to 5.5 Volts, $\mathrm{SRWn} 1=0, \mathrm{SRWnO}=1$

|  | Symbol | Parameter | 4MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | 1/t ${ }_{\text {CLCL }}$ | Oscillator Frequency |  |  | 0.0 | 8 | MHz |
| 10 | $\mathrm{t}_{\text {RLDV }}$ | Read Low to Data Valid |  | 440 |  | $2.0 \mathrm{t}_{\text {CLCL }}-60$ | ns |
| 12 | $\mathrm{t}_{\text {RLRH }}$ | RD Pulse Width | 485 |  | $2.0 \mathrm{t}_{\text {CLCL }}-15$ |  |  |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 500 |  | $2.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 485 |  | $2.0 \mathrm{t}_{\text {CLCL }}-15$ |  |  |

Table 31-17. External Data Memory Characteristics, 2.7 to 5.5 Volts, $\mathrm{SRWn} 1=1, \mathrm{SRWn0}=0$

|  | Symbol | Parameter | 4MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | $1 / \mathrm{t}_{\text {CLCL }}$ | Oscillator Frequency |  |  | 0.0 | 8 | MHz |
| 10 | $\mathrm{t}_{\text {RLDV }}$ | Read Low to Data Valid |  | 690 |  | $3.0 \mathrm{t}_{\mathrm{CLCL}}-60$ |  |
| 12 | $t_{\text {RLRH }}$ | RD Pulse Width | 735 |  | 3.0t $\mathrm{CLCL}^{-15}$ |  |  |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 750 |  | $3.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 735 |  | $3.0 \mathrm{t}_{\text {CLCL }}-15$ |  |  |

Table 31-18. External Data Memory Characteristics, 2.7 to 5.5 Volts, $\mathrm{SRWn} 1=1, \mathrm{SRWnO}=1$

|  | Symbol | Parameter | 4MHz Oscillator |  | Variable Oscillator |  | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Min. | Max. | Min. | Max. |  |
| 0 | 1/t ${ }_{\text {CLCL }}$ | Oscillator Frequency |  |  | 0.0 | 8 | MHz |
| 10 | $t_{\text {RLDV }}$ | Read Low to Data Valid |  | 690 |  | $3.0 \mathrm{t}_{\text {CLCL }}-60$ | ns |
| 12 | $\mathrm{t}_{\text {RLRH }}$ | RD Pulse Width | 735 |  | $3.0 \mathrm{t}_{\text {CLCL }}-15$ |  |  |
| 14 | $t_{\text {WHDX }}$ | Data Hold After WR High | 485 |  | $2.0 \mathrm{t}_{\text {CLCL }}-15$ |  |  |
| 15 | $\mathrm{t}_{\text {DVWH }}$ | Data Valid to WR High | 750 |  | $3.0 \mathrm{t}_{\text {CLCL }}$ |  |  |
| 16 | $\mathrm{t}_{\text {WLWH }}$ | WR Pulse Width | 735 |  | $3.0 \mathrm{t}_{\text {CLCL }}-15$ |  |  |

Figure 31-9. External Memory Timing (SRWn1 $=0, \mathrm{SRWn0}=0$


Figure 31-10. External Memory Timing (SRWn1 = $0, \mathrm{SRWn0}=1$ )


Figure 31-11. External Memory Timing (SRWn1 = 1, SRWn0 = 0)


Figure 31-12. External Memory Timing $(S R W n 1=1, S R W n 0=1)^{0}$


The ALE pulse in the last period (T4-T7) is only present if the next instruction accesses the RAM (internal or external).

## 32. Typical Characteristics

The following charts show typical behavior. These figures are not tested during manufacturing. All current consumption measurements are performed with all I/O pins configured as inputs and with internal pull-ups enabled. A sine wave generator with rail-to-rail output is used as clock source.

All Active- and Idle current consumption measurements are done with all bits in the PRR registers set and thus, the corresponding I/O modules are turned off. Also the Analog Comparator is disabled during these measurements. Table 32-1 on page 378 and Table 32-2 on page 379 show the additional current consumption compared to $\mathrm{I}_{\mathrm{CC}}$ Active and $\mathrm{I}_{\mathrm{CC}}$ Idle for every I/O module controlled by the Power Reduction Register. See "Power Reduction Register" on page 52 for details.

The power consumption in Power-down mode is independent of clock selection.
The current consumption is a function of several factors such as: operating voltage, operating frequency, loading of I/O pins, switching rate of I/O pins, code executed and ambient temperature. The dominating factors are operating voltage and frequency.
The current drawn from capacitive loaded pins may be estimated (for one pin) as $C_{L} \times V_{C C} \times f$ where $C_{L}=$ load capacitance, $\mathrm{V}_{\mathrm{CC}}=$ operating voltage and $\mathrm{f}=$ average switching frequency of $\mathrm{I} / \mathrm{O}$ pin.

The parts are characterized at frequencies higher than test limits. Parts are not guaranteed to function properly at frequencies higher than the ordering code indicates.

The difference between current consumption in Power-down mode with Watchdog Timer enabled and Power-down mode with Watchdog Timer disabled represents the differential current drawn by the Watchdog Timer.

### 32.1 Active Supply Current

Figure 32-1. Active Supply Current vs. frequency ( $0.1 \mathrm{MHz}-1.0 \mathrm{MHz}$ )


Figure 32-2. Active Supply Current vs. Frequency ( $1 \mathrm{MHz}-16 \mathrm{MHz}$ )


Figure 32-3. Active Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Internal RC Oscillator, 8MHz)


Figure 32-4. Active Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Internal RC Oscillator, 1 MHz )


Figure 32-5. Active Supply Current vs. V ${ }_{\mathrm{CC}}$ (Internal RC Oscillator, 128 kHz )


### 32.2 Idle Supply Current

Figure 32-6. Idle Supply Current vs. Low Frequency ( $0.1 \mathrm{MHz}-1.0 \mathrm{MHz}$ )


Figure 32-7. Idle Supply Current vs. Frequency ( $1 \mathrm{MHz}-16 \mathrm{MHz}$ )


Figure 32-8. Idle Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Internal RC Oscillator, 8 MHz )


Figure 32-9. Idle Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Internal RC Oscillator, 1 MHz )


Figure 32-10. Idle Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Internal RC Oscillator, 128kHz)I


### 32.2.1 Supply Current of IO modules

The tables and formulas below can be used to calculate the additional current consumption for the different I/O modules in Active and Idle mode. The enabling or disabling of the I/O modules are controlled by the Power Reduction Register. See "Power Reduction Register" on page 52 for details.

Table 32-1. Additional Current Consumption for the different I/O modules (absolute values)

| PRR bit | Typical numbers |  |  |
| :--- | :---: | :---: | :---: |
|  | $\mathbf{V}_{\mathbf{c c}}=\mathbf{2 V}, \mathbf{F}=\mathbf{1 M H z}$ | $\mathbf{V}_{\mathbf{c c}}=\mathbf{3 V}, \mathbf{F}=\mathbf{4 M H z}$ | $\mathbf{V}_{\mathbf{c c}}=\mathbf{5 V}, \mathbf{F}=\mathbf{8 M H z}$ |
| PRUSART3 | $8.0 \mu \mathrm{~A}$ | $51 \mu \mathrm{~A}$ | $220 \mu \mathrm{~A}$ |
| PRUSART2 | $8.0 \mu \mathrm{~A}$ | $51 \mu \mathrm{~A}$ | $220 \mu \mathrm{~A}$ |
| PRUSART1 | $8.0 \mu \mathrm{~A}$ | $51 \mu \mathrm{~A}$ | $220 \mu \mathrm{~A}$ |
| PRUSART0 | $8.0 \mu \mathrm{~A}$ | $51 \mu \mathrm{~A}$ | $220 \mu \mathrm{~A}$ |
| PRTWI | $12 \mu \mathrm{~A}$ | $75 \mu \mathrm{~A}$ | $315 \mu \mathrm{~A}$ |
| PRTIM5 | $6.0 \mu \mathrm{~A}$ | $39 \mu \mathrm{~A}$ | $150 \mu \mathrm{~A}$ |
| PRTIM4 | $6.0 \mu \mathrm{~A}$ | $39 \mu \mathrm{~A}$ | $150 \mu \mathrm{~A}$ |
| PRTIM3 | $6.0 \mu \mathrm{~A}$ | $39 \mu \mathrm{~A}$ | $150 \mu \mathrm{~A}$ |
| PRTIM2 | $11 \mu \mathrm{~A}$ | $72 \mu \mathrm{~A}$ | $300 \mu \mathrm{~A}$ |
| PRTIM1 | $6.0 \mu \mathrm{~A}$ | $39 \mu \mathrm{~A}$ | $150 \mu \mathrm{~A}$ |
| PRTIM0 | $4.0 \mu \mathrm{~A}$ | $24 \mu \mathrm{~A}$ | $100 \mu \mathrm{~A}$ |
| PRSPI | $15 \mu \mathrm{~A}$ | $95 \mu \mathrm{~A}$ | $400 \mu \mathrm{~A}$ |
| PRADC | $12 \mu \mathrm{~A}$ | $75 \mu \mathrm{~A}$ | $315 \mu \mathrm{~A}$ |

Table 32-2. Additional Current Consumption (percentage) in Active and Idle mode

| PRR bit | Additional Current consumption compared to <br> Active with external clock | Additional Current consumption compared to <br> Idle with external clock |
| :--- | :---: | :---: |
| PRUSART3 | $3.0 \%$ | $17 \%$ |
| PRUSART2 | $3.0 \%$ | $17 \%$ |
| PRUSART1 | $3.0 \%$ | $17 \%$ |
| PRUSART0 | $3.0 \%$ | $17 \%$ |
| PRTWI | $4.4 \%$ | $24 \%$ |
| PRTIM5 | $1.8 \%$ | $10 \%$ |
| PRTIM4 | $1.8 \%$ | $10 \%$ |
| PRTIM3 | $1.8 \%$ | $10 \%$ |
| PRTIM2 | $4.3 \%$ | $23 \%$ |
| PRTIM1 | $1.8 \%$ | $10 \%$ |
| PRTIM0 | $1.5 \%$ | $8.0 \%$ |
| PRSPI | $3.3 \%$ | $18 \%$ |
| PRADC | $4.5 \%$ | $24 \%$ |

It is possible to calculate the typical current consumption based on the numbers from Table 32-1 on page 378 for other $\mathrm{V}_{\mathrm{CC}}$ and frequency settings than listed in Table 32-2.

### 32.2.1.1 Example 1

Calculate the expected current consumption in idle mode with USARTO, TIMER1, and TWI enabled at $V_{C C}=2.0 \mathrm{~V}$ and $F=1 \mathrm{MHz}$. From Table 32-2, third column, we see that we need to add $17 \%$ for the USART0, $24 \%$ for the TWI, and $10 \%$ for the TIMER1 module. Reading from Figure $32-6$ on page 376 , we find that the idle current consumption is $\sim 0.15 \mathrm{~mA}$ at $\mathrm{V}_{\mathrm{CC}}=2.0 \mathrm{~V}$ and $\mathrm{F}=1 \mathrm{MHz}$. The total current consumption in idle mode with USART0, TIMER1, and TWI enabled, gives:
$\mathrm{ICC}_{\text {total }} \approx 0.15 \mathrm{~mA} \bullet(1+0.17+0.24+0.10) \approx 0.227 \mathrm{~mA}$

### 32.3 Power-down Supply Current

Figure 32-11. Power-down Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Watchdog Timer Disabled)


Figure 32-12. Power-down Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Watchdog Timer Enabled)


### 32.4 Power-save Supply Current

Figure 32-13. Power-save Supply Current vs. V ${ }_{\text {CC }}$ (Watchdog Timer Disabled)


Figure 32-14. Power-save Supply Current vs. V ${ }_{\text {CC }}$ (Watchdog Timer Enabled)


### 32.5 Standby Supply Current

Figure 32-15. Standby Supply Current vs. $\mathrm{V}_{\mathrm{CC}}$ (Watchdog Timer Disabled)


### 32.6 Pin Pull-up

Figure 32-16. I/O Pin Pull-up Resistor Current vs. Input Voltage ( $\mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}$ )


Figure 32-17. I/O Pin Pull-up Resistor Current vs. Input Voltage ( $\mathrm{V}_{\mathrm{CC}}=2.7 \mathrm{~V}$ )


Figure 32-18. I/O Pin Pull-up Resistor Current vs. Input Voltage ( $\mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V}$ )


Figure 32-19. Reset Pull-up Resistor Current vs. Reset Pin Voltage ( $\mathrm{V}_{\mathrm{CC}}=1.8 \mathrm{~V}$ )


Figure 32-20. Reset pull-up Resistor Current vs. Reset Pin Voltage ( $\mathrm{V}_{\mathrm{CC}}=2.7 \mathrm{~V}$ )


Figure 32-21. Reset Pull-up Resistor Current vs. Reset Pin Voltage ( $\mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V}$ )


### 32.7 Pin Driver Strength

Figure 32-22. I/O Pin output Voltage vs.Sink Current ( $\mathrm{V}_{\mathrm{CC}}=3 \mathrm{~V}$ )


Figure 32-23. I/O Pin Output Voltage vs. Sink Current $\left(\mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V}\right)$


Figure 32-24. I/O Pin Output Voltage vs. Source Current $\left(\mathrm{V}_{\mathrm{CC}}=3 \mathrm{~V}\right)$


Figure 32-25. I/O Pin Output Voltage vs. Source Current $\left(\mathrm{V}_{\mathrm{CC}}=5 \mathrm{~V}\right)$


### 32.8 Pin Threshold and Hysteresis

Figure 32-26. I/O Pin Input Threshold Voltage vs. $\mathrm{V}_{\mathrm{CC}}\left(\mathrm{V}_{\mathrm{IH}}\right.$, IO Pin Read as "1")


Figure 32-27. I/O Pin Input Threshold Voltage vs. $\mathrm{V}_{\mathrm{CC}}$ (V $\mathrm{V}_{\mathrm{LL}}$, IO Pin Read as "0")


Figure 32-28. I/O Pin Input Hysteresis


Figure 32-29. Reset Input Threshold Voltage vs. $\mathrm{V}_{\mathrm{CC}}\left(\mathrm{V}_{\mathrm{IH}}\right.$, IO Pin Read as "1")


Figure 32-30. Reset Input Threshold Voltage vs. $\mathrm{V}_{\mathrm{CC}}$ ( $\mathrm{V}_{\mathrm{IL}}$, IO Pin Read as "0")


Figure 32-31. Reset Pin Input Hysteresis vs. $\mathrm{V}_{\mathrm{Cc}}$


### 32.9 BOD Threshold and Analog Comparator Offset

Figure 32-32. BOD Threshold vs. Temperature (BOD Level is 4.3V)


Figure 32-33. BOD Threshold vs. Temperature (BOD Level is 2.7V)


Figure 32-34. BOD Threshold vs. Temperature (BOD Level is 1.8 V )


### 32.10 Internal Oscillator Speed

Figure 32-35. Watchdog Oscillator Frequency vs. $\mathrm{V}_{\mathrm{CC}}$


Figure 32-36. Watchdog Oscillator Frequency vs. Temperature


Figure 32-37. Calibrated 8MHz RC Oscillator Frequency vs. $\mathrm{V}_{\mathrm{CC}}$


Figure 32-38. Calibrated 8MHz RC Oscillator Frequency vs. Temperature


Figure 32-39. Calibrated 8MHz RC Oscillator Frequency vs. Osccal Value


### 32.11 Current Consumption of Peripheral Units

Figure 32-40. Brownout Detector Current vs. $\mathrm{V}_{\mathrm{Cc}}$


Figure 32-41. $A D C$ Current vs. $V_{C C}\left(A R E F=A V_{C C}\right)$


Figure 32-42. AREF External Reference Current vs. $\mathrm{V}_{\mathrm{CC}}$


Figure 32-43. Watchdog Timer Current vs. $\mathrm{V}_{\mathrm{CC}}$


Figure 32-44. Analog Comparator Current vs. $\mathrm{V}_{\mathrm{CC}}$


Figure 32-45. Programming Current vs. $\mathrm{V}_{\mathrm{CC}}$


### 32.12 Current Consumption in Reset and Reset Pulsewidth

Figure 32-46. Reset Supply Current vs $\mathrm{V}_{\mathrm{CC}}(0.1 \mathrm{MHz}-1.0 \mathrm{MHz}$, Excluding Current Through The Reset Pull-up)


Figure 32-47. Reset Supply Current vs. $\mathrm{V}_{\mathrm{Cc}}(1 \mathrm{MHz}-16 \mathrm{MHz}$, Excluding Current Through The Reset Pull-up)


Figure 32-48. Minimum Reset Pulse Width vs. $\mathrm{V}_{\mathrm{CC}}$


## 33. Register Summary

| Address | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Page |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x1FF) | Reserved | - | - | - | - | - | - | - | - |  |
| ... | Reserved | - | - | - | - | - | - | - | - |  |
| (0x13F) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x13E) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x13D) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x13C) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x13B) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x13A) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x139) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x138) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x137) | Reserved |  |  |  |  |  |  |  |  |  |
| (0x136) | UDR3 | USART3 I/O Data Register |  |  |  |  |  |  |  | page 218 |
| (0x135) | UBRR3H | - | - | - | - | USART3 Baud Rate Register High Byte |  |  |  | page 222 |
| (0x134) | UBRR3L | USART3 Baud Rate Register Low Byte |  |  |  |  |  |  |  | page 222 |
| (0x133) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x132) | UCSR3C | UMSEL31 | UMSEL30 | UPM31 | UPM30 | USBS3 | UCSZ31 | UCSZ30 | UCPOL3 | page 235 |
| (0x131) | UCSR3B | RXCIE3 | TXCIE3 | UDRIE3 | RXEN3 | TXEN3 | UCSZ32 | RXB83 | TXB83 | page 234 |
| (0x130) | UCSR3A | RXC3 | TXC3 | UDRE3 | FE3 | DOR3 | UPE3 | U2X3 | MPCM3 | page 233 |
| (0x12F) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x12E) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x12D) | OCR5CH | Timer/Counter5 - Output Compare Register C High Byte |  |  |  |  |  |  |  | page 160 |
| (0x12C) | OCR5CL | Timer/Counter5 - Output Compare Register C Low Byte |  |  |  |  |  |  |  | page 160 |
| (0x12B) | OCR5BH | Timer/Counter5 - Output Compare Register B High Byte |  |  |  |  |  |  |  | page 160 |
| (0x12A) | OCR5BL | Timer/Counter5 - Output Compare Register B Low Byte |  |  |  |  |  |  |  | page 160 |
| (0x129) | OCR5AH | Timer/Counter5 - Output Compare Register A High Byte |  |  |  |  |  |  |  | page 160 |
| (0x128) | OCR5AL | Timer/Counter5 - Output Compare Register A Low Byte |  |  |  |  |  |  |  | page 160 |
| (0x127) | ICR5H | Timer/Counter5 - Input Capture Register High Byte |  |  |  |  |  |  |  | page 161 |
| (0x126) | ICR5L | Timer/Counter5 - Input Capture Register Low Byte |  |  |  |  |  |  |  | page 161 |
| (0x125) | TCNT5H | Timer/Counter5 - Counter Register High Byte |  |  |  |  |  |  |  | page 158 |
| (0x124) | TCNT5L | Timer/Counter5 - Counter Register Low Byte |  |  |  |  |  |  |  | page 158 |
| (0x123) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x122) | TCCR5C | FOC5A | FOC5B | FOC5C | - | - | - | - | - | page 157 |
| (0x121) | TCCR5B | ICNC5 | ICES5 | - | WGM53 | WGM52 | CS52 | CS51 | CS50 | page 156 |
| (0x120) | TCCR5A | COM5A1 | COM5A0 | COM5B1 | COM5B0 | COM5C1 | COM5C0 | WGM51 | WGM50 | page 154 |
| (0x11F) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x11E) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x11D) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x11C) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x11B) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x11A) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x119) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x118) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x117) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x116) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x115) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x114) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x113) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x112) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x111) | Reserved | - | - | - | $\cdot$ | - | - | - | - |  |
| (0x110) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x10F) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x10E) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x10D) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x10C) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x10B) | PORTL | PORTL7 | PORTL6 | PORTL5 | PORTL4 | PORTL3 | PORTL2 | PORTL1 | PORTLO | page 100 |
| (0x10A) | DDRL | DDL7 | DDL6 | DDL5 | DDL4 | DDL3 | DDL2 | DDL1 | DDLO | page 100 |
| (0x109) | PINL | PINL7 | PINL6 | PINL5 | PINL4 | PINL3 | PINL2 | PINL1 | PINLO | page 100 |
| (0x108) | PORTK | PORTK7 | PORTK6 | PORTK5 | PORTK4 | PORTK3 | PORTK2 | PORTK1 | PORTK0 | page 99 |
| (0x107) | DDRK | DDK7 | DDK6 | DDK5 | DDK4 | DDK3 | DDK2 | DDK1 | DDK0 | page 99 |
| (0x106) | PINK | PINK7 | PINK6 | PINK5 | PINK4 | PINK3 | PINK2 | PINK1 | PINKO | page 99 |
| (0x105) | PORTJ | PORTJ7 | PORTJ6 | PORTJ5 | PORTJ4 | PORTJ3 | PORTJ2 | PORTJ1 | PORTJO | page 99 |
| (0x104) | DDRJ | DDJ7 | DDJ6 | DDJ5 | DDJ4 | DDJ3 | DDJ2 | DDJ1 | DDJo | page 99 |
| (0x103) | PINJ | PINJ7 | PINJ6 | PINJ5 | PINJ4 | PINJ3 | PINJ2 | PINJ1 | PINJO | page 99 |
| (0x102) | PORTH | PORTH7 | PORTH6 | PORTH5 | PORTH4 | PORTH3 | PORTH2 | PORTH1 | PORTH0 | page 98 |
| (0x101) | DDRH | DDH7 | DDH6 | DDH5 | DDH4 | DDH3 | DDH2 | DDH1 | DDH0 | page 99 |


| Address | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Page |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x100) | PINH | PINH7 | PINH6 | PINH5 | PINH4 | PINH3 | PINH2 | PINH1 | PINH0 | page 99 |
| (0xFF) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xFE) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xFD) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xFC) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xFB) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xFA) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF9) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF8) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF7) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF6) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF5) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF4) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF3) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF2) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xF1) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xFO) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xEF) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xEE) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xED) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xEC) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xEB) | Reserved | - | - | - | - |  | - | - | - |  |
| (0xEA) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xE9) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xE8) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xE7) | Reserved | - | - | - | - |  | - | - | - |  |
| (0xE6) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xE5) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xE4) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xE3) | Reserved | - | - | - | - |  | - | - | - |  |
| (0xE2) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xE1) | Reserved | - | - | - | - |  | - | - | - |  |
| (0xE0) | Reserved | - | - | - | - |  | - | - | - |  |
| (0xDF) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xDE) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xDD) | Reserved | - | - | - | - |  | - | - | - |  |
| (0xDC) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xDB) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xDA) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xD9) | Reserved | - | - | - | - |  | - | - | - |  |
| (0xD8) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xD7) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xD6) | UDR2 | USART2 I/O Data Register |  |  |  |  |  |  |  | page 218 |
| (0xD5) | UBRR2H | - | - | - | - | USART2 Baud Rate Register High Byte |  |  |  | page 222 |
| (0xD4) | UBRR2L | USART2 Baud Rate Register Low Byte |  |  |  |  |  |  |  | page 222 |
| (0xD3) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xD2) | UCSR2C | UMSEL21 | UMSEL20 | UPM21 | UPM20 | USBS2 | UCSZ21 | UCSZ20 | UCPOL2 | page 235 |
| (0xD1) | UCSR2B | RXCIE2 | TXCIE2 | UDRIE2 | RXEN2 | TXEN2 | UCSZ22 | RXB82 | TXB82 | page 234 |
| (0xD0) | UCSR2A | RXC2 | TXC2 | UDRE2 | FE2 | DOR2 | UPE2 | U2X2 | MPCM2 | page 233 |
| (0xCF) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xCE) | UDR1 | USART1 I/O Data Register |  |  |  |  |  |  |  | page 218 |
| (0xCD) | UBRR1H | - | - | - | - | USART1 Baud Rate Register High Byte |  |  |  | page 222 |
| (0xCC) | UBRR1L | USART1 Baud Rate Register Low Byte |  |  |  |  |  |  |  | page 222 |
| (0xCB) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xCA) | UCSR1C | UMSEL11 | UMSEL10 | UPM11 | UPM10 | USBS1 | UCSZ11 | UCSZ10 | UCPOL1 | page 235 |
| (0xC9) | UCSR1B | RXCIE1 | TXCIE1 | UDRIE1 | RXEN1 | TXEN1 | UCSZ12 | RXB81 | TXB81 | page 234 |
| (0xC8) | UCSR1A | RXC1 | TXC1 | UDRE1 | FE1 | DOR1 | UPE1 | U2X1 | MPCM1 | page 233 |
| (0xC7) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xC6) | UDRO | USARTO I/O Data Register |  |  |  |  |  |  |  | page 218 |
| (0xC5) | UBRROH | - | - | - | - | USART0 Baud Rate Register High Byte |  |  |  | page 222 |
| (0xC4) | UBRROL | USART0 Baud Rate Register Low Byte |  |  |  |  |  |  |  | page 222 |
| (0xC3) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xC2) | UCSROC | UMSEL01 | UMSELO0 | UPM01 | UPM00 | USBS0 | UCSZ01 | UCSZ00 | UCPOLO | page 235 |
| (0xC1) | UCSROB | RXCIE0 | TXCIE0 | UDRIE0 | RXEN0 | TXENO | UCSZ02 | RXB80 | TXB80 | page 234 |
| (0xC0) | UCSROA | RXC0 | TXC0 | UDREO | FEO | DORO | UPEO | U2X0 | MPCM0 | page 234 |
| (0xBF) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xBE) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xBD) | TWAMR | TWAM6 | TWAM5 | TWAM4 | TWAM3 | TWAM2 | TWAM1 | TWAM0 | - | page 264 |


| Address | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Page |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0xBC) | TWCR | TWINT | TWEA | TWSTA | TWSTO | TWWC | TWEN | - | TWIE | page 261 |
| (0xBB) | TWDR | 2-wire Serial Interface Data Register |  |  |  |  |  |  |  | page 263 |
| (0xBA) | TWAR | TWA6 | TWA5 | TWA4 | TWA3 | TWA2 | TWA1 | TWAO | TWGCE | page 263 |
| (0xB9) | TWSR | TWS7 | TWS6 | TWS5 | TWS4 | TWS3 | - | TWPS1 | TWPS0 | page 262 |
| (0xB8) | TWBR | 2-wire Serial Interface Bit Rate Register |  |  |  |  |  |  |  | page 261 |
| (0xB7) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xB6) | ASSR | - | EXCLK | AS2 | TCN2UB | OCR2AUB | OCR2BUB | TCR2AUB | TCR2BUB | page 179 |
| (0xB5) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xB4) | OCR2B | Timer/Counter2 Output Compare Register B |  |  |  |  |  |  |  | page 186 |
| (0xB3) | OCR2A | Timer/Counter2 Output Compare Register A |  |  |  |  |  |  |  | page 186 |
| (0xB2) | TCNT2 | Timer/Counter2 (8 Bit) |  |  |  |  |  |  |  | page 186 |
| (0xB1) | TCCR2B | FOC2A | FOC2B | - | - | WGM22 | CS22 | CS21 | CS20 | page 185 |
| (0xBO) | TCCR2A | COM2A1 | COM2A0 | COM2B1 | COM2B0 | - | - | WGM21 | WGM20 | page 186 |
| (0xAF) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xAE) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xAD) | OCR4CH | Timer/Counter4 - Output Compare Register C High Byte |  |  |  |  |  |  |  | page 160 |
| (0xAC) | OCR4CL | Timer/Counter4 - Output Compare Register C Low Byte |  |  |  |  |  |  |  | page 160 |
| (0xAB) | OCR4BH | Timer/Counter4 - Output Compare Register B High Byte |  |  |  |  |  |  |  | page 160 |
| (0xAA) | OCR4BL | Timer/Counter4 - Output Compare Register B Low Byte |  |  |  |  |  |  |  | page 160 |
| (0xA9) | OCR4AH | Timer/Counter4 - Output Compare Register A High Byte |  |  |  |  |  |  |  | page 159 |
| (0xA8) | OCR4AL | Timer/Counter4 - Output Compare Register A Low Byte |  |  |  |  |  |  |  | page 159 |
| (0xA7) | ICR4H | Timer/Counter4 - Input Capture Register High Byte |  |  |  |  |  |  |  | page 161 |
| (0xA6) | ICR4L | Timer/Counter4 - Input Capture Register Low Byte |  |  |  |  |  |  |  | page 161 |
| (0xA5) | TCNT4H | Timer/Counter4-Counter Register High Byte |  |  |  |  |  |  |  | page 158 |
| (0xA4) | TCNT4L | Timer/Counter4 - Counter Register Low Byte |  |  |  |  |  |  |  | page 158 |
| (0xA3) | Reserved | - | - | - | - | - | - | - | - |  |
| (0xA2) | TCCR4C | FOC4A | FOC4B | FOC4C | - | - | - | - | - | page 157 |
| (0xA1) | TCCR4B | ICNC4 | ICES4 | - | WGM43 | WGM42 | CS42 | CS41 | CS40 | page 156 |
| (0xA0) | TCCR4A | COM4A1 | COM4A0 | COM4B1 | COM4B0 | COM4C1 | COM4C0 | WGM41 | WGM40 | page 154 |
| (0x9F) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x9E) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x9D) | OCR3CH | Timer/Counter3 - Output Compare Register C High Byte |  |  |  |  |  |  |  | page 159 |
| (0x9C) | OCR3CL | Timer/Counter3 - Output Compare Register C Low Byte |  |  |  |  |  |  |  | page 159 |
| (0x9B) | OCR3BH | Timer/Counter3 - Output Compare Register B High Byte |  |  |  |  |  |  |  | page 159 |
| (0x9A) | OCR3BL | Timer/Counter3 - Output Compare Register B Low Byte |  |  |  |  |  |  |  | page 159 |
| (0x99) | OCR3AH | Timer/Counter3 - Output Compare Register A High Byte |  |  |  |  |  |  |  | page 159 |
| (0x98) | OCR3AL | Timer/Counter3 - Output Compare Register A Low Byte |  |  |  |  |  |  |  | page 159 |
| (0x97) | ICR3H | Timer/Counter3 - Input Capture Register High Byte |  |  |  |  |  |  |  | page 161 |
| (0x96) | ICR3L | Timer/Counter3 - Input Capture Register Low Byte |  |  |  |  |  |  |  | page 161 |
| (0x95) | TCNT3H | Timer/Counter3 - Counter Register High Byte |  |  |  |  |  |  |  | page 158 |
| (0x94) | TCNT3L | Timer/Counter3 - Counter Register Low Byte |  |  |  |  |  |  |  | page 158 |
| (0x93) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x92) | TCCR3C | FOC3A | FOC3B | FOC3C | - | - | - | - | - | page 157 |
| (0x91) | TCCR3B | ICNC3 | ICES3 | - | WGM33 | WGM32 | CS32 | CS31 | CS30 | page 156 |
| (0x90) | TCCR3A | COM3A1 | COM3A0 | COM3B1 | COM3B0 | COM3C1 | COM3C0 | WGM31 | WGM30 | page 154 |
| (0x8F) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x8E) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x8D) | OCR1CH | Timer/Counter1 - Output Compare Register C High Byte |  |  |  |  |  |  |  | page 159 |
| (0x8C) | OCR1CL | Timer/Counter1 - Output Compare Register C Low Byte |  |  |  |  |  |  |  | page 159 |
| (0x8B) | OCR1BH | Timer/Counter1 - Output Compare Register B High Byte |  |  |  |  |  |  |  | page 159 |
| (0x8A) | OCR1BL | Timer/Counter1 - Output Compare Register B Low Byte |  |  |  |  |  |  |  | page 159 |
| (0x89) | OCR1AH | Timer/Counter1 - Output Compare Register A High Byte |  |  |  |  |  |  |  | page 159 |
| (0x88) | OCR1AL | Timer/Counter1 - Output Compare Register A Low Byte |  |  |  |  |  |  |  | page 159 |
| (0x87) | ICR1H | Timer/Counter1 - Input Capture Register High Byte |  |  |  |  |  |  |  | page 160 |
| (0x86) | ICR1L | Timer/Counter1 - Input Capture Register Low Byte |  |  |  |  |  |  |  | page 160 |
| (0x85) | TCNT1H | Timer/Counter1 - Counter Register High Byte |  |  |  |  |  |  |  | page 158 |
| (0x84) | TCNT1L | Timer/Counter1 - Counter Register Low Byte |  |  |  |  |  |  |  | page 158 |
| (0x83) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x82) | TCCR1C | FOC1A | FOC1B | FOC1C | - | - | - | - | - | page 157 |
| (0x81) | TCCR1B | ICNC1 | ICES1 | - | WGM13 | WGM12 | CS12 | CS11 | CS10 | page 156 |
| (0x80) | TCCR1A | COM1A1 | COM1A0 | COM1B1 | COM1B0 | COM1C1 | COM1C0 | WGM11 | WGM10 | page 154 |
| (0x7F) | DIDR1 | - | - | - | - | - | - | AIN1D | AINOD | page 267 |
| (0x7E) | DIDR0 | ADC7D | ADC6D | ADC5D | ADC4D | ADC3D | ADC2D | ADC1D | ADCOD | page 287 |
| (0x7D) | DIDR2 | ADC15D | ADC14D | ADC13D | ADC12D | ADC11D | ADC10D | ADC9D | ADC8D | page 288 |
| (0x7C) | ADMUX | REFS1 | REFSO | ADLAR | MUX4 | MUX3 | MUX2 | MUX1 | MUX0 | page 281 |
| (0x7B) | ADCSRB | - | ACME | - | - | MUX5 | ADTS2 | ADTS1 | ADTS0 | page 266, 282, 287 |
| (0x7A) | ADCSRA | ADEN | ADSC | ADATE | ADIF | ADIE | ADPS2 | ADPS1 | ADPS0 | page 285 |
| (0x79) | ADCH | ADC Data Register High byte |  |  |  |  |  |  |  | page 286 |


| Address | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Page |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (0x78) | ADCL | ADC Data Register Low byte |  |  |  |  |  |  |  | page 286 |
| (0x77) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x76) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x75) | XMCRB | XMBK | - | - | - | - | XMM2 | XMM1 | XMM0 | page 38 |
| (0x74) | XMCRA | SRE | SRL2 | SRL1 | SRLO | SRW11 | SRW10 | SRW01 | SRW00 | page 36 |
| (0x73) | TIMSK5 | - | - | ICIE5 | - | OCIE5C | OCIE5B | OCIE5A | TOIE5 | page 162 |
| (0x72) | TIMSK4 | - | - | ICIE4 | - | OCIE4C | OCIE4B | OCIE4A | TOIE4 | page 161 |
| (0x71) | TIMSK3 | - | - | ICIE3 | - | OCIE3C | OCIE3B | OCIE3A | TOIE3 | page 161 |
| (0x70) | TIMSK2 | - | - | - | - | - | OCIE2B | OCIE2A | TOIE2 | page 188 |
| (0x6F) | TIMSK1 | - | - | ICIE1 | - | OCIE1C | OCIE1B | OCIE1A | TOIE1 | page 161 |
| (0x6E) | TIMSK0 | - | - | - | - | - | OCIEOB | OCIEOA | TOIEO | page 131 |
| (0x6D) | PCMSK2 | PCINT23 | PCINT22 | PCINT21 | PCINT20 | PCINT19 | PCINT18 | PCINT17 | PCINT16 | page 113 |
| (0x6C) | PCMSK1 | PCINT15 | PCINT14 | PCINT13 | PCINT12 | PCINT11 | PCINT10 | PCINT9 | PCINT8 | page 113 |
| (0x6B) | PCMSK0 | PCINT7 | PCINT6 | PCINT5 | PCINT4 | PCINT3 | PCINT2 | PCINT1 | PCINT0 | page 114 |
| (0x6A) | EICRB | ISC71 | ISC70 | ISC61 | ISC60 | ISC51 | ISC50 | ISC41 | ISC40 | page 110 |
| (0x69) | EICRA | ISC31 | ISC30 | ISC21 | ISC20 | ISC11 | ISC10 | ISC01 | ISC00 | page 110 |
| (0x68) | PCICR | - | - | - | - | - | PCIE2 | PCIE1 | PCIEO | page 112 |
| (0x67) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x66) | OSCCAL | Oscillator Calibration Register |  |  |  |  |  |  |  | page 48 |
| (0x65) | PRR1 | - | - | PRTIM5 | PRTIM4 | PRTIM3 | PRUSART3 | PRUSART2 | PRUSART1 | page 56 |
| (0x64) | PRR0 | PRTWI | PRTIM2 | PRTIM0 | - | PRTIM1 | PRSPI | PRUSARTO | PRADC | page 55 |
| (0x63) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x62) | Reserved | - | - | - | - | - | - | - | - |  |
| (0x61) | CLKPR | CLKPCE | - | - | - | CLKPS3 | CLKPS2 | CLKPS1 | CLKPSO | page 48 |
| (0x60) | WDTCSR | WDIF | WDIE | WDP3 | WDCE | WDE | WDP2 | WDP1 | WDP0 | page 65 |
| 0x3F (0x5F) | SREG | I | T | H | S | V | N | Z | C | page 13 |
| 0x3E (0x5E) | SPH | SP15 | SP14 | SP13 | SP12 | SP11 | SP10 | SP9 | SP8 | page 15 |
| 0x3D (0x5D) | SPL | SP7 | SP6 | SP5 | SP4 | SP3 | SP2 | SP1 | SP0 | page 15 |
| 0x3C (0x5C) | EIND | - | - | - | - | - | - | - | EINDO | page 16 |
| 0x3B (0x5B) | RAMPZ | - | - | - | - | - | - | RAMPZ1 | RAMPZO | page 16 |
| 0x3A (0x5A) | Reserved | - | - | - | - | - | - | - | - |  |
| 0x39 (0x59) | Reserved | - | - | - | - | - | - | - | - |  |
| 0x38 (0x58) | Reserved | - | - | - | - | - | - | - | - |  |
| 0x37 (0x57) | SPMCSR | SPMIE | RWWSB | SIGRD | RWWSRE | BLBSET | PGWRT | PGERS | SPMEN | page 323 |
| 0x36 (0x56) | Reserved | - | - | - | - | - | - | - | - |  |
| 0x35 (0x55) | MCUCR | JTD | - | - | PUD | - | - | IVSEL | IVCE | page 64, 108, 96, 301 |
| 0x34 (0x54) | MCUSR | - | - | - | JTRF | WDRF | BORF | EXTRF | PORF | page 301 |
| 0x33 (0x53) | SMCR | - | - | - | - | SM2 | SM1 | SM0 | SE | page 50 |
| 0x32 (0x52) | Reserved | - | - | - | - | - | - | - | - |  |
| 0x31 (0x51) | OCDR | OCDR7 | OCDR6 | OCDR5 | OCDR4 | OCDR3 | OCDR2 | OCDR1 | OCDR0 | page 294 |
| 0x30 (0x50) | ACSR | ACD | ACBG | ACO | ACI | ACIE | ACIC | ACIS1 | ACISO | page 266 |
| 0x2F (0x4F) | Reserved | - | - | - | - | - | - | - | - |  |
| 0x2E (0x4E) | SPDR | SPI Data Register |  |  |  |  |  |  |  | page 199 |
| 0x2D (0x4D) | SPSR | SPIF | WCOL | - | - | - | - | - | SPI2X | page 198 |
| 0x2C (0x4C) | SPCR | SPIE | SPE | DORD | MSTR | CPOL | CPHA | SPR1 | SPR0 | page 197 |
| 0x2B (0x4B) | GPIOR2 | General Purpose I/O Register 2 |  |  |  |  |  |  |  | page 36 |
| 0x2A (0x4A) | GPIOR1 | General Purpose I/O Register 1 |  |  |  |  |  |  |  | page 36 |
| 0x29 (0x49) | Reserved | - | - | - | - | - | - | - | - |  |
| 0x28 (0x48) | OCROB | Timer/Counter0 Output Compare Register B |  |  |  |  |  |  |  | page 130 |
| 0x27 (0x47) | OCROA | Timer/Counter0 Output Compare Register A |  |  |  |  |  |  |  | page 130 |
| 0x26 (0x46) | TCNT0 | Timer/Counter0 (8 Bit) |  |  |  |  |  |  |  | page 130 |
| 0x25 (0x45) | TCCROB | FOCOA | FOCOB | - | - | WGM02 | CS02 | CS01 | CSOO | page 129 |
| 0x24 (0x44) | TCCROA | COM0A1 | COMOAO | COM0B1 | COMOB0 | - | - | WGM01 | WGM00 | page 126 |
| 0x23 (0x43) | GTCCR | TSM | - | - | - | - | - | PSRASY | PSRSYNC | page 166, 189 |
| 0x22 (0x42) | EEARH | - | - | - | - |  | PROM Addre | Register High |  | page 34 |
| 0x21 (0x41) | EEARL | EEPROM Address Register Low Byte |  |  |  |  |  |  |  | page 34 |
| 0x20 (0x40) | EEDR | EEPROM Data Register |  |  |  |  |  |  |  | page 34 |
| 0x1F (0x3F) | EECR | - | - | EEPM1 | EEPM0 | EERIE | EEMPE | EEPE | EERE | page 34 |
| 0x1E (0x3E) | GPIOR0 | General Purpose I/O Register 0 |  |  |  |  |  |  |  | page 36 |
| 0x1D (0x3D) | EIMSK | INT7 | INT6 | INT5 | INT4 | INT3 | INT2 | INT1 | INTO | page 111 |
| 0x1C (0x3C) | EIFR | INTF7 | INTF6 | INTF5 | INTF4 | INTF3 | INTF2 | INTF1 | INTF0 | page 112 |
| 0x1B (0x3B) | PCIFR | - | - | - | - | - | PCIF2 | PCIF1 | PCIFO | page 113 |
| 0x1A (0x3A) | TIFR5 | - | - | ICF5 | - | OCF5C | OCF5B | OCF5A | TOV5 | page 162 |
| 0x19 (0x39) | TIFR4 | - | - | ICF4 | - | OCF4C | OCF4B | OCF4A | TOV4 | page 162 |
| 0x18 (0x38) | TIFR3 | - | - | ICF3 | - | OCF3C | OCF3B | OCF3A | TOV3 | page 162 |
| 0x17 (0x37) | TIFR2 | - | - | - | - | - | OCF2B | OCF2A | TOV2 | page 188 |
| 0x16 (0x36) | TIFR1 | - | - | ICF1 | - | OCF1C | OCF1B | OCF1A | TOV1 | page 162 |
| 0x15 (0x35) | TIFR0 | - | - | - | - | - | OCFOB | OCFOA | TOV0 | page 131 |


| Address | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Page |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0x14 (0x34) | PORTG | - | - | PORTG5 | PORTG4 | PORTG3 | PORTG2 | PORTG1 | PORTG0 | page 98 |
| $0 \times 13$ (0x33) | DDRG | - | - | DDG5 | DDG4 | DDG3 | DDG2 | DDG1 | DDG0 | page 98 |
| $0 \times 12$ (0x32) | PING | - | - | PING5 | PING4 | PING3 | PING2 | PING1 | PINGO | page 98 |
| 0x11 (0x31) | PORTF | PORTF7 | PORTF6 | PORTF5 | PORTF4 | PORTF3 | PORTF2 | PORTF1 | PORTF0 | page 97 |
| $0 \times 10$ (0x30) | DDRF | DDF7 | DDF6 | DDF5 | DDF4 | DDF3 | DDF2 | DDF1 | DDF0 | page 98 |
| 0x0F (0x2F) | PINF | PINF7 | PINF6 | PINF5 | PINF4 | PINF3 | PINF2 | PINF1 | PINF0 | page 98 |
| 0x0E (0x2E) | PORTE | PORTE7 | PORTE6 | PORTE5 | PORTE4 | PORTE3 | PORTE2 | PORTE1 | PORTE0 | page 97 |
| 0x0D (0x2D) | DDRE | DDE7 | DDE6 | DDE5 | DDE4 | DDE3 | DDE2 | DDE1 | DDE0 | page 97 |
| 0x0C (0x2C) | PINE | PINE7 | PINE6 | PINE5 | PINE4 | PINE3 | PINE2 | PINE1 | PINE0 | page 98 |
| 0x0B (0x2B) | PORTD | PORTD7 | PORTD6 | PORTD5 | PORTD4 | PORTD3 | PORTD2 | PORTD1 | PORTD0 | page 97 |
| $0 \times 0 \mathrm{~A}(0 \times 2 \mathrm{~A})$ | DDRD | DDD7 | DDD6 | DDD5 | DDD4 | DDD3 | DDD2 | DDD1 | DDD0 | page 97 |
| $0 \times 09$ (0x29) | PIND | PIND7 | PIND6 | PIND5 | PIND4 | PIND3 | PIND2 | PIND1 | PIND0 | page 97 |
| 0x08 (0x28) | PORTC | PORTC7 | PORTC6 | PORTC5 | PORTC4 | PORTC3 | PORTC2 | PORTC1 | PORTC0 | page 97 |
| $0 \times 07$ (0x27) | DDRC | DDC7 | DDC6 | DDC5 | DDC4 | DDC3 | DDC2 | DDC1 | DDC0 | page 97 |
| $0 \times 06$ (0x26) | PINC | PINC7 | PINC6 | PINC5 | PINC4 | PINC3 | PINC2 | PINC1 | PINC0 | page 97 |
| 0x05 (0x25) | PORTB | PORTB7 | PORTB6 | PORTB5 | PORTB4 | PORTB3 | PORTB2 | PORTB1 | PORTB0 | page 96 |
| 0x04 (0x24) | DDRB | DDB7 | DDB6 | DDB5 | DDB4 | DDB3 | DDB2 | DDB1 | DDB0 | page 96 |
| $0 \times 03$ (0x23) | PINB | PINB7 | PINB6 | PINB5 | PINB4 | PINB3 | PINB2 | PINB1 | PINB0 | page 96 |
| $0 \times 02$ (0x22) | PORTA | PORTA7 | PORTA6 | PORTA5 | PORTA4 | PORTA3 | PORTA2 | PORTA1 | PORTA0 | page 96 |
| $0 \times 01(0 \times 21)$ | DDRA | DDA7 | DDA6 | DDA5 | DDA4 | DDA3 | DDA2 | DDA1 | DDA0 | page 96 |
| 0x00 (0x20) | PINA | PINA7 | PINA6 | PINA5 | PINA4 | PINA3 | PINA2 | PINA1 | PINAO | page 96 |

Notes: 1. For compatibility with future devices, reserved bits should be written to zero if accessed. Reserved I/O memory addresses should never be written.
2. I/O registers within the address range $\$ 00-\$ 1 \mathrm{~F}$ are directly bit-accessible using the SBI and CBI instructions. In these registers, the value of single bits can be checked by using the SBIS and SBIC instructions.
3. Some of the status flags are cleared by writing a logical one to them. Note that the CBI and SBI instructions will operate on all bits in the I/O register, writing a one back into any flag read as set, thus clearing the flag. The CBI and SBI instructions work with registers $0 \times 00$ to $0 \times 1 \mathrm{~F}$ only.
4. When using the I/O specific commands IN and OUT, the I/O addresses $\$ 00-\$ 3 F$ must be used. When addressing I/O registers as data space using LD and ST instructions, $\$ 20$ must be added to these addresses. The
ATmega640/1280/1281/2560/2561 is a complex microcontroller with more peripheral units than can be supported within the 64 location reserved in Opcode for the IN and OUT instructions. For the Extended I/O space from $\$ 60-\$ 1 F F$ in SRAM, only the ST/STS/STD and LD/LDS/LDD instructions can be used.

## 34. Instruction Set Summary

| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| ARITHMETIC AND LOGIC INSTRUCTIONS |  |  |  |  |  |
| ADD | Rd, Rr | Add two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}$ | Z, C, N, V, H | 1 |
| ADC | Rd, Rr | Add with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}+\mathrm{C}$ | Z, C, N, V, H | 1 |
| ADIW | Rdl, K | Add Immediate to Word | Rdh:Rdl $\leftarrow$ Rdh:Rdl + K | Z, C, N, V, S | 2 |
| SUB | Rd, Rr | Subtract two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}$ | Z, C, N, V, H | 1 |
| SUBI | Rd, K | Subtract Constant from Register | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}$ | Z, C, N, V, H | 1 |
| SBC | Rd, Rr | Subtract with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z, C, N, V, H | 1 |
| SBCI | Rd, K | Subtract with Carry Constant from Reg. | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}-\mathrm{C}$ | Z, C, N, V, H | 1 |
| SBIW | Rdi, K | Subtract Immediate from Word | Rdh:RdI $\leftarrow$ Rdh:Rdl - K | Z, C, N, V, S | 2 |
| AND | Rd, Rr | Logical AND Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rr}$ | Z, N, V | 1 |
| ANDI | Rd, K | Logical AND Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{K}$ | Z, N, V | 1 |
| OR | Rd, Rr | Logical OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rdv} \mathrm{Rr}$ | Z, N, V | 1 |
| ORI | Rd, K | Logical OR Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rd}$ v K | Z, N, V | 1 |
| EOR | Rd, Rr | Exclusive OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rr}$ | Z, N, V | 1 |
| COM | Rd | One's Complement | $\mathrm{Rd} \leftarrow 0 \mathrm{xFF}-\mathrm{Rd}$ | Z, C, N, V | 1 |
| NEG | Rd | Two's Complement | Rd $\leftarrow 0 \times 00-\mathrm{Rd}$ | Z, C, N, V, H | 1 |
| SBR | Rd, K | Set Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \mathrm{vK}$ | Z, N, V | 1 |
| CBR | Rd, K | Clear Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet(0 x \mathrm{FF}-\mathrm{K})$ | Z, N, V | 1 |
| INC | Rd | Increment | $\mathrm{Rd} \leftarrow \mathrm{Rd}+1$ | Z, N, V | 1 |
| DEC | Rd | Decrement | $\mathrm{Rd} \leftarrow \mathrm{Rd}-1$ | Z, N, V | 1 |
| TST | Rd | Test for Zero or Minus | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rd}$ | Z, N, V | 1 |
| CLR | Rd | Clear Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rd}$ | Z, N, V | 1 |
| SER | Rd | Set Register | $\mathrm{Rd} \leftarrow 0 \mathrm{xFF}$ | None | 1 |
| MUL | Rd, Rr | Multiply Unsigned | $\mathrm{R} 1: \mathrm{R0} \leftarrow \mathrm{Rd} \times \mathrm{Rr}$ | Z, C | 2 |
| MULS | Rd, Rr | Multiply Signed | $\mathrm{R} 1: \mathrm{R0} 0 \leftarrow \mathrm{Rdx} \times \mathrm{Rr}$ | Z, C | 2 |
| MULSU | Rd, Rr | Multiply Signed with Unsigned | $\mathrm{R} 1: \mathrm{R0} 5 \mathrm{Rd} \times \mathrm{Rr}$ | Z, C | 2 |
| FMUL | Rd, Rr | Fractional Multiply Unsigned | $\mathrm{R} 1: \mathrm{RO} \leftarrow(\mathrm{Rd} \times \mathrm{Rr}) \ll 1$ | Z, C | 2 |
| FMULS | Rd, Rr | Fractional Multiply Signed | $\mathrm{R} 1: \mathrm{RO} \leftarrow(\mathrm{Rd} \times \mathrm{Rr}) \ll 1$ | Z, C | 2 |
| FMULSU | Rd, Rr | Fractional Multiply Signed with Unsigned | $\mathrm{R} 1: \mathrm{R} 0 \leftarrow(\mathrm{Rd} \times \mathrm{Rr}) \ll 1$ | Z, C | 2 |
| BRANCH INSTRUCTIONS |  |  |  |  |  |
| RJMP | k | Relative Jump | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 2 |
| IJMP |  | Indirect Jump to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 2 |
| EIJMP |  | Extended Indirect Jump to (Z) | PC $\leftarrow$ (EIND:Z) | None | 2 |
| JMP | k | Direct Jump | $\mathrm{PC} \leftarrow \mathrm{k}$ | None | 3 |
| RCALL | k | Relative Subroutine Call | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 4 |
| ICALL |  | Indirect Call to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 4 |
| EICALL |  | Extended Indirect Call to (Z) | PC $\leftarrow(E I N D: Z)$ | None | 4 |
| CALL | k | Direct Subroutine Call | $\mathrm{PC} \leftarrow \mathrm{k}$ | None | 5 |
| RET |  | Subroutine Return | $\mathrm{PC} \leftarrow$ STACK | None | 5 |
| RETI |  | Interrupt Return | $\mathrm{PC} \leftarrow$ STACK | 1 | 5 |
| CPSE | Rd, Rr | Compare, Skip if Equal | if ( $\mathrm{Rd}=\mathrm{Rr}$ ) $\mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2/3 |
| CP | Rd, Rr | Compare | $\mathrm{Rd}-\mathrm{Rr}$ | Z, N, V, C, H | 1 |
| CPC | Rd, Rr | Compare with Carry | $\mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z, N, V, C, H | 1 |
| CPI | Rd,K | Compare Register with Immediate | Rd-K | Z, N, V, C, H | 1 |
| SBRC | Rr, b | Skip if Bit in Register Cleared | if $(\operatorname{Rr}(\mathrm{b})=0) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2/3 |
| SBRS | $\mathrm{Rr}, \mathrm{b}$ | Skip if Bit in Register is Set | if $(\operatorname{Rr}(\mathrm{b})=1) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2/3 |
| SBIC | P, b | Skip if Bit in I/O Register Cleared | if $(P(b)=0) P C \leftarrow P C+2$ or 3 | None | 1/2/3 |
| SBIS | P, b | Skip if Bit in I/O Register is Set | if $(P(b)=1) P C \leftarrow P C+2$ or 3 | None | 1/2/3 |
| BRBS | s, k | Branch if Status Flag Set | if (SREG(s) $=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRBC | s, k | Branch if Status Flag Cleared | if (SREG(s) $=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BREQ | k | Branch if Equal | if $(\mathrm{Z}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRNE | k | Branch if Not Equal | if $(Z=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRCS | k | Branch if Carry Set | if ( $\mathrm{C}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRCC | k | Branch if Carry Cleared | if ( $\mathrm{C}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRSH | k | Branch if Same or Higher | if ( $\mathrm{C}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRLO | k | Branch if Lower | if ( $\mathrm{C}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRMI | k | Branch if Minus | if ( $\mathrm{N}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRPL | k | Branch if Plus | if ( $\mathrm{N}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRGE | k | Branch if Greater or Equal, Signed | if ( $\mathrm{N} \oplus \mathrm{V}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRLT | k | Branch if Less Than Zero, Signed | if ( $\mathrm{N} \oplus \mathrm{V}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRHS | k | Branch if Half Carry Flag Set | if $(\mathrm{H}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRHC | k | Branch if Half Carry Flag Cleared | if ( $\mathrm{H}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRTS | k | Branch if T Flag Set | if ( $\mathrm{T}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRTC | k | Branch if T Flag Cleared | if $(\mathrm{T}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRVS | k | Branch if Overflow Flag is Set | if ( $\mathrm{V}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |


| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| BRVC | k | Branch if Overflow Flag is Cleared | if ( $\mathrm{V}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRIE | k | Branch if Interrupt Enabled | if ( $\mathrm{I}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRID | k | Branch if Interrupt Disabled | if $(\mathrm{I}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BIT AND BIT-TEST INSTRUCTIONS |  |  |  |  |  |
| SBI | P, b | Set Bit in I/O Register | $\mathrm{I} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 1$ | None | 2 |
| CBI | P, b | Clear Bit in I/O Register | $\mathrm{l} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 0$ | None | 2 |
| LSL | Rd | Logical Shift Left | $\operatorname{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \mathrm{Rd}(0) \leftarrow 0$ | Z, C, N, V | 1 |
| LSR | Rd | Logical Shift Right | $\operatorname{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \operatorname{Rd}(7) \leftarrow 0$ | Z, C, N, V | 1 |
| ROL | Rd | Rotate Left Through Carry | $\operatorname{Rd}(0) \leftarrow \mathrm{C}, \mathrm{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \mathrm{C} \leftarrow \operatorname{Rd}(7)$ | Z, C, N, V | 1 |
| ROR | Rd | Rotate Right Through Carry | $\operatorname{Rd}(7) \leftarrow \mathrm{C}, \operatorname{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \mathrm{C} \leftarrow \operatorname{Rd}(0)$ | Z, C, N, V | 1 |
| ASR | Rd | Arithmetic Shift Right | $\mathrm{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \mathrm{n}=0 . .6$ | Z, C, N, V | 1 |
| SWAP | Rd | Swap Nibbles | $\operatorname{Rd}(3 . .0) \leftarrow \operatorname{Rd}(7 . .4), \operatorname{Rd}(7 . .4) \leftarrow \operatorname{Rd}(3 . .0)$ | None | 1 |
| BSET | s | Flag Set | SREG(s) $\leftarrow 1$ | SREG(s) | 1 |
| BCLR | s | Flag Clear | SREG(s) $\leftarrow 0$ | SREG(s) | 1 |
| BST | $\mathrm{Rr}, \mathrm{b}$ | Bit Store from Register to $T$ | $\mathrm{T} \leftarrow \operatorname{Rr}(\mathrm{b})$ | T | 1 |
| BLD | Rd, b | Bit load from T to Register | $\operatorname{Rd}(\mathrm{b}) \leftarrow \mathrm{T}$ | None | 1 |
| SEC |  | Set Carry | $\mathrm{C} \leftarrow 1$ | C | 1 |
| CLC |  | Clear Carry | $\mathrm{C} \leftarrow 0$ | C | 1 |
| SEN |  | Set Negative Flag | $\mathrm{N} \leftarrow 1$ | N | 1 |
| CLN |  | Clear Negative Flag | $\mathrm{N} \leftarrow 0$ | N | 1 |
| SEZ |  | Set Zero Flag | $\mathrm{Z} \leftarrow 1$ | Z | 1 |
| CLZ |  | Clear Zero Flag | $\mathrm{Z} \leftarrow 0$ | Z | 1 |
| SEI |  | Global Interrupt Enable | $1 \leftarrow 1$ | 1 | 1 |
| CLI |  | Global Interrupt Disable | $1 \leftarrow 0$ | 1 | 1 |
| SES |  | Set Signed Test Flag | $\mathrm{S} \leftarrow 1$ | S | 1 |
| CLS |  | Clear Signed Test Flag | $\mathrm{S} \leftarrow 0$ | S | 1 |
| SEV |  | Set Twos Complement Overflow. | $\mathrm{V} \leftarrow 1$ | V | 1 |
| CLV |  | Clear Twos Complement Overflow | $\mathrm{V} \leftarrow 0$ | V | 1 |
| SET |  | Set T in SREG | $\mathrm{T} \leftarrow 1$ | T | 1 |
| CLT |  | Clear T in SREG | $\mathrm{T} \leftarrow 0$ | T | 1 |
| SEH |  | Set Half Carry Flag in SREG | $\mathrm{H} \leftarrow 1$ | H | 1 |
| CLH |  | Clear Half Carry Flag in SREG | $\mathrm{H} \leftarrow 0$ | H | 1 |
| DATA TRANSFER INSTRUCTIONS |  |  |  |  |  |
| MOV | Rd, Rr | Move Between Registers | $\mathrm{Rd} \leftarrow \mathrm{Rr}$ | None | 1 |
| MOVW | Rd, Rr | Copy Register Word | $\mathrm{Rd}+1: \mathrm{Rd} \leftarrow \mathrm{Rr}+1: \mathrm{Rr}$ | None | 1 |
| LDI | Rd, K | Load Immediate | $\mathrm{Rd} \leftarrow \mathrm{K}$ | None | 1 |
| LD | Rd, X | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{X})$ | None | 2 |
| LD | Rd, $\mathrm{X}_{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{X}), \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| LD | Rd, - X | Load Indirect and Pre-Dec. | $X \leftarrow X-1, R d \leftarrow(X)$ | None | 2 |
| LD | Rd, Y | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LD | Rd, $\mathrm{Y}+$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Y}), \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| LD | Rd, - Y | Load Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1, \mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LDD | Rd, $\mathrm{Y}+\mathrm{q}$ | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Y}+\mathrm{q})$ | None | 2 |
| LD | Rd, Z | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LD | Rd, $\mathrm{Z}_{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Z}), \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| LD | Rd, -Z | Load Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1, \mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LDD | Rd, $\mathrm{Z}+\mathrm{q}$ | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Z}+\mathrm{q})$ | None | 2 |
| LDS | Rd, k | Load Direct from SRAM | $\mathrm{Rd} \leftarrow(\mathrm{k})$ | None | 2 |
| ST | $\mathrm{X}, \mathrm{Rr}$ | Store Indirect | $(\mathrm{X}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | X + , Rr | Store Indirect and Post-Inc. | $(\mathrm{X}) \leftarrow \mathrm{Rr}, \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| ST | - $\mathrm{X}, \mathrm{Rr}$ | Store Indirect and Pre-Dec. | $\mathrm{X} \leftarrow \mathrm{X}-1,(\mathrm{X}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{Y}, \mathrm{Rr}$ | Store Indirect | $(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | Y + , Rr | Store Indirect and Post-Inc. | $(\mathrm{Y}) \leftarrow \mathrm{Rr}, \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| ST | - $\mathrm{Y}, \mathrm{Rr}$ | Store Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1,(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | $\mathrm{Y}+\mathrm{q}, \mathrm{Rr}$ | Store Indirect with Displacement | $(\mathrm{Y}+\mathrm{q}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{Z}, \mathrm{Rr}$ | Store Indirect | $(\mathrm{Z}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | Z + , Rr | Store Indirect and Post-Inc. | $(\mathrm{Z}) \leftarrow \mathrm{Rr}, \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| ST | -Z, Rr | Store Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1,(\mathrm{Z}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | Z $+\mathrm{q}, \mathrm{Rr}$ | Store Indirect with Displacement | $(Z+q) \leftarrow \operatorname{Rr}$ | None | 2 |
| STS | k, Rr | Store Direct to SRAM | $(\mathrm{k}) \leftarrow \mathrm{Rr}$ | None | 2 |
| LPM |  | Load Program Memory | $\mathrm{R} 0 \leftarrow(\mathrm{Z})$ | None | 3 |
| LPM | Rd, Z | Load Program Memory | $\mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 3 |
| LPM | Rd, $\mathrm{Z}^{+}$ | Load Program Memory and Post-Inc | $\mathrm{Rd} \leftarrow(\mathrm{Z}), \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 3 |
| ELPM |  | Extended Load Program Memory | R0 $\leftarrow($ RAMPZ:Z) | None | 3 |
| ELPM | Rd, Z | Extended Load Program Memory | $\mathrm{Rd} \leftarrow(\mathrm{RAMPZ}: Z)$ | None | 3 |
| ELPM | Rd, $\mathrm{Z}_{+}$ | Extended Load Program Memory | Rd $\leftarrow($ RAMPZ:Z), RAMPZ:Z $\leftarrow$ RAMPZ:Z +1 | None | 3 |
| SPM |  | Store Program Memory | (Z) $\leftarrow$ R1:R0 | None | - |
| IN | Rd, P | In Port | $\mathrm{Rd} \leftarrow \mathrm{P}$ | None | 1 |


| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| OUT | P, Rr | Out Port | $\mathrm{P} \leftarrow \mathrm{Rr}$ | None | 1 |
| PUSH | Rr | Push Register on Stack | STACK $\leftarrow \mathrm{Rr}$ | None | 2 |
| POP | Rd | Pop Register from Stack | Rd $\leftarrow$ STACK | None | 2 |
| MCU CONTROL INSTRUCTIONS |  |  |  |  |  |
| NOP |  | No Operation |  | None | 1 |
| SLEEP |  | Sleep | (see specific descr. for Sleep function) | None | 1 |
| WDR |  | Watchdog Reset | (see specific descr. for WDR/timer) | None | 1 |
| BREAK |  | Break | For On-chip Debug Only | None | N/A |

Note: EICALL and EIJMP do not exist in ATmega640/1280/1281. ELPM does not exist in ATmega640.

## 35. Ordering Information

### 35.1 ATmega640

| Speed [MHz] ${ }^{(2)}$ | Power Supply | Ordering Code | Package ${ }^{(1)(3)}$ | Operation Range |
| :---: | :---: | :---: | :---: | :---: |
| 8 | 1.8-5.5V | ATmega640V-8AU ATmega640V-8AUR ${ }^{(4)}$ <br> ATmega640V-8CU <br> ATmega640V-8CUR ${ }^{(4)}$ | $\begin{aligned} & \hline 100 \mathrm{~A} \\ & 100 \mathrm{~A} \\ & 100 \mathrm{C} 1 \\ & 100 \mathrm{C} 1 \end{aligned}$ | Industrial ( $-40^{\circ} \mathrm{C}$ to $85^{\circ} \mathrm{C}$ ) |
| 16 | 2.7-5.5V | ATmega640-16AU <br> ATmega640-16AUR ${ }^{(4)}$ <br> ATmega640-16CU <br> ATmega640-16CUR ${ }^{(4)}$ | $\begin{aligned} & 100 \mathrm{~A} \\ & 100 \mathrm{~A} \\ & 100 \mathrm{C} 1 \\ & 100 \mathrm{C} 1 \end{aligned}$ |  |

Notes: 1. This device can also be supplied in wafer form. Contact your local Atmel sales office for detailed ordering information and minimum quantities.
2. See "Speed Grades" on page 357.
3. Pb-free packaging, complies to the European Directive for Restriction of Hazardous Substances (RoHS directive). Also Halide free and fully Green.
4. Tape \& Reel.

## Package Type

| 100A | 100-lead, Thin (1.0mm) Plastic Gull Wing Quad Flat Package (TQFP) |
| :--- | :--- |
| $\mathbf{1 0 0 C 1}$ | 100-ball, Chip Ball Grid Array (CBGA) |

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### 35.2 ATmega1280

| Speed [MHz] ${ }^{(2)}$ | Power Supply | Ordering Code | Package ${ }^{(1)(3)}$ | Operation Range |
| :---: | :---: | :---: | :---: | :---: |
| 8 | 1.8V-5.5V | ATmega1280V-8AU <br> ATmega1280V-8AUR ${ }^{(4)}$ <br> ATmega1280V-8CU <br> ATmega1280V-8CUR ${ }^{(4)}$ | $\begin{aligned} & 100 \mathrm{~A} \\ & 100 \mathrm{~A} \\ & 100 \mathrm{C} 1 \\ & 100 \mathrm{C} 1 \end{aligned}$ | Industrial ( $-40^{\circ} \mathrm{C}$ to $85^{\circ} \mathrm{C}$ ) |
| 16 | 2.7V-5.5V | ATmega1280-16AU <br> ATmega1280-16AUR ${ }^{(4)}$ <br> ATmega1280-16CU <br> ATmega1280-16CUR ${ }^{(4)}$ | $\begin{aligned} & 100 \mathrm{~A} \\ & 100 \mathrm{~A} \\ & 100 \mathrm{C} 1 \\ & 100 \mathrm{C} 1 \end{aligned}$ |  |

Notes: 1. This device can also be supplied in wafer form. Contact your local Atmel sales office for detailed ordering information and minimum quantities.
2. See "Speed Grades" on page 357.
3. Pb-free packaging, complies to the European Directive for Restriction of Hazardous Substances (RoHS directive). Also Halide free and fully Green.
4. Tape \& Reel.

## Package Type

| 100A | 100-lead, Thin (1.0mm) Plastic Gull Wing Quad Flat Package (TQFP) |
| :--- | :--- |
| $\mathbf{1 0 0 C 1}$ | 100-ball, Chip Ball Grid Array (CBGA) |

### 35.3 ATmega1281

| Speed [MHz] ${ }^{(2)}$ | Power Supply | Ordering Code | Package ${ }^{(1)(3)}$ | Operation Range |
| :---: | :---: | :---: | :---: | :---: |
| 8 | 1.8-5.5V | ATmega1281V-8AU <br> ATmega1281V-8AUR ${ }^{(4)}$ <br> ATmega1281V-8MU <br> ATmega1281V-8MUR ${ }^{(4)}$ | 64A <br> 64A <br> 64M2 <br> 64M2 | $\begin{gathered} \text { Industrial } \\ \left(-40^{\circ} \mathrm{C} \text { to } 85^{\circ} \mathrm{C}\right) \end{gathered}$ |
| 16 | 2.7-5.5V | ATmega1281-16AU <br> ATmega1281-16AUR ${ }^{(4)}$ <br> ATmega1281-16MU <br> ATmega1281-16MUR ${ }^{(4)}$ | $\begin{aligned} & \text { 64A } \\ & 64 \mathrm{~A} \\ & 64 \mathrm{M} 2 \\ & 64 \mathrm{M} 2 \end{aligned}$ |  |

Notes: 1. This device can also be supplied in wafer form. Contact your local Atmel sales office for detailed ordering information and minimum quantities.
2. See "Speed Grades" on page 357.
3. Pb-free packaging, complies to the European Directive for Restriction of Hazardous Substances (RoHS directive). Also Halide free and fully Green.
4. Tape \& Reel.

## Package Type

| 64A | 64-lead, Thin (1.0mm) Plastic Gull Wing Quad Flat Package (TQFP) |
| :--- | :--- |
| 64M2 | 64 -pad, $9 \mathrm{~mm} \times 9 \mathrm{~mm} \times 1.0 \mathrm{~mm}$ Body, Quad Flat No-lead/Micro Lead Frame Package (QFN/MLF) |

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### 35.4 ATmega2560

| Speed [MHz] ${ }^{(2)}$ | Power Supply | Ordering Code | Package ${ }^{(1)(3)}$ | Operation Range |
| :---: | :---: | :---: | :---: | :---: |
| 8 | 1.8V-5.5V | ATmega2560V-8AU <br> ATmega2560V-8AUR ${ }^{(4)}$ <br> ATmega2560V-8CU <br> ATmega2560V-8CUR ${ }^{(4)}$ | $\begin{gathered} \hline 100 \mathrm{~A} \\ 100 \mathrm{~A} \\ 100 \mathrm{C} 1 \\ 100 \mathrm{C} 1 \end{gathered}$ | Industrial ( $-40^{\circ} \mathrm{C}$ to $85^{\circ} \mathrm{C}$ ) |
| 16 | 4.5V-5.5V | ATmega2560-16AU <br> ATmega2560-16AUR ${ }^{(4)}$ <br> ATmega2560-16CU <br> ATmega2560-16CUR ${ }^{(4)}$ | $\begin{gathered} 100 \mathrm{~A} \\ 100 \mathrm{~A} \\ 100 \mathrm{C} 1 \\ 100 \mathrm{C} 1 \end{gathered}$ |  |

Notes: 1. This device can also be supplied in wafer form. Contact your local Atmel sales office for detailed ordering information and minimum quantities.
2. See "Speed Grades" on page 357.
3. Pb-free packaging, complies to the European Directive for Restriction of Hazardous Substances (RoHS directive). Also Halide free and fully Green.
4. Tape \& Reel.

## Package Type

| 100A | 100-lead, Thin (1.Omm) Plastic Gull Wing Quad Flat Package (TQFP) |
| :--- | :--- |
| $\mathbf{1 0 0 C 1}$ | 100-ball, Chip Ball Grid Array (CBGA) |

### 35.5 ATmega2561

| Speed [MHz] ${ }^{(2)}$ | Power Supply | Ordering Code | Package ${ }^{(1)(3)}$ | Operation Range |
| :---: | :---: | :---: | :---: | :---: |
| 8 | 1.8V-5.5V | ATmega2561V-8AU <br> ATmega2561V-8AUR ${ }^{(4)}$ <br> ATmega2561V-8MU <br> ATmega2561V-8MUR ${ }^{(4)}$ | $\begin{aligned} & \hline 64 \mathrm{~A} \\ & 64 \mathrm{~A} \\ & 64 \mathrm{M} 2 \\ & 64 \mathrm{M} 2 \end{aligned}$ | $\begin{gathered} \text { Industrial } \\ \left(-40^{\circ} \mathrm{C} \text { to } 85^{\circ} \mathrm{C}\right) \end{gathered}$ |
| 16 | 4.5V-5.5V | ATmega2561-16AU <br> ATmega2561-16AUR ${ }^{(4)}$ <br> ATmega2561-16MU <br> ATmega2561-16MUR ${ }^{(4)}$ | $\begin{aligned} & \text { 64A } \\ & 64 \mathrm{~A} \\ & 64 \mathrm{M} 2 \\ & 64 \mathrm{M} 2 \end{aligned}$ |  |

Notes: 1. This device can also be supplied in wafer form.Contact your local Atmel sales office for detailed ordering information and minimum quantities.
2. See "Speed Grades" on page 357.
3. Pb-free packaging, complies to the European Directive for Restriction of Hazardous Substances (RoHS directive). Also Halide free and fully Green.
4. Tape \& Reel.

## Package Type

| 64A | 64-lead, Thin (1.0mm) Plastic Gull Wing Quad Flat Package (TQFP) |
| :--- | :--- |
| 64M2 | $64-$ pad, $9 \mathrm{~mm} \times 9 \mathrm{~mm} \times 1.0 \mathrm{~mm}$ Body, Quad Flat No-lead/Micro Lead Frame Package (QFN/MLF) |

## 36. Packaging Information

### 36.1 100A



Notes:

1. This package conforms to JEDEC reference MS-026, Variation AED.
2. Dimensions D1 and E1 do not include mold protrusion. Allowable protrusion is 0.25 mm per side. Dimensions D1 and E1 are maximum plastic body size dimensions including mold mismatch.
3. Lead coplanarity is 0.08 mm maximum.

COMMON DIMENSIONS
(Unit of Measure = mm)

| SYMBOL | MIN | NOM | MAX | NOTE |
| :---: | :---: | :---: | :---: | :---: |
| A | - | - | 1.20 |  |
| A1 | 0.05 | - | 0.15 |  |
| A2 | 0.95 | 1.00 | 1.05 |  |
| D | 15.75 | 16.00 | 16.25 |  |
| D1 | 13.90 | 14.00 | 14.10 | Note 2 |
| E | 15.75 | 16.00 | 16.25 |  |
| E1 | 13.90 | 14.00 | 14.10 | Note 2 |
| B | 0.17 | - | 0.27 |  |
| C | 0.09 | - | 0.20 |  |
| L | 0.45 | - | 0.75 |  |
| e | 0.50 TYP |  |  |  |

2010-10-20

| н七ค ( $\begin{aligned} & \text { Package Drawing Contact: } \\ & \text { packagedrawings@atmel.com }\end{aligned}$ | TITLE <br> 100A, 100-lead, $14 \times 14$ mm Body Size, 1.0 mm Body Thickness, 0.5 mm Lead Pitch, Thin Profile Plastic Quad Flat Package (TQFP) | DRAWING NO. 100A | REV. <br> D |
| :---: | :---: | :---: | :---: |




COMMON DIMENSIONS
(Unit of measure $=\mathrm{mm}$ )

| SYMBOL | MIN | NOM | MAX | NOTE |
| :---: | :---: | :---: | :---: | :---: |
| A | - | - | 1.20 |  |
| A1 | 0.05 | - | 0.15 |  |
| A2 | 0.95 | 1.00 | 1.05 |  |
| D | 15.75 | 16.00 | 16.25 |  |
| D1 | 13.90 | 14.00 | 14.10 | Note 2 |
| E | 15.75 | 16.00 | 16.25 |  |
| E1 | 13.90 | 14.00 | 14.10 | Note 2 |
| B | 0.30 | - | 0.45 |  |
| C | 0.09 | - | 0.20 |  |
| L | 0.45 | - | 0.75 |  |
| e | 0.80 TYP |  |  |  |

2010-10-20
Notes:
1.This package conforms to JEDEC reference MS-026, Variation AEB.
2. Dimensions D1 and E1 do not include mold protrusion. Allowable protrusion is 0.25 mm per side. Dimensions D1 and E1 are maximum plastic body size dimensions including mold mismatch.
3. Lead coplanarity is 0.10 mm maximum.

64A, 64-lead, $14 \times 14 \mathrm{~mm}$ Body Size, 1.0mm Body Thickness, 0.8 mm Lead Pitch, Thin Profile Plastic Quad Flat Package (TQFP)

### 36.4 64M2



## 37. Errata

### 37.1 ATmega640 rev. B

- Inaccurate ADC conversion in differential mode with $200 \times$ gain
- High current consumption in sleep mode

1. Inaccurate ADC conversion in differential mode with $200 \times$ gain

With AVCC $<3.6 \mathrm{~V}$, random conversions will be inaccurate. Typical absolute accuracy may reach 64 LSB.

## Problem Fix/Workaround

None.
2. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.

## Problem Fix/Workaround

Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.2 ATmega640 rev. A

- Inaccurate ADC conversion in differential mode with 200× gain
- High current consumption in sleep mode

1. Inaccurate ADC conversion in differential mode with $200 \times$ gain

With AVCC $<3.6 \mathrm{~V}$, random conversions will be inaccurate. Typical absolute accuracy may reach 64 LSB.

## Problem Fix/Workaround

None.
2. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.

Problem Fix/Workaround
Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.3 ATmega1280 rev. B

- High current consumption in sleep mode

1. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.

## Problem Fix/Workaround

Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.4 ATmega1280 rev. A

- Inaccurate ADC conversion in differential mode with $200 \times$ gain
- High current consumption in sleep mode

1. Inaccurate ADC conversion in differential mode with $200 \times$ gain

With AVCC $<3.6 \mathrm{~V}$, random conversions will be inaccurate. Typical absolute accuracy may reach 64 LSB.

## Problem Fix/Workaround

None.
2. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.

Problem Fix/Workaround
Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.5 ATmega1281 rev. B

- High current consumption in sleep mode

1. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.
Problem Fix/Workaround
Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.6 ATmega1281 rev. A <br> - Inaccurate ADC conversion in differential mode with $200 \times$ gain <br> - High current consumption in sleep mode

1. Inaccurate ADC conversion in differential mode with $200 \times$ gain

With AVCC $<3.6 \mathrm{~V}$, random conversions will be inaccurate. Typical absolute accuracy may reach 64 LSB.
Problem Fix/Workaround
None.
2. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.

Problem Fix/Workaround
Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.7 ATmega2560 rev. F <br> - ADC differential input amplification by 46 dB (200x) not functional

1. ADC differential input amplification by 46 dB (200x) not functional Problem Fix/Workaround None.

### 37.8 ATmega2560 rev. E

No known errata.

### 37.9 ATmega2560 rev. D

Not sampled.

### 37.10 ATmega2560 rev. C

- High current consumption in sleep mode

1. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.

## Problem Fix/Workaround

Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.11 ATmega2560 rev. B

Not sampled.

### 37.12 ATmega2560 rev. A

- Non-Read-While-Write area of flash not functional
- Part does not work under 2.4 volts
- Incorrect ADC reading in differential mode
- Internal ADC reference has too low value
- IN/OUT instructions may be executed twice when Stack is in external RAM
- EEPROM read from application code does not work in Lock Bit Mode 3

1. Non-Read-While-Write area of flash not functional

The Non-Read-While-Write area of the flash is not working as expected. The problem is related to the speed of the part when reading the flash of this area.

## Problem Fix/Workaround

- Only use the first 248 K of the flash.
- If boot functionality is needed, run the code in the Non-Read-While-Write area at maximum 1/4th of the maximum frequency of the device at any given voltage. This is done by writing the CLKPR register before entering the boot section of the code.

2. Part does not work under 2.4 volts

The part does not execute code correctly below 2.4 volts.

## Problem Fix/Workaround

Do not use the part at voltages below 2.4 volts.
3. Incorrect ADC reading in differential mode

The ADC has high noise in differential mode. It can give up to 7 LSB error.
Problem Fix/Workaround
Use only the 7 MSB of the result when using the ADC in differential mode.
4. Internal ADC reference has too low value

The internal ADC reference has a value lower than specified.

## Problem Fix/Workaround

- Use AVCC or external reference.
- The actual value of the reference can be measured by applying a known voltage to the ADC when using the internal reference. The result when doing later conversions can then be calibrated.

5. IN/OUT instructions may be executed twice when Stack is in external RAM

If either an IN or an OUT instruction is executed directly before an interrupt occurs and the stack pointer is located in external ram, the instruction will be executed twice. In some cases this will cause a problem, for example:

- If reading SREG it will appear that the I-flag is cleared.
- If writing to the PIN registers, the port will toggle twice.
- If reading registers with interrupt flags, the flags will appear to be cleared.


## Problem Fix/Workaround

There are two application workarounds, where selecting one of them, will be omitting the issue:

- Replace IN and OUT with LD/LDS/LDD and ST/STS/STD instructions.
- Use internal RAM for stack pointer.

6. EEPROM read from application code does not work in Lock Bit Mode 3

When the Memory Lock Bits LB2 and LB1 are programmed to mode 3, EEPROM read does not work from the application code.

## Problem Fix/Workaround

Do not set Lock Bit Protection Mode 3 when the application code needs to read from EEPROM.

### 37.13 ATmega2561 rev. F

- ADC differential input amplification by 46dB (200x) not functional

1. ADC differential input amplification by 46 dB (200x) not functional Problem Fix/Workaround None.

### 37.14 ATmega2561 rev. E

No known errata.

### 37.15 ATmega2561 rev. D

Not sampled.

### 37.16 ATmega2561 rev. C

- High current consumption in sleep mode.

1. High current consumption in sleep mode

If a pending interrupt cannot wake the part up from the selected sleep mode, the current consumption will increase during sleep when executing the SLEEP instruction directly after a SEI instruction.

## Problem Fix/Workaround

Before entering sleep, interrupts not used to wake the part from the sleep mode should be disabled.

### 37.17 ATmega2561 rev. B

Not sampled.

### 37.18 ATmega2561 rev. A

- Non-Read-While-Write area of flash not functional
- Part does not work under 2.4 Volts
- Incorrect ADC reading in differential mode
- Internal ADC reference has too low value
- IN/OUT instructions may be executed twice when Stack is in external RAM
- EEPROM read from application code does not work in Lock Bit Mode 3

1. Non-Read-While-Write area of flash not functional

The Non-Read-While-Write area of the flash is not working as expected. The problem is related to the speed of the part when reading the flash of this area.

## Problem Fix/Workaround

- Only use the first 248 K of the flash.
- If boot functionality is needed, run the code in the Non-Read-While-Write area at maximum 1/4th of the maximum frequency of the device at any given voltage. This is done by writing the CLKPR register before entering the boot section of the code.

2. Part does not work under 2.4 volts

The part does not execute code correctly below 2.4 volts.

## Problem Fix/Workaround

Do not use the part at voltages below 2.4 volts.
3. Incorrect ADC reading in differential mode

The ADC has high noise in differential mode. It can give up to 7 LSB error.
Problem Fix/Workaround
Use only the 7 MSB of the result when using the ADC in differential mode.
4. Internal ADC reference has too low value

The internal ADC reference has a value lower than specified.

## Problem Fix/Workaround

- Use AVCC or external reference.
- The actual value of the reference can be measured by applying a known voltage to the ADC when using the internal reference. The result when doing later conversions can then be calibrated.

5. IN/OUT instructions may be executed twice when Stack is in external RAM

If either an IN or an OUT instruction is executed directly before an interrupt occurs and the stack pointer is located in external ram, the instruction will be executed twice. In some cases this will cause a problem, for example:

- If reading SREG it will appear that the I-flag is cleared.
- If writing to the PIN registers, the port will toggle twice.
- If reading registers with interrupt flags, the flags will appear to be cleared.


## Problem Fix/Workaround

There are two application workarounds, where selecting one of them, will be omitting the issue:

- Replace IN and OUT with LD/LDS/LDD and ST/STS/STD instructions.
- Use internal RAM for stack pointer.

6. EEPROM read from application code does not work in Lock Bit Mode 3

When the Memory Lock Bits LB2 and LB1 are programmed to mode 3, EEPROM read does not work from the application code.

## Problem Fix/Workaround

Do not set Lock Bit Protection Mode 3 when the application code needs to read from EEPROM.

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## 38. Datasheet Revision History

Note that the referring page numbers in this section are referring to this document. The referring revisions in this section are referring to the document revision.

### 38.1 Rev. 2549Q-02/2014

1. Updated the "Reset Sources" on page 57. Brown-out Reset: The MCU is reset when the supply voltage AVcc is below the Brown-out Reset threshold (VBOT) and the Brown-out Detector is enabled.
2. Updated the Figure 12-1 on page 58. Power-on reset is now connected to AVcc and not to Vcc.
3. Updated the content in "Brown-out Detection" on page 59. Replaced Vcc by AVcc throughout the section.
4. Updated the Figure 12-5 on page 60. Replaced Vcc by AVcc.
5. Updated "External Interrupts" on page 109. Removed the text "Note that recognition of falling or rising edge.....".
6. Updated the description of "PCMSK1 - Pin Change Mask Register 1" on page 113. The description mentions "PCIE1 bit in EIMSK". This has been changed to "PCIE1 bit in PCICR".
7. Updated "Ordering Information" in "ATmega2561" on page 411.
8. Removed Errata "Inaccurate ADC conversion in differential mode with $200 \times$ gain" from "ATmega1280 rev. B" on page 416 and from "ATmega1281 rev. B" on page 417
9. Updated "Errata" in "ATmega2560 rev. F" on page 417 and in "ATmega2561 rev. F" on page 419.
10. Updated the datasheet with new Atmel brand (new logo and addresses).

### 38.2 Rev. 2549P-10/2012

1. Replaced drawing of " 64 M 2 " on page 415.
2. Former page 439 has been deleted as the content of this page did not belong there (same page as the last page).
3. Some small correction made in the setup.

### 38.3 Rev. 25490-05/2012

1. The datasheet changed status from Preliminary to Complete. Removed "Preliminary" from the front page.
2. Replaced Figure $10-3$ on page 44 by a new one.
3. Updated the last page to include the new address for Atmel Japan site.

### 38.4 Rev. 2549N-05/2011

1. Added Atmel QTouch Library Support and QTouch Sensing Capablity Features.
2. Updated Cross-reference in "Bit 5, 2:0-WDP3:0: Watchdog Timer Prescaler 3, 2, 1 and 0" on page 65.
3. Updated Assembly codes in section "USART Initialization" on page 205.
4. Added "Standard Power-On Reset" on page 360.
5. Added "Enhanced Power-On Reset" on page 361.
6. Updated Figure 32-13 on page 381
7. Updated "Ordering Information" on page 407 to include Tape \& Reel devices.

### 38.5 Rev. 2549M-09/2010

1. Updated typos in Figure 26-9 on page 276 and in Figure $26-10$ on page 277.
2. Note is added below Table 1-1 on page 3.
3. The values for "typical characteristics" in Table 31-9 on page 365 and Table 31-10 on page 366, has been rounded.
4. Units for $t_{\text {RST }}$ and $t_{\text {BOD }}$ in Table 31-3 on page 360 have been changed from " $n s$ " to " $\mu \mathrm{s}$ ".
5. The figure text for Table 31-2 on page 359 has been changed.
6. Text in first column in Table 30-3 on page 326 has been changed from "Fuse Low Byte" to "Extended Fuse Byte".
7. The text in "Power Reduction Register" on page 52 has been changed.
8. The value of the inductor in Figure 26-9 on page 276 and Figure $26-10$ on page 277 has been changed to $10 \mu \mathrm{H}$.
9. "Port A" has been changed into "Port K" in the first paragraph of "Features" on page 268.
10. Minimum wait delay for $\mathrm{t}_{\text {WD_EEPROM }}$ in Table $30-16$ on page 340 has been changed from 9.0 ms to 3.6 ms
11. Dimension A 3 is added in " 64 M 2 " on page 415.
12. Several cross-references are corrected.
13. "COM0A1:0" on page 127 is corrected to "COMOB1:0".
14. Corrected some Figure and Table numbering.
15. Updated Section 10.6 "Low Frequency Crystal Oscillator" on page 43.

### 38.6 Rev. 2549L-08/07

1. Updated note in Table 10-11 on page 45.
2. Updated Table 10-3 on page 42, Table 10-5 on page 43, Table 10-9 on page 45.
3. Updated typos in "DC Characteristics" on page 355
4. Updated "Clock Characteristics" on page 359
5. Updated "External Clock Drive" on page 359.
6. Added "System and Reset Characteristics" on page 360.
7. Updated "SPI Timing Characteristics" on page 363.
8. Updated "ADC Characteristics - Preliminary Data" on page 365.
9. Updated ordering code in "ATmega640" on page 407.

### 38.7 Rev. 2549K-01/07

1. Updated Table 1-1 on page 3.
2. Updated "Pin Descriptions" on page 7.
3. Updated "Stack Pointer" on page 15.
4. Updated "Bit 1 - EEPE: EEPROM Programming Enable" on page 35.
5. Updated Assembly code example in "Thus, when the BOD is not enabled, after setting the ACBG bit or enabling the ADC, the user must always allow the reference to start up before the output from the Analog Comparator or ADC is used. To reduce power consumption in Power-down mode, the user can avoid the three conditions above to ensure that the reference is turned off before entering Power-down mode." on page 60.
6: Updated "EIMSK - External Interrupt Mask Register" on page 111.
6. Updated Bit description in "PCIFR - Pin Change Interrupt Flag Register" on page 113.
7. Updated code example in "USART Initialization" on page 205.
8. Updated Figure 26-8 on page 276.
9. Updated "DC Characteristics" on page 355.

### 38.8 Rev. 2549J-09/06

1. Updated "" on page 46.
2. Updated code example in "Moving Interrupts Between Application and Boot Section" on page 107.
3. Updated "Timer/Counter Prescaler" on page 180.
4. Updated "Device Identification Register" on page 296.
5. Updated "Signature Bytes" on page 328.
6. Updated "Instruction Set Summary" on page 404.

### 38.9 Rev. 2549I-07/06

1. Added "Data Retention" on page 10.
2. Updated Table 16-3 on page 126, Table 16-6 on page 127, Table 16-8 on page 128, Table 17-2 on page 145 , Table 17-4 on page 155, Table 17-5 on page 155, Table $20-3$ on page 182, Table 20-6 on page 183 and Table 20-8 on page 184.
3. Updated "Fast PWM Mode" on page 146.

### 38.10 Rev. 2549H-06/06

1. Updated "" on page 46.
2. Updated "OSCCAL - Oscillator Calibration Register" on page 48.
3. Added Table 31-1 on page 359.

### 38.11 Rev. 2549G-06/06

1. Updated "Features" on page 1.
2. Added Figure 1-2 on page 3, Table 1-1 on page 3.
3. Updated "" on page 46.
4. Updated "Power Management and Sleep Modes" on page 50.
5. Updated note for Table 12-1 on page 65.
6. Updated Figure 26-9 on page 276 and Figure 26 -10 on page 277.
7. Updated "Setting the Boot Loader Lock Bits by SPM" on page 316.
8. Updated "Ordering Information" on page 407.
9. Added Package information " 100 C 1 " on page 413.
10. Updated "Errata" on page 416.

### 38.12 Rev. 2549F-04/06

1. Updated Figure 9-3 on page 29, Figure 9-4 on page 30 and Figure 9-5 on page 30.
2. Updated Table 20-2 on page 182 and Table 20-3 on page 182.
3. Updated Features in "ADC - Analog to Digital Converter" on page 268.
4. Updated "Fuse Bits" on page 326.

### 38.13 Rev. 2549E-04/06

1. Updated "Features" on page 1.
2. Updated Table 12-1 on page 62.
3. Updated note for Table 12-1 on page 62.
4. Updated "Bit 6 - ACBG: Analog Comparator Bandgap Select" on page 266.
5. Updated "Prescaling and Conversion Timing" on page 271.
6. Updated "Maximum speed vs. VCC" on page 373.
7. Updated "Ordering Information" on page 407.

### 38.14 Rev. 2549D-12/05

1. Advanced Information Status changed to Preliminary.
2. Changed number of I/O Ports from 51 to 54 .
3. Updatet typos in "TCCROA - Timer/Counter Control Register A" on page 126.
4. Updated Features in "ADC - Analog to Digital Converter" on page 268.
5. Updated Operation in"ADC - Analog to Digital Converter" on page 268
6. Updated Stabilizing Time in "Changing Channel or Reference Selection" on page 274.
7. Updated Figure 26-1 on page 269, Figure 26-9 on page 276, Figure 26-10 on page 277.
8. Updated Text in "ADCSRB - ADC Control and Status Register B" on page 282.
9. Updated Note for Table 4 on page 42, Table 13-15 on page 82, Table 26-3 on page 281 and Table 26-6 on page 287.
10. Updated Table 31-9 on page 365 and Table 31-10 on page 366.
11. Updated "Filling the Temporary Buffer (Page Loading)" on page 315.
12. Updated "Typical Characteristics" on page 373.
13. Updated "Packaging Information" on page 412.
14. Updated "Errata" on page 416.

### 38.15 Rev. 2549C-09/05

1. Updated Speed Grade in section "Features" on page 1.
2. Added "Resources" on page 10.
3. Updated "SPI - Serial Peripheral Interface" on page 190. In Slave mode, low and high period SPI clock must be larger than 2 CPU cycles.
4. Updated "Bit Rate Generator Unit" on page 242.
5. Updated "Maximum speed vs. VCC" on page 373.
6. Updated "Ordering Information" on page 407.
7. Updated "Packaging Information" on page 412. Package 64M1 replaced by 64M2.
8. Updated "Errata" on page 416.

### 38.16 Rev. 2549B-05/05

1. JTAG ID/Signature for ATmega640 updated: $0 \times 9608$.
2. Updated Table 13-7 on page 78.
3. Updated "Serial Programming Instruction set" on page 340.
4. Updated "Errata" on page 416.
38.17 Rev. 2549A-03/05
5. Initial version.

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## 5A 35DKHz 25V Buck DC to DC Converter

## Features

■ Wide 3.6 V to 25 V Input Voltage Range

- Output Adjustable from 0.8 V to 23 V
- Maximum Duty Cycle $100 \%$
- Minimum Drop Out 0.6V
- Fixed 350 KHz Switching Frequency
- 5A Constant Output Current Capability
- Internal Optimize Power MOSFET
- High efficiency
- Excellent line and load regulation
- TTL shutdown capability
- EN pin with hysteresis function
- Built in thermal shutdown function
- Built in current limit function
- Built in output short protection function
- Available in TO-263 package


## Applications

- LCD Monitor and LCD TV
- Digital Photo Frame
- Set-up Box
- ADSL Modem
- Telecom / Networking Equipment


## General Description

The GS2678 is a 350 KHz fixed frequency PWM buck (step-down) DC/DC converter, capable of driving a 5 A load with high efficiency, low ripple and excellent line and load regulation. Requiring a minimum number of external components, the regulator is simple to use and include internal frequency compensation and a fixed-frequency oscillator.

The PWM control circuit is able to adjust the duty ratio linearly from 0 to $100 \%$. An enable function, an over current protection function is built inside. When short protection function happens, the operation frequency will be reduced from 350 KHz to 80 KHz . An internal compensation block is built in to minimize external component count.

## TO263-5L

Figure1. Package Type of GS2678

## 5A 350KHz 25V Buck DC to DC Converter

## Pin Configurations



Figure2. Pin Configuration of GS2678 (Top View)
Table 1 Pin Description

| Pin Number | Pin Name | Description |
| :---: | :---: | :--- |
| 1 | GND | Ground Pin. Care must be taken in layout. This pin should be <br> placed outside of the Schottky Diode to output capacitor <br> ground path to prevent switching current spikes from inducing <br> voltage noise into GS2678. |
| 2 | FB | Feedback Pin (FB). Through an external resistor divider <br> network, FB senses the output voltage and regulates it. The <br> feedback threshold voltage is 0.8V. |
| 3 | SW | Power Switch Output Pin (SW). SW is the switch node that <br> supplies power to the output. |
| 4 | EN | Enable Pin. Drive EN pin high to turn on the device, drive it <br> low to turn it off. Floating is default high. |
| 5 | VIN | Supply Voltage Input Pin. GS2678 operates from a 3.6V to <br> 25V DC voltage. Bypass Vin to GND with a suitably large <br> capacitor to eliminate noise on the input. |

## 5A 350KHz 25V Buck DC to DC Converter

## Function Block



Figure3. Function Block Diagram of GS2678

## Typical Application Circuit



Figure4. GS2678 Typical Application Circuit 12V-5V/5A

## 5A 350KHz 25V Buck DC to DC Converter

## Ordering Information

| Package | Temperature <br> Range | Part Number | Marking ID | Packing Type |
| :---: | :---: | :--- | :--- | :---: |
|  |  | Lead Free | Lead Free |  |
|  |  | GS2678E1 | GS2678E1 | Tube |
|  |  | GS2678TRE1 | GS2678E1 | Tape \& Reel |

GSM Pb-free products, as designated with "E1" suffix in the par number, are RoHS compliant.

## Absolute Maximum Ratings (Note1)

| Parameter | Symbol | Value | Unit |
| :--- | :---: | :---: | :---: |
| Input Voltage | $\mathrm{Vin}^{2}$ | -0.3 to 30 | V |
| Feedback Pin Voltage | $\mathrm{V}_{\mathrm{FB}}$ | -0.3 to Vin | V |
| EN Pin Voltage | $\mathrm{V}_{\mathrm{EN}}$ | -0.3 to Vin | V |
| Output Switch Pin Voltage | $\mathrm{V}_{\text {Output }}$ | -0.3 to Vin | V |
| Power Dissipation | $\mathrm{P}_{\mathrm{D}}$ | Internally limited | mW |
| Thermal Resistance (TO263) <br> (Junction to Ambient, No Heatsink, Free Air) | $\mathrm{R}_{\mathrm{JA}}$ | 50 | ${ }^{\circ} \mathrm{C} / \mathrm{W}$ |
| Operating Junction Temperature | $\mathrm{T}_{\mathrm{J}}$ | -40 to 125 | ${ }^{\circ} \mathrm{C}$ |
| Storage Temperature | $\mathrm{T}_{\text {STG }}$ | -65 to 150 | ${ }^{\circ} \mathrm{C}$ |
| Lead Temperature (Soldering, 10 sec) | $\mathrm{T}_{\text {LEAD }}$ | 260 | ${ }^{\circ} \mathrm{C}$ |
| ESD (HBM) |  | 2000 | V |

Note1: Stresses greater than those listed under Maximum Ratings may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operation is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.

## 5A 350KHz 25V Buck DC to DC Converter

## GS2EElectrical Characteristics

$\mathrm{T}_{\mathrm{a}}=25^{\circ} \mathrm{C}$;unless otherwise specified.

| Symbol | Parameter | Test Condition | Min. | Typ. | Max. | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| System parameters test circuit figure4 |  |  |  |  |  |  |
| VFB | Feedback <br> Voltage | $\begin{gathered} \text { Vin }=5 \mathrm{~V} \text { to } 25 \mathrm{~V}, \text { Vout }=5 \mathrm{~V} \\ \text { Iload }=0.5 \mathrm{~A} \text { to } 5 \mathrm{~A} \end{gathered}$ | 0.776 | 0.8 | 0.824 | V |
| Efficiency | 1 | $\begin{gathered} \text { Vin }=12 \mathrm{~V}, \text { Vout }=5 \mathrm{~V} \\ \text { Iout }=5 \mathrm{~A} \end{gathered}$ | - | 90 | - | \% |

## Electrical Characteristics (DC Parameters)

Vin $=12 \mathrm{~V}, \mathrm{GND}=0 \mathrm{~V}$, Vin \& GND parallel connect a $220 \mathrm{uf} / 50 \mathrm{~V}$ capacitor; Iout $=500 \mathrm{~mA}, \mathrm{~T}_{\mathrm{a}}=$ $25^{\circ} \mathrm{C}$; the others floating unless otherwise specified.

| Parameters | Symbol | Test Condition | Min. | Typ. | Max. | Unit |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: |
| Input operation voltage | Vin |  | 3.6 |  | 25 | V |
| Shutdown Supply Current | $\mathrm{I}_{\mathrm{STBY}}$ | $\mathrm{V}_{\mathrm{EN}}=0 \mathrm{~V}$ |  | 60 | 200 | uA |
| Quiescent Supply Current | $\mathrm{I}_{\mathrm{q}}$ | $\mathrm{V}_{\mathrm{EN}}=2 \mathrm{~V}$, <br> $\mathrm{V}_{\mathrm{FB}}=\mathrm{Vin}$ |  | 3 | 5 | mA |
| Oscillator Frequency | Fosc |  | 298 | 350 | 402 | Khz |
| Switch Current Limit | $\mathrm{I}_{\mathrm{L}}$ | $\mathrm{V}_{\mathrm{FB}}=0$ |  | 8 |  | A |
| EN Pin Threshold | $\mathrm{V}_{\mathrm{EN}}$ | High (Regulator ON) <br> Low (Regulator OFF) |  | 1.4 <br> 0.8 |  | V |
|  | $\mathrm{I}_{\mathrm{H}}$ | $\mathrm{V}_{\mathrm{EN}}=2 \mathrm{~V}(\mathrm{ON})$ |  | 1 | 15 | uA |
|  | $\mathrm{I}_{\mathrm{L}}$ | $\mathrm{V}_{\mathrm{EN}}=0 \mathrm{~V}(\mathrm{OFF})$ |  | 1 | 15 | uA |
| Max. Duty Cycle | $\mathrm{D}_{\mathrm{MAX}}$ | $\mathrm{V}_{\mathrm{FB}}=0 \mathrm{~V}$ |  | 100 |  | $\%$ |

## 5A 350KHz 25V Buck DC to DC Converter

## Test Circuit and Layout guidelines



Figure5. Standard Test Circuits and Layout Guides

Select R1 to be approximately 2 K , use a $1 \%$ resistor for best stability.
C 1 and CFF are optional; in order to increase stability and reduce the input power line noise, CIN and C1 must be placed near to PIN1 and PIN3;
For output voltages greater than approximately 10 V , an additional capacitor CFF is required. The compensation capacitor is typically between 100 pf and 33 nf , and is wired in parallel with the output voltage setting resistor, R2. It provides additional stability for high output voltage, low input-output voltages, and/or very low ESR output capacitors, such as solid tantalum capacitors. CFF $=1 /\left(31^{*} 1000 *\right.$ R2 ); This capacitor type can be ceramic, plastic, silver mica, etc. (Because of the unstable characteristics of ceramic capacitors made with Z5U material, they are not recommended.)

## 5A 350KHz 25V Buck DC to DC Converter

Schottky Diode Selection Table

| Current | Surface <br> Mount | Through Hole | VR (The same as system maximum input voltage) |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | 20V | 30V | 40V | 50V | 60 V |
| 1A |  | $\checkmark$ | 1N5817 | 1N5818 | 1N5819 |  |  |
|  |  |  |  |  |  |  |  |
| 3A |  | $\checkmark$ | 1N5820 | 1N5821 | 1N5822 |  |  |
|  |  | $\checkmark$ | MBR320 | MBR330 | MBR340 | MBR350 | MBR360 |
|  | $\checkmark$ |  | SK32 | SK33 | SK34 | SK35 | SK36 |
|  | $\checkmark$ |  |  | 30WQ03 | 30WQ04 | 30WQ05 |  |
|  |  | $\checkmark$ |  | 31DQ03 | 31DQ04 | 31DQ05 |  |
|  |  | $\checkmark$ | SR302 | SR303 | SR304 | SR305 | SR306 |
| - |  |  |  |  |  |  |  |
| 5A |  | $\checkmark$ | 1N5823 | 1N5824 | 1N5825 |  |  |
|  |  | $\checkmark$ | SR502 | SR503 | SR504 | SR505 | SR506 |
|  |  | $\checkmark$ | SB520 | SB530 | SB540 | SB550 | SB560 |
|  | $\checkmark$ |  |  | 50WQ03 | 50WQ04 | 50WQ05 |  |

## Typical System Application for 24V ~ 12V/3A Version



Figure6. GS2678 System Parameters Test Circuit ( $24 \mathrm{~V} \sim 12 \mathrm{~V} / 3 \mathrm{~A}$ )

## Typical System Application for 24V ~ 5V/5A



Figure7. GS2678 System Parameters Test Circuit (24V ~5V/5A)

5A 35DKHz 25V Buck DC to DC Converter

## Package Information

(2) TO263-5L


| Symbol | Dimensions In Millimeters |  | Dimensions In Inches |  |
| :---: | :---: | :---: | :---: | :---: |
|  | Min. | Max. | Min. | Max. |
| A | 4.06 | 4.83 | 0.160 | 0.190 |
| B | 0.76 | 1.02 | 0.030 | 0.040 |
| C | 0.36 | 0.64 | 0.014 | 0.025 |
| C2 | 1.14 | 1.40 | 0.045 | 0.055 |
| D | 8.64 | 9.65 | 0.340 | 0.380 |
| E | 9.78 | 10.54 | 0.385 | 0.415 |
| e | 1.57 | 1.85 | 0.062 | 0.073 |
| F | 6.60 | 7.11 | 0.260 | 0.280 |
| L | 15.11 | 15.37 | 0.595 | 0.605 |
| L2 | - | 1.40 | - | 0.055 |

## COM-08653



DISCLAIMER:
This document is the result of an obscene amount of work and research because I really wanted to make a nice wood case / bezel using Ponoko for a project I'm working on. Literally everything in this document might be wrong and if you waste money or injure somebody based on this document, it's your own look out. If you can and want to correct any of the information here, contact andrew@amorrow.com.

## Mechanical



## Electrical

PushButtons: Columns + Rows
This is the same information in several different ways.
I've gotten requests to visualize it differently, so why fight it?



## Arduino information

Thanks to Jim Winburn for additional Arduino related info, and for changing the groupings on the pin layouts. MUCH more clear now!

Keypad Pin 3 Arduino digital 2
Keypad Pin 1 Arduino digital 3
Keypad Pin 5 Arduino digital 4
Keypad Pin 2 Arduino digital 5
Keypad Pin 7 Arduino digital 6
Keypad Pin 6 Arduino digital 7
Keypad Pin 4 Arduino digital 8
Notice the rowPins and colPins below in the sample code example:
// example from Arduino playground: http://www.arduino.cc/playground/Code/Keypad

```
#include <Keypad.h>
const byte ROWS = 4; //four rows
const byte cOLS = 3; //three columns
char keys[ROWS][COLS] = {
    {'1','2','3'},
    {'4','5','6'},
    {'7','8','9'},
    {'*','0','#'}
};
byte rowPins[ROWS] = {5, 6, 7, 8}; //connect to the row pinouts of the keypad
byte colPins[COLS] = {2, 3, 4}; //connect to the column pinouts of the keypad
```


## Sources

http://www.sparkfun.com/products/8653
http://www.tenrod.com.au/pdf/accord/keypads.pdf
http://www.accordia.com.tw/html/tel/index.htm (AK-304-BBW)
http://www.accordia.com.tw/html/tel/imagepages/imagepg11.htm
http://www.idpcorp.net/product_expertise/silicone_keypads/stock_keypads/technical_information.html

5-5-2011<br>- changed part \# to Sparkfun COM-08653<br>- added link to tutorial

6-27-2011

- changed grouping on column/row pins for clarity - Thanks Jim
- Added source for Arduino with new grouping
- Added new visualization for pin mapping
- Split document into multiple pages


## 16 x 2 Character LCD

## FEATURES

- $5 \times 8$ dots with cursor
- Built-in controller (KS 0066 or Equivalent)
- +5V power supply (Also available for +3 V )
- 1/16 duty cycle
- B/L to be driven by pin 1, pin 2 or pin 15, pin 16 or A.K (LED)
- N.V. optional for + 3V power supply

MECHANICAL DATA

| ITEM | STANDARD VALUE | UNIT |
| :--- | :---: | :---: |
| Module Dimension | $80.0 \times 36.0$ | mm |
| Viewing Area | $66.0 \times 16.0$ | mm |
| Dot Size | $0.56 \times 0.66$ | mm |
| Character Size | $2.96 \times 5.56$ | mm |


| ABSOLUTE MAXIMUM RATING |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: |
| ITEM | SYMBOL | STANDARD VALUE |  |  |  |
|  |  |  |  |  |  |
|  |  | MIN. | TYP. | MAX. |  |
| Power Supply | VDD-VSS | -0.3 | - | 7.0 | V |
| Input Voltage | VI | -0.3 | - | VDD |  |

NOTE: VSS = 0 Volt, VDD $=5.0$ Volt

## ELECTRICAL SPECIFICATIONS

| ITEM | SYMBOL | CONDITION |  | STANDARD VALUE |  |  | UNIT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | MIN. | TYP. | MAX. |  |
| Input Voltage | VDD | $V D D=+5 \mathrm{~V}$ |  | 4.7 | 5.0 | 5.3 | V |
|  |  | $\mathrm{VDD}=+3 \mathrm{~V}$ |  | 2.7 | 3.0 | 5.3 | V |
| Supply Current | IDD | $\mathrm{VDD}=5 \mathrm{~V}$ |  | - | 1.2 | 3.0 | mA |
| Recommended LC Driving Voltage for Normal Temp. Version Module | VDD - Vo | $-20^{\circ} \mathrm{C}$ |  | - | - | - | V |
|  |  | $0^{\circ} \mathrm{C}$ |  | 4.2 | 4.8 | 5.1 |  |
|  |  | $25^{\circ} \mathrm{C}$ |  | 3.8 | 4.2 | 4.6 |  |
|  |  | $50^{\circ} \mathrm{C}$ |  | 3.6 | 4.0 | 4.4 |  |
|  |  | $70^{\circ} \mathrm{C}$ |  | - | - | - |  |
| LED Forward Voltage | VF | $25^{\circ} \mathrm{C}$ |  | - | 4.2 | 4.6 | V |
| LED Forward Current | IF | $25^{\circ} \mathrm{C}$ | Array | - | 130 | 260 | mA |
|  |  |  | Edge | - | 20 | 40 |  |
| EL Power Supply Current | IEL | $\mathrm{Vel}=110 \mathrm{VAC}: 400 \mathrm{~Hz}$ |  | - | - | 5.0 | mA |

## DISPLAY CHARACTER ADDRESS CODE:

Display Position
DD RAM Address
DD RAM Address

| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | 01 |  |  |  |  |  |  |  |  |  |  |  |  |  | OF |
| 40 | 41 |  |  |  |  |  |  |  |  |  |  |  |  |  | 4F |

Vishay

| PIN NUMBER | SYMBOL | FUNCTION |
| :--- | :---: | :---: |
| 1 | Vss | GND |
| 2 | Vdd | + 3V or +5 V |
| 3 | Vo | Contrast Adjustment |
| 4 | RS | H/L Register Select Signal |
| 5 | R/W | H/L Read/Write Signal |
| 6 | E | H $\rightarrow$ L Enable Signal |
| 7 | DB0 | H/L Data Bus Line |
| 8 | DB1 | H/L Data Bus Line |
| 9 | DB2 | H/L Data Bus Line |
| 10 | DB3 | H/L Data Bus Line |
| 11 | DB4 | H/L Data Bus Line |
| 12 | DB5 | H/L Data Bus Line |
| 13 | DB6 | H/L Data Bus Line |
| 14 | DB7 | H/L Data Bus Line |
| 15 | A/Vee | Power Supply for B/L (OV) |
| 16 | K |  |

DIMENSIONS in millimeters


This datasheet has been download from:

www.datasheetcatalog.com

Datasheets for electronics components.

## $\mathscr{P}_{\text {roduct }}$ Feature

The joysticks are used the high-quality carbon film potentiometers and could accurately control the direction and speed.

Feature
*Easyoperation *SmallResel Eror $\int_{\text {\& Lune Life }}$

## Application

*Move the camera and CCTV control system
*Measurement equipment

* Medical imaging system
$\star$ Automatic machine control
$\star$ Video editing and production equipmen
*Industrial and scientific instrument



## How to choice?

STEP 1: Choose the bezel


A

B

C

D

## STEP 2: Choose the resistance

M1: Total resistance value is 5 K , Resistance range is from 1.8 K to 2.5 k , then 2.5 k to 3.2 K .
M2: Total resistance value is 5 K . Resistance range is from 0K to 2.5 K then 2.5 K to 5 K .
M3: Total resistance value is 10 K . Resistance range is from 3.2 K to 5 K , then 5 K to 6.8 K
M4: Total resistance value is 10 K , Resistance range is from 0 K to 5 K , then 5 K to 10 K .
M5: Total resistance value is 100 K . Resistance range is from 36 K to 50 k , then 50 K to 64 K .
M6: Total resistance value is 100 K . Resistance range is from 0 K to 50 k , then from 50 K to 100 K .



## Dimensional $\mathscr{D}_{\text {rawiong }}$




## Actuator LA23

The LA23 actuator is a small and strong push or pull actuator (up to 2500 N ). The LA23 can be used in various applications where size is important.
Some of the benefits the LA23 offers you are:

- Compact design
- High lifting force
- Exchangeable cables

The standard LA23 is available for both the CARELINE ${ }^{\oplus}$, MEDLINE ${ }^{\oplus}$, and DESKLINE ${ }^{\oplus}$ product ranges.

## Features and options:

- Load in push: $2500 \mathrm{~N}, 1800 \mathrm{~N}, 1500 \mathrm{~N}, 1200 \mathrm{~N}$ or 900 N
- Load in pull: 2500N, 1800N, 1500N, 1200N or 900N
- Housing colour: Grey or black
- Protection class: IPX4, IPX6
- Motor: 12 V DC, 24 V DC
- Stroke length: 20-300 mm
- Built-in dimensions: 110-146 mm + stroke length
- Positioning options: Potential free end-stop signals Hall potentiometer or Hall PWM position Single Hall, Dual Hall
- Back fixture material: Plastic or steel
- Nut: Guided
- Safety nut: In push or pull ( 2500 N and 1800 N version only safety nut in push)
- Mechanical spline: Yes
- Built-in electrical end-stop: Yes
- Exchangeable cable: Yes
- Static safety factor: 2.5
- Noise level: Max. $58.5 \mathrm{~dB}(\mathrm{~A})$ (At nominal voltage and with no load, according to EN ISO 3743-1)
- Mechanical end stop: Yes
- Integrated Control Yes


## Usage:

- Duty cycle: $10 \%, 2$ minutes continuous use followed by 18 minutes not in use
- Usage temperature:
$-30^{\circ} \mathrm{C}$ to $+55^{\circ} \mathrm{C}$ (according to ISO 7176-9)
- Storage temperature:
$-45^{\circ} \mathrm{C}$ to $+70^{\circ} \mathrm{C}$ (according to ISO 7176-9)
- Compatibility: CB20, CB16, CB6S, CBJ 1/2, CBJC, CBJH, CBD4, CBD5 \& CBD6*, CA30, CA40, C061
* SLS must be ignored Up + Down in the CBD4, CBD5 or CBD6 when configured for LA23.
* Only the 3,6 \& 12 mm versions can be configured in the CBD4, CBD5 or CBD6.
* Only tested for single use.
- Approvals:

IEC60601-1:2005 $3^{\text {rd }}$ ed., ANSI / AAMI ES60601-1:2005
$3^{\text {rd }}$ edition, LA23 in combination with CBD4, CBD5 \&
CBD6 has no approvals.

- Fire catagory: Enclosure UL94-V0



The built-in dimension depends upon the chosen safety option and stroke length. Please see the table below to decide upon the built-in dimension.

| Safety option | Stroke length | Spindle pitch | Min. built-in Dimensions |
| :---: | :---: | :---: | :---: |
| 0 = No safety option | 20-49 | 6,9 or 12 | 160 |
| 0 = No safety option | 20-49 | 3, 5 | 168 |
| 1 = Safety nut for push | 20-49 | 6,9 or 12 | 160 |
| 1 = Safety nut for push | 20-49 | 3, 5 | 168 |
| 2 = Safety nut for pull | 20-49 | 6,9 or 12 | 172 |
| 3 = Mechanical Spline for push | 20-49 | 6,9 or 12 | 180 |
| 3 = Mechanical Spline for push | 20-49 | 3,5 | 196 |
| 4 = Mechanical Spline \& safety nut for push | 20-49 | 6,9 or 12 | 180 |
| 4 = Mechanical Spline \& safety nut for push | 20-49 | 3, 5 | 196 |
| 0 = No safety option | 50-200 | 6,9 or 12 | 110 + stroke |
| 0 = No safety option | 50-200 | 3, 5 | 118 + stroke |
| 1 = Safety nut for push | 50-200 | 6,9 or 12 | 110 + stroke |
| 1 = Safety nut for push | 50-200 | 3, 5 | 118 + stroke |
| 2 = Safety nut for pull | 50-200 | 6,9 or 12 | $122+$ stroke |
| 3 = Mechanical Spline for push | 50-200 | 6,9 or 12 | 130 + stroke |
| 3 = Mechanical Spline for push | 50-200 | 3,5 | 146 + stroke |
| 4 = Mechanical Spline \& safety nut for push | 50-200 | 6,9 or 12 | 130 + stroke |
| 4 = Mechanical Spline \& safety nut for push | 50-200 | 3, 5 | 146 + stroke |
| $0=$ No safety option | 201-300 | 6,9 or 12 | 130 + stroke |
| 0 = No safety option | 201-300 | 3, 5 | 138 + stroke |
| 1 = Safety nut for push | 201-300 | 6,9 or 12 | 130 + stroke |
| 1 = Safety nut for push | 201-300 | 3,5 | 138 + stroke |
| 2 = Safety nut for pull | 201-300 | 6,9 or 12 | 142 + stroke |
| 3 = Mechanical Spline for push | 201-300 | 6,9 or 12 | $150+$ stroke |
| 3 = Mechanical Spline for push | 201-300 | 3, 5 | 166 + stroke |
| 4 = Mechanical Spline \& safety nut for push | 201-300 | 6,9 or 12 | 150 + stroke |
| 4 = Mechanical Spline \& safety nut for push | 201-300 | 3, 5 | 166 + stroke |

It is possible to order LA23 with extended built-in dimensions if the following requirements are fulfilled

|  | Spindle pitch = 6, 9, 12 | $\begin{gathered} \text { Spindle pitch }= \\ 3,5 \end{gathered}$ | Spindle pitch $=$ 6, 9, 12 | $\begin{gathered} \text { Spindle pitch }= \\ 6,9,12 \end{gathered}$ | $\begin{gathered} \text { Spindle pitch }= \\ 3,5 \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | Safety option 0 : <br> No safety option |  | Safety option 2 : safety nut pull | Safety option 3 : <br> Spline without safety nut |  |
|  | Safety option 1 : safety nut push |  |  | Safety option 4 : <br> Spline + safety nut push |  |
| Max. built-in dimensions | $\leq 730$ - stroke | $\leq 738$ - stroke | $\leq 742$ - stroke | $\leq 750$ - stroke | $\leq 766$ - stroke |

Example:
A) 6 mm pitch no safety option, stroke 200, BID can be max. $(730-200)=530$
B) 3 mm pitch no safety option, stroke 20, BID can be max. (738-20) $=718$

Technical specifications:

| Power supply | Spindle pitch (mm) | Load max. Push or Pull ( N ) | Motor type | *Typical speed at $0 /$ full load ( $\mathrm{mm} / \mathrm{sec}$.) | *Typical current at $0 /$ full load (Amp.) | Inrush current <br> (Amp) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 12VDC | 3 | $2500 / 2500$ | A: 12 V | $3.1 / 2.5$ | $0.8 / 3.6$ | 13.4 |
| CBJ1/2, CBJH and CBJC | 3 | $2500 / 2500$ | B: 24 V | $3.2 / 2.6$ | 0.4 / 1.9 | 8.7 |
| OpenBus ${ }^{\text {TM }}$ | 3 | $2500 / 2500$ | G: 24 V | $3.3 / 2.7$ | $0.3 / 1.4$ | 6.2 |
| 12VDC | 5 | 1800 / 1800 | A: 12 V | 5.4 / 4.2 | $0.8 / 3.9$ | 13.4 |
| CBJ1/2, CBJH and CBJC | 5 | 1800/1800 | B: 24 V | $5.4 / 4.5$ | $0.4 / 1.9$ | 8.7 |
| OpenBus ${ }^{\text {TM }}$ | 5 | 1800/1800 | $\mathrm{G}: 24 \mathrm{~V}$ | $5.6 / 4.6$ | $0.3 / 1.4$ | 6.2 |
| 12VDC | 6 | $1500 / 1500$ | A: 12 V | $6.6 / 5.2$ | $0.8 / 3.6$ | 13.4 |
| CBJ1/2, CBJH and CBJC | 6 | $1500 / 1500$ | B: 24 V | $6.4 / 5.5$ | $0.4 / 1.7$ | 8.7 |
| OpenBus ${ }^{\text {TM }}$ | 6 | $1500 / 1500$ | $\mathrm{G}: 24 \mathrm{~V}$ | $6.7 / 5.5$ | $0.3 / 1.3$ | 6.2 |
| 12VDC | 9 | $1200 / 1200$ | A: 12 V | $9.9 / 7.5$ | $0.9 / 4.0$ | 13.4 |
| CBJ1/2, CBJH and CBJC | 9 | $1200 / 1200$ | B: 24 V | 9.5 / 8.1 | $0.4 / 1.9$ | 8.7 |
| OpenBus ${ }^{\text {TM }}$ | 9 | $1200 / 1200$ | $\mathrm{G}: 24 \mathrm{~V}$ | $9.9 / 8.1$ | $0.3 / 1.3$ | 6.2 |
| 12VDC | 12 | $900 / 900$ | A: 12 V | $13 / 9.6$ | $0.9 / 3.8$ | 13.4 |
| CBJ1/2, CBJH and CBJC | 12 | $900 / 900$ | B: 24 V | 12.6 / 10.4 | 0.4 / 1.9 | 8.7 |
| OpenBus ${ }^{\text {TM }}$ | 12 | $900 / 900$ | $\mathrm{G}: 24 \mathrm{~V}$ | 13.3 / 10.7 | $0.3 / 1.4$ | 6.2 |

Safety nut and steel back fixture overview

| Pitch <br> $(\mathbf{m m})$ | Load <br> (N) | Safety nut | Steel back fixture | Plastic back fixture |
| :---: | :---: | :---: | :---: | :---: |
| 12 | 900 N | Optional in push or pull | Required in pull | Only in push |
| 9 | 1200 N | Optional in push or pull | Required in pull | Only in push |
| 6 | 1500 N | Optional in push or pull | Required in pull | Only in push |
| 5 | 1800 N | Optional in push <br> (Safety nut 2500N not available in pull) | Always required | Not available |
| 3 | 2500 N | Optional in push <br> (Safety nut 2500N not available in pull) | Always required | Not available |

## Self-locking specifications

| Maximum self-lock (N) | Without short circuit | With short circuit |
| :---: | :---: | :---: |
| 12 mm pitch | 750 | 900 |
| 9 mm pitch | 750 | 1200 |
| 6 mm pitch | 1200 | 1500 |
| 5 mm pitch | 1600 | 1800 |
| 3 mm pitch | 2500 | 2500 |

## Terms of use

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## Heavy Duty Rod Actuators

Heavy Duty line was developed with high standards in mind. Our clients demand the best quality and performance for their money, inspiring us to come up with such a rugged design. This versatile unit features an aluminum shaft to shed weight, internal limit switches, an aluminum alloy inner tube, as well as a high performance 12 volt DC motor which is capable of producing between 200 and 1000 lbs . (90-453 kg ) of force.

No matter your application, installation doesn't need to be a hassle. Required components consist of just two brackets, one for either side, with quick release pins for easy fitment and set-up. Controlling the unit is just as easy via a 2-wire system, making this unit a 'plug and play' type accession unique to Firgelli Automations.

All of our units are CE and ROHS certified, and come with a 12 month mechanical warranty to keep your units serving you best. Our company values revolve around providing quick and friendly customer service that will follow your experience with our products and staff.

## Technical specifications

| Model |  | FA-200-L-12-XX |  | FA-400-L-12-XX | FA-1000-L-12-XX |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Dynamic Force |  | 200 lb |  | 400 lb | 1000 lb |
| Static Force |  | 350 lb |  | 600 lb | 1350 lb |
| Speed ("/S) @ Full Load |  | 0.8" |  | 0.4" | 0.2" |
| Duty Cycle |  | 20\% at 100\% load, 50\% at 25\% of load |  |  |  |
| IP Rating |  | 43 |  |  |  |
| Input |  | 12v DC |  |  |  |
| Max Draw |  | 5 A |  |  |  |
| Operational Temperature |  | $-26^{\circ} \mathrm{C} / 65^{\circ} \mathrm{C}\left(-15^{\circ} \mathrm{F} / 150^{\circ} \mathrm{F}\right)$ |  |  |  |
| Limit Switch |  | Built-in (factory preset) |  |  |  |
| Safety Certification |  | CE |  |  |  |
| Bracket(s) |  | MB3 or MB3U |  |  |  |
| Stroke | Retracted Length |  | Extended Length |  | Weight |
| $3^{\prime \prime}$ | 9.88" |  | 12.88" |  | 4.5 lbs |
| $6^{\prime \prime}$ | 12.88" |  | 18.88" |  | 5.1 lbs |
| 8" | 14.88" |  | 22.88" |  | 5.3 lbs |
| 9" | 15.88" |  | 24.88" |  | 5.5 lbs |
| 12" | 18.88" |  | 30.88" |  | 5.9 lbs |
| 15" | 21.88" |  | 36.88" |  | 6.1 lbs |
| 18" | 24.88" |  | 42.88" |  | 6.7 lbs |
| 24" | 30.88" |  | 54.88" |  | 7.4 lbs |
| 30" | $36.88{ }^{\prime \prime}$ |  | 66.88" |  | 8.2 lbs |

## Technical Drawings





## APPROVED DATA

## TO:

Date: Dec. $9^{\text {th }}, 2009$

File No.
Description RELAY

Part No. JZC-11F-05VDC - 1Z

| Approved | Reporter | Checker |
| :---: | :---: | :---: |
| Chen Peng | Huang Xianke | Xing Jianming |



WENGYANG INDUSTRIAL ZONE, YUEQING
325606 ZHEJIANG, CHINA Email: info@cndongya.com
Home page: www.cndongya.com

MINIATURE RELAY
FEATURES

- 5 A switching capability
- SPDT configuration
- Power consumption of 200 mW
- Standard PCB layout
- Application: Power supply, Air-conditioning, Monitor and Security alarm system, etc...
- UL File No. E145623
- TUV File No. R50121570
- CQC File No. CQC09002033531
- Meet the EU RoHS environmental protection request

SPECIFICATIONS

| Model No. | JZC-11F |
| :---: | :---: |
| Contact Arrangement | 1H (SPDT) |
| Contact Material | Ag ALLOY |
| Contact Rating (at Resistive Load) | 5A 250VAC |
|  | 5A 30VDC |
| Max. Switching Voltage | 250VAC, 30VDC |
| Max. Operating Power | 1250 VA |
| Max. Switching current | 5A |
| Min. Switching Load | 10 mA 5 VDC |
| Coil Rated Voltage | 12VDC |
| Coil Rated Current | 37.5 mA |
| Pick-up Voltage | 9VDC ( $75 \%$ Max.of nominal voltage) |
| Drop-out Voltage | $0.6 \mathrm{~V}(5 \% \mathrm{Min}$. of nominal voltage) |
| Coil Resistance | $320 \Omega \pm 10 \%\left(\right.$ at $\left.20^{\circ} \mathrm{C}\right)$ |
| Contact Resistance | Max. $100 \mathrm{~m} \Omega$ (initial) |
| Insulation Resistance | Min. 100M $\Omega$ at 500VDC |
| Dielectric Strength <br> Between Coil \& Contact <br> Between Contacts | 3000VAC 50 HZ/60 HZ (1 minute) <br> 750VAC $50 \mathrm{HZ} / 60 \mathrm{HZ}$ (1 minute) |
| Operate Time | Max.20mSec. |
| Release Time | Max. 10 mSec . |
| Ambient Temperature | $-40^{\circ} \mathrm{C} \sim+85^{\circ} \mathrm{C}$ |
| Vibration Resistance (Endurance) | 1.5 mm D.A. $10-55 \mathrm{HZ}$ |
| Shock Resistance | Min. 10G Unerror |
| Mechanical Life | 10,000,000 Operations (no load) |
| Electrical Life | 100,000 Operations (at rated load) |

DIMENSIONS（in mm）
DRILLING PLAN（in mm）WIRING DIAGRAM
（Bottom View）
（Bottom View）


产品图片


## JZC-11F MATERIAL DESCRIPTION LIST

| ITEM | DESCRIPTION | MATERIAL | UL-94 |
| :---: | :--- | :--- | :--- |
| 1 | COVER | P.B.T | $94 \mathrm{~V}-0$ |
| 2 | BOBBIN | P.B.T | $94 \mathrm{~V}-0$ |
| 3 | BASE | P.B.T | $94 \mathrm{~V}-0$ |
| 4 | STATIONARY CONTACT | Ag ALLOY |  |
| 5 | MOVABLE CONTACT | Ag ALLOY |  |
| 6 | MOVABLE CONTACT ARM. | Cu ALLOY |  |
| 7 | STATIONARY CONTACT ARM. | Cu ALLOY |  |
| 8 | COIL TERMINAL | Cu ALLOY |  |
| 9 | CORE | STEEL |  |
| 10 | HEEL PIECE | STEEL |  |
| 11 | ARMATURE | STEEL |  |
| 12 | RESIN | EPOXY |  |
| 13 | MAGNET WIRE | QA-1 |  |

PACKINGLIST:

| MODEL | QTY. | N.W. | G.W. | MEAS. |
| :---: | :---: | :---: | :---: | :---: |
| JZC-11F | 2000 PCS. | 12KGS. | 13 KGS. | $48 \times 25 \times 24 \mathrm{~mm}$ |

## HC Serial Bluetooth Products

## User Instructional Manual

## 1 Introduction

HC serial Bluetooth products consist of Bluetooth serial interface module and Bluetooth adapter, such as:
(1) Bluetooth serial interface module:

Industrial level: $\quad \mathrm{HC}-03, \mathrm{HC}-04(\mathrm{HC}-04-\mathrm{M}, \mathrm{HC}-04-\mathrm{S})$
Civil level:

> HC-05, HC-06(HC-06-M, HC-06-S)

HC-05-D, HC-06-D (with baseboard, for test and evaluation)
(2) Bluetooth adapter:

HC-M4
HC-M6
This document mainly introduces Bluetooth serial module. Bluetooth serial module is used for converting serial port to Bluetooth. These modules have two modes: master and slaver device. The device named after even number is defined to be master or slaver when out of factory and can't be changed to the other mode. But for the device named after odd number, users can set the work mode (master or slaver) of the device by AT commands.

HC-04 specifically includes:
Master device: $\quad$ HC-04-M, M=master
Slave device: HC-04-S, S=slaver
The default situation of HC-04 is slave mode. If you need master mode, please state it clearly or place an order for $\mathrm{HC}-\mathrm{O} 4-\mathrm{M}$ directly.The naming rule of HC-06 is same.

When HC-03 and HC-05 are out of factory, one part of parameters are set for activating the device. The work mode is not set, since user can set the mode of HC-03, HC-05 as they want.

The main function of Bluetooth serial module is replacing the serial port line, such as:

1. There are two MCUs want to communicate with each other. One connects to Bluetooth master device while the other one connects to slave device. Their connection can be built once the pair is made. This Bluetooth connection is equivalently liked to a serial port line connection including RXD, TXD
signals. And they can use the Bluetooth serial module to communicate with each other.
2. When MCU has Bluetooth salve module, it can communicate with Bluetooth adapter of computers and smart phones. Then there is a virtual communicable serial port line between MCU and computer or smart phone.
3. The Bluetooth devices in the market mostly are salve devices, such as Bluetooth printer, Bluetooth GPS. So, we can use master module to make pair and communicate with them.

Bluetooth Serial module's operation doesn't need drive, and can communicate with the other Bluetooth device who has the serial. But communication between two Bluetooth modules requires at least two conditions:
(1) The communication must be between master and slave.
(2) The password must be correct.

However, the two conditions are not sufficient conditions. There are also some other conditions basing on different device model. Detailed information is provided in the following chapters.

In the following chapters, we will repeatedly refer to Linvor's (Formerly known as Guangzhou HC Information Technology Co., Ltd.) material and photos.

## 2 Selection of the Module

The Bluetooth serial module named even number is compatible with each other; The salve module is also compatible with each other. In other word, the function of HC-04 and HC-06, HC-03 and HC-05 are mutually compatible with each other. HC-04 and HC-06 are former version that user can't reset the work mode (master or slave). And only a few AT commands and functions can be used, like reset the name of Bluetooth (only the slaver), reset the password, reset the baud rate and check the version number. The command set of HC-03 and HC-05 are more flexible than HC-04 and HC-06's. Generally, the Bluetooth of $\mathrm{HC}-03 / \mathrm{HC}-05$ is recommended for the user.

Here are the main factory parameters of HC-05 and HC-06. Pay attention to the differences:

| HC-05 | HC-06 |
| :--- | :--- |
| Master and slave mode can be switched | Master and slave mode can't be switched |
| Bluetooth name: HC-05 | Bluetooth name: linvor |
| Password:1234 | Password:1234 |

Master role: have no function to remember the last paired salve device. It can be made paired to any slave device. In other words, just set AT + CMODE $=1$ when out of factory. If you want HC-05 to remember the last paired slave device address like HC-06, you can set AT+CMODE=0 after paired with the other device. Please refer the command set of HC-05 for the details.

Pairing: The master device can not only make pair with the specified Bluetooth address, like cell-phone, computer adapter, slave device, but also can search and make pair with the slave device automatically.

Typical method: On some specific conditions, master device and slave device can make pair with each other automatically. (This is the default method.)

Multi-device communication: There is only point to point communication for modules, but the adapter can communicate with multi-modules.

AT Mode 1: After power on, it can enter the AT mode by triggering PIN34 with high level. Then the baud rate for setting AT command is equal to the baud rate in communication, for example: 9600.

AT mode 2: First set the PIN34 as high level, or while on powering the module set the PIN34 to be high level, the Baud rate used here is 38400 bps . Notice: All AT commands can be operated only

Master role: have paired memory to remember last slave device and only make pair with that device unless KEY (PIN26) is triggered by high level. The default connected PIN26 is low level.

Pairing: Master device search and make pair with the slave device automatically.

Typical method: On some specific conditions, master and slave device can make pair with each other automatically.

Multi-device communication: There is only point to point communication for modules, but the adapter can communicate with multi-modules.

AT Mode: Before paired, it is at the AT mode. After paired it's at transparent communication.

| when the PIN34 is at high level. Only part of the AT commands can be used if PIN34 doesn't keep the high level after entering to the AT mode. Through this kind of designing, set permissions for the module is left to the user's external control circuit, that makes the application of HC-05 is very flexible. |  |
| :---: | :---: |
| During the process of communication, the module can enter to AT mode by setting PIN34 to be high level. By releasing PIN34, the module can go back to communication mode in which user can inquire some information dynamically. For example, to inquire the pairing is finished or not. | During the communication mode, the module can't enter to the AT mode. |
| Default communication baud rate: 9600, 4800-1.3M are settable. | Default communication baud rate: 9600, 1200-1.3M are settable. |
| KEY: PIN34, for entering to the AT mode. | KEY: PIN26, for master abandons memory. |
| LED1: PIN31, indicator of Bluetooth mode. Slow flicker $(1 \mathrm{~Hz})$ represents entering to the AT mode2, while fast flicker $(2 \mathrm{~Hz})$ represents entering to the AT mode1 or during the communication pairing. Double flicker per second represents pairing is finished, the module is communicable. <br> LED2: PIN32, before pairing is at low level, after the pairing is at high level. <br> The using method of master and slaver's indicator is the same. <br> Notice: The PIN of LED1 and LED2 are connected with LED+. | LED: The flicker frequency of slave device is 102 ms . If master device already has the memory of slave device, the flicker frequency during the pairing is $110 \mathrm{~ms} / \mathrm{s}$. If not, or master has emptied the memory, then the flicker frequency is $750 \mathrm{~m} / \mathrm{s}$. After pairing, no matter it's a master or slave device, the LED PIN is at high level. <br> Notice: The LED PIN connects to LED+ PIN. |
| Consumption: During the pairing, the current is | Consumption: During the pairing, the current is |


| fluctuant in the range of $30-40 \mathrm{~mA}$. The mean | fluctuant in the range of $30-40 \mathrm{~m}$. The mean |
| :--- | :--- |
| current is about 25 mA . After paring, no matter | current is about 25 mA . After paring, no matter |
| processing communication or not, the current is | processing communication or not, the current is |
| 8mA. There is no sleep mode. This parameter is | 8 mA . There is no sleep mode. This parameter is |
| same for all the Bluetooth modules. | same for all the Bluetooth modules. |
| Reset: PIN11, active if it's input low level. It can | Reset: PIN11, active if it's input low level. It can |
| be suspended in using. | be suspended in using. |
| Level: Civil | Level: Civil |

The table above that includes main parameters of two serial modules is a reference for user selection.

HC-03/HC-05 serial product is recommended.

## 3. Information of Package

The PIN definitions of HC-03, HC-04, HC-05 and HC-06 are kind of different, but the package size is the same: $28 \mathrm{~mm} * 15 \mathrm{~mm} * 2.35 \mathrm{~mm}$.

The following figure 1 is a picture of HC-06 and its main PINs. Figure 2 is a picture of HC-05 and its main PINs. Figure 3 is a comparative picture with one coin. Figure 4 is their package size information. When user designs the circuit, you can visit the website of Guangzhou HC Information Technology Co., Ltd. (www.wavesen.com) to download the package library of protle version.


Figure $1 \quad$ HC-06


Figure 2 HC-05


Figure 3 Comparative picture with one coin
$\qquad$ ．

LINVOR BLUE T

$$
\text { LV-BC-2. } 0
$$

www．linvor．com
单位：mm
PCB Layout 请参考实物

| NO | PIN NAME | NO PIN NAME |
| :---: | :--- | :--- |
| 1 | TX | 20 USB D＋ |
| 2 | RX | 21 |
| 3 | CTS | 22 |

Figure 4 Package size information

## 4. The Using and Testing Method of HC-06 for the First Time

This chapter will introduce the using method of HC-06 in detail. User can test the module according to this chapter when he or she uses the module at the first time.

PINs description:

| PIN1 | UART_TXD, TTL/CMOS level, UART Data output |
| :--- | :--- |
| PIN2 | UART_RXD, TTL/COMS level, s UART Data input |
| PIN11 | RESET, the reset PIN of module, inputting low level can reset the module, <br> when the module is in using, this PIN can connect to air. |
| PIN12 | VCC, voltage supply for logic, the standard voltage is 3.3V, and can work <br> at 3.0-4.2V |
| PIN13 | GND |
| PIN22 | GND <br> Slave device: Before paired, this PIN outputs the period of 102ms square <br> wave. After paired, this PIN outputs high level. <br> Master device: On the condition of having no memory of pairing with a <br> slave device, this PIN outputs the period of 110ms square wave. On the <br> condition of having the memory of pairing with a slave device, this PIN <br> outputs the period of 750ms square wave. After paired, this PIN outputs <br> high level. |
| PIN26 | For master device, this PIN is used for emptying information about <br> pairing. After emptying, master device will search slaver randomly, then <br> remember the address of the new got slave device. In the next power on, <br> master device will only search this address. |

(1) The circuit 1 (connect the module to 3.3 V serial port of MCU) is showed by figure 5.


Figure 5 The circuit 1
In principle, HC-06 can work when UART_TXD, UART_RXD, VCC and GND are connected. However, for better testing results, connecting LED and KEY are recommended (when testing the master).

Where, the 3.3 V TXD of MCU connects to HC-06's UART_RXD, the 3.3V RXD of MCU connects to HC-06's UART_TXD, and 3.3 V power and GND should be connected. Then the minimum system is finished.

Note that, the PIN2:UART_RXD of Bluetooth module has no pull-up resistor. If the MCU TXD doesn't have pull-up function, then user should add a pull-up resistor to the UART_RXD. It may be easy to be ignored.

If there are two MCU which connect to master and slave device respectively, then before paired(LED will flicker) user can send AT commands by serial port when the system is power on. Please refer to HC-04 and HC-06's data sheet for detailed commands. In the last chapter, the command set will be introduced. Please pay attention to that the command of HC-04/HC-06 doesn't have terminator. For example, consider the call command, sending out AT is already enough, need not add the CRLF (carriage return line feed).

If the LED is constant lighting, it indicates the pairing is finished. The two MCUs can communicate with each other by serial port. User can think there is a serial port line between two MCUs.

## (2) The circuit 2 (connect the module to 5 V serial port of MCU) is showed by figure 6.

Figure 6 is the block diagram of Bluetooth baseboard. This kind of circuit can amplify Bluetooth module's operating voltage to $3.1-6.5 \mathrm{~V}$. In this diagram, the J1 port can not only be connected with MCU system of 3.3 V and 5 V , but also can be connected with computer serial port.


Figure 6 The circuit 2

## (3) AT command test

Before paired, the mode of HC-04 and HC-06 are AT mode.
On the condition of 9600 N 81 , OK will be received when user send the two letters AT. Please refer to the last chapter of datasheet for other commands of HC-06. Please pay attention to that sending out AT is already enough, need not add the CRLF (carriage return line feed).

The command set of Version V1.4 doesn't include parity. The version V1.5 and its later version have parity function. Moreover, there are three more commands of V1.5 than V1.4. They are:

No parity (default)
AT+PN

| Odd parity | $\mathrm{AT}+\mathrm{PO}$ |
| :--- | :--- |
| Even parity | $\mathrm{AT}+\mathrm{PE}$ |

Do not let the sending frequency of AT command of $\mathrm{HC}-06$ exceed 1 Hz , because the command of HC-06 end or not is determined by the time interval.

## (4) Pairing with adapter

User can refer to the download center of the company's website for "The Introduction of IVT" that introduces the Bluetooth module makes pair with computer adapter. That document taking HC-06-D for example introduces how the serial module makes pair with the adapter. That method is like to make pair with cell-phone. But the difference is that cell-phone need a third-party communication software to help. It's liked the kind of PC serial helper of and the hyper terminal. A software named "PDA serial helper" provided by our company is suitable for WM system. It has been proven that this serial module is supported by many smart phone systems' Bluetooth, such as, sybian, android, windows mobile and etc.

## (5) Pairing introduction

HC-06 master device has no memory before the first use. If the password is correct, the mater device will make pair with the slave device automatically in the first use. In the following use, the master device will remember the Bluetooth address of the last paired device and search it. The searching won't stop until the device is found. If master device's PIN26 is input high level, the device will lose the memory. In that occasion, it'll search the proper slave device like the first use. Based on this function, the master device can be set to make pair with the specified address or any address by user.

## (6) Reset new password introduction

User can set a new password for the HC-06 through AT+PINxxxx command. But the new password will become active after discharged all the energy of the module. If the module still has any energy, the old one is still active. In the test, for discharging all the system energy and activating the new password, we can connect the power supply PIN with GND about 20 seconds after the power is cut off. Generally, shutting down the device for 30 minutes also can discharge the energy, if there is no peripheral circuit helps discharge energy. User should make the proper way according to the specific situation.

## (7) Name introduction

If the device has no name, it's better that user doesn't try to change the master device name. The name should be limited in 20 characters.

Summary: The character of HC-06: 1 not many command 2 easy for application 3 low price. It's good for some specific application. HC-04 is very similar with HC-06. Their only one difference is HC-04 is for industry, HC-06 is for civil. Except this, they don't have difference.

The following reference about $\mathrm{HC}-04$ and $\mathrm{HC}-06$ can be downloaded from company website www.wavesen.com:

HC-06 datasheet .pdf
HC-04 datasheet .pdf
IVT BlueSoleil-2.6
Bluetooth FAQ.pdf
HC-04-D(HD-06-D)datasheet(English).pdf
HC-06-AT command software (test version)
PCB package of Bluetooth key modules
IVT software manual.pdf

PDA serial test helper.exe
(the command set introduction is included)
(the command set introduction is included)
(IVT Bluetooth drive test version)
(some commands in V1.5 is not supported by V1.4) (PCB package lib in protel)
(introduce how to operate the modern and make pair with Bluetooth module)
(serial helper used for WM system)

## 5 manual for the first use of HC-05

This chapter will introduce how to test and use the HC-05 if it's the first time for user to operate it.

## (1) PINs description

| PIN1 | UART_TXD, Bluetooth serial signal sending PIN, can connect with MCU's RXD PIN |
| :--- | :--- |
| PIN2 | UART_RXD, Bluetooth serial signal receiving PIN, can connect with the MCU's TXD PIN, <br> there is no pull-up resistor in this PIN. But It needs to be added an eternal pull-up resistor. |
| PIN11 | RESET, the reset PIN of module, inputting low level can reset the module, when the module <br> is in using, this PIN can connect to air. |
| PIN12 | VCC, voltage supply for logic, the standard voltage is 3.3V, and can work at 3.0-4.2V |
| PIN13 | GND |


| PIN31 | LED1, indicator of work mode. Has 3 modes: <br> When the module is supplied power and PIN34 is input high level, PIN31 output 1Hz square <br> wave to make the LED flicker slowly. It indicates that the module is at the AT mode, and the <br> baud rate is 38400; <br> When the module is supplied power and PIN34 is input low level, PIN31 output 2Hz square <br> wave to make the LED flicker quickly. It indicates the module is at the pairable mode. If <br> PIN34 is input high level, then the module will enter to AT mode, but the output of PIN31 is <br> still 2Hz square wave. <br> After the pairing, PIN31 output 2Hz square ware. <br> Note: if PIN34 keep high level, all the commands in the AT command set can be in <br> application. Otherwise, if just excite PIN34 with high level but not keep, only some command <br> can be used. More information has provided at chapter 2. |
| :--- | :--- |
| PIN32 | Output terminal. Before paired, it output low level. Once the pair is finished, it output high <br> level. |
| Mode switch input. If it is input low level, the module is at paired or communication mode. If <br> it's input high level, the module will enter to AT mode. Even though the module is at <br> communication, the module can enter to the AT mode if PIN34 is input high level. Then it will <br> go back to the communication mode if PIN34 is input low level again. |  |

## (2) Application circuit 1 (connect to the 3.3 V system)



Figure 7 Application 1
(3) Application circuit 2 (connect to 5V serial system or PC serial)


Figure 8 Application circuit 2

## (4) AT command test

This chapter introduces some common commands in use. The detail introduction about HC-05 command is in HC-0305 AT command set.

## Enter to AT mode:

Way1: Supply power to module and input high level to PIN34 at the same time, the module will enter to AT mode with the baud rate- 38400 .

Way2: In the first step, supply power to module; In the second step, input high level to PIN34. Then the module will enter to AT mode with the baud rate-9600. Way1 is recommended.

Command structure: all command should end up with " $\backslash \mathrm{r} \backslash n$ " (Hex: 0X0D X0A) as the terminator. If
the serial helper is installed, user just need enter "ENTER" key at the end of command.

## Reset the master-slave role command:

AT + ROLE $=0 \quad$----Set the module to be salve mode. The default mode is salve.
AT + ROLE $=1 \quad---$-Set the module to be master mode.

## Set memory command:

$\mathrm{AT}+\mathrm{CMODE}=1$
Set the module to make pair with the other random Bluetooth module (Not specified address). The default is this mode.
$\mathrm{AT}+\mathrm{CMODE}=1$
Set the module to make pair with the other Bluetooth module (specified address). If set the module to make pair with random one first, then set the module to make pair with the Bluetooth module has specified address. Then the module will search the last paired module until the module is found.

## Reset the password command

AT+PSWD=XXXX
Set the module pair password. The password must be 4-bits.

## Reset the baud rate

AT+UART== <Param>,<Param2>,<Param3>.
More information is provided at HC-0305 command set

## Example:

AT+UART=9600,0,0 ----set the baud rate to be 9600N81
Reset the Bluetooth name
AT+NAME=XXXXX

## Summary:

HC-05 has many functions and covers all functions of HC-06. The above commands are the most common ones. Besides this, HC-05 leaves lots of space for user. So HC-05 is better than HC-06 and
recommended. HC-03 is similar with HC-05. The above introduction also suits HC-03
The following reference about HC-03 and HC-05 can be downloaded from company website www.wavesen.com:

HC-03 datasheet .pdf
HC-05 datasheet .pdf
IVT BlueSoleil-2.6
Bluetooth FAQ.pdf
PCB package of Bluetooth key modules
IVT software manual.pdf

PDA serial test helper.exe
(the command set introduction is included)
(the command set introduction is included)
(IVT Bluetooth drive test version)
(PCB package lib in protel)
(introduce how to operate the modern and make pair with Bluetooth module)
(serial helper used for WM system)

HC-03/05 Bluetooth serial command set.pdf

## 6. Ordering information

The website of Guangzhou HC Information Technology Co., Ltd is www.wavesen.com The contact information is provided at the company website.

Order Way: If you want our product, you can give order to the production center of our company directly or order it in Taobao. There is a link to Taobao in our company website.

Package: 50 pieces chips in an anti-static blister package. The weight of a module is about 0.9 g . The weight of a package is about 50 g .


Please provide the product's model when you order:
HC-04-M $\quad \mathrm{HC}-04$ master module
HC-04-S HC-04 slave module
HC-06-M HC-06 master module
HC-06-S HC-06 slave module
HC-03
HC-05 HC-03/05 can be preset to be master module or slave module.


UPG

Sealed Lead-Acid Battery
Absorbant Glass Mat (AGM) technology for superior performance. Valve regulated, spill proof construction allows safe operation in any position. Approved for transport by air. D.O.T., I.A.T.A., F.A.A. and C.A.B. certified. U.L. r ecognized under


Charge Method (Constant Voltage)
Cycle Use (Repeating Use)

| Initial Current | 17.5 A or smaller |
| :--- | :--- |
| Control Voltage | $14.5-14.9 \mathrm{~V}$ |

Float Use
Control Voltage $\quad 13.6-13.8 \mathrm{~V}$

## Physical Dimensions: in (mm)



$$
\begin{aligned}
& \text { L: } 7.72 \mathrm{in}(196.1 \mathrm{~mm}) \\
& \text { W: } 6.50 \mathrm{in}(165.1 \mathrm{~mm}) \\
& \text { H: } 6.89 \mathrm{in}(175 \mathrm{~mm}) \\
& \text { TH: } 7.17 \mathrm{in}(182.1 \mathrm{~mm}) \\
& \text { Tolerances are }+/-0.04 \mathrm{in} .(+/-1 \mathrm{~mm}) \\
& \text { and }+/-0.08 \text { in. }(+/-2 \mathrm{~mm}) \text { for height } \\
& \text { dimensions. All data subject to } \\
& \text { change without notice. }
\end{aligned}
$$

Constant Current Discharge Characteristics Unit:A $\left(25^{\circ} \mathrm{C}, 77^{\circ} \mathrm{F}\right)$

| F.V/Time | 5MIN | 10 MIN | 15 MIN | 30MIN | 1HR | 2HR | 3HR | 4HR | 5HR | 8HR | 10HR | 20HR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 9.60 V | 158.2 | 115.4 | 81.2 | 49.2 | 25.7 | 15.0 | 11.0 | 8.6 | 7.1 | 5.0 | 4.5 | 2.5 |
| 10.20 V | 139.4 | 105.2 | 72.7 | 46.6 | 24.1 | 14.3 | 10.7 | 8.3 | 6.9 | 4.9 | 4.4 | 2.5 |
| 10.50 V | 134.2 | 100.0 | 68.4 | 45.3 | 23.5 | 13.9 | 10.4 | 8.2 | 6.8 | 4.8 | 4.3 | 2.5 |
| 10.80 V | 129.1 | 94.9 | 64.1 | 44.0 | 22.7 | 13.6 | 10.2 | 8.1 | 6.7 | 4.7 | 4.3 | 2.3 |
| 11.10 V | 124.0 | 89.8 | 59.9 | 42.8 | 21.8 | 13.3 | 9.8 | 7.8 | 6.5 | 4.6 | 4.1 | 2.2 |

Constant Power Discharge Characteristics Unit:W $\left(25^{\circ} \mathrm{C}, 77^{\circ} \mathrm{F}\right)$

| F.V/Time | 5MIN | 10MIN | 15MIN | 30MIN | 1HR | 2HR | 3HR | 4HR | 5HR | 8HR | 10HR | 20HR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 9.60 V | 1679.6 | 1268.4 | 862.7 | 522.0 | 297.1 | 173.1 | 127.4 | 99.2 | 81.7 | 57.7 | 52.2 | 28.1 |
| 10.20 V | 1547.1 | 1167.5 | 806.7 | 517.3 | 279.2 | 165.4 | 124.0 | 96.6 | 64.2 | 56.4 | 50.9 | 27.4 |
| 10.50 V | 1522.3 | 1134.6 | 775.5 | 513.9 | 270.2 | 161.6 | 121.0 | 94.9 | 79.1 | 56.0 | 50.0 | 27.0 |
| 10.80 V | 1502.7 | 1104.7 | 746.4 | 512.6 | 262.9 | 158.2 | 118.4 | 93.2 | 77.8 | 54.7 | 49.6 | 26.9 |
| 11.10 V | 1475.3 | 1068.3 | 712.2 | 508.7 | 259.5 | 157.7 | 117.1 | 92.8 | 77.4 | 54.3 | 48.3 | 26.1 |

Discharge Time vs. Discharge Current


Discharge Characteristics


Shelf Life \& Storage


Open Circuit Voltage vs Residual Capacity



Charge Current \& Final Discharge Voltage

| Application | Charge Voltage(V/Cell) |  |  | Max.Charge Current | Final Discharge | 1.75 | 1.70 | 1.60 | 1.30 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Temperature | Set Point | Allowable Range |  | Voltage V/Cell |  |  |  |  |
| Cycle Use | $25^{\circ} \mathrm{C}\left(77^{\circ} \mathrm{F}\right)$ | 2.45 | 2.40~2.50 | 0.35C | Discharge | $0.2 C>(A)$ | $0.2 \mathrm{C}<(\mathrm{A})<0.5 \mathrm{C}$ | $0.5 C<(A)<1.0 C$ | $(\mathrm{A})>1.0 \mathrm{C}$ |
| Standby | $25^{\circ} \mathrm{C}\left(77^{\circ} \mathrm{F}\right)$ | 2.325 | 2.30~2.35 |  | Current(A) |  |  |  |  |

## Power MOSFET

## PRODUCT SUMMARY

| PRODUCT SUMMARY |  |  |
| :--- | :---: | :---: |
| $\mathrm{V}_{\mathrm{DS}}(\mathrm{V})$ | 100 |  |
| $\mathrm{R}_{\mathrm{DS}(\mathrm{on})}(\Omega)$ | $\mathrm{V}_{\mathrm{GS}}=5.0 \mathrm{~V}$ | 0.077 |
| $\mathrm{Q}_{\mathrm{g}}(\mathrm{Max}).(\mathrm{nC})$ | 64 |  |
| $\mathrm{Q}_{\mathrm{gs}}(\mathrm{nC})$ | 9.4 |  |
| $\mathrm{Q}_{\mathrm{gd}}(\mathrm{nC})$ | 27 |  |
| Configuration | Single |  |



N-Channel MOSFET

## FEATURES

- Dynamic dV/dt Rating
- Repetitive Avalanche Rated
- Logic-Level Gate Drive
- $R_{D S(o n)}$ Specified at $V_{G S}=4 \mathrm{~V}$ and 5 V
- $175{ }^{\circ} \mathrm{C}$ Operating Temperature
- Fast Switching
- Ease of Paralleling
- Compliant to RoHS Directive 2002/95/EC


## DESCRIPTION

Third generation Power MOSFETs from Vishay provide the designer with the best combination of fast switching, ruggedized device design, low on-resistance and cost-effectiveness.
The TO-220AB package is universally preferred for all commercial-industrial applications at power dissipation levels to approximately 50 W . The low thermal resistance and low package cost of the TO-220AB contribute to its wide acceptance throughout the industry.

## ORDERING INFORMATION

| Package | TO-220AB |
| :--- | :--- |
| Lead (Pb)-free | IRL540PbF |
|  | SiHL540-E3 |
| SnPb | IRL540 |
|  | SiHL540 |


| ABSOLUTE MAXIMUM RATINGS ( $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$, unless otherwise noted) |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| PARAMETER |  |  | SYMBOL | LIMIT | UNIT |
| Drain-Source Voltage |  |  | $\mathrm{V}_{\mathrm{DS}}$ | 100 | V |
| Gate-Source Voltage |  |  | $\mathrm{V}_{\mathrm{GS}}$ | $\pm 10$ |  |
| Continuous Drain Current | $\mathrm{V}_{\mathrm{GS}}$ at 5.0 V | $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$ | ID | 28 | A |
|  |  | $\mathrm{T}_{\mathrm{C}}=100^{\circ} \mathrm{C}$ |  | 20 |  |
| Pulsed Drain Current ${ }^{\text {a }}$ |  |  | $\mathrm{I}_{\mathrm{DM}}$ | 110 |  |
| Linear Derating Factor |  |  |  | 1.0 | W/ ${ }^{\circ} \mathrm{C}$ |
| Single Pulse Avalanche Energy ${ }^{\text {b }}$ |  |  | $\mathrm{E}_{\text {AS }}$ | 440 | mJ |
| Avalanche Current ${ }^{\text {a }}$ |  |  | $\mathrm{I}_{\text {AR }}$ | 28 | A |
| Repetitive Avalanche Energy ${ }^{\text {a }}$ |  |  | $\mathrm{E}_{\text {AR }}$ | 15 | mJ |
| Maximum Power Dissipation | $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$ |  | $\mathrm{P}_{\mathrm{D}}$ | 150 | W |
| Peak Diode Recovery dV/dt ${ }^{\text {c }}$ |  |  | $\mathrm{dV} / \mathrm{dt}$ | 5.5 | V/ns |
| Operating Junction and Storage Temperature Range |  |  | $\mathrm{T}_{\mathrm{J},} \mathrm{T}_{\text {stg }}$ | -55 to + 175 | ${ }^{\circ} \mathrm{C}$ |
| Soldering Recommendations (Peak Temperature) | for 10 s |  |  | $300^{\text {d }}$ |  |
| Mounting Torque | 6-32 or M3 screw |  |  | 10 | lbf - in |
|  |  |  |  | 1.1 | $\mathrm{N} \cdot \mathrm{m}$ |

## Notes

a. Repetitive rating; pulse width limited by maximum junction temperature (see fig. 11).
b. $\mathrm{V}_{\mathrm{DD}}=25 \mathrm{~V}$, starting $\mathrm{T}_{\mathrm{J}}=25^{\circ} \mathrm{C}, \mathrm{L}=841 \mu \mathrm{H}, \mathrm{R}_{\mathrm{g}}=25 \Omega, \mathrm{I}_{\mathrm{AS}}=28 \mathrm{~A}$ (see fig. 12c).
c. $\mathrm{I}_{\mathrm{SD}} \leq 28 \mathrm{~A}, \mathrm{dl} / \mathrm{dt} \leq 170 \mathrm{~A} / \mu \mathrm{s}, \mathrm{V}_{\mathrm{DD}} \leq \mathrm{V}_{\mathrm{DS}}, \mathrm{T}_{\mathrm{J}} \leq 175^{\circ} \mathrm{C}$.
d. 1.6 mm from case.

## Vishay Siliconix

THERMAL RESISTANCE RATINGS

| PARAMETER | SYMBOL | TYP. | MAX. | UNIT |
| :--- | :---: | :---: | :---: | :---: |
| Maximum Junction-to-Ambient | $\mathrm{R}_{\text {thJA }}$ | - | 62 |  |
| Case-to-Sink, Flat, Greasd Surface | $\mathrm{R}_{\text {thcs }}$ | 0.50 | - |  |
| Maximum Junction-to-Case (Drain) | $\mathrm{R}_{\text {thJc }}$ | - | 1.0 |  |



## Notes

a. Repetitive rating; pulse width limited by maximum junction temperature (see fig. 11).
b. Pulse width $\leq 300 \mu \mathrm{~s}$; duty cycle $\leq 2 \%$.


Fig. 1 - Typical Output Characteristics, $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$


Fig. 2-Typical Output Characteristics, $\mathrm{T}_{\mathrm{C}}=175^{\circ} \mathrm{C}$


Fig. 3-Typical Transfer Characteristics


Fig. 4 - Normalized On-Resistance vs. Temperature


Fig. 5 - Typical Capacitance vs. Drain-to-Source Voltage


Fig. 6 - Typical Gate Charge vs. Gate-to-Source Voltage


Fig. 7 - Typical Source-Drain Diode Forward Voltage


Fig. 8 - Maximum Safe Operating Area

IRL540, SiHL540
Vishay Siliconix


Fig. 9 - Maximum Safe Operating Area


Fig. 10a - Switching Time Test Circuit


Fig. 10b - Switching Time Waveforms


Fig. 11 - Maximum Effective Transient Thermal Impedance, Junction-to-Case


Fig. 12a - Unclamped Inductive Test Circuit


Fig. 12b - Unclamped Inductive Waveforms


Fig. 12c - Maximum Avalanche Energy vs. Drain Current


Fig. 13a - Basic Gate Charge Waveform


Fig. 13b - Gate Charge Test Circuit

## Peak Diode Recovery dV/dt Test Circuit



Fig. 14 - For N-Channel

[^3]

| DIM. | MILLIMETERS |  | INCHES |  |
| :---: | :---: | :---: | :---: | :---: |
|  | MIN. | MAX. | MIN. | MAX. |
| A | 4.24 | 4.65 | 0.167 | 0.183 |
| b | 0.69 | 1.02 | 0.027 | 0.040 |
| b(1) | 1.14 | 1.78 | 0.045 | 0.070 |
| c | 0.36 | 0.61 | 0.014 | 0.024 |
| D | 14.33 | 15.85 | 0.564 | 0.624 |
| E | 9.96 | 10.52 | 0.392 | 0.414 |
| e | 2.41 | 2.67 | 0.095 | 0.105 |
| $\mathrm{e}(1)$ | 4.88 | 5.28 | 0.192 | 0.208 |
| F | 1.14 | 1.40 | 0.045 | 0.055 |
| $\mathrm{H}(1)$ | 6.10 | 6.71 | 0.240 | 0.264 |
| $\mathrm{~J}(1)$ | 2.41 | 2.92 | 0.095 | 0.115 |
| L | 13.36 | 14.40 | 0.526 | 0.567 |
| $\mathrm{~L}(1)$ | 3.33 | 4.04 | 0.131 | 0.159 |
| $\varnothing$ P | 3.53 | 3.94 | 0.139 | 0.155 |
| Q | 2.54 | 3.00 | 0.100 | 0.118 |

ECN: X15-0364-Rev. C, 14-Dec-15
DWG: 6031
Note

- $\mathrm{M}^{*}=0.052$ inches to 0.064 inches (dimension including protrusion), heatsink hole for HVM



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## Power MOSFET

| PRODUCT SUMMARY |  |  |
| :--- | :--- | :--- |
| $\mathrm{V}_{\mathrm{DS}}(\mathrm{V})$ | -50 |  |
| $\mathrm{R}_{\mathrm{DS}(o n)}(\Omega)$ | $\mathrm{V}_{\mathrm{GS}}=-10 \mathrm{~V}$ | 0.14 |
| $\mathrm{Q}_{\mathrm{g}}(\mathrm{Max}).(\mathrm{nC})$ | 39 |  |
| $\mathrm{Q}_{\mathrm{gs}}(\mathrm{nC})$ | 10 |  |
| $\mathrm{Q}_{\mathrm{gd}}(\mathrm{nC})$ | 15 |  |
| Configuration | Single |  |



P-Channel MOSFET

## FEATURES

- P-Channel Versatility
- Compact Plastic Package
- Fast Switching
- Low Drive Current
- Ease of Paralleling
- Excellent Temperature Stability
- Material categorization: For definitions of compliance please see www.vishay.com/doc?99912


## Note

* Lead (Pb)-containing terminations are not RoHS-compliant. Exemptions may apply.


## DESCRIPTION

The power MOSFET technology is the key to Vishay's advanced line of power MOSFET transistors. The efficient geometry and unique processing of the power MOSFET design achieve very low on-state resistance combined with high transconductance and extreme device ruggedness.
The p-channel power MOSFET's are designed for application which require the convenience of reverse polarity operation. They retain all of the features of the more common n-channel Power MOSFET's such as voltage control, very fast switching, ease of paralleling, and excellent temperature stability.
P-channel power MOSFETs are intended for use in power stages where complementary symmetry with n-channel devices offers circuit simplification. They are also very useful in drive stages because of the circuit versatility offered by the reverse polarity connection. Applications include motor control, audio amplifiers, switched mode converters, control circuits and pulse amplifiers.

## ORDERING INFORMATION

| Package | TO-22OAB |
| :--- | :--- |
| Lead $(\mathrm{Pb})$-free | IRF9Z30PbF |
|  | SiHF9Z30-E3 |
| SnPb | IRF9Z30 |
|  | SiHF9Z30 |


| ABSOLUTE MAXIMUM RATINGS ( $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$, unless otherwise noted) |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| PARAMETER |  |  | SYMBOL | LIMIT | UNIT |
| Drain-Source Voltage |  |  | $\mathrm{V}_{\text {DS }}$ | -50 | V |
| Gate-Source Voltage |  |  | $\mathrm{V}_{\mathrm{GS}}$ | $\pm 20$ |  |
| Continuous Drain Current | $\mathrm{V}_{\mathrm{GS}}$ at - 10 V | $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$ | ID | -18 | A |
|  |  | $\mathrm{T}_{\mathrm{C}}=100^{\circ} \mathrm{C}$ |  | -11 |  |
| Pulsed Drain Current ${ }^{\text {a }}$ |  |  | $\mathrm{I}_{\mathrm{DM}}$ | -60 |  |
| Linear Derating Factor |  |  |  | 0.59 | W/ ${ }^{\circ} \mathrm{C}$ |
| Inductive Current, Clamped | $\mathrm{L}=100 \mu \mathrm{H}$ |  | ILM | -60 | A |
| Unclamped Inductive Current (Avalanche Current) |  |  | $\mathrm{I}_{\mathrm{L}}$ | -3.1 | A |
| Maximum Power Dissipation | $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$ |  | $\mathrm{P}_{\mathrm{D}}$ | 74 | W |
| Operating Junction and Storage Temperature Range |  |  | $\mathrm{T}_{\mathrm{J}}, \mathrm{T}_{\text {stg }}$ | -55 to +150 | ${ }^{\circ} \mathrm{C}$ |
| Soldering Recommendations (Peak Temperature) | for 10 s |  |  | $300{ }^{\text {c }}$ |  |

## Notes

a. Repetitive rating; pulse width limited by maximum junction temperature (see fig. 14).
b. $\mathrm{V}_{\mathrm{DD}}=-25 \mathrm{~V}$, starting $\mathrm{T}_{\mathrm{J}}=25^{\circ} \mathrm{C}, \mathrm{L}=100 \mu \mathrm{H}, \mathrm{R}_{\mathrm{g}}=25 \Omega$
c. 0.063 " $(1.6 \mathrm{~mm})$ from case.

| THERMAL RESISTANCE RATINGS |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: |
| PARAMETER | SYMBOL | TYP. | MAX. | UNIT |  |
| Maximum Junction-to-Ambient | $\mathrm{R}_{\text {thJA }}$ | - | 80 |  |  |
| Maximum Junction-to-Case (Drain) | $\mathrm{R}_{\text {thJc }}$ | - | 1.7 |  |  |


| SPECIFICATIONS ( $\mathrm{T}_{J}=25^{\circ} \mathrm{C}$, unless otherwise noted) |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PARAMETER | SYMBOL | TEST CONDITIONS |  | MIN. | TYP. | MAX. | UNIT |
| Static |  |  |  |  |  |  |  |
| Drain-Source Breakdown Voltage | $\mathrm{V}_{\mathrm{DS}}$ | $\mathrm{V}_{\mathrm{GS}}=0 \mathrm{~V}, \mathrm{I}_{\mathrm{D}}=-250 \mu \mathrm{~A}$ |  | - 50 | - | - | V |
| Gate-Source Threshold Voltage | $\mathrm{V}_{\mathrm{GS}(\text { (th) }}$ | $\mathrm{V}_{\mathrm{DS}}=\mathrm{V}_{\mathrm{GS}}, \mathrm{I}_{\mathrm{D}}=-250 \mu \mathrm{~A}$ |  | -2.0 | - | -4.0 | V |
| Gate-Source Leakage | $\mathrm{I}_{\text {GSS }}$ | $\mathrm{V}_{\mathrm{GS}}= \pm 20 \mathrm{~V}$ |  | - | - | $\pm 500$ | nA |
| Zero Gate Voltage Drain Current | Idss | $\mathrm{V}_{\mathrm{DS}}=$ max. rating, $\mathrm{V}_{\mathrm{GS}}=0 \mathrm{~V}$ |  | - | - | -250 | $\mu \mathrm{A}$ |
|  |  | $\begin{gathered} \mathrm{V}_{\mathrm{DS}}=\text { max. } \\ \mathrm{T}_{\mathrm{J}}=125^{\circ} \mathrm{C}, \\ \hline 0.8 \mathrm{~V} \\ \hline \end{gathered}$ |  | - | - | - 1000 |  |
| Drain-Source On-State Resistance | $\mathrm{R}_{\mathrm{DS} \text { (on) }}$ | $V_{G S}=-10 \mathrm{~V}$ | $\mathrm{I}_{\mathrm{D}}=-9.3 \mathrm{~A}^{\mathrm{b}}$ | - | 0.093 | 0.14 | $\Omega$ |
| Forward Transconductance | $\mathrm{g}_{\mathrm{fs}}$ | $\mathrm{V}_{\mathrm{DS}}=2 \times \mathrm{V}_{\mathrm{GS}}, \mathrm{I}_{\mathrm{DS}}=-9 \mathrm{~A}^{\mathrm{b}}$ |  | 3.1 | 4.7 | - | S |
| Dynamic |  |  |  |  |  |  |  |
| Input Capacitance | $\mathrm{C}_{\text {iss }}$ | $\begin{gathered} V_{G S}=0 V, \\ V_{D S}=-25 \mathrm{~V}, \\ f=1.0 \mathrm{MHz}, \text { see fig. } 9 \end{gathered}$ |  | - | 900 | - | pF |
| Output Capacitance | $\mathrm{Cosss}^{\text {a }}$ |  |  | - | 570 | - |  |
| Reverse Transfer Capacitance | $\mathrm{C}_{\text {rss }}$ |  |  | - | 140 | - |  |
| Total Gate Charge | $\mathrm{Q}_{\mathrm{g}}$ | $\mathrm{V}_{\mathrm{GS}}=-10 \mathrm{~V}$ | $\begin{aligned} & I_{D}=-18 \mathrm{~A}, V_{D S}=-0.8 \\ & \text { max. rating. see fig. } 17 \end{aligned}$ | - | 26 | 39 | nC |
| Gate-Source Charge | $\mathrm{Q}_{\mathrm{gs}}$ |  |  | - | 6.9 | 10 |  |
| Gate-Drain Charge | $\mathrm{Q}_{\mathrm{gd}}$ |  |  | - | 9.7 | 15 |  |
| Turn-On Delay Time | $\mathrm{t}_{\mathrm{d}(\mathrm{on})}$ | $V_{D D}=-25 \mathrm{~V}, \mathrm{I}_{\mathrm{D}}=-18 \mathrm{~A} \text {, }$ <br> $R_{g}=13 \Omega, R_{D}=1.3 \Omega$, see fig. 16 <br> (MOSFET switching times are essentially independent of operating temperature) |  | - | 12 | 18 | ns |
| Rise Time | $\mathrm{t}_{\mathrm{r}}$ |  |  | - | 110 | 170 |  |
| Turn-Off Delay Time | $t_{\text {d(off }}$ |  |  | - | 21 | 32 |  |
| Fall Time | $\mathrm{t}_{\mathrm{f}}$ |  |  | - | 64 | 96 |  |
| Drain-Source Body Diode Characteristics |  |  |  |  |  |  |  |
| Continuous Source-Drain Diode Current | Is | MOSFET symbol showing the integral reverse $\mathrm{p}-\mathrm{n}$ junction diode |  | - | - | -18 | A |
| Pulsed Diode Forward Current ${ }^{\text {a }}$ | $I_{\text {SM }}$ |  |  | - | - | -60 |  |
| Body Diode Voltage | $\mathrm{V}_{\text {SD }}$ | $\mathrm{T}_{\mathrm{J}}=25^{\circ} \mathrm{C}$ | $\mathrm{I}_{\mathrm{S}}=-18 \mathrm{~A}, \mathrm{~V}_{\mathrm{GS}}=0 \mathrm{Vb}^{\text {b }}$ | - | - | -6.3 | V |
| Body Diode Reverse Recovery Time | $\mathrm{t}_{\mathrm{rr}}$ | $\mathrm{T}_{J}=25^{\circ} \mathrm{C}, \mathrm{I}_{\mathrm{F}}=-18 \mathrm{~A}, \mathrm{dl} / \mathrm{dt}=100 \mathrm{~A} / \mu \mathrm{s}^{\mathrm{b}}$ |  | 54 | 120 | 250 | ns |
| Body Diode Reverse Recovery Charge | $\mathrm{Q}_{\mathrm{rr}}$ |  |  | 0.20 | 0.47 | 1.1 | $\mu \mathrm{C}$ |

## Notes

a. Repetitive rating; pulse width limited by maximum junction temperature (see fig. 14).
b. Pulse width $\leq 300 \mu$ s; duty cycle $\leq 2 \%$.


Fig. 1 - Typical Output Characteristics


Fig. 2-Typical Transfer Characteristics


Fig. 3-Typical Saturation Characteristics


Fig. 4 - Maximum Safe Operating Area

IRF9Z30, SiHF9Z30


- $I_{D}$, Drain Current (A)

Fig. 5 - Typical Transconductance vs. Drain Current


Fig. 6 - Typical Source-Drain Diode Forward Voltage

$\mathrm{T}_{\mathrm{J}}$, Junction Temperature $\left({ }^{\circ} \mathrm{C}\right)$

Fig. 7 - Breakdown Voltage vs. Temperature


Fig. 8 - Normalized On-Resistance vs. Temperature


- $\mathrm{V}_{\mathrm{DS}}$, Drain-to-Source Voltage (V)

Fig. 9 - Typical Capacitance vs. Drain-to-Source Voltage


Fig. 10 - Typical Gate Charge vs. Gate-to-Source Voltage


Fig. 11 - Typical On-Resistance vs. Drain Current


Fig. 12 - Maximum Drain Current vs. Case Temperature


Fig. 13a - Unclamped Inductive Test Circuit


Fig. 13b - Unclamped Inductive Load Test Waveforms


Fig. 14 - Maximum Effective Transient Thermal Impedance, Junction-to-Case vs. Pulse Duration


Fig. 15 - Switching Time Test Circuit


Fig. 16 - Gate Charge Test Circuit

[^4]

| DIM. | MILLIMETERS |  | INCHES |  |
| :---: | :---: | :---: | :---: | :---: |
|  | MIN. | MAX. | MIN. | MAX. |
| A | 4.24 | 4.65 | 0.167 | 0.183 |
| b | 0.69 | 1.02 | 0.027 | 0.040 |
| b(1) | 1.14 | 1.78 | 0.045 | 0.070 |
| c | 0.36 | 0.61 | 0.014 | 0.024 |
| D | 14.33 | 15.85 | 0.564 | 0.624 |
| E | 9.96 | 10.52 | 0.392 | 0.414 |
| e | 2.41 | 2.67 | 0.095 | 0.105 |
| $\mathrm{e}(1)$ | 4.88 | 5.28 | 0.192 | 0.208 |
| F | 1.14 | 1.40 | 0.045 | 0.055 |
| $\mathrm{H}(1)$ | 6.10 | 6.71 | 0.240 | 0.264 |
| $\mathrm{~J}(1)$ | 2.41 | 2.92 | 0.095 | 0.115 |
| L | 13.36 | 14.40 | 0.526 | 0.567 |
| $\mathrm{~L}(1)$ | 3.33 | 4.04 | 0.131 | 0.159 |
| $\varnothing$ P | 3.53 | 3.94 | 0.139 | 0.155 |
| Q | 2.54 | 3.00 | 0.100 | 0.118 |

ECN: X15-0364-Rev. C, 14-Dec-15
DWG: 6031
Note

- $\mathrm{M}^{*}=0.052$ inches to 0.064 inches (dimension including protrusion), heatsink hole for HVM



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## SMALL SIGNAL NPN TRANSISTOR

PRELIMINARY DATA

| Ordering Code | Marking | Package / Shipment |
| :--- | :---: | :---: |
| 2N3904 | 2N3904 | TO-92 / Bulk |
| 2N3904-AP | 2N3904 | TO-92 / Ammopack |

- SILICON EPITAXIAL PLANAR NPN TRANSISTOR
- TO-92 PACKAGE SUITABLE FOR THROUGH-HOLE PCB ASSEMBLY
- THE PNP COMPLEMENTARY TYPE IS 2N3906


## APPLICATIONS

- WELL SUITABLE FOR TV AND HOME APPLIANCE EQUIPMENT
- SMALL LOAD SWITCH TRANSISTOR WITH HIGH GAIN AND LOW SATURATION VOLTAGE



## INTERNAL SCHEMATIC DIAGRAM



ABSOLUTE MAXIMUM RATINGS

| Symbol | Parameter | Value | Unit |
| :---: | :--- | :---: | :---: |
| $\mathrm{V}_{\mathrm{CBO}}$ | Collector-Base Voltage $\left(\mathrm{I}_{\mathrm{E}}=0\right)$ | 60 | V |
| $\mathrm{~V}_{\text {CEO }}$ | Collector-Emitter Voltage $\left(\mathrm{I}_{\mathrm{B}}=0\right)$ | 40 | V |
| $\mathrm{~V}_{\text {EBO }}$ | Emitter-Base Voltage $\left(\mathrm{I}_{\mathrm{C}}=0\right)$ | 6 | V |
| $\mathrm{I}_{\mathrm{C}}$ | Collector Current | 200 | mA |
| $\mathrm{P}_{\text {tot }}$ | Total Dissipation at $\mathrm{T}_{\mathrm{C}}=25^{\circ} \mathrm{C}$ | 625 | mW |
| $\mathrm{~T}_{\text {stg }}$ | Storage Temperature | -65 to 150 | ${ }^{\circ} \mathrm{C}$ |
| $\mathrm{T}_{\mathrm{j}}$ | Max. Operating Junction Temperature | 150 | ${ }^{\circ} \mathrm{C}$ |

## THERMAL DATA

| $R_{\text {thj-amb }} \bullet$ | Thermal | Resistance Junction-Ambient | Max | 200 | ${ }^{\circ} \mathrm{C} / \mathrm{W}$ |
| :--- | :--- | :--- | :---: | :---: | :---: |
| $\mathrm{R}_{\text {thj-case }} \bullet$ | Thermal | Resistance Junction-Case | Max | 83.3 | ${ }^{\circ} \mathrm{C} / \mathrm{W}$ |

ELECTRICAL CHARACTERISTICS ( $\mathrm{T}_{\text {case }}=25^{\circ} \mathrm{C}$ unless otherwise specified)

| Symbol | Parameter | Test Conditions | Min. | Typ. | Max. | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Icex | Collector Cut-off Current (VBE $=-3 \mathrm{~V}$ ) | $\mathrm{V}_{\text {CE }}=30 \mathrm{~V}$ |  |  | 50 | nA |
| $I_{\text {bex }}$ | Base Cut-off Current $\left(\mathrm{V}_{\mathrm{BE}}=-3 \mathrm{~V}\right)$ | $\mathrm{V}_{\text {CE }}=30 \mathrm{~V}$ |  |  | 50 | nA |
| $\mathrm{V}_{\text {(BR) }}$ CEO* | Collector-Emitter Breakdown Voltage ( $\mathrm{I}_{\mathrm{B}}=0$ ) | $\mathrm{IC}_{\mathrm{C}}=1 \mathrm{~mA}$ | 40 |  |  | V |
| $V_{\text {(BR) }}$ CBO | Collector-Base Breakdown Voltage ( $\mathrm{I}_{\mathrm{E}}=0$ ) | $\mathrm{I}_{\mathrm{C}}=10 \mu \mathrm{~A}$ | 60 |  |  | V |
| $V_{\text {(BR) }{ }^{\text {ebo }} \text { O }}$ | Emitter-Base Breakdown Voltage ( $\mathrm{Ic}=0$ ) | $\mathrm{I}_{\mathrm{E}}=10 \mu \mathrm{~A}$ | 6 |  |  | V |
| $\mathrm{V}_{\text {CE(sat)* }}$ | Collector-Emitter Saturation Voltage | $\begin{array}{ll} \mathrm{I}_{\mathrm{C}}=10 \mathrm{~mA} & \mathrm{I}_{\mathrm{B}}=1 \mathrm{~mA} \\ \mathrm{I}_{\mathrm{C}}=50 \mathrm{~mA} & \mathrm{I}_{\mathrm{B}}=5 \mathrm{~mA} \end{array}$ |  |  | $\begin{aligned} & 0.2 \\ & 0.2 \end{aligned}$ | $\begin{aligned} & \mathrm{V} \\ & \mathrm{~V} \end{aligned}$ |
| $V_{\text {bE(sat)* }}$ | Base-Emitter Saturation Voltage | $\begin{array}{ll} \mathrm{I}_{\mathrm{C}}=10 \mathrm{~mA} & \mathrm{I}_{\mathrm{B}}=1 \mathrm{~mA} \\ I_{0}=50 \mathrm{~mA} & \mathrm{I}_{\mathrm{O}}=5 \mathrm{~m} \end{array}$ | 0.65 |  | $\begin{aligned} & 0.85 \\ & 0.95 \end{aligned}$ | $\begin{aligned} & \mathrm{V} \\ & \mathrm{~V} \end{aligned}$ |
| $\mathrm{h}_{\text {FE* }}$ | DC Current Gain | $I_{C}=0.1 \mathrm{~mA}$ $\mathrm{~V}_{\mathrm{CE}}=1 \mathrm{~V}$ <br> $\mathrm{I}_{\mathrm{C}}=1 \mathrm{~mA}$ $\mathrm{~V}_{C E}=1 \mathrm{~V}$ <br> $\mathrm{I}_{\mathrm{C}}=10 \mathrm{~mA}$ $\mathrm{~V}_{\mathrm{CE}}=1 \mathrm{~V}$ <br> $\mathrm{I}_{\mathrm{C}}=50 \mathrm{~mA}$ $\mathrm{~V}_{\mathrm{CE}}=1 \mathrm{~V}$ <br> $\mathrm{I}_{\mathrm{C}}=100 \mathrm{~mA}$ $\mathrm{~V}_{\mathrm{CE}}=1 \mathrm{~V}$ | $\begin{gathered} 60 \\ 80 \\ 100 \\ 60 \\ 30 \end{gathered}$ | , | 300 |  |
| $\mathrm{f}_{T}$ | Transition Frequency | $\mathrm{I}_{\mathrm{C}}=10 \mathrm{~mA} \mathrm{~V} \mathrm{~V}_{\text {CE }}=20 \mathrm{~V}$ f $=100 \mathrm{MHz}$ | 250 | 270 |  | MHz |
| Ссво | Collector-Base Capacitance | $\mathrm{IE}_{\mathrm{E}}=0 \quad \mathrm{~V}_{\mathrm{CB}}=10 \mathrm{~V} \quad \mathrm{f}=1 \mathrm{MHz}$ |  | 4 |  | pF |
| Cebo | Emitter-Base Capacitance | $\mathrm{IC}=0 \quad \mathrm{~V}_{\mathrm{EB}}=0.5 \mathrm{~V} \quad \mathrm{f}=1 \mathrm{MHz}$ |  | 18 |  | pF |
| NF | Noise Figure | $\begin{aligned} & \mathrm{V}_{\mathrm{CE}}=5 \mathrm{~V} \quad \mathrm{IC}_{\mathrm{C}}=0.1 \mathrm{~mA} \mathrm{f}=10 \mathrm{~Hz} \\ & \text { to } 15.7 \mathrm{KHz} \quad \mathrm{R}_{\mathrm{G}}=1 \mathrm{~K} \Omega \end{aligned}$ |  | 5 |  | dB |
| $\begin{aligned} & \mathrm{t}_{\mathrm{d}} \\ & \mathrm{t}_{\mathrm{r}} \end{aligned}$ | Delay Time Rise Time | $\begin{array}{ll} \mathrm{I}_{\mathrm{C}}=10 \mathrm{~mA} & \mathrm{I}_{\mathrm{B}}=1 \mathrm{~mA} \\ \mathrm{~V}_{\mathrm{CC}}=30 \mathrm{~V} & \end{array}$ |  |  | $\begin{aligned} & 35 \\ & 35 \end{aligned}$ | $\begin{aligned} & \mathrm{ns} \\ & \mathrm{~ns} \end{aligned}$ |
| $\begin{aligned} & \mathrm{t}_{\mathrm{s}} \\ & \mathrm{t}_{\mathrm{f}} \end{aligned}$ | Storage Time Fall Time | $\begin{array}{ll} \mathrm{I}_{\mathrm{C}}=10 \mathrm{~mA} & \mathrm{I}_{\mathrm{B} 1}=-\mathrm{I}_{\mathrm{B} 2}=1 \mathrm{~mA} \\ \mathrm{~V}_{\mathrm{CC}}=30 \mathrm{~V} & \end{array}$ |  |  | $\begin{gathered} 200 \\ 50 \end{gathered}$ | $\begin{aligned} & \text { ns } \\ & \text { ns } \end{aligned}$ |

* Pulsed: Pulse duration $=300 \mu \mathrm{~s}$, duty cycle $\leq 2 \%$


## TO-92 MECHANICAL DATA

| DIM. | mm |  |  | inch |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | MIN. | TYP. | MAX. | MIN. | TYP. | MAX. |
| A | 4.32 |  | 4.95 | 0.170 |  | 0.195 |
| b | 0.36 |  | 0.51 | 0.014 |  | 0.020 |
| D | 4.45 |  | 4.95 | 0.175 |  | 0.194 |
| E | 3.30 |  | 3.94 | 0.130 |  | 0.155 |
| e | 2.41 |  | 2.67 | 0.095 |  | 0.105 |
| e1 | 1.14 |  | 1.40 | 0.045 |  | 0.055 |
| L | 12.70 |  | 15.49 | 0.500 |  | 0.609 |
| R | 2.16 |  | 1.52 | 0.045 |  | 0.094 |
| S1 | 1.14 |  | 0.56 | 0.016 |  | 0.059 |
| W | 0.41 |  | 6 degree | 4 degree |  | 0.022 |
| V | 4 degree |  |  |  | 6 degree |  |



TO-92 AMMOPACK SHIPMENT (Suffix"-AP") MECHANICAL DATA

| DIM. | mm |  |  | inch |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | MIN. | TYP. | MAX. | MIN. | TYP. | MAX. |
| A1 |  |  | 4.80 |  |  | 0.189 |
| T |  |  | 3.80 |  |  | 0.150 |
| T1 |  |  | 1.60 |  |  | 0.063 |
| T2 |  |  | 2.30 |  |  | 0.091 |
| d |  |  | 0.48 |  | 0.019 |  |
| P0 | 12.50 | 12.70 | 12.90 | 0.492 | 0.500 | 0.508 |
| P2 | 5.65 | 6.35 | 7.05 | 0.222 | 0.250 | 0.278 |
| F1,F2 | 2.44 | 2.54 | 2.94 | 0.096 | 0.100 | 0.116 |
| delta H | -2.00 |  | 2.00 | -0.079 |  | 0.079 |
| W | 17.50 | 18.00 | 19.00 | 0.689 | 0.709 | 0.748 |
| W0 | 5.70 | 6.00 | 6.30 | 0.224 | 0.236 | 0.248 |
| W1 | 8.50 | 9.00 | 9.25 | 0.335 | 0.354 | 0.364 |
| W2 |  |  | 0.50 |  |  | 0.020 |
| H | 18.50 |  | 20.50 | 0.728 |  | 0.807 |
| H0 | 15.50 | 16.00 | 16.50 | 0.610 | 0.630 | 0.650 |
| H1 |  |  | 25.00 |  |  | 0.984 |
| D0 | 3.80 | 4.00 | 4.20 | 0.150 | 0.157 | 0.165 |
| t |  |  | 0.90 |  |  | 0.035 |
| L |  |  | 11.00 |  |  | 0.433 |
| I1 | 3.00 |  |  | 0.118 |  |  |
| delta P | -1.00 |  | 1.00 | -0.039 |  | 0.039 |



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## Inicializacion de variables

## Iniciar pantalla LCD y joystick

Activar interrupciones
loop()






Leer sensores de motores
Leer sensores de actuadores lineales


velDeseada $=$ valorMotor * maxVel / JOYSTICK_MAX

> PID( "motor". id, velDeseada )

PinA $=$ PinA * cicloTrabajo * ANALOG_MAX_VALUE PinB $=$ PinB * cicloTrabajo * ANALOG_MAX_VALUE




## E. Programa en C

```
#include <Keypad.h>
#include <LiquidCrystal.h>
#include <SoftwareSerial.h>
```


## E. 1 Declaracion de constantes

```
const int NUM_ACT LIN = 4;
const int VEL_ARRAY_SIZE = 601;
// --------------------------- Constantes / pines
```

// Definir pines del modulo Bluetooth
const int PIN_BLUETOOTH_RX $=0$;
const int PIN_BLUETOOTH_TX = 1 ;
// Definir pines del LCD
const int PIN_LCD_RS = 33 ;
const int PIN ${ }^{-}$LCD $^{-}$ENABLE $=34$;
const int PIN_LCD_D4 = 35 ;
const int PIN_LCD_D5 = 36;
const int PIN_LCD_D6 = 37;
const int PIN_LCD_D7 = 38;
// Definir pines del teclado
const int PIN_KEYPAD COL1 = 6 ;
const int PIN_KEYPAD_COL2 = 7;
const int PIN_KEYPAD_COL3 = 8;
const int PIN_KEYPAD_ROW1 = 2 ;
const int PIN_KEYPAD_ROW2 = 3 ;
const int PIN_KEYPAD_ROW3 = 4 ;
const int PIN_KEYPAD_ROW4 = 5;
// definir pines del joystick
const int PIN_JOYSTICK_X = A6;
const int PIN JOYSTICK Y = A7;
const int PIN_JOYSTICK_BUTTON = 12;
// Definir pī̄es de los motores
const int PIN_MOTOR1_A = 9 ;
const int PIN_MOTOR1_B = 10 ;
const int PIN_MOTOR1_SENSOR = A11;
const int PIN_MOTOR2_A = 44;
const int PIN_MOTOR2_B $=45$;
const int PIN_MOTOR2_SENSOR = A10 ;
// Definir pines de los actuaadores lineales
const int PIN_ACT1_A = 13 ;
const int PIN_ACT1_B = 14 ;
const int PIN_ACT1 SENSOR = A12 ;
const int PIN_ACT2_A $=15$;
const int PIN_ACT2_B = 16 ;
const int PIN_ACT2_SENSOR = A13;
const int PIN_ACT3_A = 17 ;
const int PIN_ACT3_B = 18 ;
const int PIN_ACT3-SENSOR = A14 ;
const int PIN_ACT4_A = 19;
const int PIN_ACT4_B $=20$;
const int PIN_ACT4_SENSOR = A15 ;
// Encoders
const int PIN_ENCODER1_A = A0;
// int PIN ENCODER1 B = 19;
const int $\bar{P}$ IN_ENCODER2_A $=A 1$;

```
//const int PIN_ENCODER2_B = 21;
/ /
                                    Constantes globales
-------------------------------
// Posiciones minimas/ maximas / iniciales de los actuadores lineales
const int ACT1 MIN = 7; //L1
const int ACT1_MAX = 1019;//
const int ACT1_POSO = ACT1_MAX;
const int ACT2_MIN = 24;//
const int ACT2_MAX = 844;//
const int ACT2-POSO = ACT2 MIN;
const int ACT3_L1 = 242; //
const int ACT3_L2 = 685; //
const int ACT3_L3 = 1; //
const int ACT3_MIN = ACT3_L3; //
const int ACT3 MAX = ACT3 L2; //
const int ACT3_POSO = ACT\overline{3}L1;
const int ACT4_MIN = 11; //
const int ACT4_MAX = 870; //
const int ACT4_POSO = ACT4_MIN;
const int POS_FINAL_CHAIR_\overline{A}CT1 = ACT1_MAX; //
const int POS_FINAL_CHAIR_ACT2 = ACT2_MIN; //
const int POS_FINAL_CHAIR_ACT3 = ACT3_L1 ;//
const int POS_FINAL_CHAIR_ACT4 = ACT4_MIN; //
const int POS_FINAL_STR_ACT1 = ACT1_MIN; //
const int POS_FINAL_STR_ACT2 = ACT2_MAX; //
const int POS_FINAL_BIP_ACT1 = ACT1_MIN; //
const int POS_FINAL_BIP_ACT2 = ACT2_MAX; //
const int POS_FINAL_BIP_ACT3 = ACT3_MIN; //
const int POS_FINAL_BIP_ACT4 = ACT4_MAX; //
const int MANUAL_ADJUST = -1;
const int SENSOR_GAIN_MOTOR[2]={1, 1}; // 1024 / Vmax que mide el
sensor; MODIFICAR SEGUN SENSOR
const int SENSOR_GAIN_ACTUATOR[4] = {1, 1, 1, 1}; // 1024 / Vmax que mide
el sensor; MODIFICAR SEGUN SENSOR
const int PID_KP = 5; // LAS GANANCIAS DEL PID SE DEFINEN SEGUN EL MOTOR
const int PID_KI = 0.5;
const int PID_KD = 1; //
const int BACKREST ANGLE MOVEMENT = 45;
const int MAX_VEL = 100; //Ciclo de trabajo maximo para alcanzar la
velocidad maxima <----V
const int MAX_HEIGHT = 171; // altura maxima que se eleva la silla
respecto a la posicion inicial
const int JOYSTICK_THRESHOLD = 20; // < ----E
const int MOTOR_THRESHOLD = 10; // < ---- E
const int ACTUATOR_THRESHOLD = 3; //< ---- E
const float JOYSTI\overline{CK_AXE_MAXVALUE = 512.0;}
const float ANALOG_WRITE_MAX_VALUE = 255.0;
const int NUM_MOTO\overline{R}=2;
// Numero correspondiente para cada actuador lineal segun su funcion
const int BACKREST_ACT = 0;
const int FOOTREST_ACT = 1;
const int SEAT_TILT_ACT = 2;
const int SEAT HEIG\overline{TH_ACT = 3;}
const int IDDL\overline{E}=0;
const int FORWARD = 1;
```

```
const int BACKWARD = 2;
const int MOTOR_DIR1 = 1; // "clockwise";
const int MOTOR_DIR2 = 2; // "anticlockwise";
const int CODE_BLANK = 0;
const int CODE_MOVEMENT = 111 ;
const int CODE_CHAIR = 222;
const int CODE_STRETCHER = 333;
const int CODE_BIPEDESTATION = 444;
const int CODE_BACKREST = 555;
const int CODE_FOOTREST = 666;
const int CODE_SEAT_TILT = 777;
const int CODE_SEAT_HEIGTH_MAX = 888;
const int CODE_SEAT_HEIGTH_MIN = 999;
const int FUNCTIONALITY_MOVEMENT = 1;
const int FUNCTIONALITY_CHAIR = 2;
const int FUNCTIONALITY_STRETCHER = 3;
const int FUNCTIONALITY_BIPEDESTATION = 4;
const int FUNCTIONALITY_ADJUST_BACKREST = 5;
const int FUNCTIONALITY_ADJUST_FOOTREST = 6;
const int FUNCTIONALITY_ADJUST_SEAT_TILT = 7;
const int FUNCTIONALITY_ADJUST_HEIGTH_MAX = 8;
const int FUNCTIONALITY_ADJUST_HEIGTH_MIN = 9;
const int STATE_INITIAL = 0;
const int STATE_ELEVATED_CHAIR = 1;
const int STATE_FINAL = 2;
const int dt = \overline{0.05;}
const int WHEEL RADIO = 0;
const int PULSES_PER_ROTATION = 1024;
const int DAC_VE\overline{L_SEAT_TILT = 195;}
const int DEBÜG_O\overline{FF}=\overline{0};
const int DEBUG_INPUTS = 1;
const int DEBUG_JOYSTICK_MOTOR = 2;
const int DEBUG_ACTUATORS = 3;
const char CODE_ENTER = '*';
const char CODE_RESET = '#';
const String MSG__WELCOME = "Bienvenido" ;
const String MSG_ACCEPTED_CODE = "Codigo aceptado";
const String MSG_INCORRECT
const String MSG_ENTER_CODE = "Ingrese codigo" ;
const String MSG_VEL_MOD = "Velocidad reducida";
const String MSG_NO_\overline{FUNCTIONALITY_CHANGE_1 = "No hubo cambio";}
const String MSG_NO_FUNCTIONALITY_CHANGE_2 = "de funcionalidad";
const String MSG_MA\overline{X_HEIGHT_1 = "Se llego a la";}
const String MSG_MAX_HEIGHT_2 = "altura maxima";
const String MSG_MIN_HEIGHT_1 = "Se llego a la";
const String MSG_MIN_HEIGHT_2 = "altura minima";
const String MSG_DISA\overline{BLED_MOVVEMENT_1 = "Movimiento no";}
const String MSG_DISABLED_MOVEMENT__2 = "habilitado";
const String MSG_FUNCTIONALITY_1 =- "Movimiento ";
const String MSG_FUNCTIONALITY_2 = "Posicion silla";
const String MSG_FUNCTIONALITY_3 = "Posicion camilla";
const String MSG_FUNCTIONALITY_4 = "Posicion bipedestacion";
const String MSG_FUNCTIONALITY_5 = "Ajustar Espaldar ";
const String MSG_FUNCTIONALITY_6 = "Ajustar Apoyapie";
const String MSG_FUNCTIONALITY_7 = "Ajustar inclinacion";
```

```
const String MSG FUNCTIONALITY 8 = "Subir asiento";
const String MSG_FUNCTIONALITY_9 = "Bajar asiento";
const String MSG_FUNTIONALITY[9] = {MSG FUNCTIONALITY_1,
MSG_FUNCTIONALITY_2, MSG_FUNCTIONALITY_3, MSG_FUNCTIONALITY_4,
MSG_FUNCTIONALITY_5, MSG_FUNCTIONALITY_6, MSG_FUNCTIONALITY_7,
MSG_FUNCTIONALITY_8, MSG_FUNCTIONALITY_9 };
// Inicializar el lcd
LiquidCrystal lcd(PIN_LCD_RS, PIN_LCD_ENABLE, PIN_LCD_D4, PIN_LCD_D5,
PIN_LCD_D6, PIN_LCD_D7);
// Serial para comunicacion Bluetooth
SoftwareSerial BT1(PIN_BLUETOOTH_RX, PIN_BLUETOOTH_TX); // RX | TX
// Inicializar el teclado
const byte ROWS = 4; //four rows
const byte COLS = 3; //four columns
char keys[ROWS][COLS] = {
    {'1', '2', '3'},
    {'4', '5', '6'},
    {'7', '8', '9'},
    {'*', '0', '#'}
};
byte rowPins[ROWS] = {PIN_KEYPAD_ROW1, PIN_KEYPAD_ROW2, PIN_KEYPAD_ROW3,
PIN_KEYPAD_ROW4}; //conne\overline{c}t to the row pinouts of the keypad
byte colPins[COLS] = {PIN_KEYPAD_COL1, PIN_KEYPAD_COL2, PIN_KEYPAD_COL3};
//connect to the column pinouts of the keypad
Keypad keypad = Keypad( makeKeymap(keys), rowPins, colPins, ROWS, COLS);
```


## E. 2 Declaracion de variables

```
int debugType = DEBUG_ACTUATORS ; // == DEBUG_ACTUATORS; // =
```

DEBUG_INPUTS ; //= DEBUG_JOYSTICK_MOTOR ;
int code;
char text ;
char key;
float joystickX ;
float joystickY ;
int buttonState;
int backrestAngle;
int footrestAngle;
int seatAngle;
int seatHeight;
int currentFunctionality;
int desiredFunctionality;
//int requiredActuators[4];
int posFinal[4];
int posInitial[4];
int maxVel;
int contTimer;
int debugTimer;
bool inputsEnabled;
bool desiredBipedestation;
//int contMaxHeight;
//int contMinHeight;
int errorActuators[NUM_ACT_LIN * 2];
int errorMotors[NUM_MOTOR * 2];

```
int integralMotors [NUM MOTOR];
int integralActuators[NUMM_ACT_LIN];
int wheelchairState;
int posMin[4];
int posMax[4];
unsigned int lastReportedPos1 = 1;
unsigned int lastReportedPos2 = 1;
bool A_set[2];
bool B_set[2];
volati\overline{le unsigned long encoderPos[2];}
long newposition[2]={0,0};
long oldposition[2]={0,0};
int newtime[2]={0, 0};
int oldtime[2]={0, 0};
int motorR, motorL, dirR, dirL;
/ / actuators
typedef struct {
    int actStatus; //direction
    int pinA;
    int pinB;
    int pinSensor;
    int A;
    int B;
    int dutyCicle;
    int posMin;
    int posMax;
    int sensor;
    String function;
    bool isEnabled;
    int movement;
} actuator_type;
actuator_type actuators[NUM_ACT_LIN ];
actuator_type motors[NUM_MOTOR ];
// --------------------------- Funciones
```

----------------------

## E. 3 Funcion Inicializacion de variables

```
void VariablesInitialization() {
    motorR = 0;
    motorL = 0;
    dirR = 0;
    dirL = 0;
    joystickX = 0;
    joystickY = 0;
    buttonState = 0;
    contTimer = 0;
    debugTimer = 0;
    code = 0;
    text = ' ' ;
    key = NO_KEY;
    backrestAngle = 0;
    footrestAngle = 0;
    seatAngle = 0;
    seatHeight = 0;
    //contMaxHeight = 0;
```

```
    //contMinHeight = 0;
    wheelchairState = STATE INITIAL;
    A_set[0] = 0 ; A_set[1] = 0 ; B_set[0] = 0 ; B_set[1] = 0 ;
    encoderPos[0] = 0 ; encoderPos[1] = 0 ;
    errorActuators[0] = 0; errorActuators[1] = 0; errorActuators[2] = 0;
errorActuators[3] = 0; errorActuators[4] = 0; errorActuators[5] = 0;
errorActuators[6] = 0; errorActuators[7] = 0;
    errorMotors[0] = 0; errorMotors[1] = 0; errorMotors[2] = 0;
errorMotors[3] = 0;
    integralActuators[0] = 0; integralActuators[1] = 0;
integralActuators[2] = 0; integralActuators[3] = 0;
    integralMotors[0] = 0; integralMotors[1] = 0;
    maxVel = MAX_VEL;
    currentFunctionality = FUNCTIONALITY_CHAIR;//silla de ruedas
    desiredFunctionality = FUNCTIONALITY_CHAIR;
    //Motors
    motors[0] = (actuator_type) {
        IDDLE, PIN_MOTOR1_A, PIN_MOTOR1_B, PIN_MOTOR1_SENSOR, 0, 0, 0, 0, 0,
0, "Right Motor", false , 0
    };
    motors[1] = (actuator type) {
        IDDLE, PIN_MOTOR2_A, PIN_MOTOR2_B, PIN_MOTOR2_SENSOR, 0, 0, 0, 0, 0,
0, "Left Motor", false , 0
    };
    // Lineal actuators BACKREST_ACT 0 | FOOTREST_ACT 1 |
SEAT TILT ACT 2 | SEAT HEIGTH ACT 3
    ac\overline{tuators[BACKREST ACT \overline{] = (actuator type) {}}\mathbf{\}=\mp@code{M}
        IDDLE, PIN_ACT1_\overline{A}, PIN_ACT1_B, PIN_ACT1_SENSOR, 0, 0, 0, ACT1_MIN ,
ACT1_MAX , ACT1_POS\overline{0}, "Giro Espaldar", false , 0
    };
    actuators[FOOTREST_ACT ] = (actuator_type) {
        IDDLE, PIN_ACT2_A, PIN_ACT2_B, PIN__ACT2_SENSOR, 0, 0, 0, ACT2_MIN ,
ACT2_MAX , ACT2_POSO, "Giro-Apoyapie", false, 0
    };
    actuators[SEAT_TILT_ACT ] = (actuator_type) {
        IDDLE, PIN_ACT3_A, PIN_ACT3_B, PIN_ACT3_SENSOR, 0, 0, 0, ACT3_L3 ,
ACT3_L2 , ACT3_POS0, "Giro Asiento", false, 0
    };
    actuators[SEAT_HEIGTH_ACT ] = (actuator_type) {
        IDDLE, PIN_ACT4_A, PIN_ACT4_B, PIN_ÄCT4_SENSOR, 0, 0, 0, ACT4_MIN ,
ACT4_MAX , ACT4_POS0, "Levantamiento Asiento", false, 0
    };
    //requiredActuators = {BACKREST_ACT , FOOTREST_ACT , SEAT_TILT_ACT ,
SEAT_HEIGTH_ACT };
    posFinal[0] = POS_FINAL_CHAIR_ACT1; posFinal[1] = POS_FINAL_CHAIR_ACT2
; posFinal[2] = POS_FINAL_CHAIR_ACT3; posFinal[3] = POS_FINAL_CHAIR_ACT4;
    posInitial[0] = A\overline{CT1_MAX}; posInitial[1] = ACT2_MIN ; posInitial[2] =
ACT3_L1; posInitial[3] = ACT4_MIN;
    posMin[0] = ACT1_MIN; posMin[1] = ACT2_MIN ; posMin[2] = ACT3_L3;
posMin[3] = ACT4_MIN;
    posMax[0] = ACT1_MAX; posMax[1] = ACT2_MAX ; posMax[2] = ACT3_L2;
posMax[3] = ACT4_MAX;
    inputsEnabled = true;
    desiredBipedestation = false;
}
```


## E. 4 Funcion Principal

```
void setup() {
    VariablesInitialization();
    cli();
    ConfigTimer1();
    //pinMode(12, OUTPUT); // Al poner en HIGH forzaremos el modo
AT//
    // digitalWrite(12, LOW);
    /// delay (500) ; // Espera antes de encender el modulo
    // Serial.begin(9600);
    // Serial.println("Levantando el modulo HC-06");
    // Serial.println("Esperando comandos :");
    // BT1.begin(9600);
    // BT1.println("Coneccion Bluetooth establecida, presione a para
encerder led o b para apagarlo");
    // Inicializar LCD, indicandole columnas, filas
    lcd.begin(16, 2);
    // Prender el LCD
    lcd.display();
    lcd.print(MSG_WELCOME);
    delay(2000);
    lcd.clear(); lcd.print(MSG_ENTER_CODE );
    //joystick
    //pinMode(PIN JOYSTICK X , INPUT);
    //pinMode(PIN_JOYSTICK_Y , INPUT);
    //activate pull-up resistor on the push-button pin
    //pinMode(PIN_JOYSTICK_BUTTON , INPUT_PULLUP);
    // Encoders
    pinMode(PIN_ENCODER1_A , INPUT);
    //pinMode(PIN_ENCODER1_B , INPUT);
    pinMode(PIN_ENCODER2_A , INPUT);
    //pinMode(PIN_ENCODER2_B , INPUT);
    // Pull up resistor
    //digitalWrite(PIN_ENCODER1_A, HIGH); // turn on pullup resistor
    //digitalWrite(PIN ENCODER1 B, HIGH); // turn on pullup resistor
    //digitalWrite(PIN_ENCODER2_A, HIGH); // turn on pullup resistor
    // digitalWrite(PIN_ENCODER2_B, HIGH); // turn on pullup resistor
    // encoder pin on interrupt \overline{0} (pin 2)
    attachInterrupt(18, doEncoder1A, CHANGE);
    // attachInterrupt(19, doEncoder1B, CHANGE);
    attachInterrupt(20, doEncoder2A, CHANGE);
    // attachInterrupt(21, doEncoder2B, CHANGE);
    if (debugType != DEBUG_OFF) {
        BT1.begin(9600);
        BT1.println(" Serial listo ");
    }
    sei();
}
```


## E. 5 Funcion ConfigurarTimer1

```
void ConfigTimer1() {
    //set timerl interrupt at 20Hz
    TCCR1A = 0;// set entire TCCR1A register to 0
    TCCR1B = 0;// same for TCCR1B
    TCNT1 = 0;//initialize counter value to 0
    // set compare match register for 1hz increments
    //OCR1A = 31249;// = (16*10^6) / (2*256) - 1 (must be <65536)
    OCR1A = 1249; // = (16*10^6) / (20*256) - 1 (must be <65536)
    // turn on CTC mode with OCR1A
    TCCR1B |= (1 << WGM12);
    // Set CS12 bits for 256 prescaler
    TCCR1B |= (1 << CS12) ;
    // enable timer compare interrupt
    TIMSK1 |= (1 << OCIE1A);
}
```


## E. 6 Funcion Loop

```
void loop() {
    if (inputsEnabled) {
        // Verificar si se presiono una tecla
        key = keypad.getKey();
        if (key != NO_KEY)
        {
            PressedKey(key);
        }
        // Verificar si se recibieron datos por Bluetooth
        if (BT1.available())
        { //Serial.println("Recibiendo desde HC 05");
            text = BT1.read();
            PressedKey(text);
        }
    }
    // Leer datos del joystick
    ReadJoystick();
    // delay_ms(100) ;
    // Funcion encargada de realiar los movimientos
    //Movement();
}
```


## E. 7 Funcion TeclaPresionada

```
void PressedKey( char key) {
    switch (key) {
        case CODE_ENTER:
            if (code != CODE_BLANK) {
                //imprimir el \overline{codigo aceptado}
                    lcd.clear();
                lcd.print(MSG_ACCEPTED_CODE );
                EntryCode(code);
            }
            break;
        case CODE_RESET:
            code = CODE_BLANK ;
            lcd.clear();
            lcd.print(MSG_ENTER_CODE );
            break;
        default:
            code = code * 10 + (key - '0'); // code.concat(key);
            if (code > 1000)
            { lcd.clear();
                lcd.print(MSG_ENTER_CODE );
                lcd.setCursor=
                lcd.print(MSG_INCORRECT_CODE);
                code = CODE_BLANK ;
                return;
            }
            else {
                lcd.clear();
                lcd.print(MSG_ENTER_CODE );
                lcd.setCursor(0, 2);
                lcd.print(code);
            }
            //lcd.print(desiredFunctionality);
            break;
    }
}
```


## E. 8 Funcion IngresoCodigo

```
void EntryCode(int &code) {
    cli();
    int i, id;
    // Verificacion
    switch (code)
    {
        case CODE_MOVEMENT:
            if ( CheckDisplacementConditions()) {
                desiredFunctionality = FUNCTIONALITY_MOVEMENT; //desplazamiento
```

```
    errorActuators[0] = 0; errorActuators[1] = 0; errorActuators[2] =
0; errorActuators[3] = 0; errorActuators[4] = 0; errorActuators[5] = 0;
errorActuators[6] = 0; errorActuators[7] = 0;
    errorMotors[0] = 0; errorMotors[1] = 0; errorMotors[2] = 0;
errorMotors[3] = 0;
    integralActuators[0] = 0; integralActuators[1] = 0;
integralActuators[2] = 0; integralActuators[3] = 0;
            integralMotors[0] = 0; integralMotors[1] = 0;
        }
        else {
            currentFunctionality = desiredFunctionality;
            StopActuator("motor", 0);
            StopActuator("motor", 1);
            lcd.clear();
            lcd.print(MSG_DISABLED_MOVEMENT_1);
            lcd.setCursor(0, 2);
            lcd.print(MSG_DISABLED_MOVEMENT_2);
    }
    break;
    case CODE_CHAIR:
    desiredFunctionality = FUNCTIONALITY_CHAIR ; //silla
    if (wheelchairState == STATE_ELEVATED_CHAIR ) {
        desiredFunctionality = FUNCTIONALITY`_ADJUST_HEIGTH_MIN ;
//levantamiento
    actuators[BACKREST_ACT].isEnabled = true;
actuators[FOOTREST_ACT].isEnabled = true;
actuators[SEAT_TILT_ACT].isEnabled = true;
actuators[SEAT HEIG\overline{T}H ACT].isEnabled = true;
    //requiredActuators = {SEAT HEIGTH_ACT };
    posFinal[0] = POS_FINAL_CHAIR_ACT1; posFinal[1] =
POS_FINAL_CHAIR_ACT2 ; posFinal[2] = POS_FINAL_CHAIR_ACT3; posFinal[3] =
POS_FINAL_CHAIR_ACT4;
    posInitial[0] = ACT1_MIN; posInitial[1] = ACT2_MAX ;
posInitial[2] = ACT3_L2; posInitial[3] = ACT4_MAX;
    inputsEnabled = false;
    //contMinHeight = VEL_ARRAY_SIZE - 1;
    }
    else if (wheelchairState == STATE_FINAL) {
    actuators[BACKREST ACT].isEnabl\overline{ed = true;}
actuators[FOOTREST ACT].isEnabled = true;
actuators[SEAT TILT ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = false;
    //requiredActuators = {SEAT_TILT_ACT };
    posFinal[0] = POS_FINAL_CHA\overline{IR_ACT1; posFinal[1] =}
POS_FINAL_CHAIR_ACT2 ; posFinal[2] = A}CT3_L2; posFinal[3] = ACT4_MAX
    posInitial[0] = ACT1_MIN; posInitial[1] = ACT2_MAX ;
posInitial[2] = ACT3_L3; posInitial[3] = ACT4_MAX;
    }
    else {
            actuators[BACKREST ACT].isEnabled = true;
actuators[FOOTREST ACT].isEnabled = true;
actuators[SEAT_TILT_ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = false;
    //requiredActuators[0] = 1; requiredActuators[1] =1
;requiredActuators[2] = 1;requiredActuators[3] = 1;
```

```
    posFinal[0] = POS FINAL_CHAIR_ACT1; posFinal[1] =
POS_FINAL_CHAIR_ACT2 ; posFinal[2] = POS_FINAL_CHAIR_ACT3; posFinal[3] =
POS_FINAL_CHAIR_ACT4;
    posInitial[0] = ACT1_MIN; posInitial[1] = ACT2_MAX ;
posInitial[2] = ACT3_L3; posInitial[3] = ACT4_MIN;
        }
        break;
        case CODE_STRETCHER :
        // BT1.println(" Funcionlidad Camilla ") ;
        // BT1.println(wheelchairState) ;
        desiredFunctionality = FUNCTIONALITY_STRETCHER; //camilla
    if (wheelchairState == STATE_FINAL)
    {
        actuators[BACKREST_ACT].isEnabled = false;
actuators[FOOTREST_ACT].isEnabled = false;
actuators[SEAT_TILT_ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = false;
    //requiredActuators = {BACKREST_ACT , FOOTREST_ACT };
    posFinal[0] = 0; posFinal[1] = 0 ; posFinal[2] = ACT3_L2;
posFinal[3] = 0;
    posInitial[0] = 0; posInitial[1] = 0 ; posInitial[2] = ACT3_L3;
posInitial[3] = 0;
    posMin[2] = ACT3_L3;
    posMax[2] = ACT3_L2;
    }
    else
    {
        actuators[BACKREST ACT].isEnabled = true;
actuators[FOOTREST_ACT].isEnabled = true;
actuators[SEAT_TILT_ACT].isEnabled = false;
actuators[SEAT_HEIGTH_ACT].isEnabled = false;
    //requiredActuators = {BACKREST_ACT , FOOTREST_ACT };
    posFinal[0] = POS_FINAL_STR_ACT1; posFinal[1] =
POS_FINAL_STR_ACT2 ; posFinal[2] = \overline{0}; posFinal[3] = 0;
    posInitial[0] = ACT1_MAX; posInitial[1] = ACT2_MIN ;
posInitial[2] = 0; posInitial[ [3] = 0;
        }
        break;
        case CODE BIPEDESTATION :
        desiredFunctionality = FUNCTIONALITY_BIPEDESTATION; //bipedestacion
        desiredBipedestation = true;
        if (wheelchairState == STATE_INITIAL ) {
            desiredFunctionality = FUNCTIONALITY_ADJUST_HEIGTH_MAX;
//levantamiento
            actuators[BACKREST_ACT].isEnabled = true;
actuators[FOOTREST_ACT].isEnabled = true;
actuators[SEAT_TILT_ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = true;
    //requiredActuators = {SEAT_HEIGTH_ACT };
    posFinal[0] = POS_FINAL_BIP_ACT1; \overline{posFinal[1] =}
POS_FINAL_BIP_ACT2 ; posFinal[2] = A}CT3_L2; posFinal[3] = ACT4_MAX;'
    posInitial[0] = ACT1_MAX; posInitial[1] = ACT2_MIN ;
posInitial[2] = ACT3_L1; posInitial[3] = ACT4_MIN;
    posMin[2] = ACT3_L1;
    posMax[2] = ACT3_L2;
    //posFinal = {MANUAL_ADJUST };
    inputsEnabled = false;
```

```
    //contMaxHeight = 0;
    }
    else if (wheelchairState == STATE_ELEVATED_CHAIR) {
    actuators[BACKREST_ACT].isEnabled = true;
actuators[FOOTREST ACT].isEnabled = true;
actuators[SEAT TILT ACT].isEnabled = true;
actuators[SEAT_HEIGTTH_ACT].isEnabled = false;
    //requiredActuators = {SEAT TILT ACT };
    posFinal[0] = POS_FINAL_BIP_ACT1; posFinal[1] =
POS_FINAL_BIP_ACT2 ; posFinal[2] = POS_FINAL_BIP_ACT3; posFinal[3] =
POS_FINAL_BIP_ACT4;
    posInitial[0] = ACT1 MAX; posInitial[1] = ACT2 MIN ;
posInitial[2] = ACT3_L2; posInitial[3] = ACT4_MAX;
    posMin[2] = ACT3_L3 ;
    posMax[2] = ACT3_L2;
    }
    // actuators[BACKREST ACT].isEnabled = true;
actuators[FOOTREST ACT].isEnabled = true;
actuators[SEAT TILT ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = true;
    // posFinal[0] = POS_FINAL_BIP_ACT1; posFinal[1] =
POS FINAL_BIP_ACT2 ; posFinal[2] = POS_FINAL_BIP_ACT3; posFinal[3] =
POS FINAL BIP ACT4;
break;
    case CODE_BACKREST :
    desiredFunctionality = FUNCTIONALITY_ADJUST_BACKREST ; //giro
espaldar
    actuators[BACKREST ACT].isEnabled = true;
actuators[FOOTREST_ACT].isEnabled = false;
actuators[SEAT_TILT_ACT].isEnabled = false;
actuators[SEAT_HEIGTH_ACT].isEnabled = false;
    //requiredActuators = {BACKREST_ACT };
    posFinal[0] = MANUAL_ADJUST; posFinal[1] = 0 ; posFinal[2] = 0;
posFinal[3] = 0;
    posInitial[0] = ACT1_MAX; posInitial[1] = 0 ; posInitial[2] = 0;
posInitial[3] = 0;
    //posFinal = {MANUAL_ADJUST};
    break;
    case CODE FOOTREST :
    desired\overline{Functionality = FUNCTIONALITY_ADJUST_FOOTREST ; //giro}
apoyapie
    actuators[BACKREST_ACT].isEnabled = false;
actuators[FOOTREST ACT].i.sEnabled = true;
actuators[SEAT_TILT_ACT].isEnabled = false;
actuators[SEAT_HEIGTH_ACT].isEnabled = false;
    //requiredActuators = {FOOTREST ACT };
    posFinal[0] = 0; posFinal[1] = MANUAL_ADJUST ; posFinal[2] = 0;
posFinal[3] = 0;
    posInitial[0] = 0; posInitial[1] = ACT2_MIN ; posInitial[2] = 0;
posInitial[3] = 0;
    //posFinal = { MANUAL_ADJUST };
    break;
    case CODE_SEAT_TILT :
    desiredFunctionality = FUNCTIONALITY_ADJUST_SEAT_TILT; //giro
asiento
    actuators[BACKREST ACT].isEnabled = false;
actuators[FOOTREST_ACT].isEnabled = false;
```

```
actuators[SEAT_TILT_ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = false;
    //requiredActuators = {SEAT_TILT_ACT };
    posFinal[0] = 0; posFinal[1] = 0 ; posFinal[2] = MANUAL_ADJUST;
posFinal[3] = 0;
    posInitial[0] = 0; posInitial[1] = 0 ; posInitial[2] = ACT3_L3;
posInitial[3] = 0; // Verf PLSINITIAL[2]
    posMin[2] = ACT3_L1 ;
    posMax[2] = ACT3_L1 + 100; // ****SE DEBE ASIGNAR UNA INCLINACION
MAXIMA
    //posFinal = { MANUAL_ADJUST };
    break;
    case CODE_SEAT_HEIGTH_MAX :
    desiredFunctionality = FUNCTIONALITY_ADJUST_HEIGTH_MAX;
//levantamiento
    actuators[BACKREST ACT].isEnabled = false;
actuators[FOOTREST_ACT].isEnabled = false;
actuators[SEAT_TILT_ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = true;
    //requiredActuators = {SEAT_HEIGTH_ACT };
    posFinal[0] = 0; posFinal[1] = 0 ; posFinal[2] = ACT3_L2;
posFinal[3] = ACT4_MAX;
    posInitial[0] = 0; posInitial[1] = 0 ; posInitial[2] = ACT3_L1;
posInitial[3] = ACT4_MIN;
    posMin[2] = ACT3_L1 ;
    posMax[2] = ACT3_L2;
    //posFinal = {MANUAL_ADJUST };
    inputsEnabled = false;
    //contMaxHeight = 0;
    break;
    case CODE_SEAT_HEIGTH_MIN :
    desiredFunctionality = FUNCTIONALITY_ADJUST_HEIGTH_MIN ;
//levantamiento
    actuators[BACKREST ACT].isEnabled = false;
actuators[FOOTREST ACT].isEnabled = false;
actuators[SEAT_TILT_ACT].isEnabled = true;
actuators[SEAT_HEIGTH_ACT].isEnabled = true;
    //requiredActuators = {SEAT_HEIGTH_ACT };
    posFinal[0] = 0; posFinal[1] = 0 ; posFinal[2] = ACT3_L1;
posFinal[3] = ACT4_MIN;
    posInitial[0] = 0; posInitial[1] = 0 ; posInitial[2] = ACT3_L2;
posInitial[3] = ACT4_MAX;
    posMin[2] = ACT3_L1 ;
    posMax[2] = ACT3_L2;
    inputsEnabled = false;
    //contMinHeight = VEL_ARRAY_SIZE - 1;
    // if (wheelchairState != STATE_INITIAL) {
    // wheelchairState =STATE_INITIAL;
    // }
    //posFinal = {MANUAL_ADJUST };
    break;
    default:
    desiredFunctionality = currentFunctionality;
    lcd.clear();
    lcd.print(MSG_INCORRECT_CODE);
```

```
        code = CODE_BLANK ;
        return;
        break;
    }
    if (debugType != DEBUG_OFF) {
        BT1.print(currentFunctionality);
        BT1.println(desiredFunctionality);
    }
    if (currentFunctionality != desiredFunctionality) {
    lcd.setCursor(0, 2);
    if (desiredBipedestation)
    { lcd.print(MSG_FUNTIONALITY[FUNCTIONALITY_BIPEDESTATION - 1]);
        desiredBipedestation = false;
    }
    else {
        lcd.print(MSG_FUNTIONALITY[desiredFunctionality - 1]);
    }
    //int sizeReq = sizeof(requiredActuators) / sizeof(int);
    for (i = 0; i < NUM_ACT_LIN; i++) {
        id = i;
        if (actuators[id].isEnabled) {
        actuators[id].sensor = analogRead(actuators[id].pinSensor) /
SENSOR_GAIN_ACTUATOR[id]; //revisar si es necesario
        if (posFinal[i] == MANUAL_ADJUST) posFinal[i] =
actuators[id].posMax ;
        if (id == 2) {
            actuators[id].posMax = posMax[i];
            actuators[id].posMin = posMin[i];
        }
        if ( posFinal[i] - actuators[id].sensor > ACTUATOR_THRESHOLD )
        {
            // if (i == 2) { //REMOVE
            // BT1.println("we 1 :");
            // BT1.println(i);
            // }
            actuators[id].actStatus = FORWARD;
        }
        else if ( actuators[id].sensor - posFinal[i] >
ACTUATOR_THRESHOLD ) {
            // if (i == 2) { //REMOVE
            // BT1.println("we 2");
            // BT1.println(i);
            // }
            actuators[id].actStatus = BACKWARD;
        }
        else
        {
            // if (i == 2) { // REMOVE
            // BT1.println("we 3");
            // BT1.println(i);
            // }
            actuators[id].actStatus = IDDLE;
        }
        //actuators[id].isEnabled = false; // REMOVER?
        //BT1.print("Sensor :"); BT1.print(actuators[id].sensor);
```

```
            //BT1.print(" Threshold :"); BT1.print(ACTUATOR_THRESHOLD );
            //BT1.print("PosFinal :"); BT1.print(posFinal[i]); BT1.print("
,Actuador"); BT1.println(id + 1);
            }
            }
    }
    else {
            lcd.clear();
            lcd.print(MSG_NO_FUNCTIONALITY_CHANGE_1);
            lcd.setCursor(0,- 2);
            lcd.print(MSG_NO_FUNCTIONALITY_CHANGE_2);
    }
    code = CODE_BLANK ;
    sei();
}
```


## E. 9 Funcion VerificarCondicionesDesplazamiento

```
// Funcion que verifica que se cumplan ciertos requisitos antes de
arrancar los motores
bool CheckDisplacementConditions() {
    // Si se estan realizandos cambios en la inclinacion del espaldar/
apoyapie / asiento o levamtamient del asiento, se desactiva el movimiento
    if (backrestAngle < BACKREST_ANGLE_MOVEMENT )
        maxVel = MAX_VEL;
    else {
        return false;
    }
    if (seatHeight >= MAX_HEIGHT / 4)
    { // Se reduce la velocidad si el asiento esta elevado por seguridad
        maxVel = MAX_VEL / 2;
        lcd.print(MS\overline{G_VEL_MOD);}
    }
    return true;
}
```


## E. 10 Funcion LeerJoystick

```
void ReadJoystick () {
    joystickX = analogRead(PIN_JOYSTICK_X ) - 512;
    joystickY = analogRead(PIN__JOYSTICK_Y ) - 512;
    buttonState = digitalRead(\overline{PIN_JOYSTİCK_BUTTON );}
}
```


## E. 11 Funcion Movimiento

```
void Movement() {
    switch (desiredFunctionality) {
            case FUNCTIONALITY_MOVEMENT:
            if ((abs(joystic\overline{k}X) > JOYSTICK_THRESHOLD) || (abs(joystickY) >
JOYSTICK_THRESHOLD)) {
            //BT1.print("JoystickX:"); BT1.print(joystickX);
            // BT1.print(" JoystickY:"); BT1.print(joystickY);
            motorR = joystickY + joystickX;
            if (motorR > JOYSTICK_AXE_MAXVALUE) motorR =
JOYSTICK AXE_MAXVALUE;
            if (motorR < -JOYSTICK_AXE_MAXVALUE) motorR = -
JOYSTICK_AXE_MAXVALUE;
            motorL = joystickY - joystickX;
            if (motorL > JOYSTICK_AXE_MAXVALUE) motorL =
JOYSTICK_AXE_MAXVALUE;
            if (m
JOYSTICK_AXE_MAXVALUE;
            if (motorR >= MOTOR_THRESHOLD) {
                dirR = MOTOR_DIR1 ;
            }
            else if (motorR <= -MOTOR_THRESHOLD) {
                dirR = MOTOR_DIR2 ;
            }
            else {
                dirR = IDDLE;
            }
            if (motorL >= MOTOR_THRESHOLD) {
                dirL = MOTOR_DIR1;
            }
            else if (motorL <= -MOTOR_THRESHOLD) {
            dirL = MOTOR_DIR2;
            }
            else {
                dirL = IDDLE;
            }
            StartMotor(0, abs(motorR), dirR);
            StartMotor(1, abs(motorL), dirL);
        }
        else {
            if (debugType == DEBUG_INPUTS) {
                //BT1.println("No detēction Joystick");
                //debugTimer= 0;
            }
            StopActuator("motor", 0);
            StopActuator("motor", 1);
        }
        break;
    default :
        //int sizeReq = sizeof(requiredActuators) / sizeof(int);
        int id;
```

```
    int dirMovement;
    int i;
    int actuatorsDisabled = 0;
    int aux;
    for (i = 0; i < NUM_ACT_LIN; i++) {
    id = i;
    //BT1.print("act hab ") ;BT1.print(id) ; BT1.print(" :") ;
BT1.println(actuators[id].isEnabled) ;
    if (actuators[id].isEnabled) {
            // if ((desiredFunctionality ==
FUNCTIONALITY_ADJUST_HEIGTH_MAX) && (id == SEAT_TILT_ACT)) {
                // - if ((actuators[id].posMāx - \overline{actuators[id].sensor}
<= ACTUATOR THRESHOLD ) ) {// || (dirMovement == IDDLE)) {
        // //actuators[id].isEnabled = false;
        // BT1.print("WTF TIO Pos:") ;
BT1.print(actuators[id].sensor) ;
    // BT1.print(" PosMAX:") ;
BT1.print(actuators[id].posMax)
                // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
        // StopActuator("linearActuator", id);
        // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
            // BT1.print(" Actuador:") ; BT1.println(id) ;
            // }
            // else {
            // actuators[id].movement = FORWARD;
            // analogWrite(actuators[SEAT_TILT_ACT].pinA,
actuators[SEAT_TILT_ACT].A * DAC_VEL_SEAT_TILT );
            // analogWrite(actuators[SEAT_TILT_ACT].pinB,
actuators[SEAT_TILT_ACT].B * DAC_VEL_SEAT_TILT );
            // actuators[id].isEnabled = true; // verificar si
queda
                // }
                // }
                // else if ((desiredFunctionality ==
FUNCTIONALITY_ADJUST_HEIGTH_MIN) && (id == SEAT_TILT_ACT)) {
            // if ((actuators[id].sensor - actuators[id].posMin
<= ACTUATOR THRESHOLD ) ) { //|| (dirMovement == IDDLE)) {
            // // actuators[id].isEnabled = false;
            // StopActuator("linearActuator", id);
            // }
            // else {
                // actuators[id].movement = BACKWARD;
                // analogWrite(actuators[SEAT_TILT_ACT].pinA,
actuators[SEAT_TILT_ACT].A * DAC_VEL_SEAT_TILT );
            // analogWrite(actuators[SEAT_TILT_ACT].pinB,
actuators[SEAT_TILT_ACT].B * DAC_VEL_SEAT_TILT );
            // actuators[id].isEnabled = true;// verificar si
queda
    // }
    // }
    //else {
    // BT1.print("Joystick Y: ") ; BT1.println(joystickY) ;
    //BT1.println(" Movimiento ") ;
```

```
            //actuators[id].sensor = analogRead(actuators[id].pinSensor) /
SENSOR_GAIN_ACTUATOR[id]; //revisar si es necesario
            dirMovement = actuators[id].actStatus;
            if (joystickY > JOYSTICK_THRESHOLD) {
                aux = (dirMovement == FORWARD ? posFinal[id] -
actuators[id].sensor : actuators[id].sensor - posFinal[id]) ;
                            // BT1.print("posFinal: ");
BT1.print(posFinal[id]);
                        // BT1.print(" pos: ");
BT1.print(actuators[id].sensor);
            // BT1.print(" aux: "); BT1.print(aux);
            // BT1.print(" actuador: "); BT1.println(id);
            if (( aux <= ACTUATOR_THRESHOLD ) ) { //|| (dirMovement ==
IDDLE)) {
                            //if ((actuators[id].posMax - actuators[id].sensor <=
ACTUATOR_THRESHOLD ) ) { //|| (dirMovement == IDDLE)) {
                            // BT1.print("Huhu Pos:") ;
BT1.print(actuators[id].sensor) ;
                        // BT1.print(" PosMIN:") ;
BT1.print(actuators[id].posMin) ;
            // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
                                    StopActuator("linearActuator", id);
                                    // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
                            // BT1.print(" Actuador:") ; BT1.println(id)
;
        if (id == SEAT_HEIGTH_ACT ) {
        // BT1.print("Pos:") ;
BT1.print(actuators[id].sensor) ;
            // BT1.print(" PosMAX:") ;
BT1.print(actuators[id].posMax) ;
                            // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
            // BT1.print(" Actuador:") ;
BT1.println(id) ;
            // BT1.println("what pasa max");
            if (inputsEnabled == false) {
                            if (actuators[id].sensor - actuators[id].posMin <=
ACTUATOR_THRESHOLD ) {
                lcd.clear();
                lcd.print(MSG_MIN_HEIGHT_1);
                lcd.setCursor(0, 珑);
                lcd.print(MSG_MIN_HEIGHT_2);
            }
            else if (actuators[id].posMax - actuators[id].sensor
<= ACTUATOR_THRESHOLD ) {
                        lcd.clear();
                        lcd.print(MSG_MAX_HEIGHT_1);
                        lcd.setCursor(0, \overline{2);}
            lcd.print(MSG_MAX_HEIGHT_2);
            }
            inputsEnabled = true;
        }
        if (wheelchairState == STATE_INITIAL) {
            if (desiredBipedestation) {
                BT1.println(" Bipedestacion activada :v") ;
```

```
                        currentFunctionality = desiredFunctionality;
                    wheelchairState = STATE_ELEVATED_CHAIR;
                    code = CODE_BIPEDESTATION;
                        EntryCode( code);
                    }
                    // currentFunctionality = FUNCTIONALITY_BIPEDESTATION;
                }
            }
            }
            else {
            //actuators[id].isEnabled = true;
            //dirMovement = FORWARD;
            StartActuator(dirMovement, id );
    }
    }
    else if (joystickY < -1 * JOYSTICK_THRESHOLD) {
    if (dirMovement == FORWARD) dirMovement = BACKWARD;
    else if (dirMovement == BACKWARD) dirMovement = FORWARD;
    aux = (dirMovement == FORWARD ? actuators[id].sensor -
posInitial[i] : posInitial[i] - actuators[id].sensor ) ;
    // BT1.print("posFinal: ");
BT1.print(posFinal[id]);
    // BT1.print(" pos: ");
BT1.print(actuators[id].sensor);
    // BT1.print(" aux: "); BT1.print(aux);
    // BT1.print(" actuador: "); BT1.println(id);
    if ((aux <= ACTUATOR_THRESHOLD ) ) { //|| (dirMovement ==
IDDLE)) {
    // BT1.print("WHAT PASA CTM Pos:") ;
BT1.print(actuators[id].sensor) ;
    // BT1.print(" PosMIN:") ;
BT1.print(actuators[id].posMin) ;
    // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
    StopActuator("linearActuator", id);
    // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
    // BT1.print(" Actuador:") ; BT1.println(id) ;
    if (id == SEAT_HEIGTH_ACT ) {
        // BT1.print("Pos:") ;
BT1.print(actuators[id].sensor) ;
    // BT1.print(" PosInitial:") ;
BT1.print(posInitial[i]) ;
    // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
    // BT1.print(" Actuador:") ;
BT1.println(id) ;
        // BT1.println("what pasa min");
        if (inputsEnabled == false) {
            if (actuators[id].sensor - actuators[id].posMin <=
ACTUATOR_THRESHOLD ) {
                lcd.clear();
                        lcd.print(MSG_MIN_HEIGHT_1);
                        lcd.setCursor(0, \overline{2);}
                                lcd.print(MSG_MIN_HEIGHT_2);
    }
```

```
                            else if (actuators[id].posMax - actuators[id].sensor
<= ACTUATOR_THRESHOLD ) {
                lcd.clear();
                        lcd.print(MSG_MAX_HEIGHT_1);
                        lcd.setCursor(0, \overline{2);}
                        lcd.print(MSG_MAX_HEIGHT_2);
                }
                inputsEnabled = true;
            }
            if (wheelchairState == STATE_ELEVATED_CHAIR) {
                    BT1.println("what PASAAAAA min");
                        currentFunctionality = desiredFunctionality;
                        wheelchairState = STATE_INITIAL;
                        code = CODE_CHAIR;
                        EntryCode(code );
                        // currentFunctionality = FUNCTIONALITY_BIPEDESTATION;
            }
    }
        }
        else {
            //actuators[id].isEnabled = true;
            //BT1.print("Dir: "); BT1.print(dirMovement); BT1.print("
Actuator: "); BT1.println(id);
    //BT1.print("Dir: "); BT1.print(dirMovement); BT1.print("
Actuator: "); BT1.println(id);
    // BT1.print("HEHE Pos:") ; BT1.print(actuators[id].sensor)
;
    // BT1.print(" PosMIN:") ; BT1.print(actuators[id].posMin)
;
    // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
    StartActuator(dirMovement, id);
    // BT1.print(" Dir mov:") ;
BT1.print(actuators[id].actStatus) ;
            // BT1.print(" Actuador:") ; BT1.println(id) ;
            }
        }
        else {
```

            StopActuator("linearActuator", id);
        \}
        / / \}
    \}
        else \{
            actuatorsDisabled \(=\) actuatorsDisabled +1 ;
            if (actuatorsDisabled == NUM_ACT_LIN) \{
            //lcd.print("Se llego a la funcionalidad deseada"); TODO
    VERIFICAR DONDE VA
// FUNCTIONALITY_CHAIR FUNCTIONALITY_STRETCHER
FUNCTIONALITY_BIPEDESTATION
//FUNCTIONALITY_ADJUST_BACKREST
FUNCTIONALITY_ADJUST_FOOTREST FUNCTIONALITY_ADJUST_SEAT_TILT
FUNCTIONALITY_ADJUST_HEIGTH_MAX FUNCTIONALITY_ADJUST_HEIGTH_MIN
if ((currentFunctionality == FUNCTİONALITY _ADJUS $\bar{T}$ _HEIGTH_MAX
) $\& \&(d e s i r e d F u n c t i o n a l i t y ~==~ F U N C T I O N A L I T Y ~ B I P E D E S T A T \bar{I} O N) ~) ~(~\{~$
currentFunctionality $=$ desire $\bar{d} F u n c t i o n a l i t y ;$

```
            wheelchairState = STATE_FINAL;
            }
            else if ((currentFunctionality == FUNCTIONALITY_BIPEDESTATION
) && (desiredFunctionality == FUNCTIONALITY_CHAIR ) ) {
                    currentFunctionality = FUNCTIONALITY_ADJUST_HEIGTH_MAX;
                    wheelchairState = STATE_ELEVATED_CHAIR;
                    code = CODE_CHAIR;
            EntryCode( code);
            }
            else if ((currentFunctionality ==
FUNCTIONALITY_ADJUST_HEIGTH_MIN ) && (desiredFunctionality ==
FUNCTIONALITY_CHAIR ) ) {
                    currentFunctionality = desiredFunctionality;
                            // wheelchairState = STATE INITIAL; TODO VERIFICAR SI
ESTO VA
            }
            else
            {
                        currentFunctionality = desiredFunctionality;
                    }
                    }
            }
        }
    }
}
```


## E. 12 Funcion ArrancarMotor

```
void StartMotor(int idMotor, int motorValue, int dir) {
    float velDeseada;
    VelocityCalculation(idMotor);
    // Activar los pines A/B de los motores segun sea el sentido de giro
deseado.
    switch (dir) {
        case MOTOR_DIR1 :
            motors[idMotor].actStatus = MOTOR_DIR1;
            motors[idMotor].A = 1;
            motors[idMotor].B = 0;
            break;
        case MOTOR_DIR2 :
            motors[idMOtor].actStatus = MOTOR_DIR2;
            motors[idMotor].A = 0;
            motors[idMotor].B = 1;
            break;
    }
    velDeseada = motorValue / JOYSTICK_AXE_MAXVALUE * maxVel ; // se lleva
    if (debugType == DEBUG_JOYSTICK_MOTOR) {
        BT1.print(" A: "); BT1.print(motors[idMotor].A );
        BT1.print(" B: "); BT1.print(motors[idMotor].B );
        BT1.print(" motorValue: "); BT1.print(motorValue);
        BT1.print(" velDeseada: "); BT1.print(velDeseada);
    }
```

```
    PID("motor", idMotor, velDeseada);
    //Habilita motor
    motors[idMotor].isEnabled = true;
    // Salidas
    analogWrite(motors[idMotor].pinA, floor(motors[idMotor].A *
motors[idMotor].dutyCicle / 100.0 * ANALOG_WRITE_MAX_VALUE) ); //
analogWrite values from 0 to 255
    analogWrite(motors[idMotor].pinB, floor(motors[idMotor].B *
motors[idMotor].dutyCicle / 100.0 * ANALOG_WRITE_MAX_VALUE) );
    if (debugType == DEBUG_JOYSTICK_MOTOR) {
        BT1.print(" Duty:"); BT1.print(motors[idMotor].dutyCicle);
        BT1.print(" motor:"); BT1.print(idMotor);
        BT1.print(" A: "); BT1.print(motors[idMotor].A *
motors[idMotor].dutyCicle / 100.0 * ANALOG_WRITE_MAX_VALUE);
        BT1.print(" B: "); BT1.println(motors[i\overline{dMotor].B `}
motors[idMotor].dutyCicle / 100.0 * ANALOG_WRITE_MAX_VALUE);
    }
}
```


## E. 13 CalculoVelocidad

```
void VelocityCalculation(int id) {
    newposition[id] = encoderPos[id];
    newtime[id] = millis();
    motors[id].sensor = (newposition[id] - oldposition[id]) * 2 * PI *
WHEEL_RADIO / (PULSES_PER_ROTATION * (newtime[id] - oldtime[id]) /
1000);
    oldposition[id] = newposition[id];
    oldtime[id] = newtime[id];
}
```


## E. 14 Funcion PID

```
void PID(String type, int id, int velDeseada) {
    int dutyCicle = 0;
    int proportional, integral, derivative;
    actuator_type actuator;
    if (type-== "motor") {
    errorMotors[id * 2 + 1] = velDeseada - motors[id].sensor ; //error
actual
    proportional = PID_KP * errorMotors[id * 2 + 1];
    integralMotors[id] += PID_KI * (errorMotors[id * 2 + 1] * dt);
    derivative = PID_KD * (errrorMotors[id * 2 + 1] - errorMotors[id * 2]
) / dt; //
    dutyCicle = proportional + integralMotors[id] + derivative;
    if (dutyCicle > 100)
        dutyCicle = 100;
    else if (dutyCicle < -100)
        dutyCicle = -100;
```

```
            // Ciclo de trabajo de los motores
            motors[id].dutyCicle = dutyCicle ;
            if (debugType == DEBUG_JOYSTICK_MOTOR) {
                BT1.print(" errorPasado: "); BT1.print(errorMotors[id * 2]);
                BT1.print(" errorActual: "); BT1.print(errorMotors[id * 2 + 1]);
                BT1.print(" velDes: "); BT1.print(velDeseada);
                BT1.print(" velReal: "); BT1.print(motors[id].sensor);
                BT1.print(" proportional: "); BT1.print(proportional);
                BT1.print(" integral: "); BT1.print(integralMotors[id]);
                BT1.print(" derivative: "); BT1.print(derivative);
        }
        errorMotors[id * 2] = errorMotors[id * 2 + 1];
    }
    else if (type == "actuator") {
    errorActuators[id * 2 + 1] = velDeseada - motors[id].sensor ; //error
actual
    proportional = PID_KP * errorActuators[id * 2 + 1];
    integralActuators[id] += PID_KI * (errorActuators[id * 2 + 1] * dt);
    derivative = PID_KD * (errorActuators[id * 2 + 1] -
errorActuators[id * 2] ) / dt; //
    dutyCicle = proportional + integralActuators[id] + derivative;
    if (dutyCicle > 100)
        dutyCicle = 100;
    else if (dutyCicle < -100)
        dutyCicle = -100;
        // Ciclo de trabajo de los motores
        motors[id].dutyCicle = dutyCicle ;
    }
}
```


## E. 15 Funcion ArrancarActuador

```
void StartActuator(int dirMovement, int idActuator) {
```

void StartActuator(int dirMovement, int idActuator) {
if (actuators[idActuator].isEnabled) {
if (actuators[idActuator].isEnabled) {
if (dirMovement == FORWARD ) {
if (dirMovement == FORWARD ) {
actuators[idActuator].A = 0;
actuators[idActuator].A = 0;
actuators[idActuator].B = 1;
actuators[idActuator].B = 1;
}
}
if (dirMovement == BACKWARD ) {
if (dirMovement == BACKWARD ) {
actuators[idActuator].A = 1;
actuators[idActuator].A = 1;
actuators[idActuator].B = 0;
actuators[idActuator].B = 0;
}
}
// Se prohibe el uso del actuador lineal encargado de levantamiento,
// Se prohibe el uso del actuador lineal encargado de levantamiento,
ya que su uso se controla en la interrucpion.
ya que su uso se controla en la interrucpion.
if (idActuator != SEAT_HEIGTH_ACT)
if (idActuator != SEAT_HEIGTH_ACT)
{
{
digitalWrite(actuators[idActuator].pinA, actuators[idActuator].A
digitalWrite(actuators[idActuator].pinA, actuators[idActuator].A
);
);
digitalWrite(actuators[idActuator].pinB, actuators[idActuator].B );
digitalWrite(actuators[idActuator].pinB, actuators[idActuator].B );
}
}
}
}
}

```

\section*{E. 16 Funcion DetenerActuador}
```

void StopActuator(String type, int id) {
if (type == "motor") {
motors[id].isEnabled = false;
motors[id].A = 0;
motors[id].B = 0;
analogWrite(motors[id].pinA, motors[id].A );
analogWrite(motors[id].pinB, motors[id].B );
}
else if (type == "linearActuator") {
actuators[id].isEnabled = false;
actuators[id].A = 0;
actuators[id].B = 0;
actuators[id].actStatus = IDDLE;
digitalWrite(actuators[id].pinA, actuators[id].A );
digitalWrite(actuators[id].pinB, actuators[id].B );
}
}

```

\section*{E. 17 Funcion DetenerActuadores}
```

void StopActuators() {
StopActuator("linearActuator", BACKREST_ACT);
StopActuator("linearActuator", FOOTREST_ACT);
StopActuator("linearActuator", SEAT_TILT_ACT);
StopActuator("linearActuator", SEAT_HEIGT\overline{H_ACT);}
}

```

\section*{E. 18 Funcion InterrupcionTimer1}
```

//timer1 interrupt 2Hz
ISR(TIMER1_COMPA_vect) {
int i;
actuator_type actuator;
// Funciōn encargada de mover los actuadores/motores
Movement();
// Cada 4 veces que entre al la interrupcion o 4 * 0.05 s = 0.2 s (10
hz), Leer los sensores
if (contTimer == 4) {
contTimer = 0;
for (i = 0; i < NUM_ACT_LIN; i++) {
//actuators[i].sensor = analogRead(actuators[i].pinSensor) /
SENSOR_GAIN_ACTUATOR[i]; //revisar si es necesario
//BT1.print(i);BT1.print(": ");
//BT1.println(actuators[i].sensor );
actuator = actuators[i];
if (actuator.isEnabled) {
actuator.sensor = analogRead(actuator.pinSensor) /
SENSOR_GAIN_ACTUATOR[i];

```
```

    if ( (actuator.sensor < actuator.posMin) || (actuator.sensor
    actuator.posMax)) {
StopActuator("linearActuator", i);
}
switch (i)
{
case BACKREST ACT :
backrestAngle = (actuator.posMax - actuator.sensor) * 90 /
(actuator.posMax - actuator.posMin);
break;
case FOOTREST ACT :
footrestAngle = (actuator.sensor - actuator.posMin ) * 90 /
(actuator.posMax - actuator.posMin);
break;
case SEAT_TILT_ACT :
seatAng\overline{le}=\overline{(actuator.posMax - actuator.sensor) * 90 /}
(actuator.posMax - actuator.posMin);
break;
case SEAT_HEIGTH_ACT :
seatHeight = actuator.sensor - actuator.posMin;
break;
}
actuators[i] = actuator;
}
}
}
contTimer = contTimer + 1;
if (debugTimer < 75) {
debugTimer = debugTimer + 1 ;
}
else {
switch (debugType) {
case (DEBUG_ACTUATORS):
for (int id = 0; id < 4 ; id ++) {
if (actuators[id].isEnabled) {
if (actuators[id].movement == FORWARD ) {
BT1.print("Pos:") ; BT1.print(actuators[id].sensor) ;
BT1.print(" PosMAX:") ; BT1.print(actuators[id].posMax) ;
BT1.print(" Dir mov:") ; BT1.print(actuators[id].actStatus)
;
BT1.print(" Mov:") ; BT1.print(actuators[id].movement) ;
BT1.print(" Actuador:") ; BT1.print(id + 1) ;
BT1.println(" ADELANTE") ;
}
else if (actuators[id].movement == BACKWARD ) {
BT1.print("Pos:") ; BT1.print(actuators[id].sensor) ;
BT1.print(" PosMIN:") ; BT1.print(actuators[id].posMin) ;
BT1.print(" Dir mov:") ; BT1.print(actuators[id].actStatus)
;
BT1.print(" Mov:") ; BT1.print(actuators[id].movement) ;
BT1.print(" Actuador:") ; BT1.print(id + 1) ;
BT1.println(" ATRAS") ;

```
```

                }
                else {
                    BT1.print("Pos:") ; BT1.print(actuators[id].sensor) ;
                    BT1.print(" PosMIN:") ; BT1.print(actuators[id].posMin) ;
                    BT1.print(" PosMAX:") ; BT1.print(actuators[id].posMax) ;
                    BT1.print(" Dir mov:") ; BT1.print(actuators[id].actStatus)
    ;
BT1.print(" Mov:") ; BT1.print(actuators[id].movement) ;
BT1.print(" Actuador:") ; BT1.print(id + 1) ;
BT1.println(" PARADO") ;
}
}
// printDebug = true;
}
debugTimer = 0;
break;
case DEBUG_INPUTS:
// BT1.print("JoystickX:"); BT1.print(joystickX);
// BT1.print(" JoystickY:"); BT1.print(joystickY);
// BTl.print(" Codigo Ingresado:"); BTl.println(code);
// debugTimer = 0;
BT1.print("Codigo Ingresado:"); BT1.print(code);
BT1.print(" JoystickX:"); BT1.print(joystickX);
BT1.print(" JoystickY:"); BT1.println(joystickY);
BT1.print(" motorDer:") ; BT1.print(motorR) ;
BT1.print(" motorIzq:"); BT1.print(motorL);
BT1.print(" dirDer:"); BT1.print(dirR);
BT1.print(" dirIzq:"); BT1.println(dirL);
debugTimer = 0;
break;
}
}
}

```

\section*{E. 19 Function InterrupcionEncoder1}
```

// Interrupt on A changing state
void doEncoder1A() {
// Test transition
// A_set[0] = digitalRead(PIN_ENCODER1_A) == HIGH;
if (motors[0].actStatus == MOTOR_DIR1) {
// and adjust counter + if A leads B
encoderPos[0] += 1;
}
else if (motors[0].actStatus == MOTOR_DIR2) {
encoderPos[0] += -1;
}
}
}

```
```

// Interrupt on B changing state
//void doEncoder1B() {
// Test transition
// B_set[0] = digitalRead(PIN_ENCODER1_B) == HIGH;
// and adjust counter + if B follows A
// encoderPos[0] += (A_set[0] == B_set[0]) ? +1 : -1;
/ / }
// Interrupt on A changing state

```

\section*{E. 20 Función InterrupcionEncoder2}
```

void doEncoder2A() {
// Test transition
// A_set[0] = digitalRead(PIN_ENCODER1_A) == HIGH;
if (motors[1].actStatus == MOTOR_DIR1) {
// and adjust counter + if A leads B
encoderPos[1] += 1;
}
else if (motors[1].actStatus == MOTOR_DIR2) {
encoderPos[1] += -1;
}
}
// Interrupt on B changing state
//void doEncoder2B() {
// Test transition
// B_set[1] = digitalRead(PIN_ENCODER2_B) == HIGH;
// and adjust counter + if B follows A
// encoderPos[1] += (A_set[1] == B_set[1]) ? +1 : -1;
/ / }

```

\section*{Código de la simulación del movimiento del actuador lineal 3}
```

clc
clear all
close all
% Cambio minimo en la velocidad posible
minVel = 0.0392;
% Velocidad del actuador lineal 4
v4=255* minVel; %10 m/s
% Velocidad del actuador lineal 3
v3_real= 195* minVel; %7.644 m/s
%Bfx= 0;
%Bfy= 389.6583+ deltaH;
% Posicion inicial del punto de union Actuador lineal3 / asiento
x3_0= 389.6583;
y3_0 = 389.6583;
% Longitud inicial del actuador lineal 3
13_0=(x3_0^2 + y3_0^2 )^0.5;
% Radio de giro del asiento
r= 170;
O3x =389.6583 ;
O4x =170 ;
y0 = y3_0;
13= 13_0;
10=13-0;
alpha3_0 = atand((y3_0)/(x3_0) ) ;
% Intervalo de tiempo
deltaT= 0.05;
% Distancia que sube la silla cada Intervalo de tiempo
deltaH = v4 *deltaT;
t= 0:deltaT:20;
cont = 1;
% Simulacion del movimiento
for i= t
% Movimiento 1 ( solo act lin 4 )
l_aux = ((y0+ deltaH).^2+ (03x-04x)^2).^0.5;
teta = atan2d((y0+ deltaH),(O3x-O4x));
alpha_aux = acosd(( l_aux.^2+ lo^2- r^2 )./(2*l_aux *l0));
alpha2= teta - alpha_aux;
alpha1= asind((y0+deltaH- l0*sind(alpha2))./r);
% Nueva posicion del punto B
Bfx_1= 03x-l0* cosd(alpha2);
Bfy_1= l0* sind(alpha2);

```
```

    %deltaX_1= Bfx_1 ;
    %deltaY_1= -Bfy 
    % movimiento 2 ( solo act lin 3 )
    x0 = Bfx_1;
    y0 = Bfy_1;
    10=( (0\overline{3x}-\textrm{x}0).^^2 + y0.^2).^0.5;
% Cambio en la longitud del actuador lineal 3
deltaL = v3_real*deltaT;
% Nueva longitud del act lin 3
lf = l0 + deltaL;
alpha_aux = acosd((lf^2 + l_aux^2 - r^2)/(2*lf*l_aux));
alpha2 = teta - alpha_aux;
x0_aux(cont)= x0;
y0_aux(cont)= y0;
% Posicion final
xf(cont)= O3x- lf * cosd(alpha2);
yf(cont) = lf * sind(alpha2);
% Actualizando las posiciones para la siguiente iteracion.
x0= xf(cont);
y0= yf(cont);
l0 =( (O3x- x0).^2 + y0.^2).^^0.5;
cont = cont +1;
end
% Posicion X
figure(); plot(t,x0_aux);
hold on; plot(t,xf);
% Posicion Y
figure(); plot(t,y0_aux);
hold on; plot(t,yf);
%figure();plot(t,559.6583- yf);
%figure();plot(t,xf);

```

\title{
G. Cálculo del centro de masa de la silla de ruedas
}

\section*{G. 1 Centro de gravedad de la estructura}

Símbolos:
\(x_{1}:\) Ubicación del eje \(X\) del centro de gravedad de la pieza 1
\(x_{4}:\) Ubicación del eje \(X\) del centro de gravedad de la pieza 4
\(x_{5}:\) Ubicación del eje \(X\) del centro de gravedad de la pieza 5
\(W_{1}\) : Peso de la pieza 1(Ensamble Pata trasera)
\(W_{2}\) : Peso de la pieza 2(Ensamble Transversal)
\(W_{3}:\) Peso de la pieza 3(Ensamble Transversal)
\(W_{4}:\) Peso de la pieza 4(Ensamble pata frontal)
\(W_{5}\) : Peso de la pieza 5(Ensamble soporte frontal)
\(W_{6}\) : Peso de la pieza 6(Ensamble Transversal)

En la figura 1 se muestran las diversas que forman la estructura de la silla de ruedas y en la figura 2. se muestra la distribución de los pesos de cada uno de estos componentes


Figura 1 Piezas de la estructura


Figura 2 Distribución de los pesos según cada pieza - Estructura
La ecuación 1 representa el cálculo del centro de masa de este componente.
\[
\begin{equation*}
x_{\text {estructura }}=\frac{x_{1} W_{1}+x_{2} W_{2}+x_{3} W_{3}+x_{4} W_{4}+x_{5} W_{5}+x_{6} W_{6}}{W_{1}+W_{2}+W_{3}+W_{4}+W_{5}+W_{6}} \tag{1}
\end{equation*}
\]

Con,
\[
\begin{array}{cc}
x_{1}=39.982 \mathrm{~mm} & W_{1}=0.329 \mathrm{~kg} \\
x_{2}=0 \mathrm{~mm} & W_{2}=0.18 \mathrm{~kg} \\
x_{3}=0 \mathrm{~mm} & W_{3}=0.18 \mathrm{~kg} \\
x_{4}=385.48 \mathrm{~mm} & W_{4}=0.386 \mathrm{~kg} \\
x_{5}=385.48 \mathrm{~mm} & W_{5}=0.197 \mathrm{~kg} \\
x_{6}=368.919 \mathrm{~mm} & W_{6}=0.18 \mathrm{~kg}
\end{array}
\]

Se calcula:
\[
x_{\text {estructura }}=209.569 \mathrm{~mm}
\]

\section*{G. 2 Centro de gravedad de la silla- función silla}

Símbolos:
\(W_{1}:\) Peso de la pieza 1(Cabecera)
\(W_{2}:\) Peso de la pieza 2(Espaldar)
\(W_{3}:\) Peso de la pieza 3(Asiento)
\(W_{4}:\) Peso de la pieza 4(Apoyapié)
\(W_{5}:\) Peso de la pieza 5 (Actuador lineal para el giro del apoyapié)
\(W_{6}:\) Peso de la pieza 6(Actuador lineal para el giro del espaldar)
\(x_{1}:\) Ubicación del eje X del centro de gravedad de la pieza 1
\(x_{2}:\) Ubicación del eje X del centro de gravedad de la pieza 2
\(x_{3}:\) Ubicación del eje X del centro de gravedad de la pieza 3
\(x_{4}:\) Ubicación del eje X del centro de gravedad de la pieza 4
\(x_{5}:\) Ubicación del eje X del centro de gravedad de la pieza 5
\(x_{6}:\) Ubicación del eje X del centro de gravedad de la pieza 6

En la figura 3. se muestra la distribución de los pesos de cada uno de los componentes que forman el ensamble Silla en la funcionalidad Silla de ruedas.


Figura 3 Distribución de los pesos Silla - Función silla

La ecuación 2 representa el cálculo del centro de masa de la estructura cuando la silla está en funcionalidad Silla de ruedas.
\[
\begin{equation*}
x_{\text {silla }}=\frac{x_{1} W_{1}+x_{2} W_{2}+x_{3} W_{3}+x_{4} W_{4}+x_{5} W_{5}+x_{6} W_{6}}{W_{1}+W_{2}+W_{3}+W_{4}+W_{5}+W_{6}} \tag{2}
\end{equation*}
\]

Con,
\[
\begin{array}{cc}
x_{1}=0 & W_{1}=0.963 \\
x_{2}=0 & W_{2}=1.739 \\
x_{3}=185.612 & W_{3}=2.735 \\
x_{4}=351 & W_{4}=2.08 \\
x_{5}=120 & W_{5}=2 \\
x_{6}=182.588 & W_{6}=2
\end{array}
\]

Se calcula:
\[
x_{\text {silla }}=156.478 \mathrm{~mm}
\]

El valor hallado es la distancia desde el punto A (unión espaldar/ asiento) hacia el centro de gravedad. Se procede a calcular la distancia del centro de gravedad de la persona desde el punto de referencia 0 (parte baja de la pata posterior).
\[
x_{\text {silla } / o}=x_{\text {silla }}+273.407-180=249.88 \mathrm{~mm}
\]

\section*{G. 3 Centro de gravedad de la silla- función camilla}

En la figura 4. se muestra la distribución de los pesos de cada uno de los componentes que forman el ensamble Silla en la funcionalidad Silla de ruedas.


Figura 4 Distribución de los pesos Silla - Función camilla
La ecuación 3 representa el cálculo del centro de masa de la estructura cuando la silla está en funcionalidad Camilla.
\[
\begin{equation*}
x_{\text {camilla }}=\frac{x_{1} W_{1}+x_{2} W_{2}+x_{3} W_{3}+x_{4} W_{4}+x_{5} W_{5}+x_{6} W_{6}}{W_{1}+W_{2}+W_{3}+W_{4}+W_{5}+W_{6}} \tag{3}
\end{equation*}
\]

Con,
\[
\begin{array}{ll}
x_{1}=97.503 & W_{1}=0.325 \\
x_{2}=398.696 & W_{2}=0.587 \\
x_{3}=872.254 & W_{3}=0.923 \\
x_{4}=1313.8 & W_{4}=0.702 \\
x_{5}=806.642 & W_{5}=2 \\
x_{6}=806.642 & W_{6}=2
\end{array}
\]

Se calcula:
\[
x_{\text {camilla }}=818.673 \mathrm{~mm}
\]

El valor hallado es la distancia desde el punto C (inicio de la cabecera) hacia el centro de gravedad. Se procede a calcular la distancia del centro de gravedad de la persona desde el punto de referencia O (parte baja de la pata posterior).
\[
x_{\text {camilla/o }}=x_{\text {camilla }}-686.642+273.407-180=225.438 \mathrm{~mm}
\]

\section*{G.4 Centro de gravedad de la persona- función silla}

Símbolos:
\(l_{\text {cabeza }}\) : Longitud de la cabeza de la persona
\(l_{\text {torso }}\) : Longitud del torso de la persona
\(l_{\text {muslo }}\) : Longitud del muslo de la persona
\(l_{\text {piernas }}\) : Longitud de las piernas de la persona
\(W_{\text {cabeza }}\) : Peso de la cabeza ( \(6.81 \%\) del peso total)
\(W_{\text {torso }}\) : Peso del torso y brazos (52.45\% del peso total)
\(W_{\text {muslo }}\) : Peso de los muslos ( \(28.32 \%\) del peso total)
\(W_{\text {piernas }}\) : Peso de las piernas ( \(11.4 \%\) del peso total)

En la figura 5. se muestra la distribución de los pesos de cada uno de las partes de la persona.


Figura 5 Distribución de los pesos Persona - Función silla

La ecuación 4 representa el cálculo del centro de masa de la persona en cuando la silla está en funcionalidad Silla de ruedas.
\[
\begin{equation*}
x_{\text {personaSilla }}=\frac{x_{\text {cabeza }} W_{\text {cabeza }}+x_{\text {torso }} W_{\text {torso }}+x_{\text {muslo }} W_{\text {muslo }}+x_{\text {piernas }} W_{\text {piernas }}}{W_{\text {cabeza }}+W_{\text {torso }}+W_{\text {muslo }}+W_{\text {piernas }}} \tag{4}
\end{equation*}
\]

Con,
\[
\begin{array}{cl}
x_{\text {cabeza }}=34 & W_{\text {cabeza }}=13.62 \\
x_{\text {torso }}=94 & W_{\text {torso }}=104.9 \\
x_{\text {muslo }}=273.785 & W_{\text {muslo }}=56.64 \\
x_{\text {piernas }}=453.57 & W_{\text {piernas }}=22.8
\end{array}
\]

Se calcula:
\[
x_{\text {personasilla }}=182.725 \mathrm{~mm}
\]

El valor hallado es la distancia desde el punto U (unión torso/muslos) hacia el centro de gravedad. Se procede a calcular la distancia del centro de gravedad de la persona desde el punto de referencia O (parte baja de la pata posterior).
\[
x_{\text {personasilla/o }}=x_{\text {personasilla }}+273.407-180=276.132 \mathrm{~mm}
\]

\section*{G. 5 Centro de gravedad de la persona- función camilla}

En la figura 6. se muestra la distribución de los pesos de cada uno de las partes de la persona.


Figura 6 Distribución de los pesos Persona - Función camilla
La ecuación 5 representa el cálculo del centro de masa de la persona en cuando la silla está en funcionalidad Camilla.
\[
\begin{equation*}
x_{\text {personaCamilla }}=\frac{x_{\text {cabeza }} W_{\text {cabeza }}+x_{\text {torso }} W_{\text {torso }}+x_{m u s l o} W_{m u s l o}+x_{p i e r n a s} W_{p i e r n a s}}{W_{\text {cabeza }}+W_{\text {torso }}+W_{\text {muslo }}+W_{\text {piernas }}} \tag{5}
\end{equation*}
\]

Con,
\[
\begin{array}{cl}
x_{\text {cabeza }}=93.5 & W_{\text {cabeza }}=13.62 \\
x_{\text {torso }}=427.139 & W_{\text {torso }}=104.9 \\
x_{\text {muslo }}=889.128 & W_{\text {muslo }}=56.64 \\
x_{\text {piernas }}=1110.978 & W_{\text {piernas }}=22.8
\end{array}
\]

Se calcula:
\[
x_{\text {personaCamilla }}=182.725 \mathrm{~mm}
\]

El valor hallado es la distancia desde el punto U (unión torso/muslos) hacia el centro de gravedad. Se procede a calcular la distancia del centro de gravedad de la persona desde el punto de referencia O (parte baja de la pata posterior).
\[
x_{\text {personaCamilla/o }}=x_{\text {personaCamilla }}+273.407-180=276.132 \mathrm{~mm}
\]

\section*{G. 6 Centro de gravedad total-silla}

Símbolos:
OA: Distancia del punto O al punto A
\(A C\) Distancia del punto A al punto C
\(x_{\text {persona }}\) : Ubicación en el eje X del centro de gravedad de la persona.
\(x_{\text {silla }}\) : Ubicación en el eje X del centro de gravedad de la silla.
\(x_{\text {estructura }}\) : Ubicación en el eje \(X\) del centro de gravedad de la estructura.
\(W_{\text {persona }}\) : Peso de la persona
\(W_{\text {silla }}\) : Peso de la silla
\(W_{\text {estructura }}\) : Peso de la estructura
En la figura 7. se muestra la distribución de los pesos de cada uno de las partes que forman el ensamble total (estructura, silla y persona) cuando la silla se encuentra en funcionalidad Silla.


Figura 7 Distribución de los pesos total - Función silla

La ecuación 6 representa el cálculo del centro de masa de la persona en cuando la silla está en funcionalidad Silla de ruedas.
\[
\begin{equation*}
x_{\text {total }}=\frac{x_{\text {persona }} W_{\text {persona }}+x_{\text {silla }} W_{\text {silla }}+x_{\text {estructura }} W_{\text {estructura }}+x_{\text {motores }} W_{\text {motores }}+x_{\text {baterias }} W_{\text {baterias }}}{W_{\text {persona }}+W_{\text {silla }}+W_{\text {estructura }}+W_{\text {motores }}+W_{\text {baterias }}} \tag{6}
\end{equation*}
\]

Con,
\[
\begin{array}{cc}
x_{\text {estructura }}=209.569 \mathrm{~mm} & W_{\text {estructura }}=5 \mathrm{~kg} \\
x_{\text {silla }}=249.885 \mathrm{~mm} & W_{\text {silla }}=12 \mathrm{~kg} \\
x_{\text {persona }}=276.1 \mathrm{~mm} & W_{\text {persona }}=40 \mathrm{~kg} \\
x_{\text {bateria } 1}=88 \mathrm{~mm} & W_{\text {bateria } 1}=29 \mathrm{~kg} \\
x_{\text {motores }}=-110 \mathrm{~mm} & W_{\text {motores }}=14.4 \mathrm{~kg}
\end{array}
\]

Se calcula:
\[
x_{\text {totalSilla }}=160 \mathrm{~mm}
\]

\section*{G. 7 Centro de gravedad total-camilla}

Símbolos:

\author{
\(l_{\text {espaldar }}\) : Longitud del espaldar
}

En la figura 8. se muestra la distribución de los pesos de cada uno de las partes que forman el ensamble total (estructura, silla y persona) cuando la silla se encuentra en funcionalidad Camilla.


Figura 8 Distribución de los pesos total - Función camilla
La ecuación 7 representa el cálculo del centro de masa del ensamble total cuando la silla está en funcionalidad Camilla.
\[
x_{\text {total }}=\frac{x_{\text {persona }} W_{\text {persona }}+x_{\text {camilla }} W_{\text {camilla }}+x_{\text {estructura }} W_{\text {estructura }}+x_{\text {motores }} W_{\text {motores }}+x_{\text {baterias }} W_{\text {baterias }}}{W_{\text {persona }}+W_{\text {silla }}+W_{\text {estructura }}+W_{\text {motores }}+W_{\text {baterias }}}
\]

Con,
\[
\begin{array}{cc}
x_{\text {estructura }}=209.569 \mathrm{~mm} & W_{\text {estructura }}=5 \mathrm{~kg} \\
x_{\text {camilla }}=225.438 \mathrm{~mm} & W_{\text {silla }}=12 \mathrm{~kg} \\
x_{\text {bateria } 1}=88 \mathrm{~mm} & W_{\text {bateria } 1}=29 \mathrm{~kg} \\
x_{\text {motores }}=-110 \mathrm{~mm} & W_{\text {motores }}=14.4 \mathrm{~kg} \\
x_{\text {personaCamilla }}=21.9 \mathrm{~mm} & W_{\text {persona }}=20 \mathrm{~kg}
\end{array}
\]

Se calcula:
\(x_{\text {totalCamilla }}=55.7 \mathrm{~mm}\)

\section*{G. 8 Centro de gravedad total-bipedestacion}

Símbolos:
\begin{tabular}{ll}
\(x_{\text {motor }}:\) & Posición de los motores \\
\(x_{\text {estructura }}:\) & Posición de la estructura \\
\(x_{\text {baterias }}:\) & Posición de las baterías \\
\(x_{\text {silla }}:\) & Posición de la silla \\
\(x_{\text {persona }}:\) & Posición de la persona \\
\(W_{\text {motor }}:\) & Peso de los motores \\
\(W_{\text {estructura }}:\) & Peso de la estructura \\
\(W_{\text {baterias }}:\) & Peso de las baterías \\
\(W_{\text {silla }}:\) & Peso de la silla \\
\(W_{\text {persona }}:\) & Peso de la persona
\end{tabular}

En la figura 9 se muestra la distribución de los pesos de cada uno de las partes que forman el ensamble total (estructura, componentes, silla y persona) cuando la silla se encuentra en funcionalidad Bipedestación. A continuación se calculara la ubicación del centro de gravedad en esta posición.


Figura 9 Distribución de los pesos total - Función camilla
La ecuación 7 representa el cálculo del centro de masa del ensamble total cuando la silla está en funcionalidad Bipedestación.
\[
\begin{equation*}
x_{\text {total }}=\frac{x_{\text {persona }} W_{\text {persona }}+x_{\text {camilla }} W_{\text {camilla }}+x_{\text {estructura }} W_{\text {estructura }+x_{\text {motores }} W_{\text {motores }}+x_{\text {baterias }} W_{\text {baterias }}}}{W_{\text {persona }}+W_{\text {silla }}+W_{\text {estructura }}+W_{\text {baterias }}+W_{\text {motores }}} \tag{I.8}
\end{equation*}
\]

Con,
\[
\begin{gathered}
x_{\text {estructura }}=209 \mathrm{~mm} \\
x_{\text {silla }}=375 \mathrm{~mm} \\
x_{\text {persona }}=475 \mathrm{~mm}
\end{gathered}
\]
\[
W_{\text {estructura }}=5 \mathrm{~kg}
\]
\[
\begin{gathered}
W_{\text {silla }}=12 \mathrm{~kg} \\
W_{\text {persona }}=40 \mathrm{~kg}
\end{gathered}
\]
\[
\begin{array}{ll}
x_{\text {bateria1 }}=88 \mathrm{~mm} & W_{\text {bateria1 }}=29 \mathrm{~kg} \\
x_{\text {motor }}=-110 \mathrm{~mm} & W_{\text {motor }}=14.4 \mathrm{~kg}
\end{array}
\]

Se calcula:
\[
x_{\text {totalBipedestacion }}=254 \mathrm{~mm}
\]```


[^0]:    Note: 1. Only for ATmega1281/2561. For ATmega640/1280/2560 these functions are placed on MISO/MOSI pins.

[^1]:    Notes: 1. Z0: should be zero for all SPM commands, byte select for the (E)LPM instruction.
    2. See "Addressing the Flash During Self-Programming" on page 314 for details about the use of Z-pointer during Self-Programming.
    3. The Z-register is only 16 bits wide. Bit 16 is located in the RAMPZ register in the I/O map.

[^2]:    Notes: 1. Z0: should be zero for all SPM commands, byte select for the (E)LPM instruction.
    2. See "Addressing the Flash During Self-Programming" on page 314 for details about the use of $Z$-pointer during Self-Programming.
    3. The Z-register is only 16 bits wide. Bit 16 is located in the RAMPZ register in the I/O map.

[^3]:    Vishay Siliconix maintains worldwide manufacturing capability. Products may be manufactured at one of several qualified locations. Reliability data for Silicon Technology and Package Reliability represent a composite of all qualified locations. For related documents such as package/tape drawings, part marking, and reliability data, see www.vishay.com/ppg?91300.

[^4]:    Vishay Siliconix maintains worldwide manufacturing capability. Products may be manufactured at one of several qualified locations. Reliability data for Silicon Technology and Package Reliability represent a composite of all qualified locations. For related documents such as package/tape drawings, part marking, and reliability data, see www.vishay.com/ppg?91459.

